# OPENCL编程模型

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### **Outline**

- OpenCL概述
- · OpenCL抽象模型
  - 平台模型
  - 执行模型
  - 内存模型
  - 编程模型
- · OpenCL程序设计
  - 程序设计流程
  - 编程实例

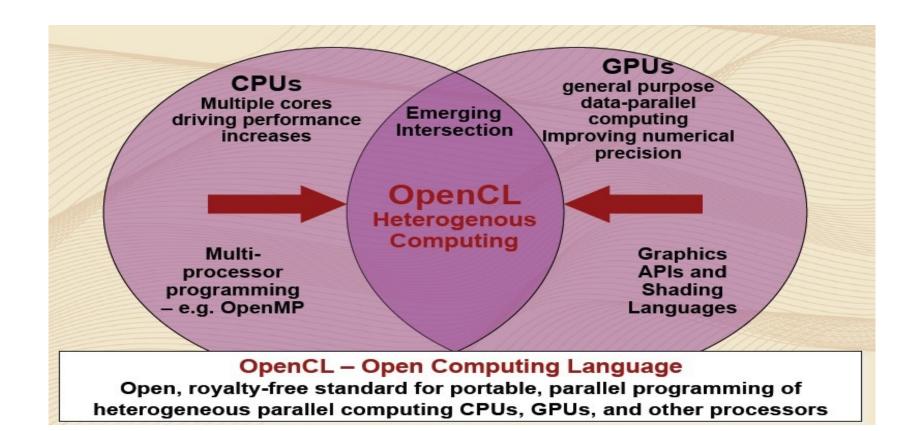
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## 异构计算系统

- A modern computer system has:
  - CPU(s)
  - GPU(s)
  - DSP processors
  - ... other?
- Need to make the best use of all the available resources from within a single program:
  - One program that runs well (i.e., reasonably close to "hand-tuned" performance) on a heterogeneous mixture of processors.

# 异构计算系统编程



# 什么是OpenCL?

- OpenCL (Open Computing Language) 是一个为异构平台编写程序的框架,此异构平台可由CPU,GPU或其他类型的处理器组成。
- Low-level programming API for data parallel computation

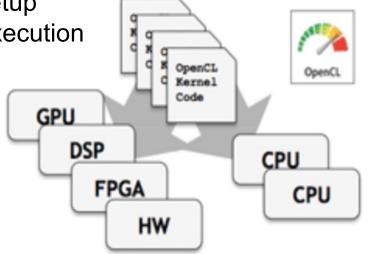
Platform API: device query and pipeline setup

Runtime API: resources management + execution

Cross-platform API

Windows, MAC, Linux, Mobile, Web...

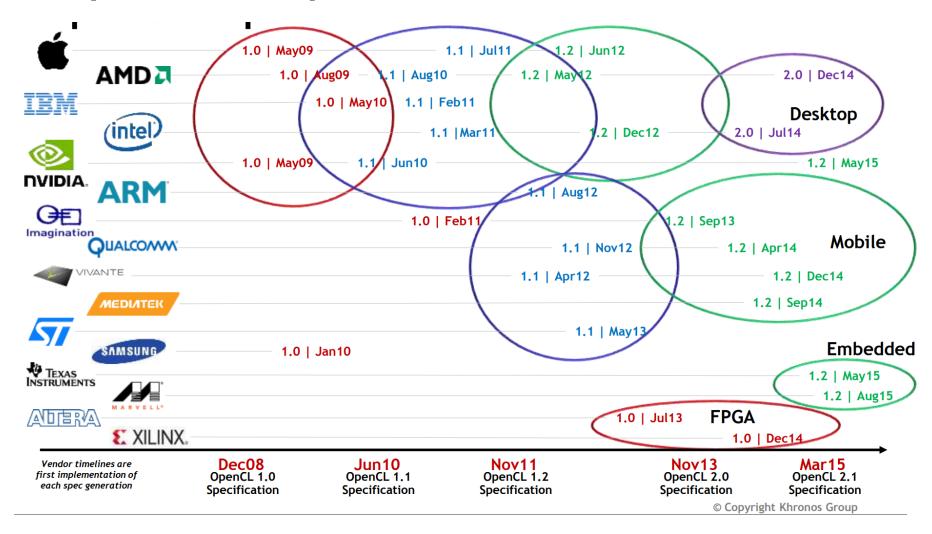
- Portable device targets
  - CPUs, GPUs, FPGAs, DSPs, etc...
  - One code tree can be executed on CPUs, GPUs, DSPs, FPGA and hardware
- Implementation based on C99
- Maintained by Kronos Group (<u>www.khronos.org</u>)
- Current version: 3.0 with C++ support (classes & templates)



