



上海科技大学
ShanghaiTech University

EE140 Introduction to Communication Systems

Lecture 12

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ShanghaiTech University, Fall 2022

Signal Space

- Norm $|x(t)|$. Energy \mathcal{E}_x

$$\|x(t)\| = \left(\int_{-\infty}^{\infty} |x(t)|^2 dt \right)^{1/2} = \sqrt{\mathcal{E}_x}$$

- Inner produce:

$$\langle x_1(t), x_2(t) \rangle = \int_{-\infty}^{\infty} x_1(t) x_2^*(t) dt$$

- $x_1(t), x_2(t)$ are **orthogonal** if their inner product is 0.
- Signals are **orthonormal** if they are orthogonal and their norm are unity
- Signals are **linearly independent** if no signal can be represented as a linear combination of remaining signals
- $\|x_1(t) + x_2(t)\| \leq \|x_1(t)\| + \|x_2(t)\|$
- $\|\langle x_1(t), x_2(t) \rangle\| \leq \|x_1(t)\| \|x_2(t)\|$

Gram-Schmidt for Signals

Given an finite energy signal waveforms $\{s_m(t)\}_{m=1}^M$, generate an orthonormal waveforms $\{\phi_1(t), \phi_2(t), \dots, \phi_M(t)\}$

① $\phi_1(t) = s_1(t)/\sqrt{E_1}.$

② $c_{21} = \langle s_2(t), \phi_1(t) \rangle$

③ $\gamma_2(t) = s_2(t) - c_{21}\phi_1(t)$, and let $E_2 = \int_{-\infty}^{\infty} \gamma_2^2(t)dt$

④ $\phi_2(t) = \gamma_2(t)/\sqrt{E_2}$

⑤ $\phi_k(t) = \gamma_k(t)/\sqrt{E_k}$

where

$$\gamma_k(t) = s_k(t) - \sum_{i=1}^{k-1} c_{ki}\phi_i(t)$$

$$c_{ki} = \langle s_k(t), \phi_i(t) \rangle$$

$$E_k = \int_{-\infty}^{\infty} \gamma_k^2(t)dt$$

Orthonormal Expansions in \mathcal{L}_2

- Given an orthonormal basis $\{\phi_1(t), \phi_2(t), \dots\}$ of \mathcal{L}_2 , any \mathcal{L}_2 signal $x(t)$ can be represented as $x(t) = \sum_j x_j \phi_j(t)$.
 - $x_j = \langle x(t), \phi_j(t) \rangle = \int x(t) \phi_j^*(t) dt$.
 - $\mathbf{x} = (x_1, x_2, \dots)$
 - Equality holds in the sense of MSE=0.
 - $\int_{-\infty}^{\infty} |x(t) - \sum_j x_j \phi_j(t)|^2 dt = 0$
 - $E_x = \int_{-\infty}^{\infty} |x(t)|^2 dt = \sum_j |x_j|^2 = \|\mathbf{x}\|^2$
 - $\langle x(t), s(t) \rangle = \langle \mathbf{x}, \mathbf{s} \rangle$
- Example
 - Fourier Series
 - Sampling Theorem

Orthonormal Expansions in \mathcal{L}_2

- Fourier Series

Given a function $v(t)$, with duration $[-T/2, T/2]$

- Define $\theta_k(t) = e^{2\pi i k t / T} \text{rect}(t/T)$
- $\{\theta_k(t)\}_k$ are orthogonal functions, with $\|\theta_k(t)\|^2 = T$
- Let $\phi_k(t) = \theta_k(t) / \sqrt{T}$:

$$\phi_k(t) = \frac{1}{\sqrt{T}} e^{2\pi i k t / T} \text{rect}(t/T)$$

- By FSE, $v(t)$ can be written as

$$v(t) = \sum_k \alpha_k \phi_k(t)$$

where $\alpha_k = \langle v(t), \phi_k(t) \rangle$.

Orthonormal Expansions in \mathcal{L}_2

- Sampling Theorem

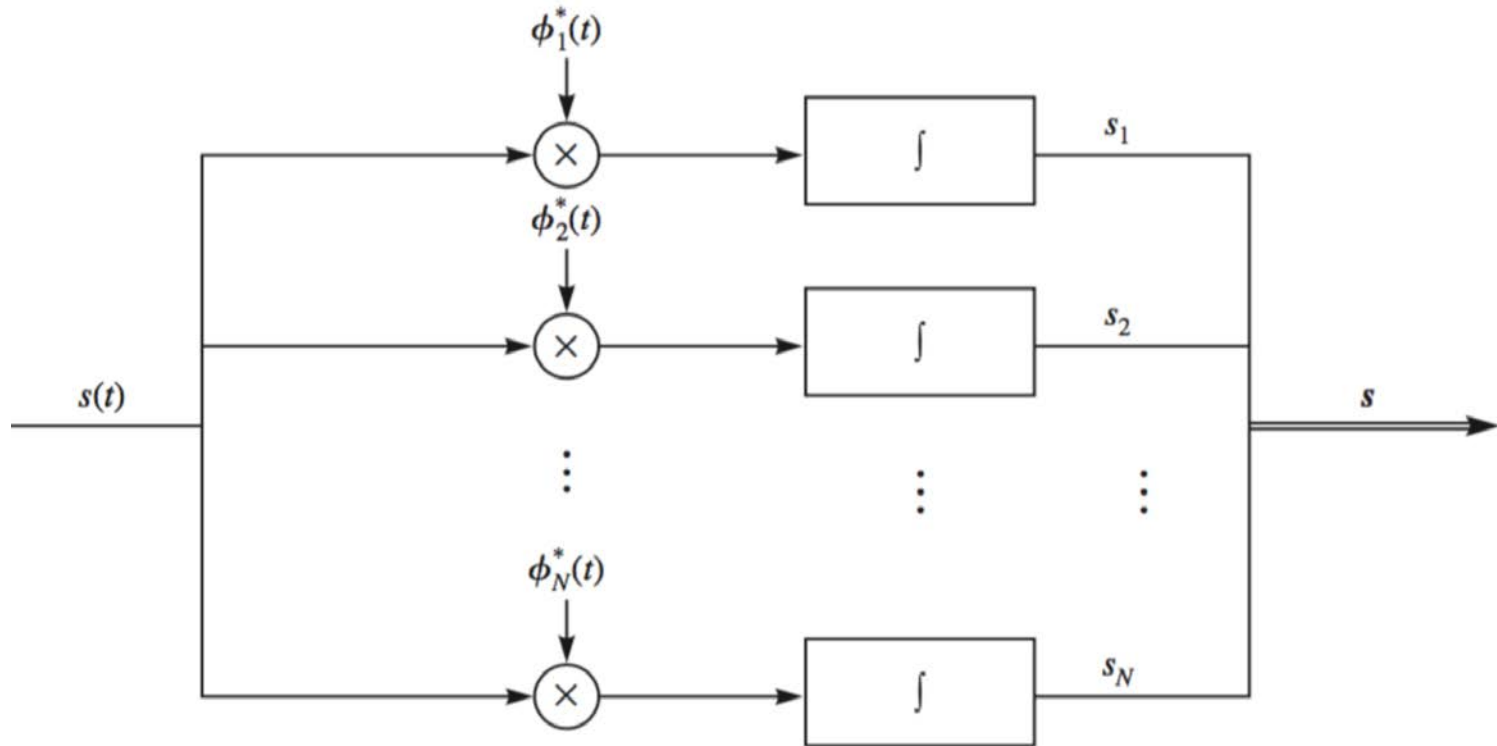
Sampling Theorem: If $x(t)$ is a band-limited in W , then

$$x(t) = \sum_{n=-\infty}^{\infty} x(nT) \operatorname{sinc} \left[\left(\frac{t}{T} - n \right) \right]$$

where $\operatorname{sinc}(x) = \frac{\sin(\pi x)}{\pi x}$, $T = 1/2W$.

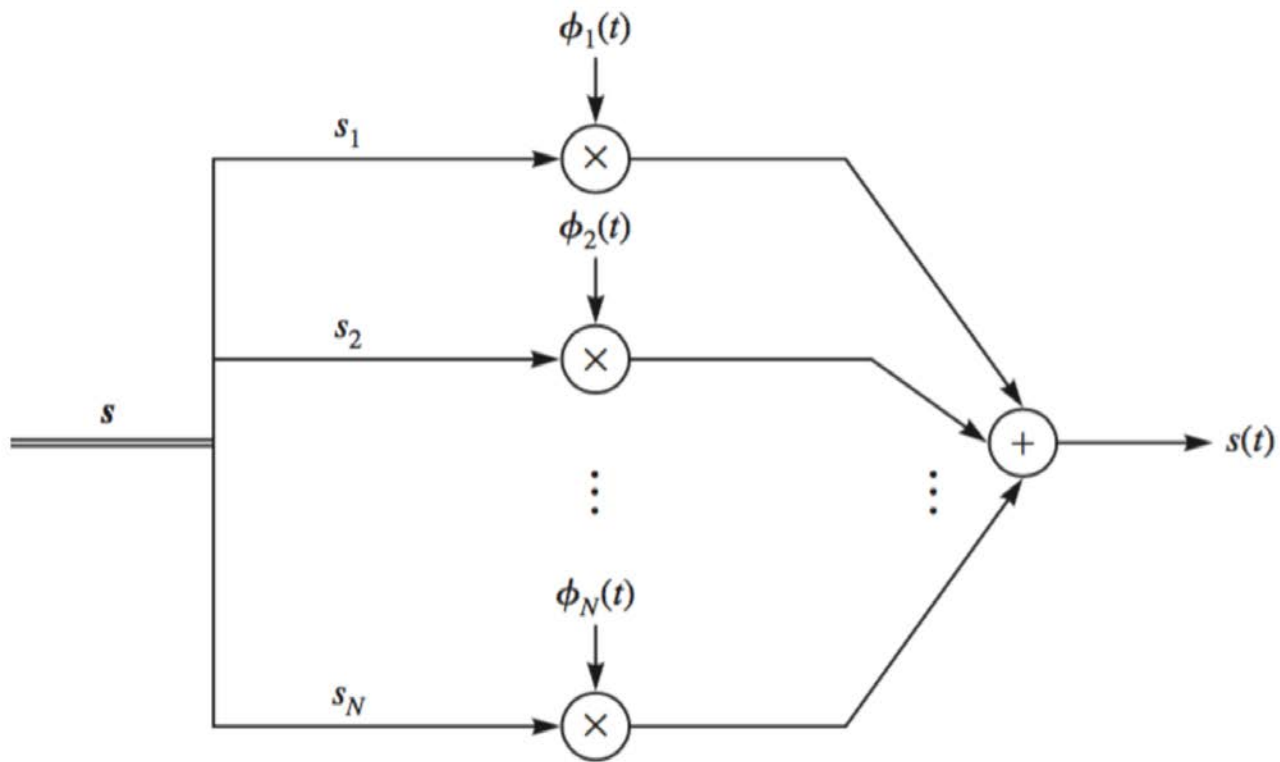
- $x(nT)$ is coefficient (x_k)
- $\{\operatorname{sinc}[(\frac{t}{T} - n)]\}_{n=-\infty}^{\infty}$ are orthogonal function
- Equality holds in sense of $E_e = \int_{-\infty}^{\infty} |x(t) - \hat{x}(t)|^2 = 0$
- View signal $x(t)$ as vectors $x(nT)$

