Name: -Pratik suntl Garkward
TE-A-39
A.I

* Assignment NO-03-x

Explain the Min-max and Alpha beta pruning algorithm
For adversal search with an Example.

If The Min-Max algorithm computes minimax decision From the current state.

21 14 13 used as sevening technique in game problems.

31 The Min-max algorithm complete dipth- 1/1st Explosur-

* augosthim:

1) Sturt Noce 15 Mux plade with current bound condiquation.

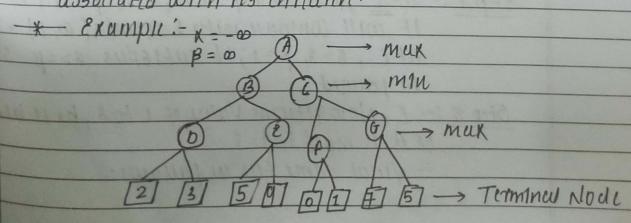
2) Expand Nodes down to some depth of look ahead in the game.

31 Apply Evaluation function at Each of lear Nodes.

H Buck up Values dor tuch non-leaf nodes until computed

6) At Min Nocks, the bucked up Value is minimum of Values associated with children.

6 At Mux Nodes, backed up Value Is muximum of Values
associated with its children.

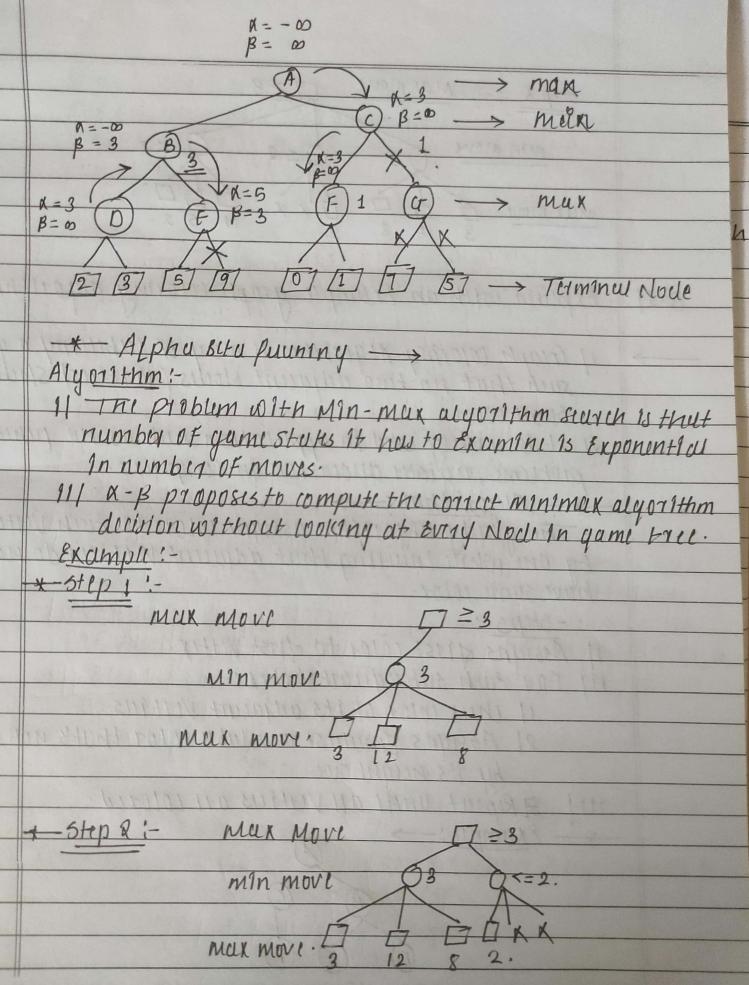


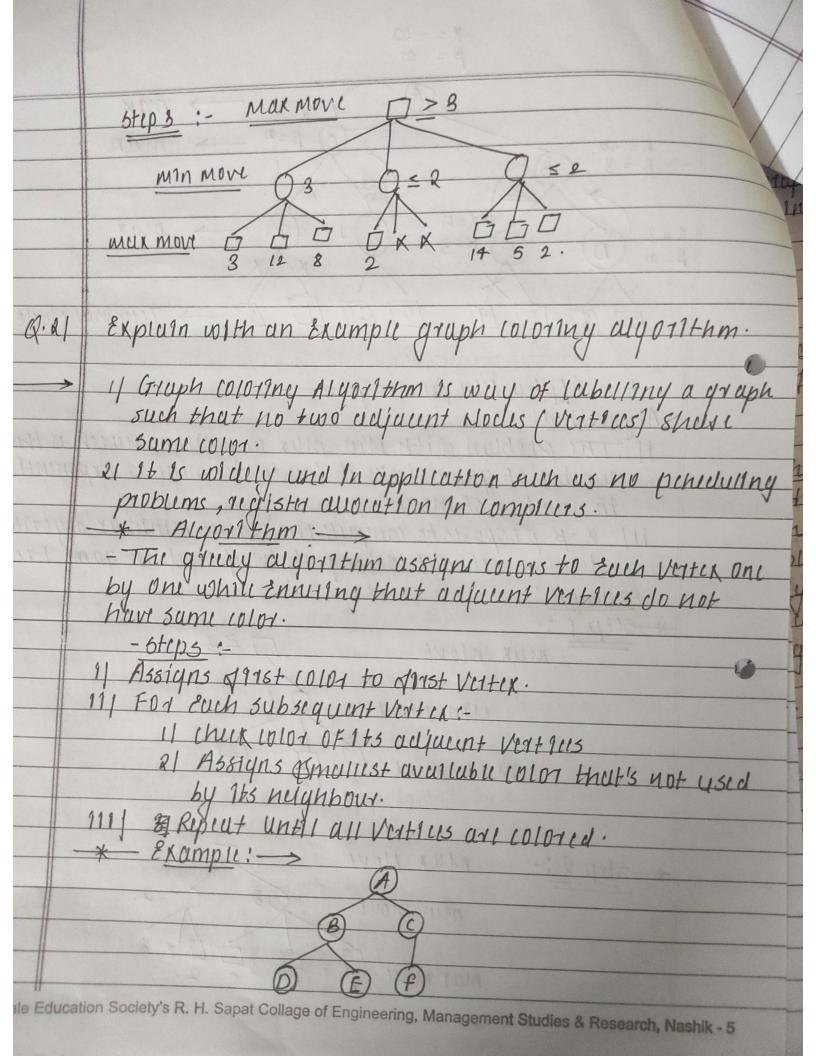
Step1: - At Alist step, max plujer will start first move From Noch A where A = -0, B = 0. Hese Value puss to B where again a = - 00, B=+00, B pusses Stepi: - At Node D, Value of & will be colleworld formax Value of x 15 compared with 2 and 3 (2,3) = 3 will be Value & at node D 4 Node Value will also 3. steps: - Now Alyo, bucktruck to Node B, Value of & will change as 1ts turn of min, B= + 0; compage with available Node Value; min (0,3) = 3, hence at Node B and now x=-0, B=3 Step 45- At Node E, max will taxe its turn, and value of x will change. current Value of x will compared with 5 30 max (-0, 51=5; hinu at Node €, x=5, β=3; while x>=B, Now Valu of E will be 5 5tp 5:- At Next step, alyorithm bucktive til, At Noch C, N = 8, B = +0, sumi Values will puned to 5+1p6:- At F, a yuin compute with lift (hild; max (32)=3
and then compare with right child which is I; max (3,1)=3 NOW NUCL OF F WILL be 1 Step 7: at [; x=3 and B= +0, here & will be changed. 1+ will compute with 1 50 min(0,1)=1 Now clt C, N=3, B=1; 1+ SU+18fils A>=B 50, C+ W111 Step 8: - C, Now yeturns Value of 1 to A, here best value

ation Society's R. H. Sapat Collage of Fnging

OF A 15 MUK(3,1)=3

- final game true as fallaros:





Stips: - Sturt with A -> Absign colors. Step 1:- Move to Band c -> Assign color & (Adjuly to A)

