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Roll No: C-111

Artificial Intelligence <u>Experiment 1</u>

1. AI for Personalized Marketing

Problem Statement: Design an AI system to create personalized marketing strategies and content tailored to individual customer preferences and behaviors.

Description: This problem involves creating an AI system that uses customer data to develop customized marketing strategies. The AI analyzes individual customer preferences, behaviors, and past interactions to tailor marketing messages, offers, and advertisements. The goal is to increase engagement, improve conversion rates, and maximize return on investment by delivering highly relevant content to each customer.

PEAS:

- Performance Measure: Engagement rates, conversion rates, and return on investment.
- Environment: Customer data, marketing channels, and product information.
- Actuators: Personalized ads, email campaigns, and product recommendations.
- **Sensors:** Customer interaction data, purchase history, and browsing patterns.

2. Predictive Text Input

Problem Statement: Develop an AI-powered predictive text input system to enhance typing efficiency and accuracy on mobile and desktop devices.

Description: This problem centers on developing an AI system that predicts and suggests text as users type on their devices. By analyzing typing patterns, user context, and linguistic data, the system aims to enhance typing efficiency and accuracy. The AI predicts the next word or phrase, offers autocomplete suggestions, and corrects errors in real-time, making typing faster and reducing user frustration.

- Performance Measure: Typing speed, accuracy, and user satisfaction.
- Environment: Text input fields, user typing patterns, and language models.

- Actuators: Predictive text suggestions, auto-corrections, and word completions.
- Sensors: User keystrokes, text input history, and contextual information.

3. AI for Drug Discovery

Problem Statement: Create an AI system to accelerate drug discovery by identifying potential drug candidates and predicting their effectiveness.

Description: This problem involves creating an AI system to accelerate the process of discovering new drugs. The AI analyzes vast amounts of chemical and biological data to identify potential drug candidates and predict their effectiveness against diseases. By leveraging machine learning algorithms, the system helps researchers focus on the most promising compounds, reducing the time and cost associated with drug development.

PEAS:

- **Performance Measure:** Discovery rate of viable drug candidates, accuracy of predictions, and reduction in research time.
- Environment: Chemical compound databases, biological data, and research literature.
- **Actuators:** Drug candidate recommendations, research prioritization, and experimental suggestions.
- Sensors: Molecular data, biological assays, and computational models.

4. Autonomous Vehicle Navigation

Problem Statement: Design an AI system to enable autonomous vehicles to navigate safely and efficiently through diverse and dynamic environments.

Description: This problem focuses on designing an AI system for autonomous vehicles that can navigate safely and efficiently through various environments. The system must handle diverse traffic conditions, interpret road signs and signals, and avoid obstacles such as pedestrians and other vehicles. The AI integrates data from sensors like cameras, LIDAR, and radar to make real-time driving decisions and ensure safe navigation.

- **Performance Measure:** Safety, navigation accuracy, efficiency, and compliance with traffic laws.
- **Environment:** Road networks, traffic conditions, and environmental factors.
- Actuators: Steering, acceleration, braking, and signaling systems.
- Sensors: Cameras, LIDAR, radar, GPS, and IMU.

5. AI for Financial Market Prediction

Problem Statement: Develop an AI model to predict stock market trends and assist in making informed investment decisions.

Description: This problem involves developing an AI model to predict financial market trends, aiding investors in making informed decisions. The AI analyzes historical market data, economic indicators, and financial news to forecast stock prices and market movements. The goal is to provide actionable insights that can help in making strategic investment decisions and managing financial risk.

PEAS:

- **Performance Measure:** Prediction accuracy, profitability of investment strategies, and risk management.
- Environment: Financial data, market trends, and economic indicators.
- Actuators: Investment recommendations, trading signals, and portfolio adjustments.
- Sensors: Stock prices, trading volume, and financial news.

6. Speech Recognition for Virtual Assistants

Problem Statement: Create an AI system for accurate speech recognition to enhance the functionality of virtual assistants in understanding and responding to user commands.

Description: This problem entails creating an AI system that accurately recognizes and processes human speech to enhance virtual assistants. The system must understand spoken commands, queries, and conversations in real-time, converting them into actionable responses or tasks. The AI improves user interaction by accurately interpreting speech and providing relevant responses or executing commands.

- **Performance Measure:** Recognition accuracy, response time, and user satisfaction.
- Environment: Audio input, user commands, and contextual information.
- Actuators: Voice responses, action execution, and command processing.
- Sensors: Microphones, speech processing algorithms, and language models.

7. AI-Driven Healthcare Diagnosis

Problem Statement: Design an AI system to assist healthcare professionals in diagnosing diseases based on patient symptoms and medical history.

Description: This problem focuses on developing an AI system to support healthcare professionals in diagnosing medical conditions. By analyzing patient symptoms, medical history, and diagnostic data, the AI provides diagnostic suggestions and treatment recommendations. The aim is to enhance diagnostic accuracy, support decision-making, and improve patient outcomes.