Map Signal to Vector



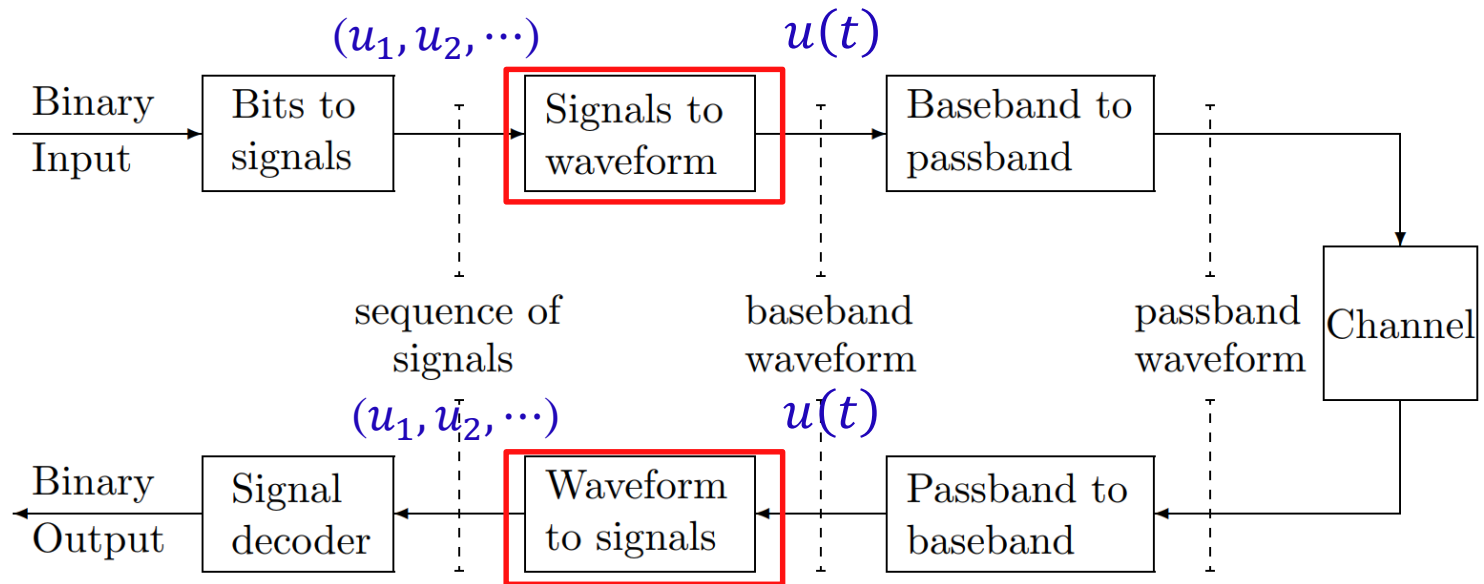
$$s_k = \int_{-\infty}^{\infty} s(t) \phi_k^*(t) dt = \langle s(t), \phi_k(t) \rangle$$

Map Vector to Signal



$$s(t) = \sum_{k=1}^N s_k \phi_k(t)$$

Modulation and Demodulation



$$u(t) = \sum_k u_k p(t - kT_s) \rightarrow (u_1, u_2, \dots)$$

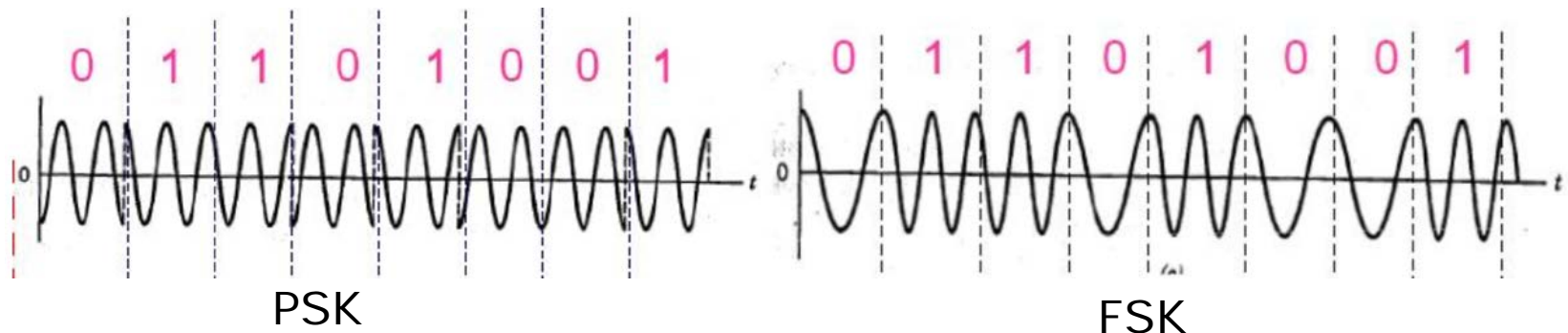
- Step 1: Binary digits \rightarrow A sequence of signal¹(numbers)
- Step 2: A sequence of numbers \rightarrow Waveforms
- Step 3: baseband \rightarrow pass band

Example: Modulation and Demodulation

1. Map information bits into (u_1, u_2, \dots) .
 - $(0010\dots) \rightarrow (+1, +1, -1, +1, \dots)$
 2. Map sequence of numbers (or signals) into a waveform
 - $(u_1, u_2, \dots) \rightarrow u_1 p(t), u_2 p(t - T_s), u_3 p(t - 3T_s), \dots$
 - $u(t) = \sum_k u_k p(t - kT_s), u_k = \{+1, -1\}, T_s$ is the signal interval
 - $\{p(t - kT_s)\}$: baseband pulse waveform (explain later)
 3. Map a baseband waveform $u(t)$ into a passband waveform
 $x(t) = \text{Re}(u(t)e^{j2\pi f_c t}) = u(t) \cos 2\pi f_c t$ (DSB-AM).
- Various Modulations: we can change
 - Amplitude: mapping $\{00, 01, 10, 11\} \rightarrow -3, -1, +1, +3$ (4PAM)
 - Phase: $0 \rightarrow p_1(t), 1 \rightarrow p_2(t) = p_1(t)e^{j\pi}$ (phase-shift-keying)
 - Frequency: $0 \rightarrow p_1(t) = a \cos(2\pi \Delta f t), 1 \rightarrow p_2(t) = a \cos(4\pi \Delta f t)$ (frequency-shift-keying)

Example: Modulation and Demodulation

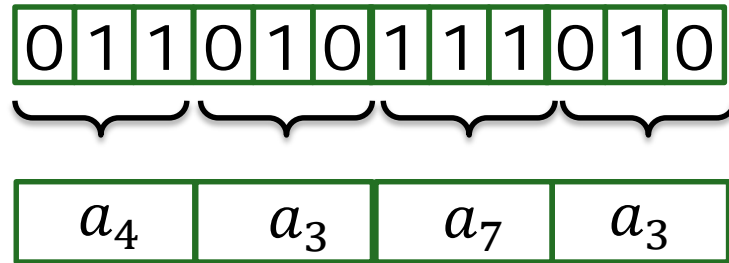
1. Map information bits into (u_1, u_2, \dots) .
 - $(0010\dots) \rightarrow (+1, +1, -1, +1, \dots)$
 2. Map sequence of numbers (or signals) into a waveform
 - $(u_1, u_2, \dots) \rightarrow u_1 p(t), u_2 p(t - T_s), u_3 p(t - 3T_s), \dots$
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 3. Map a baseband waveform $u(t)$ into a passband waveform $x(t) = \text{Re}(u(t)e^{j2\pi f_c t})$.
- Various Modulations: we can change



Pulse Amplitude Modulation (PAM)

Bit rate: R bps

Bit interval: $T_b = 1/R$ s.



b bits blocks ($b=3$)