# OpenCL的发展历史

- 2008年 Apple 提出OpenCI规范
  2008年 Khronos Group 发布OpenCL 1.0标准
- 2013年 Khronos Group
   OpenCl 2.0标准发布
- 2013年 截至2013年11月,以AMD、Apple、ATI、Intel、Nvidia等为代表的公司,已经发布了多款支持OpenCI的产品。
- 2017年 Khronos Group
   OpenCl 2.2标准发布
- 2020年 Khronos Group
   OpenCl 3.0标准发布

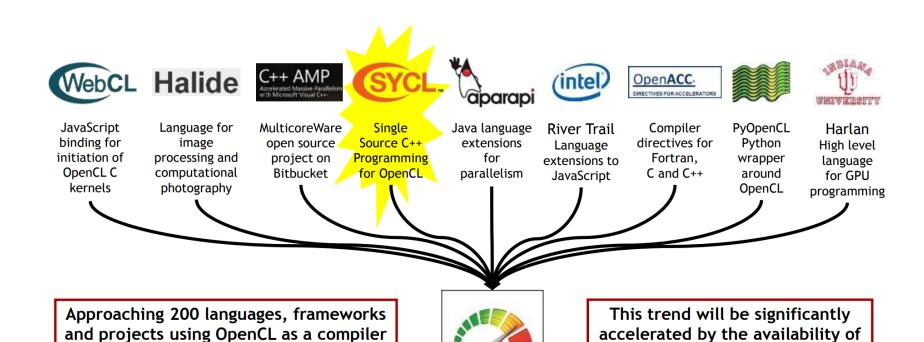
# OpenCL Implementations



# OpenCL Front-End APIs

target to access vendor optimized,

heterogeneous compute runtimes



OpenCL

SPIR-V which is specifically

designed to be a compiler target

# OpenCL: A high-level view

- OpenCL applications:
  - A host program running on the PC
  - One or more Kernels that are queued up to run on CPUs, GPUs, and "other processors".
- OpenCL is understood in terms of these models
  - Platform model
  - Execution model
  - Memory model
  - Programming model

### **Outline**

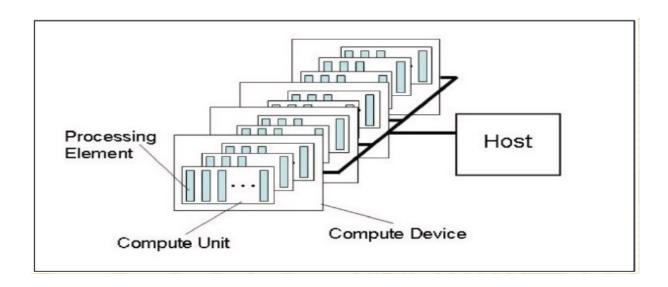
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# OpenCL平台模型

- Each OpenCL implementation (i.e. an OpenCL library from AMD, NVIDIA, etc.) defines platforms which enable the host system to interact with OpenCL-capable devices
  - Currently each vendor supplies only a single platform per implementation
- OpenCL uses an "Installable Client Driver" model
  - The goal is to allow platforms from different vendors to co-exist
  - Current systems' device driver model will not allow different vendors' GPUs to run at the same time

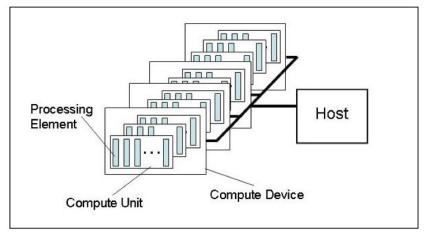
# OpenCL平台模型

- One Host + one or more Compute Devices
  - Each Compute Device is composed of one or more Compute Units
    - Each Compute Unit is further divided into one or more Processing Elements
      - PE executes code as SIMD or SPMD



