

**COSC 2123/1285 Algorithms and Analysis
Tutorial
Maths and Algorithm Analysis Workshop**

Questions

1 Simplify the following summation:

$$\sum_{j=2}^{n+1} 10j$$

Answer:

$$\begin{aligned}\sum_{j=2}^{n+1} 10j &= 10 \sum_{j=2}^{n+1} j \\ &= 10 \left[\sum_{j=2}^{n+1} j + 1 - 1 \right] \\ &= 10 \left[\sum_{j=0}^{n+1} j - 1 \right] \\ &= 10 \left[\sum_{j=0}^n j + (n+1) - 1 \right] \\ &= 10 \left[\sum_{j=0}^n j + n \right] \\ &= 10 \left[\frac{n(n+1)}{2} + n \right] \\ &= 10 \left[\frac{n^2 + 3n}{2} \right] \\ &= 5n^2 + 15n\end{aligned}$$

2 Consider the following recurrence relation. Simplify it to an expression of n variable only, i.e., no recurrence terms in your final expression.

$$C(n) = C(n-1) + 2, C(3) = 1$$

Answer:

$$\begin{aligned}C(n) &= C(n-1) + 2 \\ &= C(n-2) + 2 + 2 \\ &= C(n-3) + 2 + 2 + 2 \\ &= C(n-3) + 3 \cdot 2 \\ &= \dots \\ &= C(n-i) + i \cdot 2 \\ &= \dots \\ &= C(3) + (n-3) \cdot 2 \\ &= 1 + 2n - 6 \\ &= 2n - 5\end{aligned}$$

3 Consider the following mystery function:

Algorithm 1 `RecurMystery($A[0 \dots n - 1]$)`

Input: an array A of size n

Output: T

```

1: if  $n == 1$  then
2:   return 1
3: end if
4:  $T_1 = \text{RecurMystery}(A[0 \dots n - 2])$ 
5:  $T = T_1 + T_1$ 
6: return  $T$ 

```

a) What is this function computing?

Answer: Setup a recurrence for this. Let the mystery function be denoted by $F(n)$. Each iteration the output of $F(n)$ is twice the output of $F(n - 1)$. Therefore the recurrence is:

$$F(n) = 2F(n - 1), F(1) = 1$$

Solving this (this is not examinable, but for your interest):

$$\begin{aligned}
 F(n) &= 2F(n - 1) \\
 &= 2^2 F(n - 2) \\
 &= 2^3 F(n - 3) \\
 &= \dots \\
 &= 2^i F(n - i) \\
 &= \dots \\
 &= 2^{n-1} F(1) \\
 &= 2^{n-1}
 \end{aligned}$$

Hence this function calculates 2^{n-1} .

b) What is the basic operation in this recursive algorithm?

Answer: Addition or comparison. Assignment is possible, but addition is generally more expensive than assignment, and also the 2 assignments can be collapsed to one line and become one assignment.

c) Write the recurrence relation for the number of basic operations executed by this algorithm, including the base/termination case.

Answer: Using addition:

$$C(n) = C(n - 1) + 1, C(1) = 0$$

d) Simplify the recurrence relation.

Answer: Solving this:

$$\begin{aligned}C(n) &= C(n-1) + 1 \\&= [C(n-2) + 1] + 1 \\&= C(n-2) + 2 \\&= [C(n-3) + 1] + 2 \\&= C(n-3) + 3 \\&= \dots \\&= C(n-i) + i \\&= \dots \\&= C(1) + (n-1) \\&= 0 + n - 1 \\&= n - 1\end{aligned}$$