

Union-find: (Disjoint Set Union)

If we have a set of N elements which are partitioned into further subsets, and if we have to keep track of connectivity of each element in a particular subset or connectivity of subsets with each^{oth}er, we use Union-find.

Union(A, B) - Connect two elements A and B

Find(A, B) - find, if there is any path connecting two elements A and B.

Example:

Ans

0	1	2	3	4	5	6	7	8	9
0	1	2	3	4	5	6	7	8	9

Union(2, 1)

Ans

0	1	1	3	4	5	6	7	8	9
0	1	2	3	4	5	6	7	8	9

Union(4, 3)

Union(8, 4)

Union(9, 3)

Union(6, 5)

Ans

0	1	1	3	3	5	5	3	3	3
0	1	2	3	4	5	6	7	8	9

We now have 5 subsets.

{0}

{1, 2}

{3, 4, 8, 9}

{5, 6}

{7}

Each subset is a connected component.

find(0, 7) → false

find(8, 9) → true

We have,

Arr	0	1	1	3	3	5	5	7	3	3
	0	1	2	3	4	5	6	7	8	9

Union(5, 2)

Arr	0	1	1	3	3	1	1	7	3	3
	0	1	2	3	4	5	6	7	8	9

Code Snippets:

```
int find (int arr a[], int a, int b)
{
    if (arr[a] == arr[b])
        return 1;
    else
        return 0;
}
```



```

void union (int arr[], int n, int a, int b)
{
    int temp = arr[a];
    int i;
    for (i = 0; i < n; i++)
    {
        if (arr[i] == temp)
            arr[i] = arr[b];
    }
}

```

* Kruskal's Algorithm *

- Used to find the minimum spanning tree

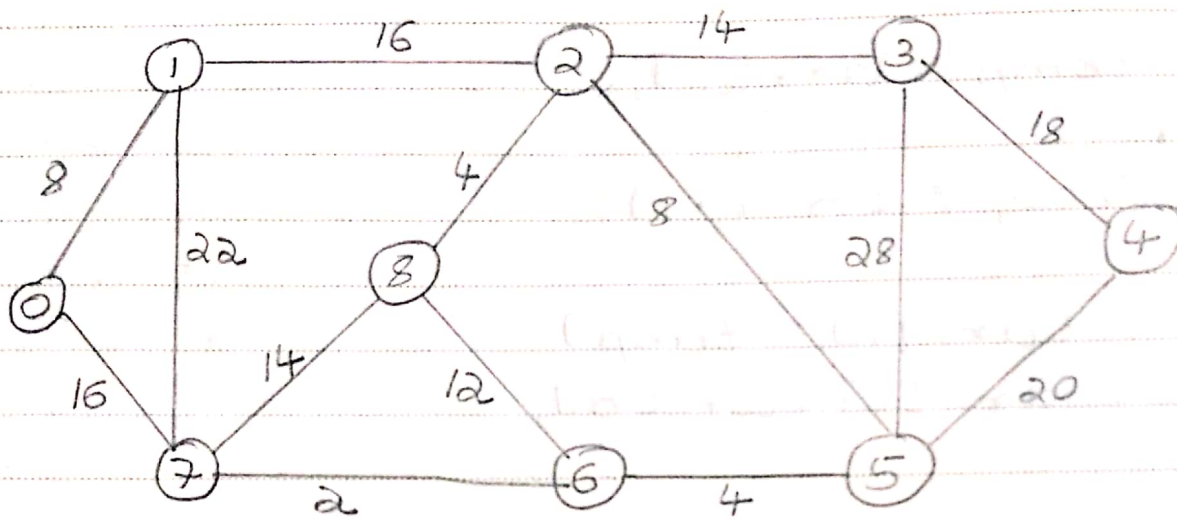
Method :

- Sort all the edges
- Keep adding edges one by one until
 - There is no cycle
 - all the nodes are connected

Every spanning tree of n vertices will have $n-1$ edges.

Alternate: Prim's.

