# 计算机体系架构 第八周作业

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作业内容: 7.1, 7.3, 7.6, 7.11, 7.25, 7.32, 7.36。

### Problem 7.1

如果不考虑成本的话,更大的 SRAM 会占据更大额度片上面积,进而影响其他的模块,并且缺少 RAM 接口,难以兼容以 RAM 接口为主的 DMA 模块,速度成为瓶颈。

#### Problem 7.3

对数据来说,时间连续,但是空间性不强,哈希表是一个很好的例子。

```
Get: HashMap map
data = map[key]
for op in {find_index, find_value, ...}:
    op.answer = op(data)
```

#### Problem 7.6

对指令来说,时间和空间性都不强,也就是指令的跳跃性很强。

```
Get: UserInput keyboard
data, choice = read(keyboard)
case(choice):
    type_1:
        multiply(data[0], data[1])
    type_2:
        add(*((int*)data[2]), 0xff3e)
    type_3:
        copy(from=data, to=new_data)
```

#### Problem 7.11

在不进行优化的情况下,通过时钟函数进行计时,代码如下,y优化比例大概有2.5倍。

```
#include <stdio.h>
#include <time.h>

int main()
{
    static int array[10000][100000];
    clock_t start, end;
    start = clock();
```

```
double t1, t2;
    for (int i = 0; i < 10000; ++i)
        for (int j = 0; j < 100000; ++j)
            array[i][j] = 2 * array[i][j];
    }
    end = clock();
    t1 = end - start;
    printf("the cost of for-row is %d\n", end - start);
    start = clock();
    for (int i = 0; i < 100000; ++i)
        for (int j = 0; j < 10000; ++j)
       {
            array[j][i] = 2 * array[j][i];
        }
    }
    end = clock();
    t2 = end - start;
    printf("the cost of for-col is d\n", end - start);
    printf("the ratio is %f", t1 / t2);
   return 0;
}
>>> the cost of for-row is 7779271
>>> the cost of for-col is 19284270
>>> the ratio is 0.403400
```

## **Problem 7.25**

```
addr:
2: empty and load to way1(1-8)
3: hit
11: empty and load to way2(9-16)
16: hit
21: miss and load to way1(17-24)
13: hit
64: miss and load to way1(57-64)
48: miss and load to way2(41-48)
19: miss and load to way1(17-24)
11: miss and load to way2(9-16)
3: miss and load to way1(1-8)
```

```
22: miss and load to way2(17-24)
```

4: hit

27: miss and load to way2(25-32)

6: hit

11: miss and load to way2(9-16)

buffer:

way1: 1 2 3 4 5 6 7 8 way2: 9 10 11 12 13 14 15 16

# **Problem 7.32**

需要计算整体未命中的惩罚:  $Total = I_{penalty} + D_{penalty}$ 。 因此分别为:

$$t1 = (6+1) \times (4\% + 6\%/2) = 0.49$$

$$t2 = (6+4) \times (2\% + 4\%/2) = 0.4$$

$$t3 = (6+4) \times (2\% + 3\%/2) = 0.35$$

# **Problem 7.36**

18 32K x 16-bit

每个数据是 4 \* 16-bit ,偏移为 2 位,那么共有 18 \* 8 K 个数据,每组 3 路,共 48 K 组,那么组索引需要  $\log_2 64K=16$  位,因此标记有 14 位。

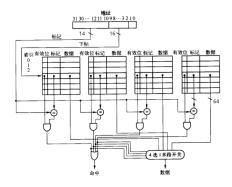


图 1: 7.36