

Combinatorics

Thomas Fleming

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Lecture 29

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I skipped taking notes on a few lectures. See the paper in the onedrive for info on the szemerdi lemma.

Lemma 0.1 (Triangle Removal Lemma). Denote $K_3(G)$ to be the number of triangles in G .

Then, for all $0 < \varepsilon < 1$ there is $\delta > 0$ so that if G is a graph of order n with $K_3(G) < \delta n^3$ there are edges e_1, e_2, \dots, e_p with $p < \varepsilon n^2$ and $K_3(G - e_1 - e_2 - \dots - e_p) = 0$.

For an example, we may take the graph $K_{n,n}$ and add an edge within one of the partite sets. Then, we see the only triangles are the ones incident to this new edge. Hence, we need only remove one edge to obtain a triangle-less graph. More generally, adding cross edges between two subsets of size εn within one of the partite sets will yield $\frac{\varepsilon^2}{2} n^3$ triangles which may be removed by simply deleting all of the edges which we added.

Proof. Fix ε and apply the regularity lemma with $\frac{\varepsilon}{4}$ and $m > \frac{2}{\varepsilon}$. Then, there is $n_0(\frac{\varepsilon}{4}, m)$ so that for all G of order $n > n_0$ we find $V(G) = \bigcup_{i=1}^m V_i$ with $|V_i| < \frac{\varepsilon}{4} n$ and $|V_1| = \dots = |V_m|$. Moreover, $m < k < M$ and all but $\frac{\varepsilon}{4} k^2$ of the pairs (V_i, V_j) , $1 \leq i < j \leq k$ are $\frac{\varepsilon}{4}$ -regular.

Then, we delete all edges incident to V_0 . We see this is at most $\sum_{x \in V_0} d(x) < \frac{\varepsilon}{4} n^2$ edges being removed. Now, fix $1 \leq i \leq k$ are delete all edges within V_i . Again, we see this is at most $\binom{|V_i|}{2}$ edges, hence doing this for all sets V_i , $1 \leq i \leq k$ we see we remove at most $k \binom{|V_i|}{2} < k \frac{n^2}{2k^2} = \frac{n^2}{2k} < \frac{n^2}{2m} \leq \frac{\varepsilon}{4} n^2$.

Next, we delete all edges within the non $\frac{\varepsilon}{4}$ -regular pairs. Again, we see this is at most $|V_i||V_j| < \frac{n^2}{k^2}$ pairs. Summing over all non $\frac{\varepsilon}{4}$ -regular pairs yields $\frac{\varepsilon}{4} n^2$ more edges to remove.

Finally, we delete all edges in the remaining pairs with $d(V_i, V_j) < \frac{\varepsilon}{2}$. We see $e(V_i, V_j) \leq \frac{\varepsilon}{2} \frac{n^2}{k^2}$. Summing over all $\binom{k}{2}$ pairs (V_i, V_j) , we see once again we remove less than $\frac{\varepsilon}{4} n^2$ edges.

Hence, in total we have removed strictly less than εn^2 edges. Now, we verify

that we have removed all triangles. Suppose there remains a triangle. Of course, this triangle will not be incident to V_0 . Moreover, this triangle will have edges contained in 3 distinct V_i s. This is because we removed all edges within each V_i . Suppose WLOG, that the edges are in V_1, V_2, V_3 . Then, as we have removed edges between irregular and low density pairs, $d(V_1, V_2), d(V_2, V_3), d(V_1, V_3) \geq \frac{\varepsilon}{2}$.