Example:

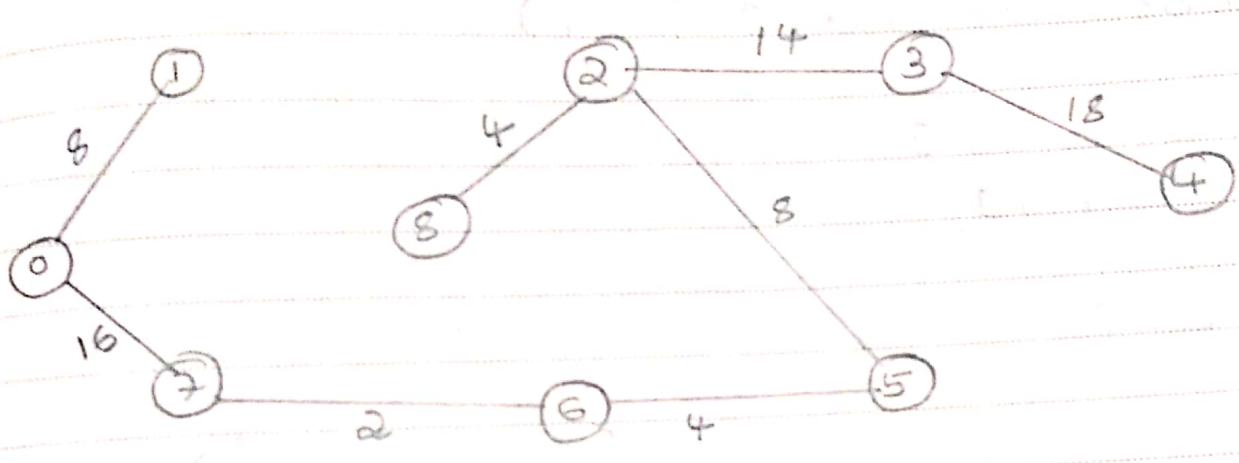


Sorted Edges:

$(6,7) \rightarrow 2$	$(2,8) \rightarrow 4$	$(5,6) \rightarrow 4$
$(0,1) \rightarrow 8$	$(2,5) \rightarrow 8$	$(6,8) \rightarrow 12$
$(7,8) \rightarrow 14$	$(0,7) \rightarrow 16$	$(1,2) \rightarrow 16$
$(4,5) \rightarrow 20$	$(1,7) \rightarrow 22$	$(3,5) \rightarrow 28$
		$(2,3) \rightarrow 14$
		$(3,4) \rightarrow 18$

Steps	(u,v)	i = find(u) j = find(v)	O/p (u,v)	Union(i,j) arr _j								
				0	1	2	3	4	5	6	7	8
init	0,7	6,7		0	1	2	3	4	5	6	7	8
1	(6,7)	6, 7	(6,7)	0	1	2	3	4	5	7	7	8
2	(2,8)	2, 8	(2,8)	0	1	8	3	4	5	7	7	8
3	(5,6)	5, 7	(5,6)	0	1	8	3	4	7	7	7	8
4	(0,1)	0, 1	(0,1)	1	1	8	3	4	7	7	7	8
5	(2,5)	8 7	(2,5)	1	1	7	3	4	7	7	7	7
6	(6,8)	7 7	discard									
7	(2,3)	7 3	(2,3)	1	1	3	3	4	3	3	3	3
8	(7,8)	3 3	discard									
9	(0,7)	1 3	(0,7)	3	3	3	3	4	3	3	3	3
10	(1,2)	3 3	discard									
11	(3,4)	3 4		4	4	4	4	4	4	4	4	4

Spanning Tree:



Minimum Cost = $8 + 16 + 2 + 4 + 8 + 4 + 14 + 18$
 $= 74.$

* Union-find : Root Method *

Arr

0	1	2	3	4	5
0	1	2	3	4	5

Union(1, 0)
 ↙ root of 0
 ↘ root of 1
 ↖ copy

Union(0, 2)

Union(1, 4)

