

# Statistics 1 Unit 2

Group 8

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## Contents

<b>1</b>	<b>Task 15</b>	<b>1</b>
1.1	a) . . . . .	1
1.2	b) . . . . .	1
<b>2</b>	<b>Task 16</b>	<b>2</b>
2.1	a) . . . . .	2
2.2	b) . . . . .	2
<b>3</b>	<b>Task 17</b>	<b>2</b>
<b>4</b>	<b>Task 18</b>	<b>3</b>
<b>5</b>	<b>Task 19</b>	<b>3</b>
<b>6</b>	<b>Task 20</b>	<b>4</b>
<b>7</b>	<b>Task 21</b>	<b>4</b>
7.1	Solve . . . . .	5
7.2	QR.Solve . . . . .	5
7.3	SVD . . . . .	5
7.4	QR + rearrangement . . . . .	5
<b>8</b>	<b>Task 23</b>	<b>6</b>
8.1	a . . . . .	6
8.2	b . . . . .	6
8.3	c . . . . .	6
<b>9</b>	<b>Task 26</b>	<b>6</b>

10 Task 28	7
11 Task 31	8
12 Task 32	9
13 Task 33	10
14 Task 34	11
15 Task 35	12
16 Task 36	12
17 Task 37	13
18 Task 38	13
19 Task 39	14

## 1 Task 15

### 1.1 a)

From Sylvester's Criterion it is clear that  $\alpha$  being the first leading principal minor must be positive.

Since matrix  $B$  is a positive definite square matrix,  $x'Bx > 0$  for all  $x$ . Let's take  $x = (0, x_1 \dots x_n)$ . In this case,  $x'Bx = y'Ay$ , where  $y = (x_1 \dots x_n)$  and it must be  $> 0$ . Since  $y'Ay > 0$  for arbitrary  $y = (x_1 \dots x_n)$ ,  $A$  is positive definite by definition.

### 1.2 b)

$$C = \begin{pmatrix} D & C^T \\ C & E \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ CD^{-1} & 1 \end{pmatrix} \begin{pmatrix} D & 0 \\ 0 & S \end{pmatrix} \begin{pmatrix} 1 & D^{-1}C^T \\ 0 & 1 \end{pmatrix}$$

Substitution  $D = L_d L_d^T$  and  $S = L_s L_s^T$  gives us the Cholesky factorization:

$$B = \begin{pmatrix} D & C^T \\ C & E \end{pmatrix} = \begin{pmatrix} L_d & 0 \\ CL_d^{-1} & L_s \end{pmatrix} \begin{pmatrix} L_d^t & L_d^{-1}C^T \\ 0 & L_s^t \end{pmatrix} = L_b L_b^t$$

In our case:  $D=\alpha$ ,  $C=a$ ,  $E=A$ ,  $C^t = a^t$  then  $L_d = \sqrt{\alpha}$ ,  $L_d^{-1} = \frac{1}{\sqrt{\alpha}}$ ,  $L_d^t = \sqrt{\alpha}$ ,  $D^{-1} = \frac{1}{\alpha}$

$$S = A - \frac{1}{\alpha}aa^t$$

$$S = L_s L_s^t - \text{Choleski - Factorisation}$$

Hence,

$$B = \begin{pmatrix} \alpha & a^T \\ a & A \end{pmatrix} = \begin{pmatrix} \sqrt{\alpha} & 0 \\ \frac{1}{\alpha}a & L_s \end{pmatrix} \begin{pmatrix} \sqrt{\alpha} & \frac{1}{\alpha}a^T \\ 0 & L_s^T \end{pmatrix}$$

## 2 Task 16

### 2.1 a)

Here we use the same idea. Since matrix  $B$  is a positive definite square matrix,  $x'Bx > 0$  for all  $x$ . Let's take  $x = (x_1 \dots x_n, 0)$ . In this case,  $x'Bx = y' Ay$ , where  $y = (x_1 \dots x_n)$  and it must be  $> 0$ . Since  $y' Ay > 0$  for arbitrary  $y = (x_1 \dots x_n)$ ,  $A$  is positive definite by definition. For  $\alpha$  we can use  $x = (0 \dots 0, 1)$ . Applying the same idea, fact that  $\alpha > 0$  is easily proved.

