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## Algorithms

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# 18. Binary Search

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# Binary Search

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- ⑱ m = sqrt2 (Binary search) 100%

2. In infinite array 100%

2. Works

3. What is the worst case time complexity?

"Simple" approach works well

Problem: no fibonaci in 21 months from 74 days + 74

Output: output nothing as needed previous flag is set

2. (returning to main)

3. Using radio button with

4. Using checkboxes instead of radio buttons

### ① Binary Search :

```

int search (vector<int> &nums, int target) {
    int start = 0;
    int end = nums.size() - 1;
    while (start <= end) {
        int mid = start + (end - start) / 2;
        if (target == nums[mid]) return mid;
        else if (target < nums[mid]) end = mid - 1;
        else start = mid + 1;
    }
    return -1; // If target not found return -1.
}

```

$T.C: O(\log n)$  /  $S.C: O(1)$

### ② Binary Search on reverse sorted array :

```

while (start <= end) {
    int mid = start + (end - start) / 2;
    if (target == nums[mid]) return mid;
    else if (target < nums[mid]) start = mid + 1;
    else end = mid - 1;
}
return -1;

```

$T.C: O(\log n)$

### ③ Order not known BS or, Agnostic BS:-

Let's use have array 'nums'

First check if array element is in ascending or, in descending  
then apply Binary search as previous two question

if ( $\text{nums}[0] > \text{nums}[1]$ ) {

Apply descending order Binary search

Set step

Apply ascending order Binary search

}

(4) First and last occurrence of an element.

Ex:- Input:  $A[ ] = [-1, 1, 2, 2, 2, 5, 6, 12, 12]$ , target = 2

Output: 2, 4  
 ↑  
 first occurrence  
 last occurrence

Code:- vector<int> searchRange (vector<int> &nums, int target)

vector<int> res;

int start = 0, first = -1, last = -1, end = nums.size() - 1;

while (start <= end) {

int mid = start + (end - start) / 2;

if (target == nums[mid]) {

first = mid;

end = mid - 1;

} else if (target < nums[mid]) end = mid - 1;

else start = mid + 1;

start = 0, end = nums.size() - 1;

while (start <= end) {

int mid = start + (end - start) / 2;

if (target == nums[mid]) {

last = mid;

start = mid + 1;

} else if (target < nums[mid]) end = mid - 1;

else start = mid + 1;

res.push\_back(first);

res.push\_back(last);

return res;

(5) Count of an element in a sorted array.

Ex:- Input:  $A[ ] = [-1, 1, 2, 2, 2, 5, 6, 12, 12]$ , target = 2

Output: 3

for that ele.

Code:- in previous question after finding first & last answer will be equal to  $\boxed{(last - first + 1)}$  if  $(last - first > 0)$   
else return 0;

### ⑥ Find minimum in Rotated sorted array.

for example an array  $\text{nums} = [0, 1, 2, 3, 4, 5, 6, 7]$  might become

$[4, 5, 6, 7, 0, 1, 2]$  or  $[0, 1, 2, 4, 5, 6, 7] \rightarrow$  unique element  
 $\hookdownarrow$  rotated by 4 times

so, in input we have given, rotated sorted array, and we have to find minimum element.

ex:- input  $\text{nums} = [3, 4, 5, 1, 2]$

O/P: 1

Code:-

```
int findMin (vector<int> & nums) {
    int n = nums.size();
    int start = 0, end = n - 1;
    if (n == 1) return nums[0];
    while (start <= end) {
        int mid = start + (end - start) / 2;
        int prev = (mid + n - 1) % n;
        int next = (mid + 1) % n;
        if (nums[mid] < nums[prev] && nums[mid] < nums[next])
            return nums[mid];
        else if (nums[end] < nums[mid]) start = mid + 1;
        else end = mid - 1;
    }
    return -1;
}
```

### ⑦ Search in Rotated sorted array

ex:-  $\text{nums} = [4, 5, 6, 7, 0, 1, 2]$ , target = 6

O/P:- 2 (index)

⑧  $\text{nums} = [4, 5, 6, 7, 0, 1, 2]$ , target = 3

O/P:- -1

Note:- all elements are distinct

Code:-

```

int search (vector<int> & nums, int target) {
    int n = nums.size();
    int i=0, j=n-1;
    while (i < j) {
        int mid = i + (j-i)/2;
        if (nums[mid] == target) return mid;
        if (nums[i] <= nums[mid]) {
            if (target >= nums[i] && target <= nums[mid])
                j = mid - 1;
            else i = mid + 1;
        } else {
            if (target >= nums[mid] && target <= nums[j])
                i = mid + 1;
            else j = mid - 1;
        }
    }
    return -1;
}