## OpenCL平台模型

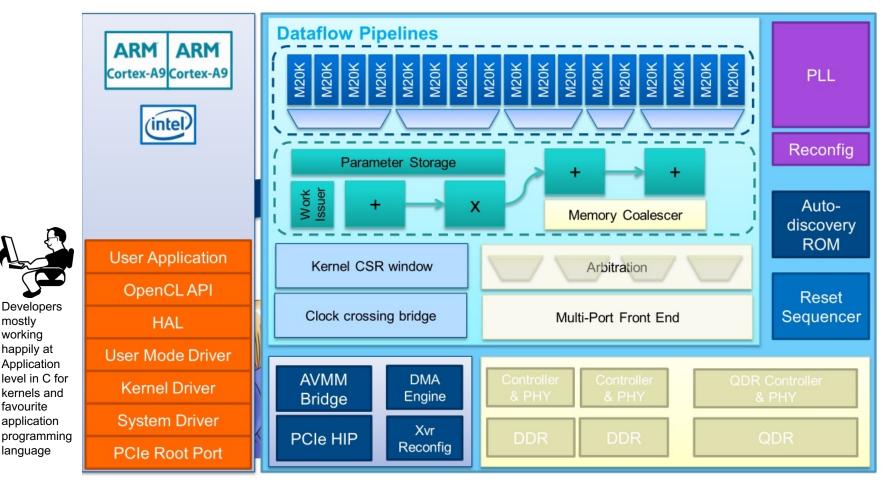
- The host is whatever the OpenCL library runs on
  - x86 CPUs for both NVIDIA and AMD
- Devices are processors that the library can talk to
  - CPUs, GPUs, and generic accelerators
- For AMD
  - All CPUs are combined into a single device (each core is a compute unit and processing element)
  - Each GPU is a separate device



# Altera OpenCL平台模型

mostly

working



The OpenCL compiler takes care of much of the worries including machinebased profiling, tuning and optimization

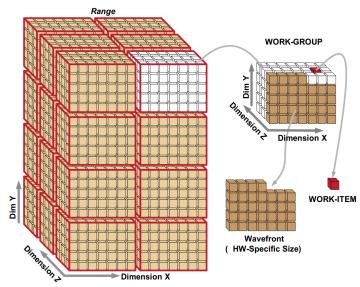


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# OpenCL执行模型

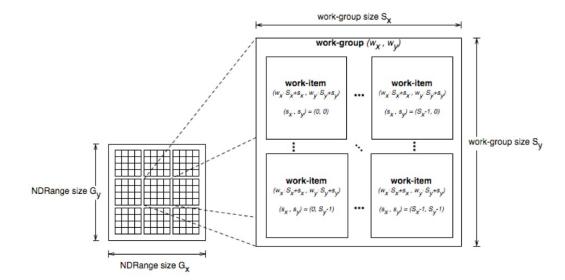
- A kernel is logical unit of instructions to be executed on a compute device.
- Kernels are executed in multidimensional index space:
   NDRange
- For every element of the index space a work-item is executed
- The index space is tiled into workgroups
- Work items within a workgroup are synchronized using barriers or fences



AMD OpenCL User Guide 2015

# OpenCL执行模型

- Host defines a command queue and associates it with a context (devices, kernels, memory, etc).
- Host enqueues commands to the command queue
- Work-items can uniquely identify themselves based on:
  - A global id (unique within the index space)
  - A work-group ID and a local ID within the work-group



### OpenCL线程组织结构

- Threads can determine their global ID in each dimension
  - get\_global\_id(dim)
  - get\_global\_size(dim)
- Or they can determine their work-group ID and ID within the workgroup
  - get\_group\_id(dim)
  - get\_num\_groups(dim)
  - get\_local\_id(dim)
  - get\_local\_size(dim)
- get\_global\_id(0) = column, get\_global\_id(1) = row
- get\_num\_groups(0) \* get\_local\_size(0) == get\_global\_size(0)

# OpenCL执行模型——与CUDA映射

OpenCL terminology aims for generality

OpenCL Terminology	CUDA Terminology	
Compute Unit	Streaming Processor (SM)	
Processing Element	Processor Core	
Kernel	Kernel	
Host program	Host program	
Wavefront (AMD)	Warp	
Work-item	Thread	
Work-group	Thread Block	
NDRange	Grid	