Arr

0	1	2	3	4	5
2	0	4	3	4	5

Union(3, 4)

Code Snippets:

```

int root (int arr[], int i)
{
    while (arr[i] != i)
        i = arr[i];

    return i;
}

```

```

int find (int u, int v, int arr[])
{
    if ((root (arr, u)) == (root (arr, v)))
        return 1;
    else
        return 0;
}

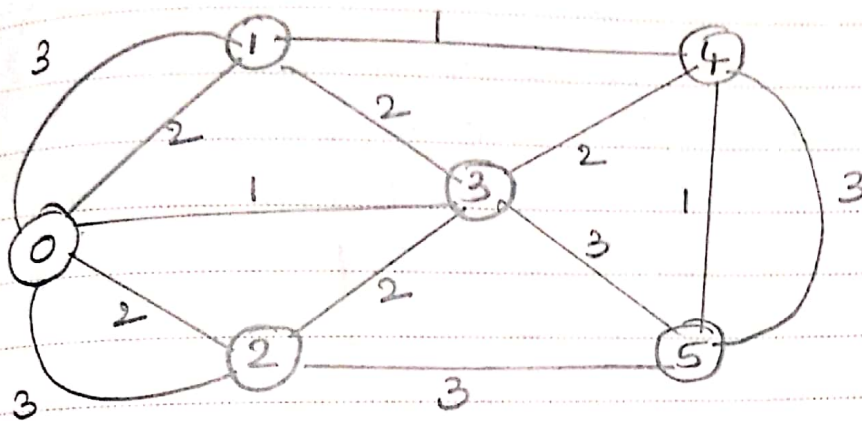
```

```

int union (int arr[], int u, int v)
{
    int rootu = root (arr, u)
    int rootv = root (arr, v)
    arr[rootu] = rootv;
}

```

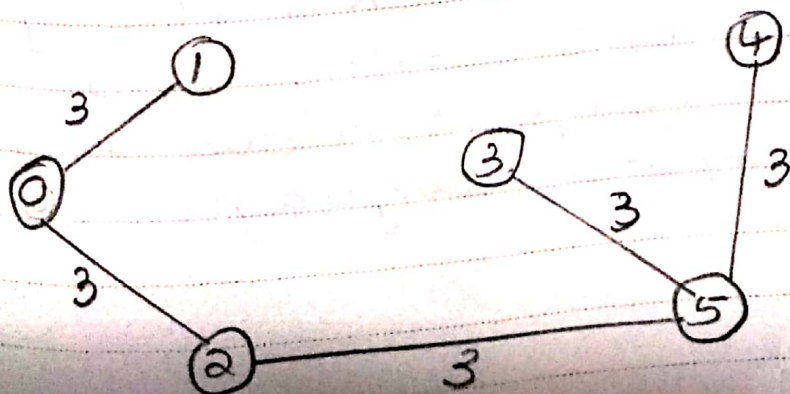

Kruskal's Example: Compute MAX Spanning Tree.



Sorted Edges: (descending)

(2,5) \rightarrow 3
 (0,1) \rightarrow 3 (0,2) \rightarrow 3 (3,5) \rightarrow 3 (4,5) \rightarrow 3
 (0,1) \rightarrow 2 (0,2) \rightarrow 2 (1,3) \rightarrow 2 (2,3) \rightarrow 2
 (3,4) \rightarrow 2 (0,3) \rightarrow 1 (1,4) \rightarrow 1 (4,5) \rightarrow 1

Steps	(u,v)	$i = \text{find}(u)$ $j = \text{find}(v)$	Output	Union(i,j)
init				0 1 2 3 4 5
1	(0,1)	0 1	(0,1)	0 1 2 3 4 5
2	(0,2)	1 2	(0,2)	1 1 2 3 4 5
3	(2,5)	2 5	(2,5)	2 2 2 3 4 5
4	(3,5)	3 5	(3,5)	5 5 5 3 4 5
5	(4,5)	4 5	(4,5)	5 5 5 5 4 5



Tree Cost
= 15.