PEAS:

- **Performance Measure:** Diagnostic accuracy, response time, and integration with medical records.
- Environment: Patient symptoms, medical history, and diagnostic guidelines.
- Actuators: Diagnostic suggestions, treatment recommendations, and alerts.
- Sensors: Patient data, medical records, and symptom descriptions.

8. AI for Real-Time Language Translation

Problem Statement: Develop an AI system to provide real-time language translation for seamless communication in multilingual contexts.

Description: This problem involves designing an AI system that translates languages in real-time, facilitating communication in multilingual settings. The system must handle both spoken and written text, providing accurate and contextually relevant translations. It aims to break down language barriers, enabling seamless communication across different languages.

PEAS:

- **Performance Measure:** Translation accuracy, latency, and user satisfaction.
- Environment: Text or speech input, languages, and context.
- Actuators: Translated text or speech output.
- Sensors: Language models, translation algorithms, and speech recognition tools.

9. AI for E-commerce Product Recommendations

Problem Statement: Create an AI system to recommend products to users based on their browsing behavior and purchase history.

Description: This problem centers on creating an AI system that generates personalized product recommendations for e-commerce platforms. The system uses user browsing history, purchase patterns, and product attributes to suggest items that match individual preferences. The goal is to enhance the shopping experience and increase sales by providing relevant product suggestions.

- **Performance Measure:** Recommendation accuracy, user engagement, and sales increase.
- Environment: User browsing history, product catalog, and purchase data.
- Actuators: Product recommendations, personalized ads, and content suggestions.
- Sensors: User behavior data, purchase history, and product attributes.

10. AI for Dynamic Pricing in Retail

Problem Statement: Design an AI system to optimize retail pricing dynamically based on market conditions, demand, and competitor pricing.

Description: This problem involves developing an AI system to adjust retail prices dynamically based on market conditions, demand, and competitor pricing. The AI analyzes real-time data to optimize pricing strategies, aiming to maximize revenue, maintain competitiveness, and respond to market fluctuations effectively.

- **Performance Measure:** Revenue maximization, price accuracy, and market competitiveness.
- Environment: Market conditions, demand data, and competitor pricing.
- Actuators: Price adjustments, promotional offers, and discount strategies.
- Sensors: Sales data, competitor prices, and customer demand patterns.

Problem Statement	Observab le	Determinist ic	Episodi c	Static	Discrete	Agents
1. Al for Personalized Marketing	Partially	Stochastic	Sequenti al	Dynami c	Discrete	Single (Marketing AI)
2. Predictive Text Input	Observable	Stochastic	Sequenti al	Static	Discrete	Single (Text Input AI)
3. Al for Drug Discovery	Observabl e	Stochastic	Sequenti al	Dynami c	Continuo us	Single (Discovery AI)
4. Autonomous Vehicle Navigation	Fully Observable	Stochastic	Sequential	Dynami c	Continuo	Single (Autonomous Vehicle)

5. Al for Financial Market Prediction	Partially	Stochastic	Sequential	Dynami c	Continuo us	Single (Prediction AI)
6. Speech Recognition for Virtual Assistants	Observable	Stochastic	Sequential	Static	Discrete	Single (Virtual Assistant)
7. Al-Driven Healthcare Diagnosis	Observable	Stochastic	Sequential	Static	Discrete	Single (Healthcare AI)
8. Al for Real- Time Language Translation	Observable	Stochastic	Sequential	Static	Discrete	Single (Translation AI)
9. Al for E- commerce Product Recommendatio ns	Observable	Stochastic	Sequential	Static	Discrete	Single (Recommendati on AI)
10. Al for Dynamic Pricing in Retail	Observable	Stochastic	Sequential	Dynami c	Discrete	Single (Pricing AI)

Conclusion:-

Understanding the distinctions between fully and partially observable, static and dynamic, and deterministic and stochastic environments is essential in various fields. Fully observable environments provide complete information, simplifying decision-making, while partially observable environments require inference and probabilistic reasoning. Static environments remain unchanged unless acted upon, whereas dynamic environments continuously evolve due to external factors. Deterministic systems yield predictable outcomes given initial conditions, whereas stochastic systems incorporate randomness, leading to variable outcomes. These concepts guide the design and implementation of algorithms and systems in AI, robotics, game development, and beyond, ensuring they effectively handle the nature and complexity of the environments they operate in.