# OpenCL执行模型——与CUDA映射

Work Items Indexing

OpenCL Terminology	CUDA Terminology	
get_num_groups()	gridDim	
get_local_size()	blockDim	
get_group_id()	blockldx	
get_local_id()	threadIdx	
get_global_id()	blockldx * blockDim + threadldx	
get_global_size()	gridDim * blockDim	

# OpenCL执行模型——与CUDA映射

Threads Synchronization

OpenCL Terminology	CUDA Terminology	
barrier()	syncthreads()	
No direct equivalent*	threadfence()	
mem_fence()	threadfence_block()	
No direct equivalent*	threadfence_system()	
No direct equivalent*	syncwarp()	
Read_mem_fence()	No direct equivalent*	
Write_mem_fence()	No direct equivalent*	

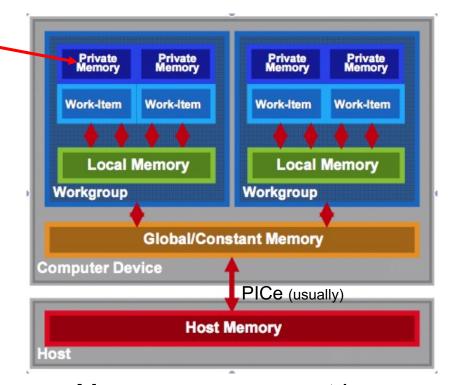
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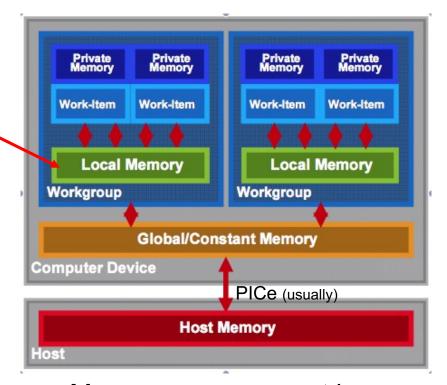
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- Private memory: available per work item
- Local memory: shared in workgroup
- NB: No synchronization between workgroups
- Synchronization possible between work items in a common workgroup
- Global/constant memory accessible by any workitems (no guarantee to be synchronized)
- Host memory: access through the CPU



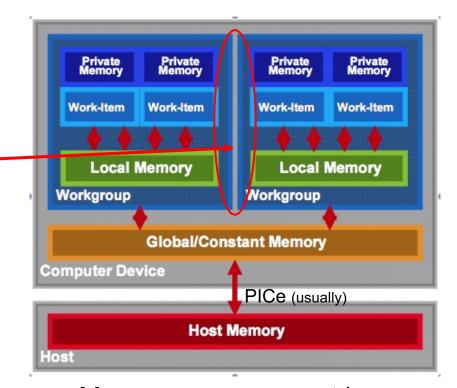
- Memory management is explicit...
- Data moved from: host->global->local and back

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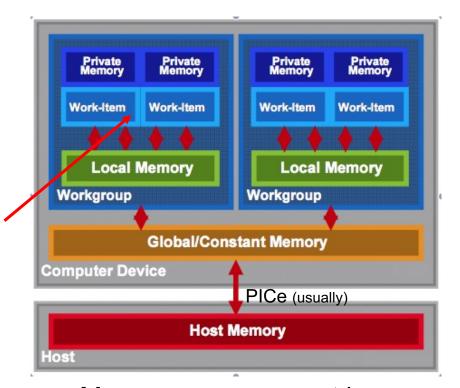
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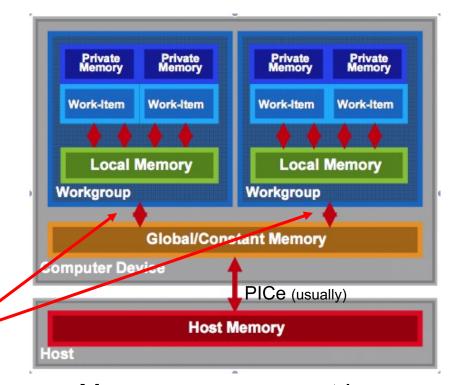
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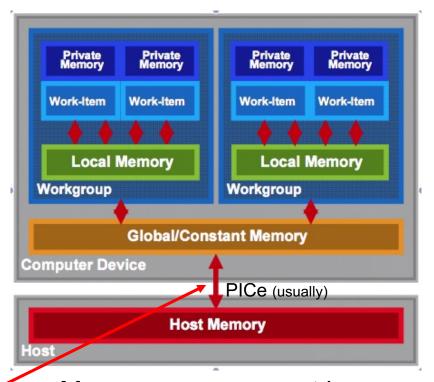
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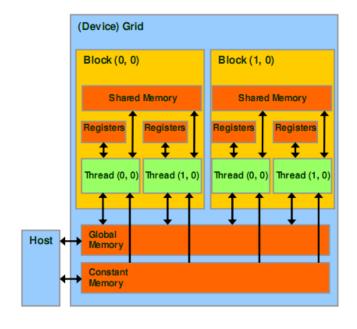