```

### (8) Searching in Nearly Sorted array

Given a sorted array arr[] of size N, some elements of arr are moved to either of the adjacent positions, i.e arr[i] may be present at arr[i+1] or arr[i-1]. i.e arr[i] can only be swapped with either arr[i+1] or arr[i-1]. The task is to search for an element in this array. If element not present then return -1, else its index.

Ex:- arr[] = {10, 3, 40, 20, 50, 80, 70}; key = 40,

O/P: 2

Code:-

```

int search (int arr[], int key) {
    int start=0, end=arr.size()-1;
    while (start <= end) {
        int mid = start + (end-start)/2;
        if (arr[mid] == key) return mid;
        if (mid > start && arr[mid-1] == key) return mid-1;
        if (mid < end && arr[mid+1] == key) return mid+1;
        if (arr[mid] > key) end = mid-2;
        else start = mid+2;
    }
    return -1;
}

```

⑨ find floor of an element in sorted array:

Given a sorted array and a value 'x'; the floor of 'x' is the largest element in the array smaller than or equal to x. Find the floor of x.

ex:- ① I/P: arr[] = {1, 4, 8, 10, 10, 12, 19}, x=5  
O/P: 2

② I/P: arr[] = {1, 4, 8, 10, 10, 12, 19}, x=20

Code: O/P: 19

```
int findFloor(vector<int> arr, int x) {
    int start=0, end=arr.size()-1, res=-1;
    while (start <= end) {
        int mid = start + (end - start)/2;
        if (x == arr[mid]) return mid;
        else if (x < arr[mid]) end = mid-1;
        else {
            res = arr[mid];
            start = mid+1;
        }
    }
    return res;
}
```

⑩ find ceil of an element in sorted array:

Ceil value of 'x' is the smallest element in the array larger than or equal to 'x'.

ex:- ① I/P: arr[] = {1, 4, 8, 10, 10, 12, 19}, x=5, O/P: 8

Code:-

```
while (start <= end) {
    int mid = start + (end - start)/2;
    if (x == arr[mid]) return mid;
    else if (x > arr[mid]) start = mid+1;
    else {
        res = arr[mid];
        end = mid-1;
    }
}
return res;
```

(11.) Next Alphabetical element :-

I/P:- letters = [ "D", "J", "K" ], target = "B", O/P:- "D"

Opp :- leftExp = ["h", "n", "5"], target = "t", opp:- "h"  
Code :- (letter also wrap around)

char nextGreaterLetter (vector<char> letters, char target) {

```
int start=0, end = letters.size() - 1; res = letters[start];
```

while (start < end) do

```
int mid = start + (end - start) / 2;
```

here we don't need to return matched if (target == letters[mid]) start = mid + 1;

else if (target < letters[mid]) {

res = letters [mid];

{ end = mid - 1;

else

~~Start = mid + 1;~~

```
return res;
```

Find the pos of an element in infinite sorted array:

```
int binarySearch (int arr[], int l, int r, int x) {
```

while ( $k < \delta$ ) {

int mid = l + (r - 1) / 2;

if ( $\text{arr}[\text{mid}] \neq x$ ) return mid;

if  $(\text{dim } M) = n$  then  $M$  is  
is it true that  $\text{dim } M = n$

IF ( $\text{abs}(\text{mid}) > x$ )  
    mid = mid \* -1

else l=mid+1;

5 return -1;

int FindPos(int arr[], int key) of zip search - O(n)

int d=0, h=1;

~~int val = arr[0];~~

while ( var <key> ). {

$$d = h_3 \cdot t - k(m - b) \leq$$

$$h = 2 * h;$$

1. Strong 2.

return binarySearch(first, last, key);

~~T C~~ bridge + logn  
~~TC~~ = logn

(13) Find the index of 1st 1's in a Binary sorted array

Code:-

```
int FirstIndex (int a[], int n) {
    int start = 0;
    int end = n - 1;
    while (start <= end) {
        int mid = start + (end - start) / 2;
        if (a[mid] == 1) {
            if (a[mid - 1] == 0)
                end = mid - 1;
            else
                return mid;
        }
        else
            start = mid + 1;
    }
    return -1;
}
```