### 2.2 b)

The idea of Choleski decomposition is absolutely the same as in 15(b). In this case:  $D=A$ ,  $C = a^t$ ,  $E=\alpha$ ,  $C^t = a$

This gives us:

$$B = \begin{pmatrix} A & a \\ a^T & \alpha \end{pmatrix} = \begin{pmatrix} L_A & 0 \\ a^t L_A^{-1} & \sqrt{\alpha - a^t A^{-1} a} \end{pmatrix} \begin{pmatrix} L_A^T & L_A^{-1} \\ 0 & \sqrt{\alpha - a^t A^{-1} a} \end{pmatrix}$$

## 3 Task 17

```
> f17 <- function(k) {
+ mat_A <- matrix(c(10^(-2*k),1,1,1), 2, 2)
+ mat_M <- matrix(c(1, -1/(10^(-2*k))), 0,1), 2, 2)
+ vec_B <- matrix(c(1+(10^(-2*k))), 2), 2, 1)
+ S <- backsolve((mat_M %*% mat_A), (mat_M %*% vec_B))
+ S
}
```

```

+ }
> a<-seq(1:10)
> sapply(a, f17)

      [,1] [,2] [,3] [,4] [,5]      [,6]      [,7]      [,8] [,9] [,10]
[1,]      1      1      1      1      1 0.9998669 0.9992007 2.220446      0      0
[2,]      1      1      1      1      1 1.0000000 1.0000000 1.000000      1      1

```

When,  $\epsilon$  decreases ( $k > 6$ ) we are getting greater error.

## 4 Task 18

```

> matrix_2norm <- function(x){
+   norm(x, type="2")
+ }
> x <- matrix(c(1,2,3,4),2,2)
> max(svd(x)$d)

```

```
[1] 5.464986
```

```
> matrix_2norm(x)
```

```
[1] 5.464986
```

It is clear that 2-norm is just the maximum value of the singular value decomposition of the matrix.

## 5 Task 19

```

> cond_num <- function(p){
+   r <- matrix_2norm(p)
+   s <- matrix_2norm(solve(p))
+   s * r
+ }
> max(svd(x)$d) / min(svd(x)$d)

```

```
[1] 14.93303
```

```
> cond_num(x)
```

```
[1] 14.93303
```

We observe that the condition number of a matrix is the quotient of the lowest and the highest value of the SVD.

This follows from:  $M = U\Sigma V'$  and  $M^{-1} = V\Sigma^{-1}U'$ . The values of the diagonal matrix  $\Sigma^{-1}$  are just the reciprocals of the diagonal values of  $\Sigma$ . Since the 2-norm is the maximum value of the SVD, the 2-norm of  $\Sigma^{-1}$  has to be 1 over the smallest singular value of the matrix.

## 6 Task 20

$Ax=b$

$$\Delta x = A^{-1}\Delta b$$

$$\|b\| = \|Ax\| \|x\|$$

Using norm's property:

$$\|Ax\| \leq \|A\| \|x\|$$

$$\Rightarrow \|\Delta x\| = \|A^{-1}\Delta b\| \leq \|A^{-1}\| \|\Delta b\| \wedge \|b\| = \|Ax\| \leq \|A\| \|x\|$$

Multiply these two inequalities (the norm > 0)

$$\|\Delta x\| \|b\| \leq \|A^{-1}\| \|\Delta b\| \|A\| \|x\|$$

Divide by:  $\|b\| \|x\|$  and get

$$\frac{\|\Delta x\|}{\|x\|} \leq k(A) \frac{\|\Delta b\|}{\|b\|}$$

## 7 Task 21

```
> # Matrix is filled by columns by default
> A <- function(epsilon) {
+   matrix(data = c(1, 1-epsilon, 1+epsilon, 1), nrow = 2, ncol = 2)}
```

Now, given a small  $\epsilon$  we will try to solve the system. "normal" solve command, qr.solve, qr.decomposition "by hand" and the SVD method.

As suggested we will use square root of the machine precision as our  $\epsilon$ .

```
> epsilon <- unname(sqrt(unlist(.Machine[1])))
> epsilon
```

```
[1] 1.490116e-08
```

```
> b <- c(1+epsilon +epsilon^2, 1)
```

First, let's check condition number using task 19 approach:

```
> svdA <- svd(A(epsilon))
> cond.nr <- max(svdA$d) / min(svdA$d)
> cond.nr
```

```
[1] 2.547621e+16
```

As we can see the condition number is very big. This means that small changes of a parameter will dramatically change the solution.