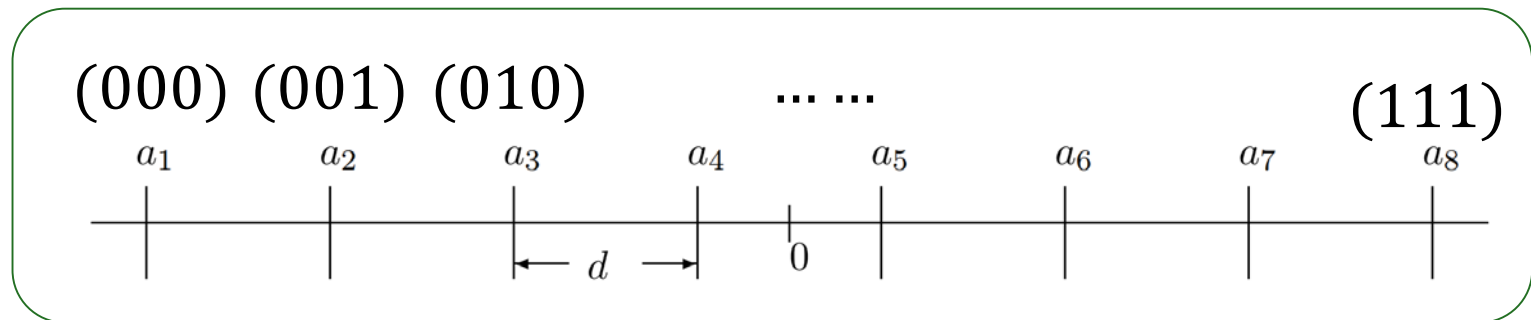
$$M = 2^b$$



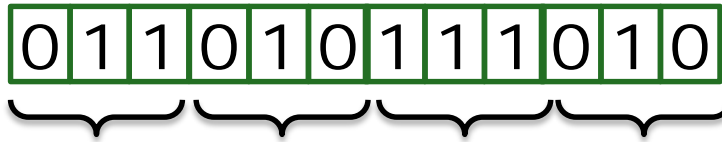
signal constellation $\mathcal{A} = \{a_1, a_2, \dots, a_M\}$ of real numbers

Signal (Symbol) Rate: $R_s = R/b$ signals/s

Signal (Symbol) interval: $T_s = \frac{1}{R_s} = bT_b$ s



Pulse Amplitude Modulation (PAM)



$$u(t) = \sum_k u_k p(t - kT_s)$$

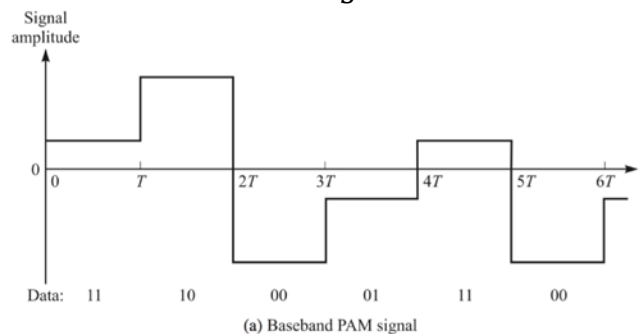
b bits blocks (b=3)

$$M = 2^b$$

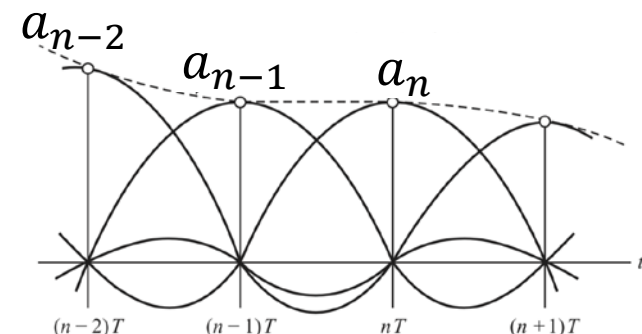


signal constellation $\mathcal{A} = \{a_1, a_2, \dots, a_M\}$, $a_j \in R$

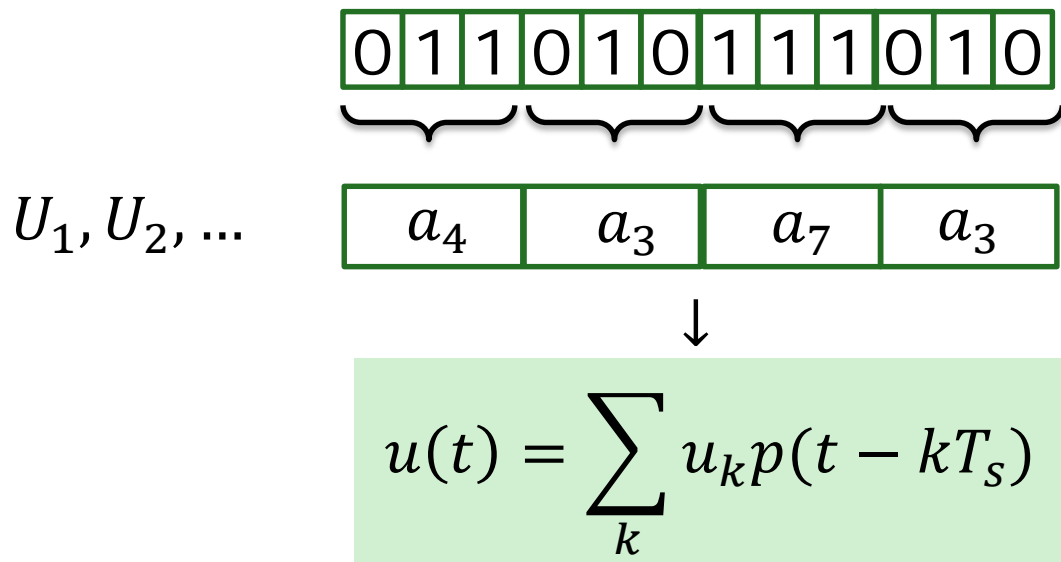
$$T_b = \frac{1}{R}, T_s = \frac{1}{R_s} = \frac{b}{R} = bT_b \text{ (interval between signals)}$$



$$\{00, 01, 10, 11\} \rightarrow -3, -1, +1, +3$$

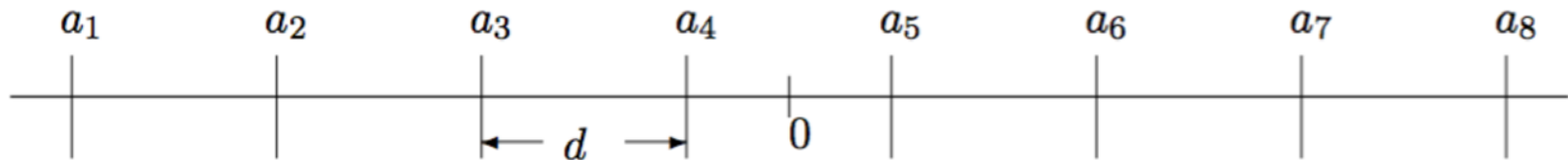


Pulse Amplitude Modulation (PAM)



- Assume: U_1, U_2, \dots , iid $\sim P(U_k = a_m) = p_m$ (k is time index, $m \in \{1, \dots, M\}$ is index in \mathcal{A})
- Energy $E_m = \int |a_m p(t)|^2 dt$
- Average signal energy $E_{\text{avg}} = \sum_{m=1}^M p_m E_m$
- Average energy per bit $E_{\text{bavg}} = \frac{E_{\text{avg}}}{b}$, power $P_{\text{avg}} = E_{\text{bavg}}/T_b$
- Binary PAM: $b = 1$; M-PAM : $M = 2^b$

PAM Signal Constellation



- Assume incoming bits are equiprobable RVs.
- Each signal $U_k = a_m$ is equiprobable in \mathcal{A}

$$\mathcal{A} = \left\{ \frac{-d(M-1)}{2}, \frac{-d(M-3)}{2}, \dots, \frac{-d}{2}, \frac{d}{2}, \dots, \frac{d(M-3)}{2}, \frac{d(M-1)}{2} \right\}$$

- The choose of values in \mathcal{A} is similar to finding representation points in quantization problem
- Average power per signal: E_{avg} ?

PAM Signal Constellation

$$\mathcal{A} = \left\{ \frac{-d(M-1)}{2}, \frac{-d(M-3)}{2}, \dots, \frac{-d}{2}, \frac{d}{2}, \dots, \frac{d(M-3)}{2}, \frac{d(M-1)}{2} \right\}$$

For symbol $U_k = a_m \in \mathcal{A} \Rightarrow$ send $u_k(t) = a_m p(t - kT)$,

$$E_m = \int_{-\infty}^{\infty} a_m^2 p^2(t - kT) dt = a_m^2 E_p$$

Thus,

$$\begin{aligned} E_{\text{avg}} &= \sum_{m=1}^M \left(\frac{1}{M} \cdot a_m^2 E_p \right) \\ &= 2 \frac{E_p}{M} \left(\frac{d}{2} \right)^2 (1^2 + 3^2 + \dots + (M-1)^2) \\ &= \frac{d^2}{2} \frac{E_p}{M} \times \frac{M(M^2 - 1)}{6} \\ &= \frac{d^2 (M^2 - 1) E_p}{12} \end{aligned}$$

PAM Signal Constellation

The signal energy, i.e., the mean square signal value assuming equiprobable signals is:

$$E_{\text{avg}} = \frac{d^2(M^2 - 1)E_p}{12}$$

- E_{avg} increases as d^2 and M^2
- d is determined by the noise
- Errors in reception are primarily due to noise exceeding $d/2$
- For many channels, the noise is independent of the signal, which explains the standard equal spacing between signal constellation values.

Quadrature Amplitude Modulation

$$u(t) = \sum_k u_k p(t - kT) = \sum_k u_k(t)$$

- Segment the incoming binary symbols into ***b*-bits** blocks.
- Map the ***b*-bits** into a **signal constellation** $\mathcal{A} = \{a_1, \dots, a_M\}$, $a_j \in \mathbb{C}$
- $u(t)$ is a complex baseband waveform
- Convert $u(t)e^{j2\pi f_c t}$ to a real waveform

$$u(t) = \sum_k u_k p(t - kT_s)$$

$$\begin{aligned} x(t) &= 2\operatorname{Re}[u(t)e^{j2\pi f_c t}] \\ &= 2\operatorname{Re}[u(t)] \cos(2\pi f_c t) - 2\operatorname{Im}[u(t)] \sin(2\pi f_c t) \\ &= u(t)e^{j2\pi f_c t} + u^*(t)e^{-j2\pi f_c t} \end{aligned}$$

- QAM solves the frequency waste problem of DSB
- When $u(t)$ is real, QAM reduces to DSB PAM
- The factor of 2 is an arbitrary scale factor, and can be left out.

