# The Method of Chaining

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MA5249 Project Presentation Part 1

### Outline

- Introduction
- 2 Finite Maxima
- Covering and Packing
- Wasserstein Law of Large Numbers

### Introduction to the Task

**Motivation:** Given a random process  $\{X_t\}_{t\in\mathcal{T}}$ , we would like to control its regularity.

**Main task:** Control the regularity of the supremum of a random process  $\{X_t\}_{t\in\mathcal{T}}$  by establishing **lower bounds** and **upper bounds** on

$$\mathbb{E}\left[\sup_{t\in\mathcal{T}}X_t\right]\quad\text{ and }\quad P\left(\sup_{t\in\mathcal{T}}X_t\geq x\right).$$

**Example:** Operator norm of random matrices

$$||M||_{\mathsf{op}} := \sup_{||v||_2 = 1} ||Mv||_2 = \sup_{v \in \mathbb{S}^n} M_v.$$

### Definition 1.1.

Let  $\sigma \in \mathbb{R}_{>0}$ . A random variable X is  $\sigma^2$ -subgaussian if for all  $\lambda \in \mathbb{R}$ , we have

$$\mathbb{E}[e^{\lambda X}] \leq e^{\frac{\lambda^2 \sigma^2}{2}}.$$

### Finite Maxima

**First step:** Control the maximum of finitely many variables  $\{X_1, X_2, \dots, X_k\}$ .

Naive approach: Bound the supremum by the sum

$$\max_{t \in T} X_t \le \sum_{t \in T} |X_t|.$$

**Key idea:** Use concavity to improve the dependence on |T|. For example,

$$\mathbb{E}\left[\max_{t\in T}X_{t}\right] \leq \mathbb{E}\left[\max_{t\in T}|X_{t}|^{2}\right]^{1/2} \leq |T|^{1/2}\max_{t\in T}\left{\mathbb{E}\left[\left|X_{t}\right|^{2}\right]^{1/2}\right\}.$$

# Finite Maximal Inequality

### Lemma 2.1. (Maximal Inequality)

Let  $\{X_t\}_{t\in T}$  be a random process. Suppose that  $X_t$  is  $\sigma^2$ -subgaussian for each  $t\in T$ . Then we have

$$\mathbb{E}\bigg[\max_{t\in T} X_t\bigg] \leq \sqrt{2\sigma^2\log|T|}.$$

### Proof.

For any  $\lambda \in \mathbb{R}_{>0}$ , Jensen's inequality gives

$$\mathbb{E}\bigg[\max_{t\in\mathcal{T}}X_t\bigg] \leq \frac{1}{\lambda}\log\bigg(\mathbb{E}\Big[e^{\lambda\max_{t\in\mathcal{T}}X_t}\Big]\bigg) \leq \frac{1}{\lambda}\log\bigg(\sum_{t\in\mathcal{T}}\mathbb{E}\Big[e^{\lambda X_t}\Big]\bigg) \leq \frac{\log|\mathcal{T}|}{\lambda} + \frac{\lambda\sigma^2}{2}.$$

The conclusion follows by optimizing in  $\lambda$  and choosing  $\lambda = \frac{\sqrt{2\log |T|}}{\sigma}$ .

Obtain elementary but workable bounds.

## Generalising to the Infinite Case

To have a meaningful theory when  $|T| = +\infty$ , we need to exploit the relationship between the variables in  $\{X_t\}_{t \in T}$ .

**Key idea:** Leverage on the **rich structure** of the index set T.

**Example:** Homogeneous discrete Markov chain

$$P(X_r = k \mid X_i = x_i, \text{ for } i = 0, \dots, r - 1, r + 1, \dots, n)$$
  
= $P(X_r = k \mid X_{r-1} = x_{r-1}, X_{r+1} = x_{r+1}).$ 

The random variable  $X_r$  shares the **strongest dependence** with its **adjacent** neighbours.

# Lipschitz Processes and $\epsilon$ -nets

To generalise this concept to arbitrary processes, we introduce the notion of a **Lipschitz process**.

## Definition 3.1. (Lipschitz process)

A random process  $\{X_t\}_{t\in\mathcal{T}}$  is **Lipschitz** for a metric d on  $\mathcal{T}$  if there exists a random variable C such that

$$|X_s - X_t| \le Cd(s,t)$$
 for all  $s,t \in T$ .

**Approach:** Approximate a Lipschitz process by a **finite** set N, and then estimate N using the inequalities defined in the previous section.

