

Combinatorics

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Lecture 32

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Definition 0.1 (Cut Norm of Matrices). Let A , be an $m \times n$ (possibly complex) matrix and define the **cut norm** of A to be

$$\|A\|_{\square} = \sup \left\{ \left| \sum_{i \in S, j \in T} a_{ij} \right| : S \subseteq [m], T \subseteq [n] \right\}.$$

Remark. If $A \geq 0$ is a nonnegative real matrix, we find

$$\|A\|_{\square} = |A|_1 = \sum_{i \in [m], j \in [n]} a_{ij}.$$

Similarly, for a nonpositive real matrix we find the cut norm to again be the modulus of the sum of entries.

Moreover, the cut norm is in fact a norm, as it is always nonnegative, it is only zero in the case of a zero matrix, it behaves linearly with real multiplication, and with a bit of derivation we find it obeys the triangle inequality.

Example. $\left\| \begin{bmatrix} -1 & 1 \\ 1 & -1 \end{bmatrix} \right\|_{\square} = 1$ as any rectangle yields a sum 0 and the square consisting of just a_{11} yields a sum 1.

$\left\| \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \right\|_{\square} = 2$ taking either a_{11}, a_{12} , or a_{11}, a_{21} , or simply summing over the whole matrix.

$$\|J_n - 2I_n\|_{\square} = \left\| \begin{bmatrix} -1 & 1 & \dots & 1 \\ 1 & -1 & \dots & 1 \\ \vdots & \vdots & \ddots & \vdots \\ 1 & 1 & \dots & -1 \end{bmatrix} \right\|_{\square} = n(n-2)$$

by taking the whole matrix. It is simple to show that if $|S|$ or $|T| \leq n-2$, then the sum over their entries must be strictly less than $n(n-2)$. Then, this leaves only four possibilities, the possible permutations of sets of size $n-1$ and n . If $|S| = n-1$ and $|T| = n$ (WLOG) we see the sum of entries is at most

$(n-1)(n-2) < n(n-2)$. Lastly, if $|T| = |S| = n-1$, then we have exactly one row and one column missing, so the sum of their entries will be

$$n(n-2) - r_i - c_j + 1 \leq n(n-2).$$

Hence, we have the claim is shown. \diamond

Proposition 0.1. The $2n \times 2n$ matrix $J_n \otimes (J_2 - 2I_2) = A$ has

$$\|A\|_{\square} = \|J_n \otimes (J_2 - 2I_2)\|_{\square} = \left\| \begin{bmatrix} -J_n & J_n \\ J_n & -J_n \end{bmatrix} \right\|_{\square} = n^2.$$

Proof. Note that for any individual row of length ℓ , we find the row sum $r_i \leq \begin{cases} \ell, & \ell \leq n \\ \ell - 2(\ell - n), & \ell > n \end{cases}$ and similarly for the column sums. Denoting $|S| = a$, $|T| = b$ for sets S, T which maximize the element sums, we first, note one of $a, b > n$ else the sum would be less than n^2 . Hence, we find $\|A\|_{\square} \leq \inf\{a(n-2(b-n)), b(n-2(b-n))\}$. Moreover, we find both a and $b > n$, hence we can assume WLOG $a \geq b > n$ and the solution follows by minimizing the two quadratic upper bounds. \square

Remark. We wish to examine the cut norm of a hadamard matrix. We will show a hadamard matrix H has $\|H\|_{\square} \leq n^{\frac{3}{2}} = n\sqrt{n}$.

The key to this proof is to let x, y be the indicator vectors for the sets S, T on which the maximum is obtained respectively. Then we find $\|H\|_{\square} = |\langle Hy, x \rangle| \leq \sigma_1(H) \|x\|_2 \|y\|_2$ (this is true for any value). Applying the fact that $\sigma_1(H) = \sqrt{n}$ and $\|x\|_2, \|y\|_2 \leq \sqrt{n}$ as H is hadamard and x, y are indicator vectors of length n and from this we obtain the earlier upper bound.

We can generalize the first steps of this argument to any matrix A in the following way:

Proposition 0.2. For an arbitrary $m \times n$ matrix A , we find

$$\|A\|_{\square} \leq \sigma_1(A) \sqrt{mn}.$$

Lecture 33

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Let A be a $m \times n$ matrix with $\vec{x} \in \mathbb{R}^n$, $\vec{y} \in \mathbb{R}^m$ and $|\vec{x}|_{\infty} \leq 1$ and $|\vec{y}|_{\infty} \leq 1$. Then, we consider $\max |\langle A\vec{x}, \vec{y} \rangle| = \|A\|_{\pi}$.

Proposition 0.3. We claim

$$\|A\|_{\square} \leq \|A\|_{\pi}.$$

Proof. If S, T are submatrices inducing $\|A\|_{\square}$. That is

$$\left| \sum_{i \in T, j \in S} a_{i,j} \right| = \|A\|_{\square}.$$

Letting \vec{x}, \vec{y} be indicator vectors for S, T respectively, we see this sum is simply

$$\left| \sum_{i \in T, j \in S} a_{ij} \right| = |\langle Ax, y \rangle| \leq \max |\langle Ax, y \rangle|.$$

It is also possible to set an upper bound, $\|A\|_\pi \leq 4\|A\|_\square$. □