## 7.1 Solve

```
> # solve(A(epsilon), b)
> # Error in solve.default(A(epsilon), b) :
> # system is computationally singular: reciprocal condition number = 5.55112e
```

## 7.2 QR.Solve

```
> # qr.solve(A(epsilon), b)
> # Error in qr.solve(A(epsilon), b) : singuläre Matrix 'a' in 'solve'
```

## 7.3 SVD

```
> # svdA <- svd(A(epsilon))
> # D <- diag(x=svdA$d)
> # D
> # a <- (tcrossprod(svdA$u %*% D, svdA$v))
> #solve(a,b)
> # Error in solve.default(a, b) :
> # system is computationally singular: reciprocal condition number = 2.77556e
```

## 7.4 QR + rearrangement

QR decomposition using rearrangement  $Rx = Q'b$  gives result, but with really big error:

```
> AQR <- qr(A(epsilon))
> qr.Q(AQR)

      [,1]      [,2]
[1,] -0.7071068 -0.7071068
[2,] -0.7071068  0.7071068

> backsolve(qr.R(AQR), crossprod(qr.Q(AQR),b))

      [,1]
[1,]  0.5
[2,]  0.5
```

Since the relative error is bounded by  $\mathcal{K}(A) \frac{\|\Delta b\|}{\|b\|}$ , larger  $\epsilon$  leads to smaller condition numbers  $\mathcal{K}(A)$  and hence more accurate solution.

## 8 Task 23

### 8.1 a

To prove the task just use the formula below:

$$Au = (uv^T)u = u(v^T u) = (v^T u)u.$$

### 8.2 b

The other eigenvalues of A. Since  $\text{rank}(A) = 1$  all other eigenvalues equals are zeros.

### 8.3 c

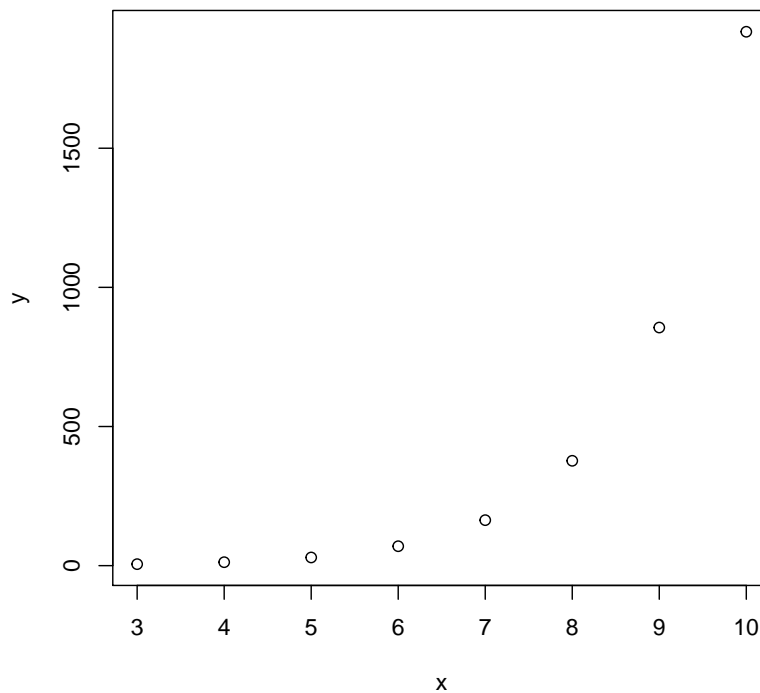
Let  $U \in C_{m \times m}$  be a unitary matrix so that  $Uu = \|u\|_2 e_1$ , and  $V \in C_{n \times n}$  be a unitary matrix so that  $Vv = \|v\|_2 e_1$ . We can substitute it into  $uv^*$  shows SVD.  $D = \|u\|_2 \|v\|_2 e_1 e_1^*$

## 9 Task 26

From task 19 follows that the ratio is a condition number of our matrix. To show the dynamic it is convenient to use a plot: x-axis the dimension, y-axis the condition number.

```
> f26 <- function(n) {
+   M <- mat.or.vec(n, n)
+   M[upper.tri(M)] <- -1
+   diag(M) <- 1
+   condition_number <- max(svd(M)$d) / min(svd(M)$d)
+   return (condition_number)
+ }
> x <- 3:10
> y <- c(rep(0, 8))
> for (i in (1:8)) {
+   y[i] <- f26(x[i])
+ }
>
```

```
> plot(x,y)
```



As we can see, condition number exponentially increases with increasing the dimension of matrix.

