

PRINT BITS IN AN INTEGER

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EXTRACT THE BITS ONE BY ONE
FROM THE INTEGER

GET WHETHER THEY ARE 0 OR 1 -
PRINT THAT OUT TO SCREEN

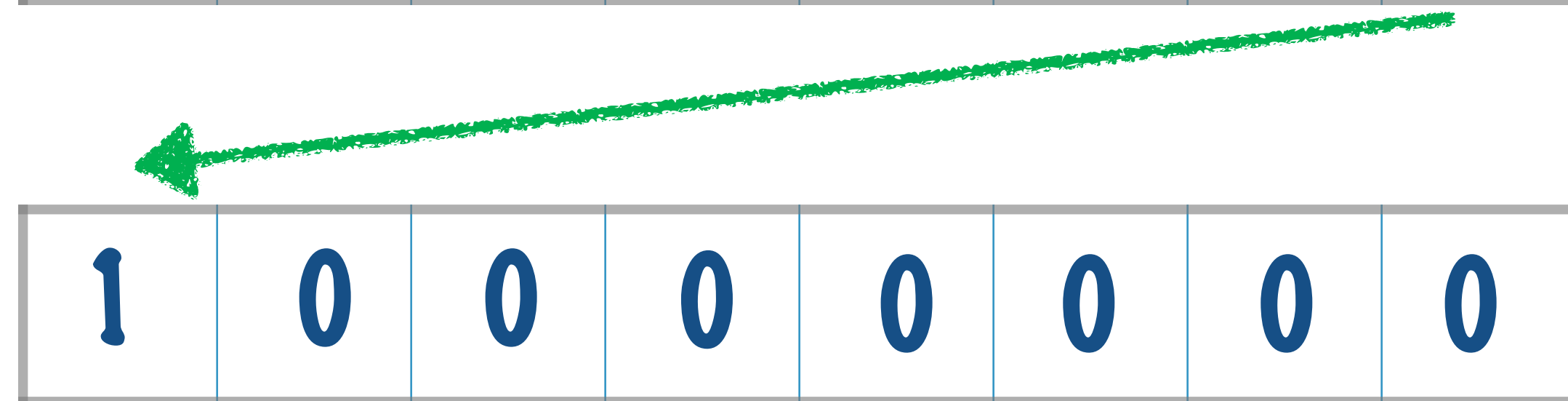
WE'VE ALREADY SEEN HOW TO GET
THE VALUE OF AN INTEGER AT THE
NTH POSITION

THE TRICK HERE IS TO START AT THE
LEFTMOST POSITION SO WE PRINT
THE BITS IN A LOGICAL ORDER

PRINT BITS IN AN INTEGER

HOW DO WE SET THE
LEFT MOST BIT OF AN
INTEGER TO 1?