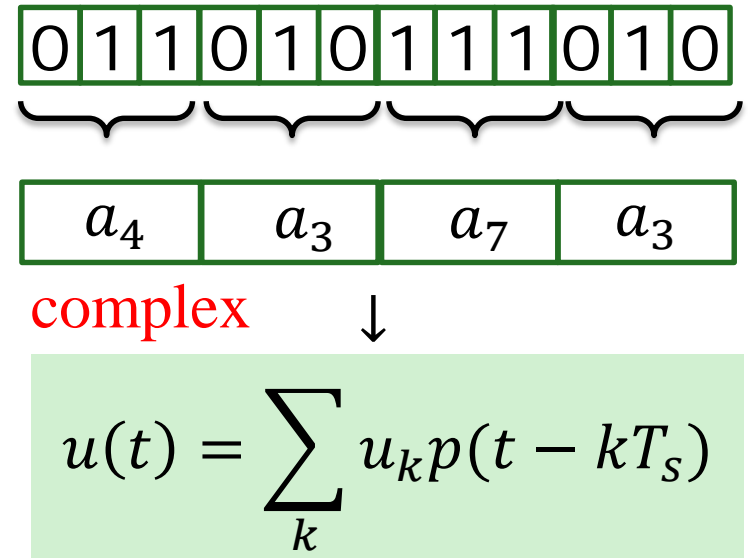
QAM Constellation

- Bit rate: R bps
- Segment binary b bits, map into complex numbers $u_k \in \mathcal{A}$, $|\mathcal{A}| = M = 2^b$.
- Signal (symbol) rate: $R_s = R/b$.
- Standard QAM Constellation
 - Cartesian product: $\mathcal{A} \times \mathcal{B}$, e.g.,

$$\mathcal{A} = \{1, 2\}, \mathcal{B} = \{3, 4\}$$

$$\mathcal{A} \times \mathcal{B} = \{(1, 3), (1, 4), (2, 3), (2, 4)\}$$

$$\mathcal{B} \times \mathcal{A} = \{(3, 1), (3, 2), (4, 1), (4, 2)\}$$



QAM Constellation

- Standard QAM Constellation
 - A standard (\sqrt{M}, \sqrt{M}) -QAM signal set is the **Cartesian product** of two \sqrt{M} -PAM set, i.e.,

$$\mathcal{A} = \{(a' + ja'') | a' \in \mathcal{A}', a'' \in \mathcal{A}'\}$$

where $\mathcal{A}' = \left\{ \frac{-d(\sqrt{M}-1)}{2}, \dots, \frac{-d}{2}, \frac{d}{2}, \dots, \frac{d(\sqrt{M}-1)}{2} \right\}$

- It is a square array signal points located as below for $M = 16$

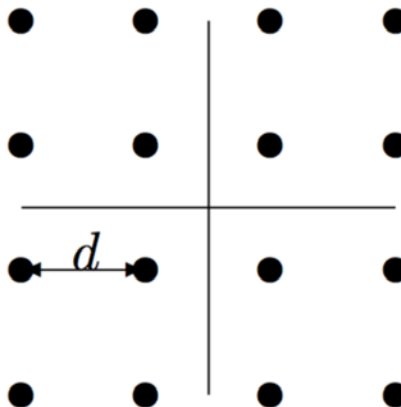


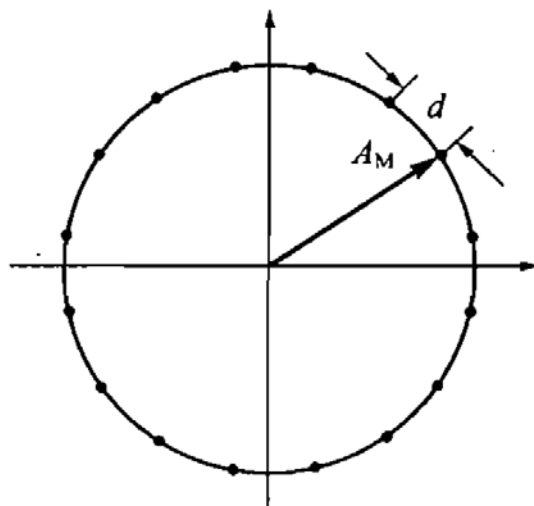
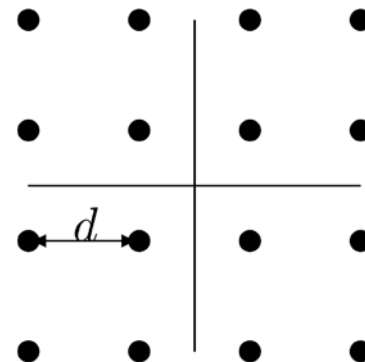
Figure: Standard 16QAM constellation [Gallagar'Book]

QAM Constellation

- Standard QAM Constellation
 - Average energy per 2D signal:

$$E_s = \frac{d^2(M'^2-1)E_p}{6} = \frac{d^2(M-1)E_p}{6}$$

- Q: How to arrange the signal points
 - Choose constellations that minimize E_s given d and M .
 - MPSK



(b) 16PSK

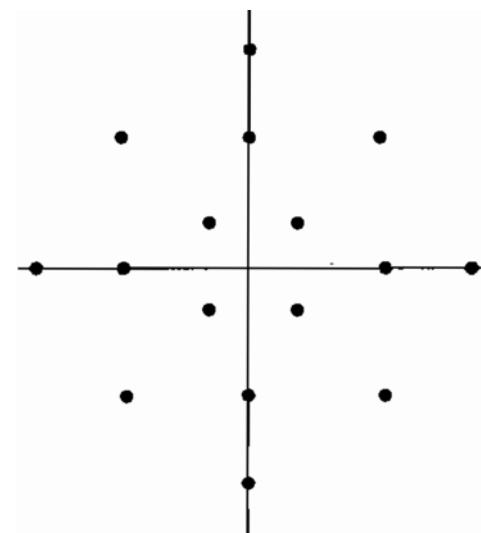
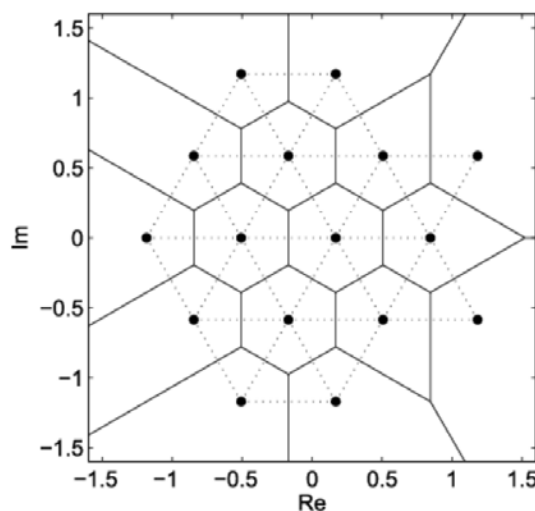
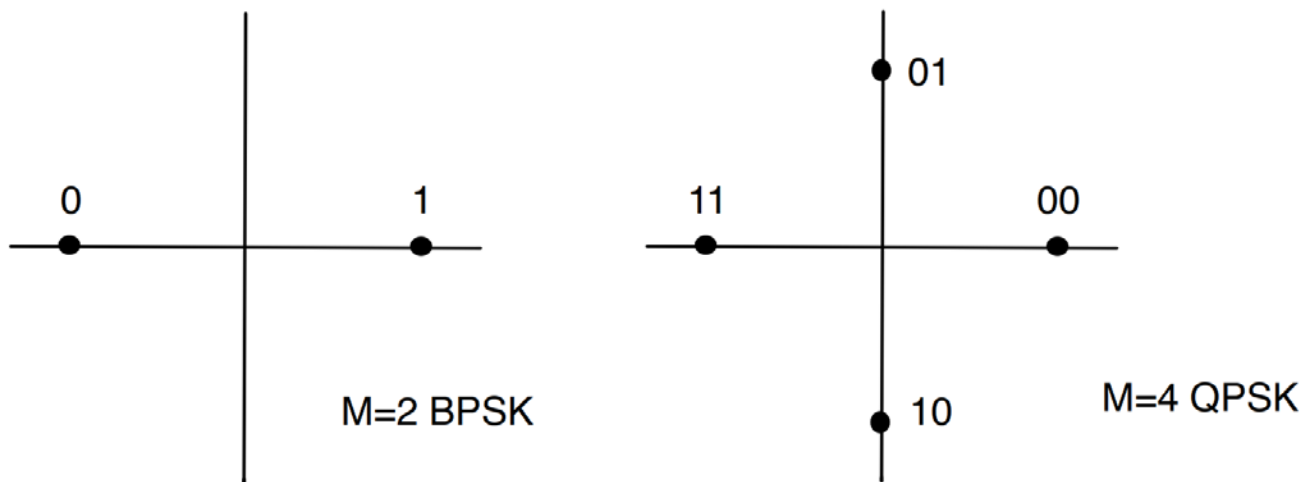


