

# Combinatorics

Thomas Fleming

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## Contents

### Lecture 16

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First, we examine some more random graphs. For a random graph  $G$ , it is a trivial result of probability theory that the number of four cycles is precisely  $\frac{1}{2} \sum_{u,v \in V(G); v \neq u} \binom{\hat{d}(u,v)}{2}$ . Then, applying our estimation  $\hat{d}(u,v) = \frac{n}{4} + o(n)$  yields  $\binom{n}{2}$  possible pairs  $u, v$  and  $\hat{d} \approx \frac{n}{4}$ , hence the number of four cycles is

$$\frac{1}{n} \binom{\frac{n}{4}}{2} \binom{n}{2} = \frac{n^4}{128} + o(n^4).$$

Now, we examine the  $k$ -walks.

**Definition 0.1** (Walks). A  $k$ -walk is a  $k$ -path  $v_1, v_2, v_3, \dots, v_k$ .  
A closed  $k$ -walk is a  $k$ -cycle,  $v_1, v_2, \dots, v_k, v_1$ .

**Remark.** Walks need not have all vertices distinct, hence a graph of order 2 where one simply oscillates between the vertices to produce a degenerate  $2n$ -walk. Similarly, one can traverse a triangle to induce a 4-walk as well. Overall this yields 14 possible 4-walks on a graph of order 4.

Now, we examine the number of closed 4-walks on a random graph of order  $n$ . We see nondegenerate 4-walks are just 4-cycles of which we know there to be  $\frac{n^4}{128}$  with 8 possible permutations of directions and starting point yields  $8 \cdot \frac{n^4}{128}$ . Similarly, we note that  $4 \cdot \sum_{v \in V} \binom{d_i(v)}{2} = 4n \binom{n}{2} = \frac{1}{2}n^3 + o(n^3) = o(n^4)$  degenerate graphs on 3 vertices exist. Lastly, the number of degenerate graphs on 2-vertices is clearly,  $2 \cdot e(g) = o(n^4)$ . Hence, the number of 4-walks is just  $\frac{n^4}{16} + o(n^4)$ .

**Proposition 0.1.**  $\text{tr} \left( A(G)^k \right) = \sum_{i=1}^n \lambda_i^k$  is the number of closed  $k$ -walks in a graph  $G$  of order  $n$ .

From this, we arrive at  $6k_3(G) = \text{tr}(A^3) = \sum_{i=1}^3 \lambda_i^3$ .  
 We also see the number of closed walks of order 4 is

$$\begin{aligned} CW_4 &= \sum_{i=1}^n \lambda_i^4 \\ \frac{n^4}{16} + o(n^4) &= \lambda_1^4 + \sum_{i=2}^n \lambda_i^4 \\ \Rightarrow \sum_{i=2}^n \lambda_i^4 &= o(n^4). \end{aligned}$$

Similarly, we find  $\sigma_2(G) = o(n)$  and  $O(\sqrt{n})$ .

**Definition 0.2** (Local Density). The **local density** of a graph is simply  $e(U)$  for some graph  $U \subseteq V$ .

**Remark.** Local density is highly variable. For instance in  $K_{n,n}$  we find  $U$  being one of the partite sets yields 0 local density and  $U$  being a set of half the vertices in each partite set yields  $\frac{1}{4}e(G)$  local density.

**Proposition 0.2.** Suppose  $G$  is a random graph of order  $n$  and let  $U$  be a set with  $|U| > 502 \log(n)$ . Then,  $\left| e(U) - \frac{1}{2} \binom{|U|}{2} \right| < \binom{|U|}{2} \left( \frac{3.5 \log n}{|U|} \right)^{\frac{1}{2}}$ .

**Proposition 0.3.** There exists a function  $f : \mathbb{N} \rightarrow \mathbb{N}$  such that almost every graph of order  $n$  has clique number  $f(n)$  or  $f(n+1)$ .

This function is approximated by

$$f(n) \approx 2 \log_2(n).$$

**Remark.** There is clearly also such a function for the independence number.

Furthermore, more investigation yields  $\chi(G) \approx \frac{n}{2 \log_2(n)}$  for almost all graphs  $G$ .

## Lecture 17: Semi-circle Law

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Recall that for eigenvalues  $\lambda_1, \lambda_2, \dots, \lambda_n$  we have  $\lambda_1 = \frac{n}{2} + \sqrt{n \log(n)} = o(n)$ .  
 Additionally, we know  $\sigma_1 = \lambda_1$  and  $\sigma_2, \sigma_3, \dots, \sigma_n$  correspond to  $|\lambda_2|, |\lambda_3|, \dots, |\lambda_n|$ .  
 Further, it is known by Furedi and Kowlos that  $\sigma_2 = O(\sqrt{n})$ .

**Theorem 0.1.** For a randomly chosen graph of order  $n$ , with eigenvalues  $\lambda_2 \geq \lambda_3 \geq \dots \geq \lambda_n$ . Define  $W_n(x) : \mathbb{R} \rightarrow \mathbb{Z}^+$  to be the number of eigenvalues  $\lambda_i$ , such that  $\frac{\lambda_i}{\sqrt{n} \leq x}$ , divided by  $n$ . Then, we find the function which  $W_n(x)$  tends to pointwise,  $W(x)$  has  $W(x) = \begin{cases} \frac{2}{\pi} \sqrt{1-x^2}, & |x| \leq 1 \\ 0, & |x| > 1 \end{cases}$

Here recall that  $\sqrt{1-x^2}$  is an upper half semicircle of radius 1 and the factor  $\frac{2}{\pi}$  compresses it into an ellipse. This fact essentially characterizes the distribution of eigenvalues of a random graph. That is, plurality of eigenvalues will be 0 and we find the number of eigenvalues of a given magnitude decreases as  $\lambda \rightarrow \sqrt{n}$ . We note that the leading  $\frac{2}{\pi}$  is to normalize the area such that this is a probability density function. Then, we note  $E[x^2 W(x)] = \int_{-1}^1 \frac{2}{\pi} x^2 \sqrt{1-x^2} dx = \frac{1}{4}$ . Hence, we find  $\frac{1}{n^2} \sum_{i=2}^n \lambda_i^2 \approx \frac{1}{4}$ .

It is a well known result that  $\sum_{i=1}^n |\lambda_i| = \sum_{i=1}^{\infty} \sigma_i \leq \frac{1}{2} n^{\frac{3}{2}} \leq 2(n-1)$ . Applying our integral formula from earlier yields  $\sum_{i=1}^{\infty} |\lambda_i| = \int_{-1}^1 |x| \sqrt{1-x^2} = 2 \int_0^1 x \sqrt{1-x^2}$ .