START WITH 1



SHIFT IT LEFT $\text{sizeof(int)} * 8 - 1$ TIMES

sizeof WILL RETURN THE SIZE OF AN
INTEGER IN BYTES

MULTIPLY BY 8 TO GET THE SIZE IN
BITS

SUBTRACT 1 TO ACCOUNT
FOR THE RIGHT MOST
POSITION BEING AT INDEX 0

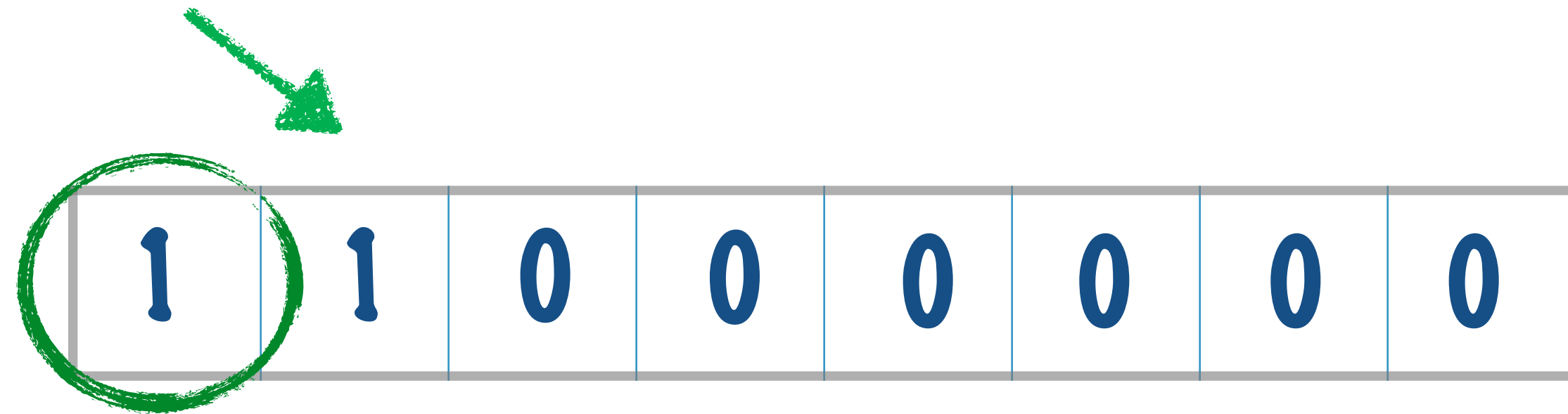
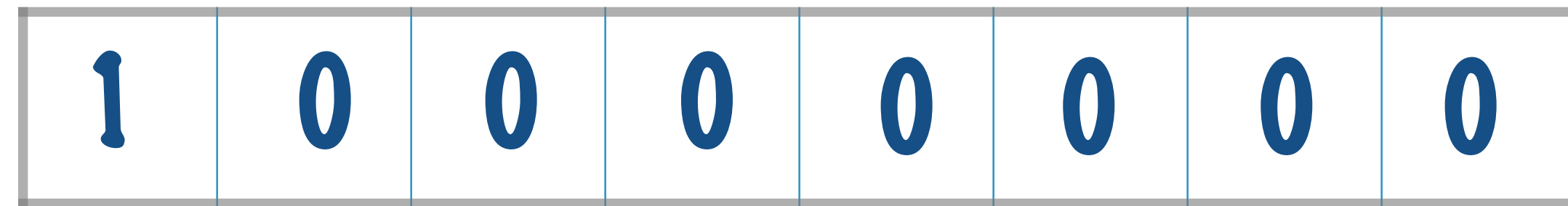
PRINT BITS IN AN INTEGER

AFTER WE GET THE 1-BIT IN THE **LEFT MOST** POSITION

IT SHOULD BE **SHIFTED RIGHT**

ONE SUBTLE DETAIL IN SHIFTING RIGHT

THIS IS BECAUSE RIGHT SHIFT **PRESERVES THE SIGN** OF THE INTEGER - A 1 IN THE RIGHT MOST POSITION IS A NEGATIVE NUMBER



IF THE LEFT MOST BIT IS 1, THE **FILL BIT** TENDS TO BE 1 WHEN SHIFTED RIGHT

USE AN **UNSIGNED INTEGER** AS A CHECK BIT TO AVOID THIS ISSUE!

PRINT BITS IN AN INTEGER

```
void print_bits(int num) {  
    unsigned int check_bit = 1 << (sizeof(int) * 8 - 1);  
  
    while (check_bit != 0) {  
        int result = num & check_bit;  
        if (result == check_bit) {  
            printf("%d ", 1);  
        } else {  
            printf("%d ", 0);  
        }  
  
        check_bit = check_bit >> 1;  
    }  
  
    printf("\n");  
}
```

THE RIGHT MOST
POSITION IN THE CHECK
BIT SHOULD BE 1 - NOTE
THE UNSIGNED INT!