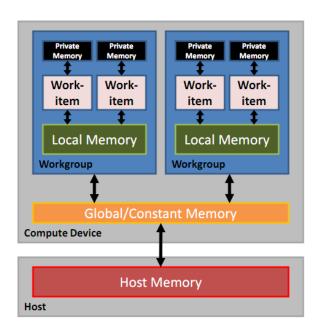
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- Memory management is explicit
  - Must move data from host memory to device global memory, from global memory to local memory, and back
- Work-groups are assigned to execute on compute-units
  - No guaranteed communication/coherency between different workgroups (no software mechanism in the OpenCL specification)

# OpenCL内存模型——与CUDA映射

OpenCL Terminology	CUDA Terminology
Global Memory	Global Memory
Constant Memory	Constant Memory
Local Memory	Shared Memory
Private Memory	Local Memory





# OpenCL内存模型——与CUDA映射

#### Resources Qualifiers

Description	OpenCL Terminology	CUDA Terminology
Kernel global function	kernel	global
Kernel local function	nothing*	device
Readonly memory	constant	device
Global memory	global	device
Private memory	local	shared

## CUDA与OpenCL系键指标对应汇总

#### <u>CUDA</u>

- Thread
- Thread-block
- Global memory
- Constant memory
- Shared memory
- Local memory
- \_\_global\_\_\_ function
- \_\_device\_\_ function
- constant variable
- device variable
- \_\_shared\_\_ variable

#### **OpenCL**

- Work-item
- Work-group
- Global memory
- Constant memory
- Local memory
- Private memory
- kernel function
- no qualification needed
- constant variable
- \_\_global variable
- \_\_local variable

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# OpenCL编程模型

- 数据并行化
  - Work-items in a work-group run the same program
  - Update data structures in parallel using the work-item ID to select
  - data and guide execution.
- 任务并行化
  - One work-item per work group ... for coarse grained task-level parallelism.
  - Native function interface: trap-door to run arbitrary code from an OpenCL command-queue.
  - Multiple different kernels can be executed in parallel

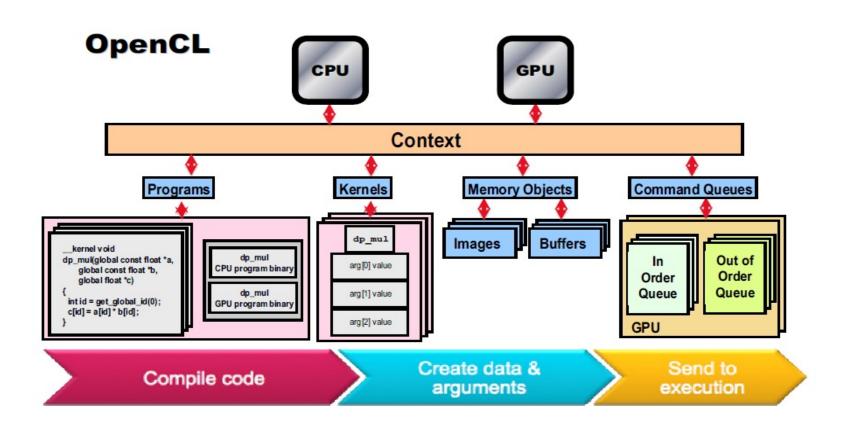
### OpenCL 数据并行化

- Kernels executed across a global domain of workitems
  - Global dimensions define the range of computation
  - One work-item per computation, executed in parallel
  - Work-items are grouped in local work-groups
    - Local dimensions define the size of the work-groups
    - Executed together on one device
    - Share local memory and synchronization
  - Caveats
    - Global work-items must be independent: No global synchronization
    - Synchronization can be done within a work-group

