## Combinatorics

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## Contents

## Lecture 16: Random Graphs

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First, we examine some more random graphs. For a random graph G, it is a trivial result of probability theory that the number of four cycles is precisely  $\frac{1}{2}\sum_{u,v\in V(G);v\neq u}{\hat{d}(u,v)\choose 2}$ . Then, applying our estimation  $\hat{d}(u,v)=\frac{n}{4}+o(u)$  yields  $\binom{n}{2}$  possible pairs u,v and  $\hat{d}\approx\frac{n}{4}$ , hence the number of four cycles is

$$\frac{1}{n} \binom{\frac{n}{4}}{2} \binom{n}{2} = \frac{n^4}{128} + o\left(n^4\right).$$

Now, we examine the k-walks.

**Definition 0.1** (Walks). A k-walk is a k-path  $v_1, v_2, v_3, \ldots, v_k$ . A closed k-walk is a k-cycle,  $v_1, v_2, \ldots, v_k, v_1$ .

**Remark.** Walks need not have all vertices distinct, hence a graph of order 2 where one simply oscillates between the vertices to produce a degenerate 2*n*-walk. Similairly, one can traverse a triangle to induce a 4-walk as well. Overall this yields 14 possible 4-walks on a graph of order 4.

Now, we examine the number of closed 4-walks on a random graph of order n. We see nondegenerate 4-walks are just 4-cycles of which we know there to be  $\frac{n^4}{128}$  with 8 possible permutations of directions and starting point yields  $8 \cdot \frac{n^4}{128}$ . Similairly, we note that  $4 \cdot \sum_{v \in V} \binom{d_i(v)}{2} = 4n\binom{n}{2} = \frac{1}{2}n^3 + o\left(n^3\right) = o\left(n^4\right)$  degenerate graphs on 3 vertices exist. Lastly, the number of degenerate graphs on 2-vertices is clearly,  $2 \cdot e\left(g\right) = o\left(n^4\right)$ . Hence, the number of 4 - walks is just  $\frac{n^4}{16} + o\left(n^4\right)$ .

**Proposition 0.1.**  $\operatorname{tr}\left(A\left(G\right)^{k}\right) = \sum_{i=1}^{n} \lambda_{i}^{k}$  is the number of closed k-walks in a graph G of order n.

From this, we arrive at  $6k_3(G) = \operatorname{tr}(A^3) = \sum_{i=1}^3 \lambda_i^3$ . We also see the number of closed walks of order 4 is

$$CW_4 = \sum_{i=1}^n \lambda_i^4$$
$$\frac{n^4}{16} + o(n^4) = \lambda_1^4 + \sum_{i=2}^n \lambda_i^4$$
$$\Rightarrow \sum_{i=2}^n \lambda_i^4 = o(n^4).$$

Similarly, we find  $\sigma_2(G) = o(n)$  and  $O(\sqrt{n})$ .

**Definition 0.2** (Local Density). The **local density** of a graph is simply e(U) for some graph  $U \subseteq V$ .

**Remark.** Local density is highly variable. For instance in  $K_{n,n}$  we find U being one of the partite sets yields 0 local density and U being a set of half the vertices in each partite set yields  $\frac{1}{4}e(G)$  local density.

**Proposition 0.2.** Suppose G is a random graph of order n and let U be a set with  $|U| > 502 \log{(n)}$ . Then,  $\left| e\left( U \right) - \frac{1}{2} \binom{|U|}{2} \right| < \binom{|U|}{2} \left( \frac{3.5 \log{n}}{|U|} \right)^{\frac{1}{2}}$ .

**Proposition 0.3.** There exists a function  $f : \mathbb{N} \to \mathbb{N}$  such that almost every graph of order n as clique number f(n) or f(n+1).

This function is approximated by

$$f(n) \approx 2 \log_2(n)$$
.

Remark. There is clearly also such a function for the independence number.

Furthermore, more investigation yields  $\chi(G) \approx \frac{n}{2\log_2(n)}$  for almost all graphs G.

## Lecture 17: Semi-circle Law

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Recall that for eigenvalues  $\lambda_1, \lambda_2, \ldots, \lambda_n$  we have  $\lambda_1 = \frac{n}{2} + \sqrt{n \log(n)} = o(n)$ . Additionally, we know  $\sigma_1 = \lambda_1$  and  $\sigma_2, \sigma_3, \ldots, \sigma_n$  correspond to  $|\lambda_2|, |\lambda_3|, \ldots, |\lambda_n|$ . Further, it is known by Furedi and Kowlos that  $\sigma_2 = O(\sqrt{n})$ .

**Theorem 0.1.** For a randomly chosen graph of order n, with eigenvalues  $\lambda_2 \geq \lambda_3 \geq \ldots \geq \lambda_n$ . Define  $W_n(x) : \mathbb{R} \to \mathbb{Z}^+$  to be the number of eigenvalues  $\lambda_i$ , such that  $\frac{\lambda_i}{\sqrt{n} \leq x}$ , divided by n. Then, we find the function which

$$W_{n}\left(x\right)$$
 tends to pointwise,  $W\left(x\right)$  has  $W\left(x\right)=\left\{ egin{array}{ll} \frac{2}{\pi}\sqrt{1-x^{2}}, & |x|\leq1\\ 0, & |x|>1 \end{array} \right.$ 

Here recall that  $\sqrt{1-x^2}$  is an upper half semicircle of radius 1 and the factor  $\frac{2}{\pi}$  compresses it into an ellipse. This fact essentially characterizes the distribution of eigenvalues of a random graph. That is, plurality of eigenvalues will be 0 and we find the number of eigenvalues of a given magnitude decreases as  $\lambda \to \sqrt{n}$ . We note that the leading  $\frac{2}{\pi}$  is to normalize the area such that this is a probability density function. Then, we note  $E\left[x^2W\left(x\right)\right] = \int_{-1}^{1} \frac{2}{\pi}x^2\sqrt{1-x^2}dx = \frac{1}{4}$ . Hence, we find  $\frac{1}{n^2}\sum_{i=2}^{n}\lambda_i^2 \approx \frac{1}{4}$ .

It is a well known result that  $\sum_{i=1}^{n} |\lambda_i| = \sum_{i=1}^{\infty} \sigma_i \leq \frac{1}{2} n^{\frac{3}{2}} \leq 2(n-1)$ . Applying our integral formula from earlier yields  $\sum_{i=1}^{\infty} |\lambda_i| = \int_{-1}^{1} |x| \sqrt{1-x^2} = 2 \int_{0}^{1} x \sqrt{1-x^2}$ .