图 8-5 改进的 16QAM 方案 21

QAM Constellation

- PSK: Special Case of QAM
 - $u_k \in \mathcal{A}$

$$\mathcal{A} = \{e^{j\frac{2\pi(m-1)}{M}}, \text{ for } m = 1, \dots, M\}$$



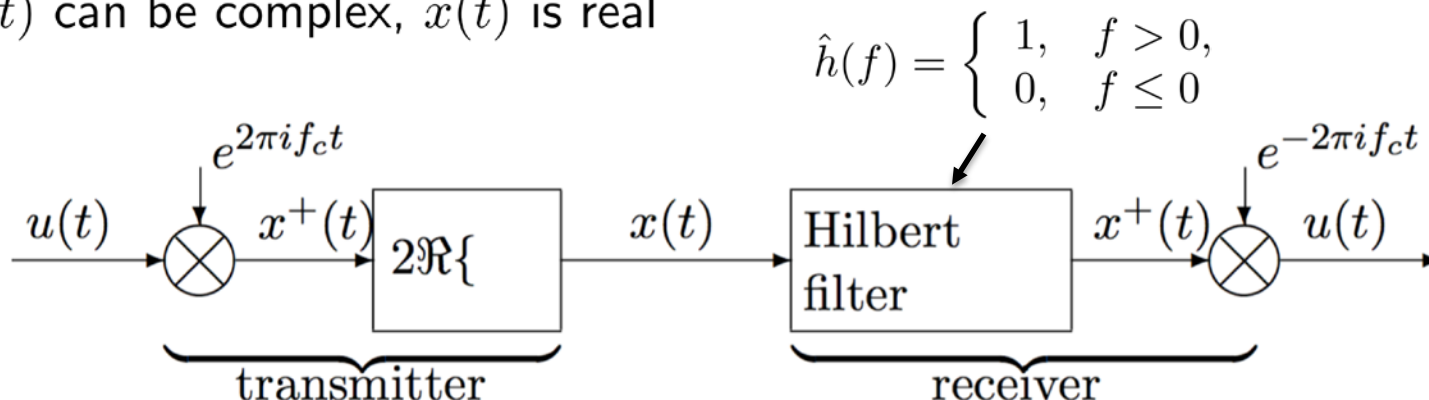
- Signal points have same amplitude
- PSK is rarely used for large M (signal points are very closed)
- Combining *PSK* and *PAM*

QAM: baseband to passband

Shift complex $u(t)$ to f_c , then add complex conjugate to form real $x(t)$

$$\begin{aligned}x(t) &= u(t)e^{j2\pi f_c t} + u^*(t)e^{-j2\pi f_c t} \\&= x^+(t) + (x^+(t))^* \\&= 2\text{Re}[u(t)e^{j2\pi f_c t}]\end{aligned}$$

- Baseband Bandwidth: $B_b < f_c$
 - $u(t)$ can be complex, $x(t)$ is real



- This is nice for analysis, but not usually so for implementation

QAM: baseband to passband

- Easier way

$$u(t) = \sum_k u_k p(t - kT)$$

$$\begin{aligned} x(t) &= 2\operatorname{Re}[u(t)e^{j2\pi f_c t}] \\ &= 2\operatorname{Re}[u(t)] \cos(2\pi f_c t) - 2\operatorname{Im}[u(t)] \sin(2\pi f_c t) \end{aligned}$$

Assume $p(t)$ is real

$$\operatorname{Re}[u(t)] = \sum_k \operatorname{Re}[u_k] p(t - kT)$$

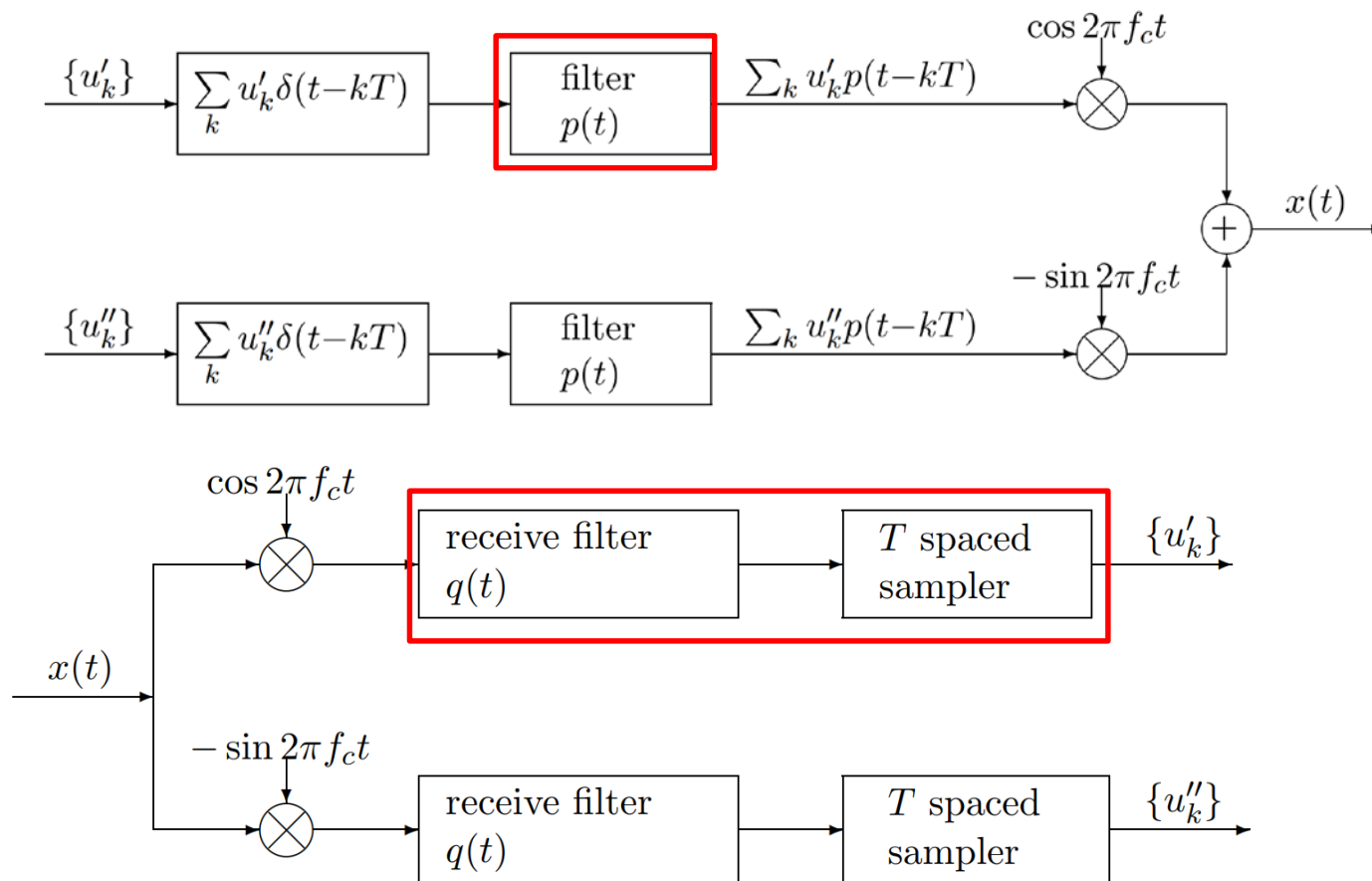
$$\operatorname{Im}[u(t)] = \sum_k \operatorname{Im}[u_k] p(t - kT)$$

with $u'_k = \operatorname{Re}[u_k]$ and $u''_k = \operatorname{Im}[u_k]$

$$x(t) = 2 \left(\sum_k u'_k p(t - kT) \right) \cos 2\pi f_c t - 2 \left(\sum_k u''_k p(t - kT) \right) \sin 2\pi f_c t$$

QAM: baseband to passband

- Easier way

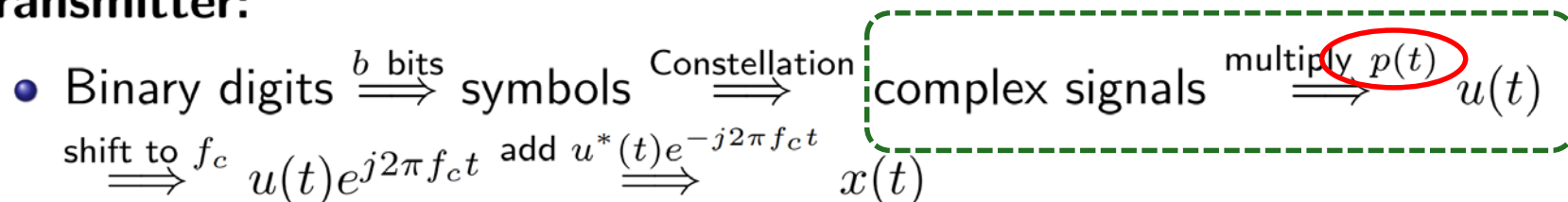


$q(t)$ should be chosen to such that $\hat{g}(f) = \hat{p}(f)\hat{q}(f)$ satisfy the Nyquist criterion. (Explained Later)

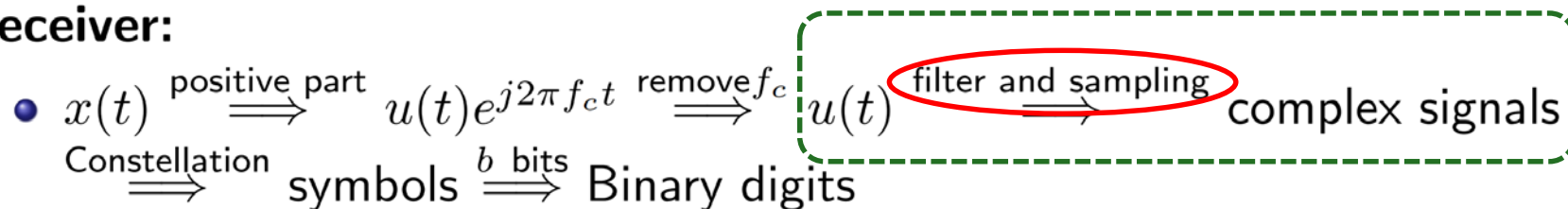
QAM: Implementation

PAM is a special case of QAM

Transmitter:



Receiver:

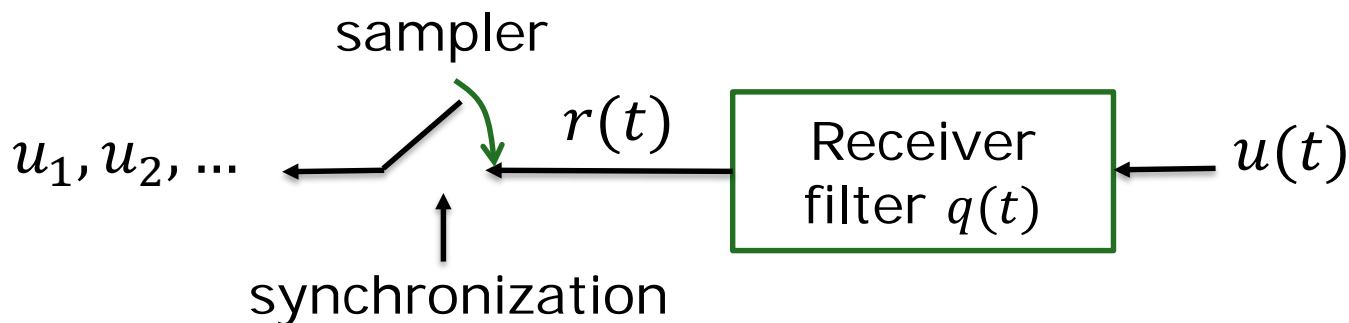


Q: How to choose pulse $p(t)$ and receive filter $q(t)$, so that there is no inter-symbol interference (ISI)??