ONCE THE RIGHT MOST 1 IS
MOVED OUT TO THE LEFT AND IS
SHIFTED OUT THEN CHECK_BIT
WILL BE 0

MOVE THE 1 TO THE RIGHT

COUNT THE NUMBER OF 1 BITS

COUNT THE NUMBER OF 1 BITS

THE SIMPLE WAY IS VERY SIMILAR
TO THE PRINT BITS EXAMPLE

CHECK EACH BIT, INCREMENT A
COUNT IF THE BIT IS 1

THIS METHOD INVOLVES CHECKING
EVERY BIT IN AN INTEGER

IT'S COMPLEXITY IS $O(\text{NUMBER OF BITS})$

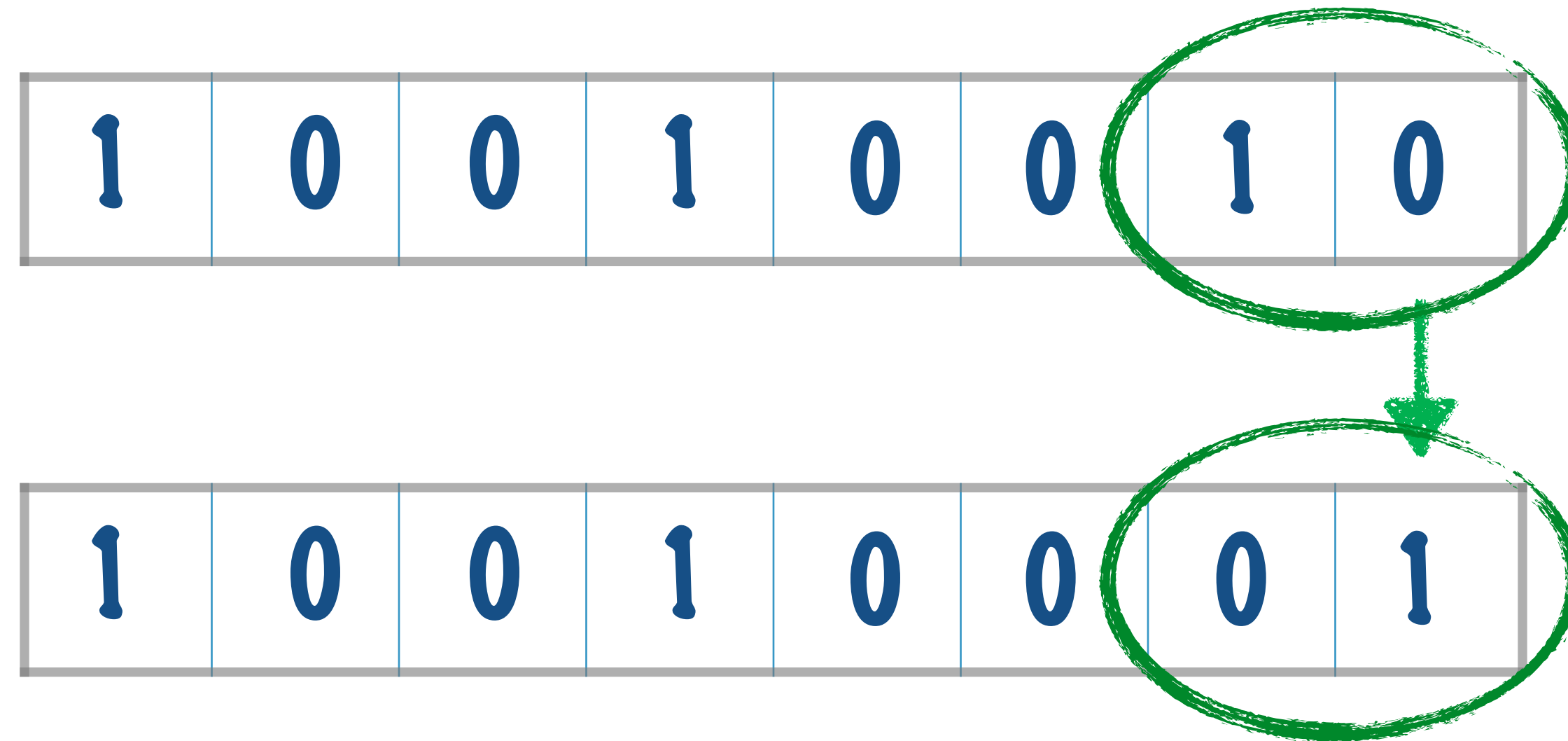
THERE IS ANOTHER, MORE
OPTIMIZED TECHNIQUE