### OpenCL任务并行化

- Expressing Task-Parallelism in OpenCL
  - Execute as a single work-item
  - A kernel written in OpenCL C
- clEnqueueTask
  - Imagine "sea of different tasks" executing concurrently
  - A task "owns the core" (i.e., a workgroup size of 1)

## OpenCL Program Flow



# OpenCL vs CUDA

#### API Terminology

OpenCL Terminology	CUDA Terminology
clGetContextInfo()	cuDeviceGet()
clCreateCommandQueue()	No direct equivalent*
clBuildProgram()	No direct equivalent*
clCreateKernel()	No direct equivalent*
clCreateBuffer()	cuMemAlloc()
clEnqueueWriteBuffer()	cuMemcpyHtoD()
clEnqueueReadBuffer()	cuMemcpyDtoH()
clSetKernelArg()	No direct equivalent*
clEnqueueNDRangeKernel()	kernel<<<>>>()
clReleaseMemObj()	cuMemFree()

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# OpenCL程序设计流程

1 查询平台	clGetPlatformIDs()
2 查询设备	clGetDeviceIDs()
3 创建上下文:将平台设备与上下 文关联起来	clCreateContext() 或 clCreateContextFromType()
4 创建命令队列	clCreateCommandQueue()
5 读取、编译内核	clCreateProgramWithSource()或 clCreateProgramWithBinary() clBuildProgram()
6 打包生成内核	clCreateKernel()
7 创建缓存对象或图像对象, 为内核参数分配内存	clCreateBuffer().clCreateSubBuffer() clCreateImage2D().clCreateImage3D()
8 设置内核参数,将上面分 配的内存发送到设备上	clSetKernelArg()
9 执行内核	clEnqueueTask()或 clEnqueueNDRangeKernel()
10 读取设备上的处理结果	clEnqueueReadBuffer()
11 释放创建的资源: 创建的内存、 命令队列、 内核、 打包的程序、 上下文	clReleaseMemObject(), clReleaseCommandQueue(), clReleaseKernel(), clReleaseProgram(), clReleaseContext()

### Selecting a Platform

```
cl_int clGetPlatformIDs (cl_uint num_entries,
cl_platform_id *platforms,
cl_uint *num_platforms)
```

- This function is usually called twice
  - The first call is used to get the number of platforms available to the implementation
  - Space is then allocated for the platform objects
  - The second call is used to retrieve the platform objects

## Selecting Devices

 Once a platform is selected, we can then query for the devices that it knows how to interact with

```
clGetDeviceIDs<sup>4</sup> (cl_platform_id platform,
cl_device_type device_type,
cl_uint num_entries,
cl_device_id *devices,
cl_uint *num_devices)
```

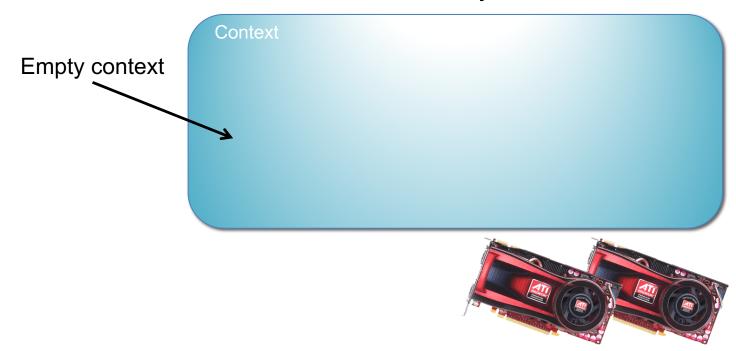
- We can specify which types of devices we are interested in (e.g. all devices, CPUs only, GPUs only)
- This call is performed twice as with clGetPlatformIDs
  - The first call is to determine the number of devices, the second retrieves the device objects

#### Contexts

- A context refers to the environment for managing OpenCL objects and resources
- To manage OpenCL programs, the following are associated with a context
  - Devices: the things doing the execution
  - Program objects: the program source that implements the kernels
  - Kernels: functions that run on OpenCL devices
  - Memory objects: data that are operated on by the device
  - Command queues: mechanisms for interaction with the devices
    - Memory commands (data transfers)
    - Kernel execution
    - Synchronization