The Nyquist Criterion

$$u(t) = \sum_k u_k p(t - kT)$$

- Assume no noise in channel
- Received baseband waveform $u(t)$, retrieve the signals $\{u_1, u_2, \dots\}$



- $r(t) = \int_{-\infty}^{\infty} u(\tau)q(t - \tau)d\tau$ is sampled as $r(0), r(T), \dots$
- Objective: choose $p(t)$ and $q(t)$ so that $r(kT) = u_k$ (No ISI)

The Nyquist Criterion

$$\begin{aligned} r(t) &= \int u(\tau)q(\tau - t)d\tau = \int \sum_k u_k p(\tau - kT)q(t - \tau)d\tau \\ &= \sum_k u_k g(t - kT) \quad \text{where } g(t) = p(t) * q(t) \end{aligned}$$

- While ignoring noise, $r(t)$ is determined by $g(t)$
- Ideal Nyquist: a wave form $g(t)$ is **ideal Nyquist** with period T if **$g(0) = 1$ and $g(kT) = 0$ for $k \neq 0$** (same property as sinc function)

$$r(jT) = \sum_k u_k g(jT - kT) \stackrel{g(t) \text{ is ideal Nyquist}}{=} u_j$$

- If $g(t)$ is ideal Nyquist, then $r(kT) = u_k$ for all k . If $g(t)$ is not ideal Nyquist, then $r(kT) \neq u_k$ for some u_k (intersymbol interference)

Recall: Aliasing Problem

Given a signal $u(t)$, its sampling approximation

$$s(t) = \sum_k u(kT) \operatorname{sinc}\left(\frac{t}{T} - k\right),$$

and the Fourier transform of $s(t)$ satisfies:

$$\hat{s}(f) = \mathcal{F}(s(t)) = \sum_k \hat{u}(f + \frac{k}{T}) \operatorname{rect}(fT)$$

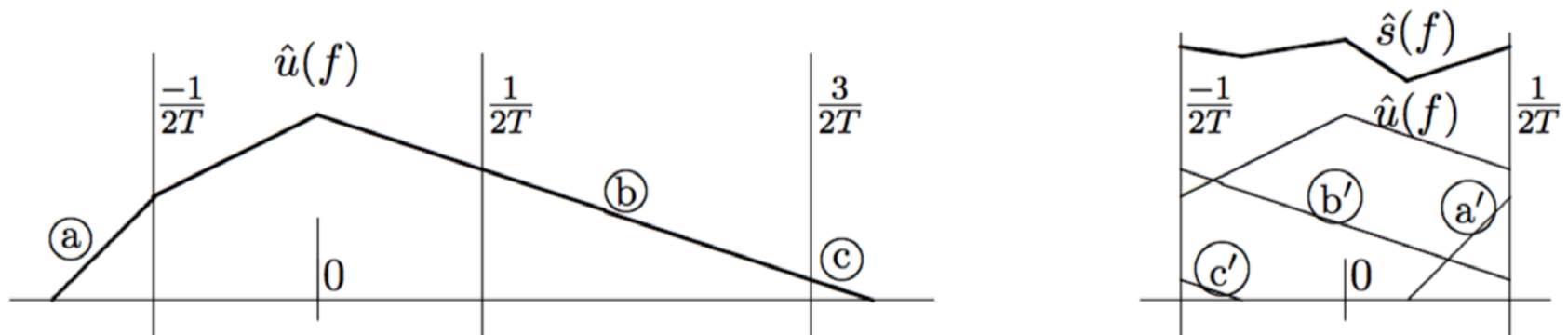


Figure: Aliasing when $1/(2T) < B_b$ [Gallagar'Book]

The Nyquist Criterion

Let $s(t)$ be signal reconstructed by samples of $g(t)$:

$$s(t) = \sum_k g(kT) \operatorname{sinc}\left(\frac{t}{T} - k\right) \quad (11)$$

$g(t)$ is ideal Nyquist: $g(kT) = \delta(k)$, substitute it into (11), we have

$$s(t) = \operatorname{sinc}\left(\frac{t}{T}\right) \longrightarrow \hat{s}(f) = T \operatorname{rect}(fT) \quad (12)$$

From (11) and the aliasing theorem

$$\hat{s}(f) = \sum_k \hat{g}\left(f + \frac{k}{T}\right) \operatorname{rect}(fT) \quad (13)$$

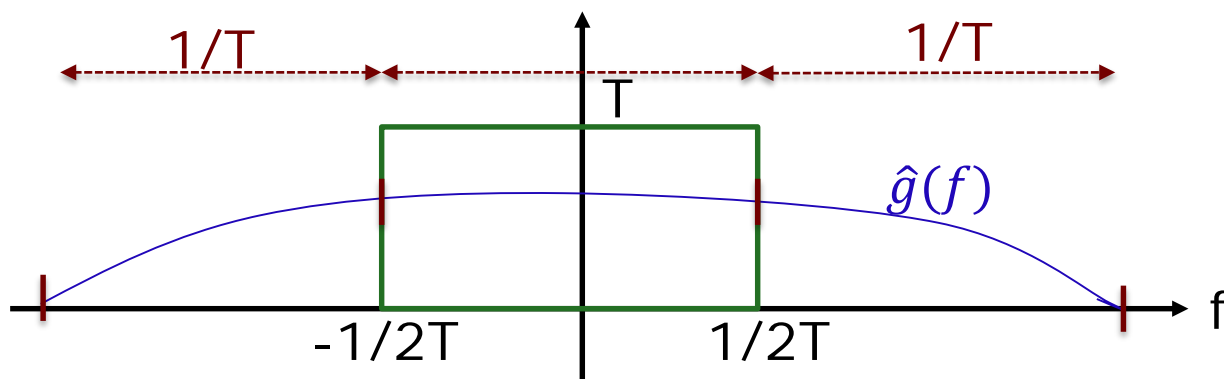
Thus from (12) and (13), we have $g(t)$ is ideal Nyquist iff

$$\sum_k \hat{g}\left(f + \frac{k}{T}\right) \operatorname{rect}(fT) = T \operatorname{rect}(fT)$$

The Nyquist Criterion

- Nyquist Criterion

$$\sum_k \hat{g}(f + \frac{k}{T}) \text{rect}(fT) = T \text{rect}(fT)$$

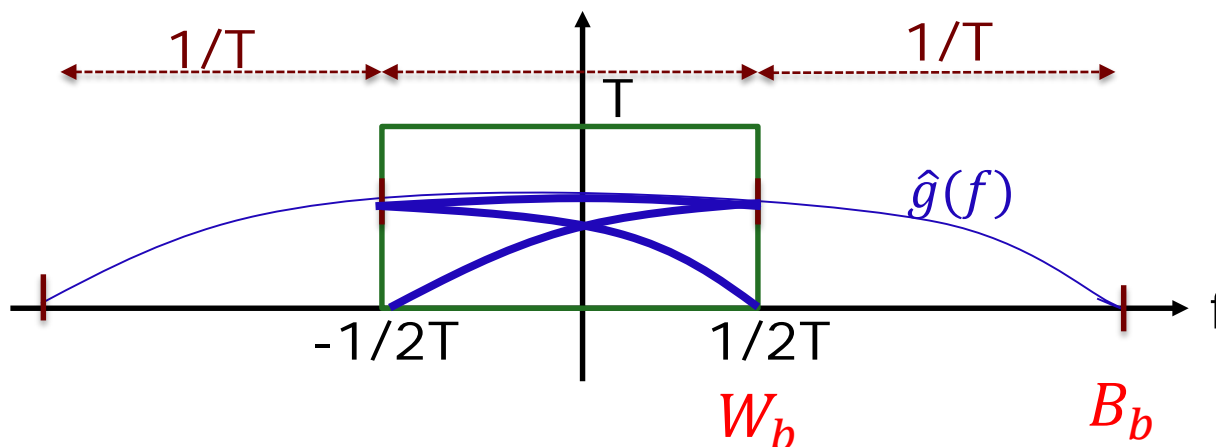


The Nyquist Criterion

- Nyquist Criterion

$$\sum_k \hat{g}(f + \frac{k}{T}) \text{rect}(fT) = T \text{rect}(fT)$$

- Signal Interval of $g(t)$: T
- Nyquist bandwidth: $W_b = \frac{1}{2T}$; Signal rate $R_s = \frac{1}{T} = 2W_b$
- Actual baseband bandwidth: B_b ($\hat{g}(f) = 0, |f| > B_b$)