## 10 Task 28

Singular Valued Decomposition constructs orthonormal bases for the range and null space of a matrix. The columns of  $U$  which correspond to non-zero singular values of  $A$  are an orthonormal set of basis vectors for the range of  $A$ . The columns of  $V$  which correspond to zero singular values form an orthonormal basis for the null space of  $A$ .

```
> # The columns of output matrix contain orthonormal bases for the range of A
> get_bases_for_the_range <- function (A)
+ {
+   A_svd = svd(A)
+   +
+   d = A_svd$d
```



```

+   u = A_svd$u
+   v = A_svd$v
+
+   column_indices <- which(d >= .Machine$double.eps)
+
+   B = u[,column_indices]
+
+   return(B)
+ }
> # The columns of output matrix contain orthonormal bases for the null space of A
> get_bases_for_the_null_space <- function (A)
+ {
+   A_svd = svd(A)
+
+   d = A_svd$d
+   u = A_svd$u
+   v = A_svd$v
+
+   column_indices <- which(d < .Machine$double.eps)
+
+   B = v[,column_indices]
+
+   return(B)
+ }
> # Test
> A <- matrix (c(3,7,8,7,10,9,8,9,12, 15) , nrow = 5, ncol = 2)
> get_bases_for_the_range(A)

      [,1]      [,2]
[1,] -0.3131156  0.7326549483
[2,] -0.3605707 -0.4066120946
[3,] -0.4080145 -0.4879538258
[4,] -0.4744268  0.2445079999
[5,] -0.6167582  0.0004828064

> get_bases_for_the_null_space(A)

[1,]
[2,]

```

## 11 Task 31

Required function is represented below:

```
> library(matrixcalc)
> Duplication <- function (n) {
+   # Arbitrary matrix
+   A <- matrix((1:n ^ 2), n, n)
+   # Make it symmetric
+   A[lower.tri(A)] = t(A)[lower.tri(A)]
+   vec(A)
+   vech(A)
+   D <- mat.or.vec(length(vec(A)), length(vech(A)))
+   for (i in 1:n ^ 2) {
+     col_num <- match(vec(A), vech(A))
+     row_num <- 1:n ^ 2
+     D[row_num[i], col_num[i]] <- 1
+   }
+   return(D)
+ }
> Duplication(2)

      [,1] [,2] [,3]
[1,]     1     0     0
[2,]     0     1     0
[3,]     0     1     0
[4,]     0     0     1
```

## 12 Task 32

Let's compute Singular Value Decomposition of the duplication matrix  $D_n$  with  $n = 2$  and  $n = 3$ :

```
> svd(Duplication(2))

$d
[1] 1.414214 1.000000 1.000000

$u
      [,1] [,2] [,3]
[1,] 0.0000000 0 1
```

```
[2,] -0.7071068    0    0
[3,] -0.7071068    0    0
[4,]  0.0000000   -1    0
```

\$v

```
      [,1] [,2] [,3]
[1,]    0    0    1
[2,]   -1    0    0
[3,]    0   -1    0
```

```
> svd(Duplication(3))
```

\$d

```
[1] 1.414214 1.414214 1.414214 1.000000 1.000000 1.000000
```

\$u

```
      [,1]      [,2]      [,3] [,4] [,5] [,6]
[1,] 0.0000000 0.0000000 0.0000000    1    0    0
[2,] 0.0000000 -0.7071068 0.0000000    0    0    0
[3,] 0.0000000 0.0000000 -0.7071068    0    0    0
[4,] 0.0000000 -0.7071068 0.0000000    0    0    0
[5,] 0.0000000 0.0000000 0.0000000    0   -1    0
[6,] -0.7071068 0.0000000 0.0000000    0    0    0
[7,] 0.0000000 0.0000000 -0.7071068    0    0    0
[8,] -0.7071068 0.0000000 0.0000000    0    0    0
[9,] 0.0000000 0.0000000 0.0000000    0    0   -1
```

\$v

```
      [,1] [,2] [,3] [,4] [,5] [,6]
[1,]    0    0    0    1    0    0
[2,]    0   -1    0    0    0    0
[3,]    0    0   -1    0    0    0
[4,]    0    0    0    0   -1    0
[5,]   -1    0    0    0    0    0
[6,]    0    0    0    0    0   -1
```

It can be seen that if  $n = 2$ , then last 2 diagonal elements of matrix  $D$ , which are singular values, are equal to 1. If  $n = 3$ , then the last 3 diagonal elements of matrix  $D$  are equal to 1. Other elements are 1.414214 what is just the square root of 2.