#### Contexts

- When you create a context, you will provide a list of devices to associate with it
  - For the rest of the OpenCL resources, you will associate them with the context as they are created



#### Contexts

- This function creates a context given a list of devices
- The properties argument specifies which platform to use (if NULL, the default chosen by the vendor will be used)
- The function also provides a callback mechanism for reporting errors to the user

#### Command Queues

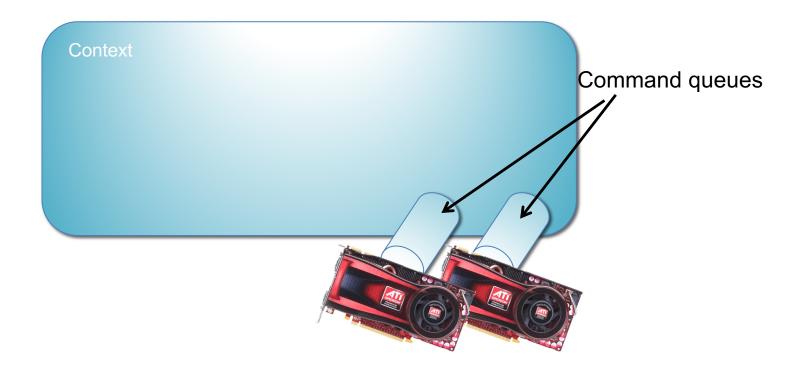
- A command queue is the mechanism for the host to request that an action be performed by the device
  - Perform a memory transfer, begin executing, etc.
- A separate command queue is required for each device
- Commands within the queue can be synchronous or asynchronous
- Commands can execute in-order or out-of-order

### **Command Queues**

- A command queue establishes a relationship between a context and a device
- The command queue properties specify:
  - If out-of-order execution of commands is allowed
  - If profiling is enabled

#### **Command Queues**

- Command queues associate a context with a device
  - Despite the figure below, they are not a physical connection



### Memory Objects

- Memory objects are OpenCL data that can be moved on and off devices
  - Objects are classified as either buffers or images
- Buffers
  - Contiguous chunks of memory stored sequentially and can be accessed directly (arrays, pointers, structs)
  - Read/write capable
- Images
  - Opaque objects (2D or 3D)
  - Can only be accessed via read\_image() and write\_image()
  - Can either be read or written in a kernel, but not both

### Creating buffers

```
cl_mem clCreateBuffer (cl_context context,
cl_mem_flags flags,
size_t size,
void *host_ptr,
cl_int *errcode_ret)
```

- This function creates a buffer (cl\_mem object) for the given context
  - Images are more complex and will be covered in a later lecture
- The flags specify:
  - the combination of reading and writing allowed on the data
  - if the host pointer itself should be used to store the data
  - if the data should be copied from the host pointer

## **Transferring Data**

- OpenCL provides commands to transfer data to and from devices
  - clEnqueue{Read|Write}{Buffer|Image}
  - Copying from the host to a device is considered writing
  - Copying from a device to the host is reading
- The write command both initializes the memory object with data and places it on a device
  - The validity of memory objects that are present on multiple devices is undefined by the OpenCL spec (i.e. are vendor specific)
- OpenCL calls also exist to directly map part of a memory object to a host pointer

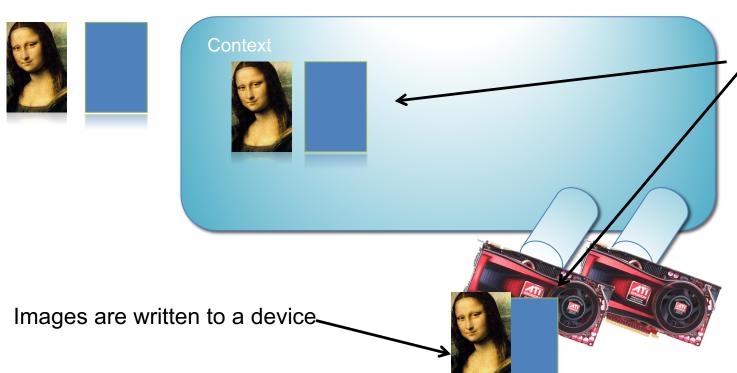
## **Transferring Data**

```
cl_int clEnqueueWriteBuffer (cl_command_queue command_queue, cl_mem buffer, cl_bool blocking_write, size_t offset, size_t cb, const void *ptr, cl_uint num_events_in_wait_list, const cl_event *event_wait_list, cl_event *event)
```