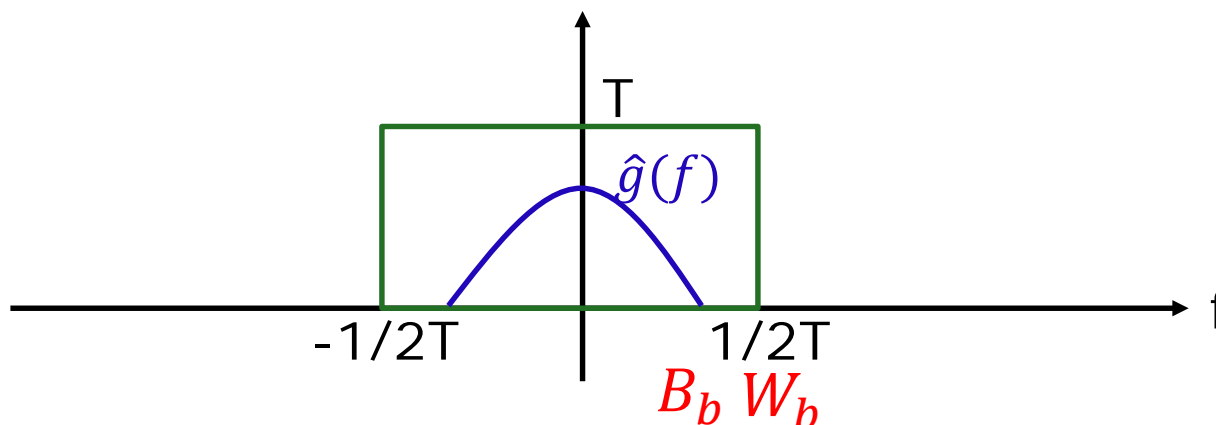


The Nyquist Criterion

- Nyquist Criterion

$$\sum_k \hat{g}(f + \frac{k}{T}) \text{rect}(fT) = T \text{rect}(fT)$$

- Nyquist bandwidth: $W_b = \frac{1}{2T}$; Signal rate $R_s = \frac{1}{T} = 2W_b$
- Actual baseband bandwidth: B_b ($\hat{g}(f) = 0, |f| > B_b$)
- Case1: $B_b < W_b$, ISI is not avoidable



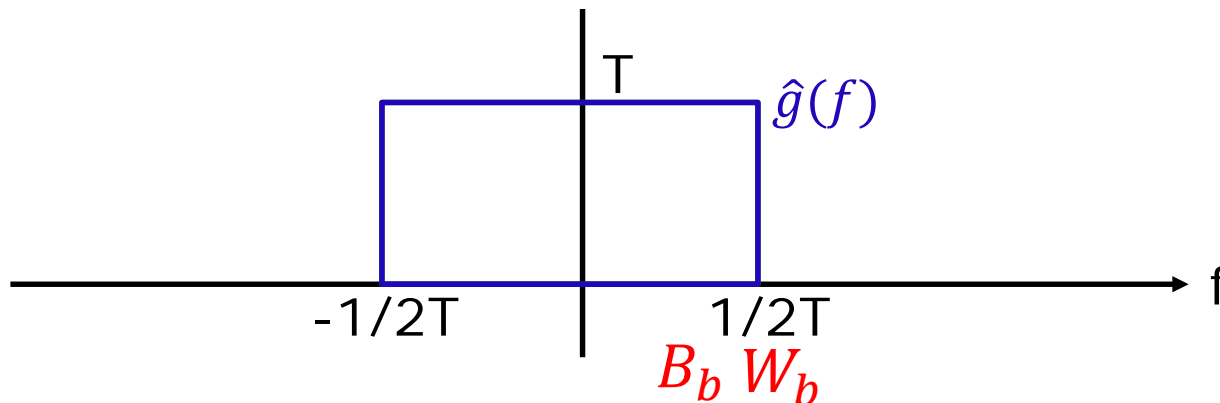
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- Actual baseband bandwidth: B_b ($\hat{g}(f) = 0, |f| > B_b$)
- Case2: $B_b = W_b$

$$\hat{g}(f) = T \text{rect}(fT) \leftrightarrow g(t) = \text{sinc}(\frac{t}{T})$$



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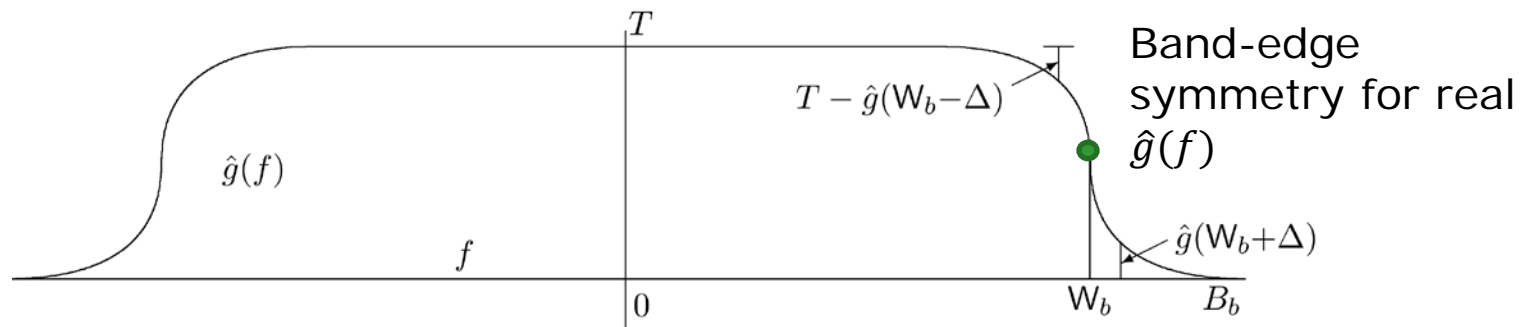
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- Actual baseband bandwidth: B_b ($\hat{g}(f) = 0, |f| > B_b$)
- Case3: $B_b > W_b$
 - If B_b is much larger than W_b , then W_b can be increased (T can be decreased), thus increasing the rate R_s at which signal can be transmitted.
 - $g(t)$ should be chosen B_b exceed W_b by a relatively small amount.
 - $W_b \leq B_b < 2W_b$: Keep $\hat{g}(f)$ almost baseband limited to $1/(2T)$

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- Actual baseband bandwidth: B_b ($\hat{g}(f) = 0, |f| > B_b$)
- Case3: $W_b \leq B_b < 2W_b$
 - Rolloff factor: $\alpha = \frac{B_b}{W_b} - 1$
 - Tradeoff between rolloff and smoothness (slow time decay and bandwidth)



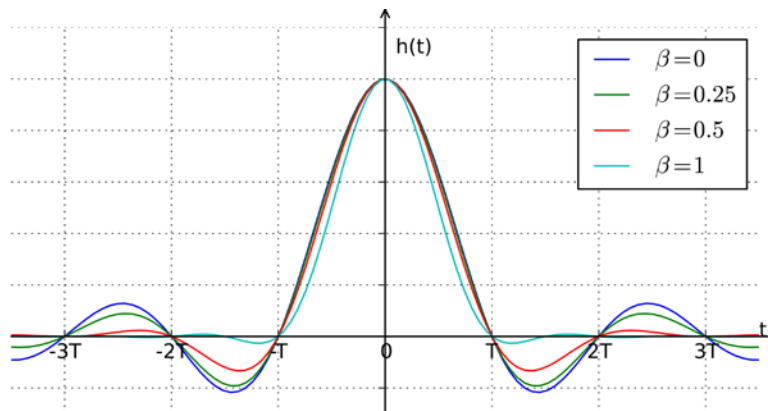
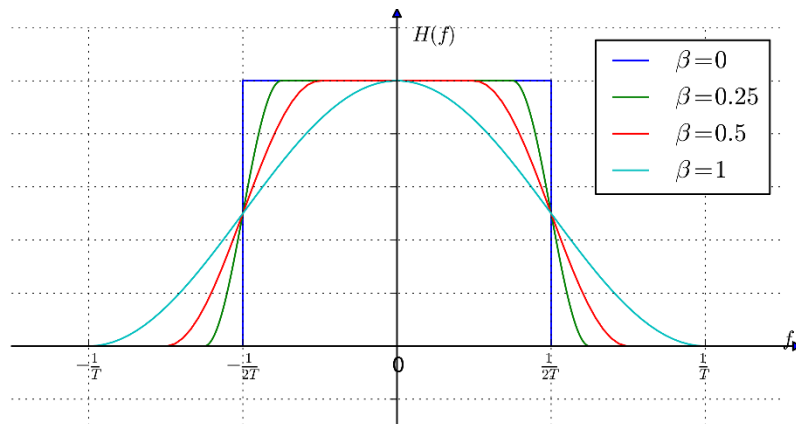
The Nyquist Criterion

- Nyquist Criterion
 - PAM filters in practice often have raised cosine transform

$$\hat{g}_\alpha(f) = \begin{cases} T, & 0 \leq |f| \leq \frac{1-\alpha}{2T}; \\ T \cos^2 \left[\frac{\pi T}{2\alpha} \left(|f| - \frac{1-\alpha}{2T} \right) \right], & \frac{1-\alpha}{2T} \leq |f| \leq \frac{1+\alpha}{2T}; \\ 0, & |f| \geq \frac{1+\alpha}{2T}. \end{cases}$$

- The inverse transform of $\hat{g}_\alpha(f)$

$$\hat{g}_\alpha(t) = \text{sinc} \left(\frac{t}{T} \frac{\cos(\pi\alpha \frac{t}{T})}{1 - 4\alpha^2 t^2 / T^2} \right)$$