## Definition 3.2. ( $\epsilon$ -net and covering number)

A set N is called an  $\epsilon$ -net for (T,d) if for every  $t\in T$ , there exists  $\pi(t)\in N$  such that  $d(t,\pi(t))\leq \epsilon$ . The smallest cardinality of an  $\epsilon$ -net for (T,d) is called the **covering number** 

$$N(T, d, \epsilon) := \inf \{ |N| : N \text{ is an } \epsilon\text{-net for } (T, d) \}.$$

# Lipschitz Maximal Inequality

## Lemma 3.3. (Lipschitz maximal inequality)

Let  $\{X_t\}_{t\in T}$  be a Lipschitz random process. Suppose that  $X_t$  is  $\sigma^2$ -subgaussian for each  $t\in T$ . Then

$$\mathbb{E}\left[\sup_{t\in\mathcal{T}}X_{t}\right]\leq\inf_{\epsilon>0}\Big\{\epsilon\mathbb{E}[C]+\sqrt{2\sigma^{2}\log\mathcal{N}(\mathcal{T},d,\epsilon)}\Big\}.$$

### Proof.

Fix  $\epsilon \in \mathbb{R}_{>0}$ . Choose an  $\epsilon$ -net N satisfying  $|N| = N(T, d, \epsilon)$  and perform the following decomposition:

$$\sup_{t \in T} X_t \leq \sup_{t \in T} \left\{ X_t - X_{\pi(t)} \right\} + \sup_{t \in T} X_{\pi(t)} \leq C\epsilon + \max_{t \in N} X_t.$$

Taking expectation and applying the maximal inequality gives

$$\mathbb{E}\left[\sup_{t\in\mathcal{T}}X_{t}\right]\leq\inf_{\epsilon>0}\left\{\epsilon\mathbb{E}[C]+\sqrt{2\sigma^{2}\log\mathcal{N}(\mathcal{T},d,\epsilon)}\right\}.$$

Let  $X_1, X_2, \cdots$  be i.i.d random variables taking values in the interval [0,1] and let  $f:[0,1]\to\mathbb{R}$  be a bounded function. By the law of large numbers,

$$\mathbb{E}\left[\left|\sum_{i=1}^n \frac{f(X_i)}{n} - \mu_f\right|\right] \lesssim n^{-1/2}$$

where  $\mu_f := \mathbb{E}[f(X_1)]$ .

Question: What is the optimal bound that also is uniform in f?

Problem setting:

$$X_f:=\sum_{i=1}^n rac{f(X_i)}{n}-\mu_f$$
 and  $\mathcal{F}:=ig\{f\in \mathsf{Lip}([0,1]): 0\leq f\leq 1ig\}.$ 

Goal: Establish an upper bound for the quantity

$$\mathbb{E}\left[\sup_{f\in\mathcal{F}}X_f\right].$$

### Two preliminary observations:

- **①** The process  $\{X_f\}_{f\in\mathcal{F}}$  is Lipschitz (with Lipschitz constant 2) with respect to the supremum norm on  $\mathcal{F}$ .
- 2 Each  $X_f$  is  $\frac{1}{n}$ -subgaussian.

By the Lipschitz maximal inequality,

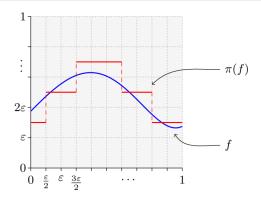
$$\mathbb{E}\left[\sup_{f\in\mathcal{F}}X_{f}\right]\leq\inf_{\epsilon>0}\left\{2\epsilon+\sqrt{\frac{2}{n}\log N(\mathcal{F},\|\cdot\|_{\infty},\epsilon)}\right\}.$$

It remains to find a good bound on the covering number.

#### Lemma 3.8.

There exists a constant  $c \in \mathbb{R}$  such that

$$N(\mathcal{F}, \|\cdot\|_{\infty}, \epsilon) \leq \mathrm{e}^{c/\epsilon} ext{ for } \epsilon < rac{1}{2}, \qquad N(\mathcal{F}, \|\cdot\|_{\infty}, \epsilon) = 1 ext{ for } \epsilon \geq rac{1}{2}.$$



Source: R. van Handel (2016, p.127)

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The equation

$$\mathbb{E}\left[\sup_{f\in\mathcal{F}}X_f\right]\leq\inf_{\epsilon>0}\left\{2\epsilon+\sqrt{\frac{2}{n}\log N(\mathcal{F},\|\cdot\|_{\infty},\epsilon)}\right\}.$$

reduces to

$$\mathbb{E}\left[\sup_{f\in\mathcal{F}}X_f\right]\leq\inf_{\epsilon>0}\left\{2\epsilon+\sqrt{\frac{2c}{\epsilon n}}\right\}\lesssim n^{-1/3}$$

which is unfortunately not **sharp**. In the second part of the presentation, we will see how sharper bounds can be obtained once we have improved our tool further.