WITH A COMPLEXITY OF
 $O(\text{NUMBER OF 1S})$

COUNT THE NUMBER OF 1 BITS

CONSIDER A NUMBER REPRESENTED
IN BINARY FORM

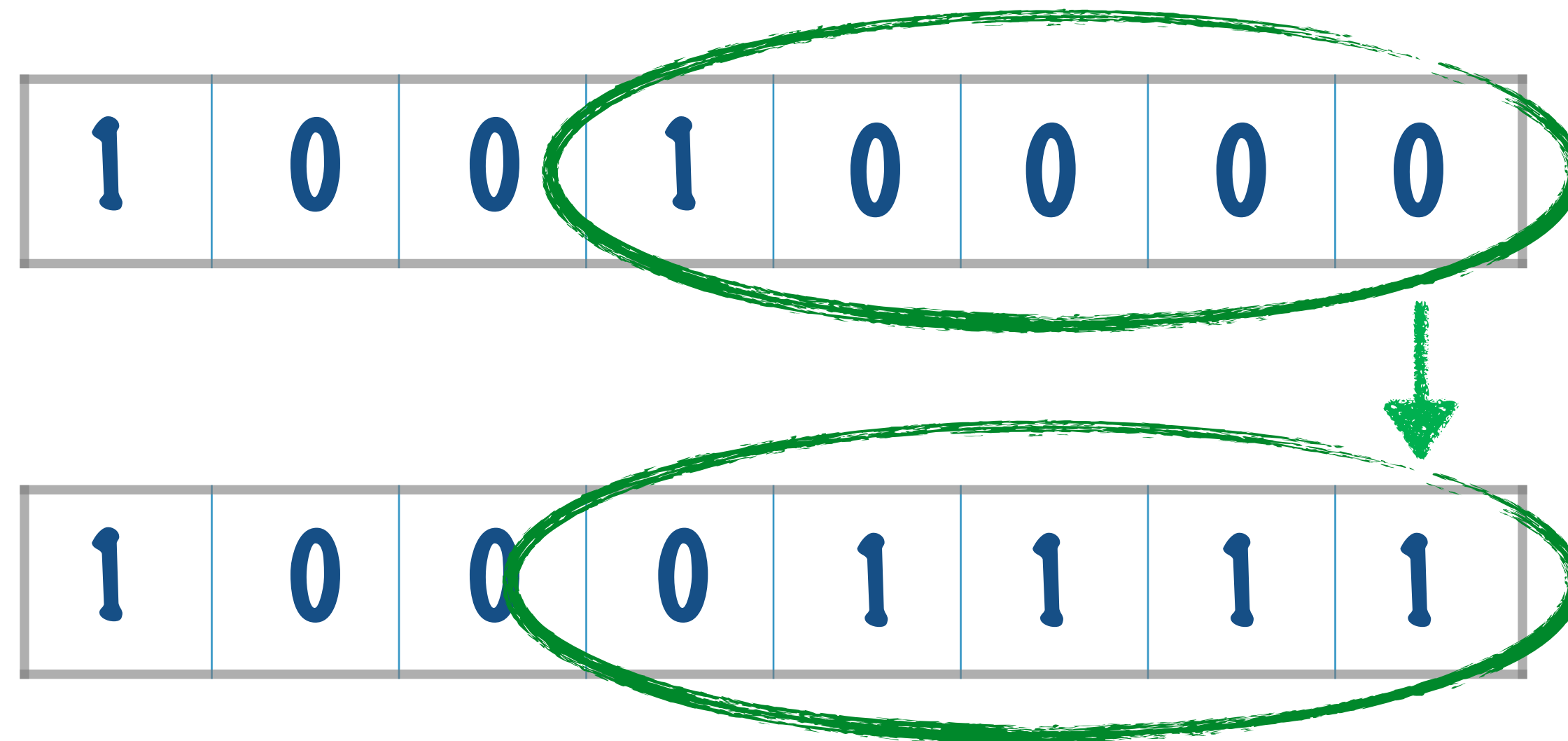
NOW IF YOU
SUBTRACT 1
FROM THIS
NUMBER THIS IS
THE RESULT



ALL THE BITS FROM THE RIGHT UP TO
THE VERY FIRST 1 ARE **TOGGLED**

COUNT THE NUMBER OF 1 BITS

CONSIDER ANOTHER EXAMPLE



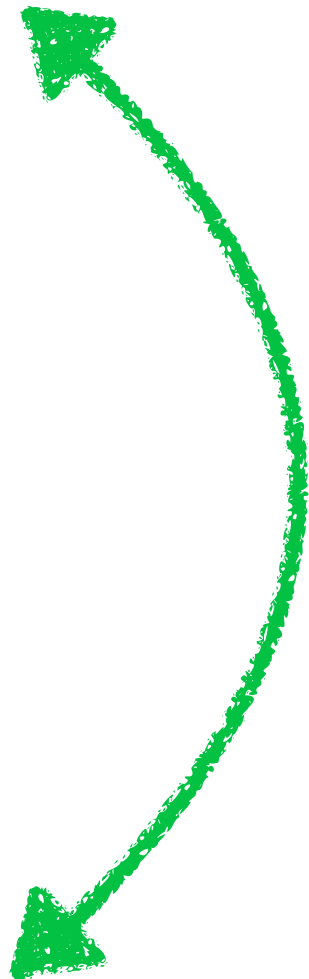
SUBTRACT 1

ALL THE BITS FROM THE RIGHT UP TO
THE VERY FIRST 1 ARE **TOGGLED**

COUNT THE NUMBER OF 1 BITS

NOW IF WE BITWISE AND THESE TWO n & $(n - 1)$

n	1	0	0	1	0	0	0	0
$n - 1$	1	0	0	0	1	1	1	1
	1	0	0	0	0	0	0	0



THE ORIGINAL NUMBER
HAD TWO 1-BITS

THE RESULT HAS JUST ONE

COUNT THE NUMBER OF 1 BITS

CONTINUE THE BITWISE AND TILL THE
NUMBER IS ALL ZEROES

A COUNTER CAN KEEP TRACK OF HOW
OFTEN THIS OPERATION IS PERFORMED

THE COUNT WILL GIVE YOU THE
NUMBER OF 1-BITS IN THE INTEGER

WITH A COMPLEXITY OF
 $O(\text{NUMBER OF 1S})$

COUNT THE NUMBER OF 1 BITS

```
int count_1s_optimized(int num) {  
    int count = 0;  
    while (num != 0) {  
        num = num & (num - 1);  
        count++;  
    }  
  
    return count;  
}
```

INITIALIZE THE COUNT
TO TRACK THE NUMBER
OF 1-BITS



CONTINUE THE BITWISE AND
WITH NUM - 1



THE COUNT FINALLY HOLDS THE
NUMBER OF 1-BITS IN THIS
NUMBER