- This command initializes the OpenCL memory object and writes data to the device associated with the command queue
  - The command will write data from a host pointer (ptr) to the device
- The blocking\_write parameter specifies whether or not the command should return before the data transfer is complete
- Events can specify which commands should be completed before this one runs

### **Transferring Data**

- Memory objects are transferred to devices by specifying an action (read or write) and a command queue
  - The validity of memory objects that are present on multiple devices is undefined by the OpenCL spec (i.e. is vendor specific)



The images are redundant here to show that they are both part of the context (on the host) and physically on the device

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 The "hello world" program of data parallel programming is a program to add two vectors

$$C[i] = A[i] + B[i]$$
 for  $i=1$  to  $N$ 

- For the OpenCl solution, there are two parts
  - Host code
  - Kernel code

Address space qualifier

kernel qualifier

Global thread index

Vector addition

Setup kernel grid

Allocate host resources

Create device context

Allocate device resources

Populate device memory

```
pint main(int argc, char **argv) {
    // set and log Global and Local work size dimensions
   localWorkSize = 256;
    globalWorkSize = RoundUp(localWorkSize, iNumElements);
    // Allocate and initialize host arrays
    srcA = malloc(sizeof(cl float) * globalWorkSize);
    srcB = malloc(sizeof(cl float) * globalWorkSize);
    dst = malloc(sizeof(cl float) * globalWorkSize);
    FillArray(srcA, iNumElements);
    FillArray(srcB, iNumElements);
    // Create the context
    clGetPlatformIDs (1, &platform, NULL);
    clGetDeviceIDs (platform, CL DEVICE TYPE GPU, 1, &device, NULL);
    context = clCreateContext(0, 1, &device, NULL, NULL, &ciErrl);
    // Create a command-queue
    cqCommandQueue = clCreateCommandQueue(context, device, 0, &ciErr1);
    // Allocate the OpenCL buffer memory objects
    dSrcA = clCreateBuffer(context, CL MEM READ ONLY, sizeof(cl float) * globalWorkSize, ...);
    dSrcB = clCreateBuffer(context, CL MEM READ ONLY, sizeof(cl float) * globalWorkSize, ...);
    dDst = clCreateBuffer(context, CL MEM WRITE ONLY, sizeof(cl float) * globalWorkSize, ...);
    // Copy data to GPU device
    clEnqueueWriteBuffer(cgCommandQueue, dSrcA, CL FALSE, 0, sizeof(cl float) * globalWorkSize, srcA, ...);
    clEnqueueWriteBuffer(cqCommandQueue, dSrcB, CL FALSE, 0, sizeof(cl float) * qlobalWorkSize, srcB, ...);
```

Build kernel program

Set kernel arguments

Launch kernel execution

Read destination buffer

```
// Build the program
progSrc = oclLoadProgSource("VectorAdd.cl", ...);
prog = clCreateProgramWithSource(context, 1, progSrc, ...);
clBuildProgram(prog, 0, NULL, NULL, NULL, NULL);
// Create the kernel
kernel = clCreateKernel (prog, "VectorAdd", &ciErr1);
// Set kernel the Arguments
clSetKernelArg(kernel, 0, sizeof(cl mem), (void*)&dSrcA);
clSetKernelArg(kernel, 1, sizeof(cl mem), (void*)&dSrcB);
clSetKernelArg(kernel, 2, sizeof(cl mem), (void*)&dDst);
clSetKernelArg(kernel, 3, sizeof(cl int), (void*)&iNumElements);
// Launch kernel
clEnqueueNDRangeKernel(cqCommandQueue, kernel, 1, NULL, &globalWorkSize, &localWorkSize,...);
// Synchronous/blocking read of results, and check accumulated errors
clEngueueReadBuffer(cgCommandQueue, dDst, ...);
// Cleanup
Cleanup();
return 0;
```

## 参考资料

- · OpenCL异构并行计算:原理、机制与优化实践陈轶,吴长江,刘文志 著
- Khronos OpenCL Registr, https://www.khronos.org/ registry/OpenCL/