(14) Find Peak Element

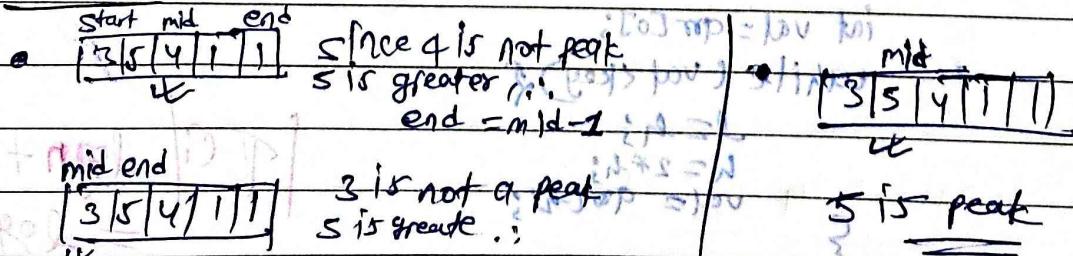
Given an array, find a peak element (print any one, if many are found). A peak element is one such that it is either greater than or equal to its neighbours. For the first & last element, it is enough to look at its only one neighbour.

Ex: ① I/P: arr = {3, 5, 4, 2, 1}, O/P: 5

② I/P: arr = {2, 6, 3, 7, 8, 9}, O/P: 6

Note:- here 9 is also a peak element.

Approach:-



Code :-

```

int findPeakElement (int arr[], int n) {
    int start=0, end=n-1;
    while (start<end) {
        int mid = (start+end)/2;
        if (mid==0)
            return arr[0]>arr[1] ? arr[0] : arr[1];
        if (mid==n-1)
            return arr[n-1]>arr[n-2] ? arr[n-1] : arr[n-2];
        if (arr[mid]>=arr[mid-1] && arr[mid]>=arr[mid+1])
            return arr[mid];
        if (arr[mid]<arr[mid-1]) end=mid-1;
        else start=mid+1;
    }
    return arr[start];
}

```

(5)

Find maxm element in Bitonic array.

Bitonic array :- monotonic  $\uparrow \rightarrow$  monotonic  $\downarrow$ 

monotonic  $\uparrow$  array is increasing sequence  
 in which  $arr[i] \neq arr[i+1]$ , and vice versa

ex:-  $arr[]: [1|3|8|12|4|2]$ , 0|1|: 12  
 increasing      decreasing

Code :-

Bitonic array maxm = Find Peak Element

so, this problem is same as find peak element.

Note:- here only one peak element are present in

Bitonic array and in previous question there can be more than one peak element

(16) Search an element in Bitonic array.

Ex:- arr[] =  $\begin{bmatrix} 0 & 1 & 2 & 3 & 4 & 5 \\ 2 & 3 & 8 & 12 & 4 & 2 \end{bmatrix}$ , key = 4

O/P:- 4.

Approach:-

① Using Find peak element / maxm element in Bitonic array

Find maxm element

② check if maxm element is not the target element

then A apply binary search on left side of maxm element

B apply reverse BS on right side of maxm element

(17) N<sup>th</sup> root of a number using binary search

- Given two no. N and m find the N<sup>th</sup> root of m.

- N<sup>th</sup> root of m is defined as a no. X, when raised to power N equals m i.e.  $X^N = M$  or,  $X = \sqrt[N]{M}$

Ex:- N=3, m=27

O/P:- 3

Ex:- N=2, m=16

O/P:- 4

Ex:- N=5, m=243

O/P:- 3

Code:-

double multiply(double number, int n) of

double ans = 1.0;

for (int i=1; i<=n; i++) {

    ans = ans \* number;

return ans;

double getNthRoot (int n, int m) of

double low = 1, high = m, eps = 1e-7;

n=2, m=16

mid = 8

$8^2 = 64 > m$

$\therefore$  high = 8

now mid = 5, compare

$5^2 = 25 < 16$

$\therefore$  high = 5

low, mid = 3.5

while ((high - low) > eps) of

    double qmid = (low + high) / 2.0;

    if (multiply(mid, n) < m) low = mid;

    else high = mid;

    // some steps here

    // some steps here

return low;

T.C:  $N \times \log(M \times 10^d)$