Step 3:- Move to D, E, F -> Assign color 1 (Adjuly to A) 5+1p4 !color 1 101012 Define constraint surasfuction Problem. Stuti types of Q. 81 consistencies. 1/ A constraint satisfaction problem has various states and goal test, a traditional problem has converted into Stundard structure and very simple representation. es A construint sotisfaction problem is defined by set OF Vortables Kigka. . Xn and set of construents Cycla, in 3) Euch Variable X, has non Empty domain Di of possible Values. H A State OF Problem is defined by an assignment of Values to some or all of Variables | XI=VI, Xj=Vj... | 5/ An assignment that does not voilate any constraints is called consistent. 6/ A compute assignment to one in which every variable is mountained and solution to csp. * Types of constraints :-> (onsistentils

I Node consistency: 1 - A Variable is hode - consistent if all values in 173 domuin satisfy 1ts unusy constraints. 2- Example: Top x can only have Even xlumbers, but Its domain 15 \$ 1, x, 3, 4 4 removing make Its noch consistent 11) Are consistincy: 1) A variable 13 are-consistent with another variable if Every Value In 175 domuin hus Valle consusponding Value In other Variable's domain. 2) - Example: - 9 F x = 3 mon, Tuey and y = 2 mon, wed y, but construint suys x = y then remove mon from x make 9+ UTL-constent 111/ Path consistency: 1) Extends A11-consistency by Enriving pary of Vortubles Suffisty construint relative to 3rd variable. 2) - Example: 9 F K, y, Z are Vatiables, 4 K and \$ 2 are consistent through y 1V/ K - constency: DA 13P 15 K- consistent of Every valled assignment to R-1 Vayrubles can be Extended to Kth Vutiable. 2) - 2 - consistency = A11 consistency. - 3 - constency = Path constency. - strong K - Consistency = Problem 1s Jos dujsk. Examples of Csp: 11 Sudoku Mup Coloring.

AH	solve Landwing Crypt Arithmetic Problem.
→	$ \begin{array}{cccccccccccccccccccccccccccccccccccc$
•	$\begin{array}{c} 0 \rightarrow 0 \\ 1 \rightarrow M \end{array}$
	$2 \rightarrow 1$ $3 \rightarrow X$
	4 -> K
	$\begin{array}{c} 5 \rightarrow E \\ 6 \rightarrow N \end{array}$
	$\begin{array}{c} 1 \rightarrow 0 \\ 8 \rightarrow R \end{array}$
	$q \rightarrow 5$
)	MORE = 1085
	MONEY = 10652
	$\frac{1}{10000000000000000000000000000000000$
	· 图1000 · 1000