The Nyquist Criterion

- Q: Restrict $\hat{g}(f)$ real and nonnegative, $\hat{g}(f) > 0$, how to choose $p(t)$ and $q(t)$, s.t. $\hat{p}(f)\hat{q}(f) = \hat{g}(f)$
- Choose $\hat{q}(f) = \hat{p}^*(f) \rightarrow |\hat{p}(f)| = |\hat{q}(f)| = \sqrt{\hat{g}(f)}$ and $q(t) = p^*(-t)$.
 - $p(t)$: square root of Nyquist;
 - $q(t)$: matched filter to $p(t)$
 - Same bandwidth for $\hat{p}(f)$, $\hat{q}(f)$ and $\hat{g}(f)$ (slightly larger than $1/2T$), truncated in time to allow finite delay.
- **Orthonormal Shifts:**
 - Let $p(t) \in \mathcal{L}_2$ and $\hat{g}(f) = |\hat{p}(f)|^2$ satisfies the Nyquist criterion for T . Then $\{p(t - kT); k \in \mathbb{Z}\}$ is a set of orthonormal functions. Conversely, if $\{p(t - kT); k \in \mathbb{Z}\}$ is a set of orthonormal functions, then $|\hat{p}(f)|^2$ satisfies the Nyquist criterion.
 - Proof:

The Nyquist Criterion

$$g(t) = \int p(\tau)q(t - \tau)d\tau = \int p(\tau)p^*(\tau - t)d\tau$$

$$g(t) = \int p(\tau)p^*(\tau - t)d\tau$$

$$g(kT) = \int p(\tau)p^*(\tau - kT)d\tau$$

- Shift τ by jT for any integer j ($\tau = \tau - jT$)

$$\begin{aligned} g(kT) &= \int p(\tau - jT)p^*(\tau - (k + j)T)d\tau \\ &= \int p(t - jT)p^*(t - (k + j)T)dt \end{aligned}$$

- If $g(t)$ is ideal Nyquist, then $g(kT) = 1$ for $k = 0$ and 0 otherwise. Thus,

$$g(kT) = \int p(t - jT)p^*(t - (k + j)T)dt = \begin{cases} 1 & \text{for } k = 0 \\ 0 & \text{for } k \neq 0 \end{cases}$$

$\Rightarrow \{p(t - kT)\}$ is a set of **orthonormal** functions

The Nyquist Criterion

- if $\{p(t - kT); k \in \mathbb{Z}\}$ is a set of orthonormal functions, then let $\hat{q}(f) = \hat{p}^*(f)$, that is $q(t) = p^*(-t)$. We have

$$g(t) = \int p(\tau)q(t - \tau)d\tau = \int p(\tau)p^*(\tau - t)d\tau$$

$$g(kT) = \int p(\tau)p^*(\tau - kT)d\tau$$

- Shift τ by jT for any integer j ($\tau = \tau - jT$)

$$g(kT) = \int p(\tau - jT)p^*(\tau - (k + j)T)d\tau$$

$$= \int p(t - jT)p^*(t - (k + j)T)dt$$

$$g(kT) = \int p(t - jT)p^*(t - (k + j)T)dt = \begin{cases} 1 & \text{for } k = 0 \\ 0 & \text{for } k \neq 0 \end{cases}$$

$\Rightarrow g(t)$ is ideal Nyquist

The Nyquist Criterion

- Orthonormal Shifts \rightarrow recover u_k from vector space perspective

$$\hat{s}(t) = \sum_{k=1}^N s_k \phi_k(t)$$
$$s_k = \int_{-\infty}^{\infty} s(t) \phi_k^*(t) dt = \langle s(t), \phi_k(t) \rangle$$

- Since $\{p(t - kT)\}$ is an orthonormal basis and $u(t) = \sum u_k p(t - kT)$ is the orthonormal expansion, thus

$$u_k = \langle u(t), p(t - kT) \rangle$$

- Retrieving u_k corresponding to projecting $u(t)$ onto $p(t - kT)$
- Note this projection is done by $u(t) * q(t)$ and then sampling at time kT

Summarize

- Bit sequence to symbol sequence
 - PAM: b-tuple of bits \rightarrow one of $M = 2^b$ signal points in \mathbb{R}^1
 - QAM: b-tuple of bits \rightarrow one of $M = 2^b$ signal points in \mathbb{C}^1
 - Signal constellation: small average energy with a large distance between points
 - Standard Mapping
- Symbol sequences to baseband waveform
 - $u(t) = \sum_k u_k p(t - kT)$
 - How to choose $p(t)$: $\hat{g}(f) = \hat{p}(f)\hat{q}(f)$ must satisfy Nyquist criterion to avoid ISI

$$\text{if } g(kT) = 1 \text{ for } k = 0, \text{ and } 0 \text{ for } k \in \mathbb{Z} \neq 0$$

$$\sum_k \hat{g}(f + \frac{k}{T}) \text{rect}(fT) = T \text{rect}(fT)$$

Summarize

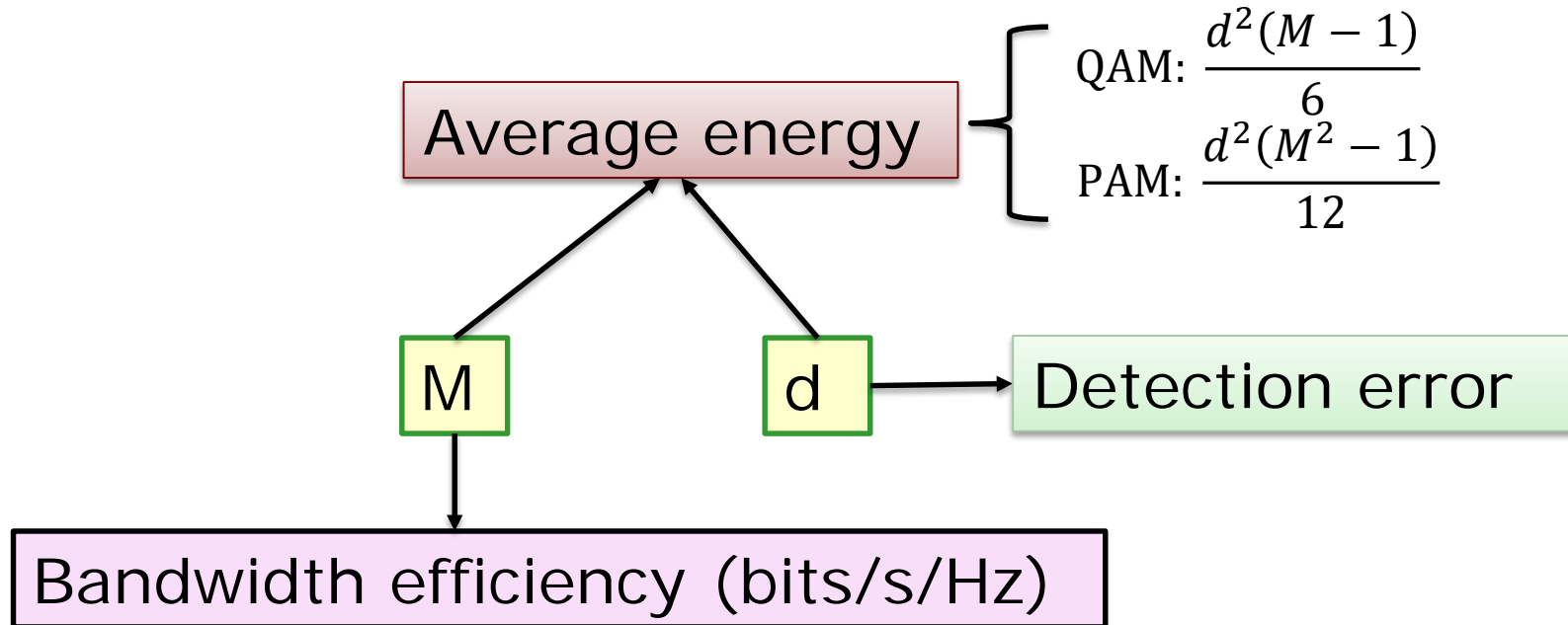
- Symbol sequences to baseband waveform

$$\sum_k \hat{g}(f + \frac{k}{T}) \text{rect}(fT) = T \text{rect}(fT)$$

- $\hat{g}(f)$: almost baseband limited to $1/2T$; smooth such that $g(t)$ goes to zero quickly
 - If $\hat{g}(f) > 0$: $|\hat{p}(f)| = |\hat{q}(f)| = \sqrt{\hat{g}(f)}$. In this case, $\{p(t - kT); k \in \mathbb{Z}\}$ is a set of orthonormal functions
- Baseband to passband: $x(t) = 2R(u(t)e^{j2\pi f_c t})$
$$x(t) = 2 \left(\sum_k u'_k p(t - kT) \right) \cos 2\pi f_c t - 2 \left(\sum_k u''_k p(t - kT) \right) \sin 2\pi f_c t$$
- Tradeoff: Bandwidth Efficiency (M), Energy (M,d), Detection Error (d)

Summarize

- Tradeoff: Bandwidth Efficiency (M), Energy (M,d), Detection Error (d)



$$\text{Bandwidth} = 1/(2T) = R_s/2 = R/(2\log M)$$

$M \nearrow$ Bandwidth \searrow (same R), Bandwidth efficiency \nearrow



Thanks for your kind attention!

Questions?

Summarize

- Bandwidth Efficiency

- $\mu = \frac{R}{B} = \frac{\log_2 MR_s}{B} = \log_2 M \frac{1}{B_b T}$ (# of bits/s/Hz)

- $\mu_s = \frac{1}{B_b T} \times a$ (# of real symbols/s/Hz)

- PAM:

- BB $B_b \approx \frac{1}{2T}$, PB $B_b \approx \frac{1}{T}$ (DSB, half of bandwidth with information)

- $\mu_s = 2$ at BB.

- $\mu = 2\log_2 M$ at BB.

- QAM:

- BB $B_b \approx \frac{1}{2T}$, PB $B_b \approx \frac{1}{T}$ (DSB)

- $\mu_s = 2$ at PB.

- $\mu = \log_2 M$ at PB.