## 13 Task 33

Function returning the elimination matrix  $L_n$  for given  $n$  and example of its application are as follows:

```
> library(matrixcalc)
> Elimination <- function (n) {
+ A <- matrix((1:n^2),n,n)
+ vec(A)
+ vech(A)
+ D <- mat.or.vec(length(vech(A)),length(vec(A)))
+ for (i in 1:n^2) {
+ col_num <-match(vech(A), vec(A))
+ row_num <- 1:(n*(n+1)/2)
+ D[row_num[i], col_num[i]] <- 1}
+ return(D)}
> Elimination(2)
```

```
      [,1] [,2] [,3] [,4]
[1,]     1     0     0     0
[2,]     0     1     0     0
[3,]     0     0     0     1
```

## 14 Task 34

Let's compute Singular Value Decomposition of the elimination matrix  $L_n$  for  $n = 2$ :

```
> svd(Elimination(2))
```

```
$d
[1] 1 1 1
```

```
$u
      [,1] [,2] [,3]
[1,]     1     0     0
[2,]     0     1     0
[3,]     0     0     1
```

```
$v
      [,1] [,2] [,3]
```

```

[1,]    1    0    0
[2,]    0    1    0
[3,]    0    0    0
[4,]    0    0    1

```

It can be seen that matrix  $U$  is an identity matrix and diagonal matrix  $D$  is an identity matrix too. In addition, the elimination matrix itself equals to  $V^t$  matrix.

## 15 Task 35

Function returning the commutation matrix  $K_{mn}$  for given  $m$  and  $n$  and example of its application are as follows:

```

> library(matrixcalc)
> commutation_matrix <- function (r, c = r)
+ {
+   H <- H.matrices(r, c)
+   p <- r * c
+   K <- matrix(0, nrow = p, ncol = p)
+   for (i in 1:r) {
+     for (j in 1:c) {
+       Hij <- H[[i]][[j]]
+       K <- K + (Hij %x% t(Hij))
+     }
+   }
+   return(K)
+ }
> k <- commutation_matrix(3,2)
> k
      [,1] [,2] [,3] [,4] [,5] [,6]
[1,]    1    0    0    0    0    0
[2,]    0    0    0    1    0    0
[3,]    0    1    0    0    0    0
[4,]    0    0    0    0    1    0
[5,]    0    0    1    0    0    0
[6,]    0    0    0    0    0    1

```

## 16 Task 36

The singular values of the commutation matrix are:

```
> svd(k)$d
```

```
[1] 1 1 1 1 1 1
```

## 17 Task 37

The Moore-Penrose inverse can be expressed in terms of SVD,  $A = UDV'$ , since  $A^+ = VD^{-1}U'$  where each element in  $D^{-1}$  is taken as a reciprocal of corresponding element in matrix  $D$ , if it is greater, than given tolerance, or 0 otherwise.

```
> f37 <- function(A, tol = 1e-10){
+   SVD <- svd(A)
+   D <- NULL
+   for(i in 1:length(SVD$d)){
+     if(SVD$d[i] < tol){
+       SVD$d[i] <- 0
+       D <- c(D, as.numeric(SVD$d[i]))
+     }
+     else{
+       D <- c(D, as.numeric(1/(SVD$d[i])))
+     }
+   }
+   SVD$v %*% diag(D)
+   A_plus <- crossprod(t(crossprod(t(SVD$v), diag(D))), t(SVD$u))
+   return(A_plus)
+ }
```

## 18 Task 38

Let's use the property that the trace is invariant under cyclic permutations when inner matrix is square.

```
> wcptrace <- function(A, w) {
+   if(dim(A)[1] == dim(A)[2]) {
+     a <- numeric(dim(A)[1])
+     for(i in 1:dim(A)[1]){
+       a[i] <- sum(A[i, ]^2)
+     }
+   }
```

```

+     return(sum(a * w))
+   } else {
+     cat("The input matrix is not square, therefore there is not much to improve")
+     sum(diag(t(A) %*% diag(w) %*% A))
+   }
+ }
> mat <- matrix(data = rexp(200, rate = 10), nrow = 100, ncol = 100)
> w <- sample(1:100)
> system.time( replicate(10000, wcptrace(mat, w)))

      user  system elapsed
      2.36   0.00   2.38

> system.time( replicate(10000, sum(diag(t(mat) %*% diag(w) %*% mat))))

      user  system elapsed
     15.16   0.00   15.24

```

## 19 Task 39

```

> #Using Cholesky Decomposition is more efficient for our task
>
> dmvnorm2 <- function (x, m, V) {
+   mat_chol <- chol(V)
+   delta <- x - m
+   y <- det(mat_chol)^2 * (2*pi)^nrow(mat_chol)
+   e <- exp(-t(delta) %*% solve(mat_chol) %*% t(solve(mat_chol)) %*% delta / 2)
+   diag(y^(-1/2) * e)
+ }
> #TEST
> library(mvtnorm)
> M <- matrix(c(1,0.5,0.5,0.5,1,0.5,0.5,0.5,1),nrow=3)
> dmvnorm(1:3, 1:3, M)

[1] 0.08979356

> #Our code test
> dmvnorm2(1:3, 1:3, M)

[1] 0.08979356

```