Now, recall for 3 mutually $\frac{\varepsilon}{4}$ -regular pairs V_1, V_2, V_3 of equal cardinality we had $k_3(V_1, V_2, V_3) \geq (1 - 2\frac{\varepsilon}{4})(d_{1,2} - \frac{\varepsilon}{4})(d_{1,3} - \frac{\varepsilon}{4})(d_{2,3} - \frac{\varepsilon}{4})|V_1|^3$. Hence, we see as $|V_1| > (1 - \varepsilon)\frac{n}{k} > (1 - \varepsilon)\frac{n}{M}$, and letting $\delta = (1 - \frac{\varepsilon}{2})\left(\frac{\varepsilon}{4}\right)^3 \frac{(1-\varepsilon)^3}{M^3}$, we see

$$\begin{aligned} K_3(V_1, V_2, V_3) &\geq \left(1 - \frac{\varepsilon}{2}\right) \left(\frac{\varepsilon}{4}\right)^3 |V_1|^3 \\ &> \delta n^3. \end{aligned}$$

Hence, there will only be triangles remaining if the initial assumptions are violated so the claim is shown. \square

Remark. We have shown this for $n > n_0$. If instead we wish to show for $n \leq n_0$ we can choose $\delta = \frac{1}{\binom{n_0}{3}}$ which will exclude the possibilities of any triangles before removing any edges, hence the lemma is trivially true in this case.

Lecture 30: The (6,3) Problem

Wed 10 Nov 2021 16:24

Definition 0.1 (Hypergraphs). A r -uniform graph or r -graph, G is a pair of two sets. V , being a vertex set and E , being a set of edges are subsets of cardinality r . That is, a graph where a single edge connects r (not necessarily distinct) vertices instead of 2 vertices as in a graph (with or without loops).

Proposition 0.1. Suppose G is a 3-uniform graph of order n such that every set of 6 vertices induces at most 2 edges. Then $e(G) = o(n^2)$.

Proof. First note that for a fixed vertex v , we find $d(v) \leq 4$ else the second assumption is violated. Hence, induce a graph from G by removing all vertices sharing a common pair. Denote this graph G' of order n' . Then, we see $e(G) - e(G') \leq 4(n - n')$ by the restriction on the degrees.

Next, induce a graph from G' by considering only the edges participating in a triple (that is pairs of two triangles sharing a single vertex) and "demote" it to a simple graph by converging every triple into a triangle. Call this G'_2 and note that we converted a triple into 3 edges, so by the preceding lemma, we see $e(G'_2) = o(n'^2)$ and $e(G') = \frac{1}{3}e(G'_2) = o(n'^2)$. Then, working backwards, we see $e(G) \leq 4(n - n') + o(n'^2) = o(n^2)$. \square

Definition 0.2 (Super Regular Pair). We define a graph $G(A, B)$ to be an (ε, δ) -**super regular pair** for some $\varepsilon, \delta \in (0, 1)$ if the following conditions hold

- For all $X \subseteq A$ and $Y \subseteq B$ with $|X| < \varepsilon |A|$ and $|Y| < \varepsilon |B|$ we have $e(X, Y) > \delta |X| |Y|$.
- For every vertex $a \in A$, $d(a) > \delta |B|$ and similarly, for all $b \in B$, $d(b) > \delta |A|$.

Proposition 0.2. If (A, B) is an ε regular pair with $d(A, B) > 3\varepsilon$, then there are A', B' so that $|A'| > (1 - \varepsilon) |A|$ and $|B'| > (1 - \varepsilon) |B|$ such that (A', B') is an $(2\varepsilon, d - 2\varepsilon)$ -super regular pair.

Proof. Define $A_0 = \{a \in A : d(a) \leq (d - \varepsilon) |B|\}$. Then, we see $|A_0| < \varepsilon |A|$ by assumption. Define $A' = A \setminus A_0$. Then, $d_B(a) > (d - \varepsilon) |B|$, so once we have removed at most ε elements from B as well, we see $d_{B'}(a) > (d - 2\varepsilon) |B|$. Similarly, we find $d_{A'}(b) > (d - 2\varepsilon) |A|$ after removing an exceptional set B_0 . Then, by the slicing lemma, we find an $\alpha > \varepsilon$ so that $|A'| > \alpha |A|$ and $|B'| > \alpha |B|$. Then, we see (A', B') is an ε' regular pair with $\varepsilon' = \sup\{\frac{\varepsilon}{\alpha}, 2\varepsilon\}$. If $\alpha = 1 - \varepsilon$, we find $\varepsilon' = 2\varepsilon$. \square