Name: Parth Das Sap: 60004220185 Roll No:- C-111 Al Experiment 2

Aim:- Implement BFS and DFS algorithm

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PI Experiment-2	
Aim To implement RES & DES to seach quality	te
Theory - 12 Boradth First search (BES):-	
BFI is a graph towersal algorithm that starts at a source and exploses all its neighbours at the present depth is	vel
a queue to keep toack of the noder to be explosed next.	Mes
Algorithm.	-
I Initialize a queue and enqueue the starting node. mark of	he
2- while the queue is not empty:	5
· Dequeue a node from front of the queue.	7.
· too each unrelated neighbour of the doqueued node:	-
> mark the reighbour as visited > Engineere the neighbour in the genue.	-
3. Repeat until the queue it empty	
property of the state of which a property of the party	-
Observation -	
is largered Explanation:	
of their distance form the starting node:	8
grow the starting node to any node in an unweighted of	edges -
The queue ensures nodes are explosed in the order to are discovered, providing a systematic level-order to time complexity: 0 (V+E) V vertue E edges	heu -
spour complexity: o (.v)	1 600

```
23 Depth Fiart Seasch (DES)
DIS is a graph to arrived algorithm that starts at source node and
explores at far as pshible along each branch before booktoning.
It was a Stock (os or auriers) to keep toack of neder to be explored
 1. Initiative stack and push starting node onto the stack. Mark the
 Starting node as writed.
 2. While the stark is not empty:
     . App a node from the top of the flork.
     . For each unvisited neighbour of the popped node
         -> mark the neighbour as weited
         -> Push the neighbour and the elect.
 3. Repeat winter stack is emply.
  Objectivation:
 is seep Explosation: DTS exploses as deep as possible along a path
 before backtracking. It does not explore nodes tour by love but
 stather follows a single poth to its end.
 2) Rath Discovery - Dis can be used to find a path between nodes
 and it well- nited for problems requiring exploration of all
 possible paths
 33 the of Stock: The Stock ensures that nodes are recurred in
 The reverse under of their discovery, so cilitating deep
 exploration and backtooking Icholviel v vertice
                                S.C. - O(V) E Edger
 conclusion.
 BOTH BET & DES are exercical graph algorithms, each with its own
 strength & appropriate use cases. BET is award see level-coder
explanation and shortest path in unweighted graphs while asserces
gin deep explosiation and postifieding. Understanding their mechanisms
 helps in selecting eight algorithm baled on problem sequitement
```

Code:-

BFS.cpp

```
#include <iostream>
#include <vector>
#include <limits>
#include <queue>
using namespace std;
int path_cost;
// BFS implementation with path cost calculation
bool bfs(const vector<vector<int>>& graph, const vector<vector<int>>& cost, int start, int goal, int numNodes, vector<bool>& visited) {
    vector<int> costToNode(numNodes, numeric_limits<int>::max()); //
Initialize costs to a large value
    queue<int> q;
```

```
q.push(start);
    visited[start] = true;
    costToNode[start] = 0;
   while (!q.empty()) {
        int current = q.front();
        q.pop();
        cout << "Visited node: " << current << endl;</pre>
        // Check if we have reached the goal
        if (current == goal) {
            path_cost = costToNode[goal];
            return true;
        }
        // Process all adjacent nodes
        for (int i = 0; i < numNodes; i++) {</pre>
            if (graph[current][i] && !visited[i]) {
                // Update cost to reach the adjacent node
                if (costToNode[current] + cost[current][i] <</pre>
costToNode[i]) {
                    costToNode[i] = costToNode[current] +
cost[current][i];
                q.push(i);
                visited[i] = true;
        }
    return false;
int main() {
   int numNodes = 6;
   int start = 0;
    int goal = 4;
   vector<vector<int>> graph(numNodes, vector<int>(numNodes, 0)); //
Adjacency matrix
    vector<vector<int>> cost(numNodes, vector<int>(numNodes,
numeric_limits<int>::max())); // Cost matrix
   vector<bool> visited(numNodes, false); // Visited array
    graph[0][1] = 1;
    graph[0][2] = 1;
    graph[1][3] = 1;
    graph[2][3] = 1;
    graph[3][4] = 1;
   graph[4][5] = 1;
    cost[0][1] = 22;
    cost[0][2] = 24;
    cost[1][3] = 11;
    cost[2][3] = 3;
```

```
cost[3][4] = 15;
cost[4][5] = 19;
if (bfs(graph, cost, start, goal, numNodes, visited)) {
    cout << "Path to goal node " << goal << " found." << endl;
    cout << "Path cost is " << path_cost << endl;
} else {
    cout << "No path to goal node " << goal << "." << endl;
}

cout << "Non-visited nodes:" << endl;
for (int i = 0; i < numNodes; i++) {
    if (!visited[i]) {
        cout << "Node " << i << endl;
}

return 0;
}</pre>
```

Output:-

Visited node: 0
Visited node: 1
Visited node: 2
Visited node: 3
Visited node: 4
Path to goal node 4 found.
Path cost is 48
Non-visited nodes:
Node 5

Code:-

DFS.cpp

```
#include <iostream>
#include <vector>
#include <limits>

using namespace std;
#define MAX_NODES 100
int path_cost;
bool dfs(const vector<vector<int>>& graph, const vector<vector<int>>& cost, int current, int goal, vector<bool>& visited, int numNodes, vector<int>& costToNode) {
    // Mark the current node as visited
    visited[current] = true;
    cout << "Visited node: " << current << endl;
    if (current == goal) {
        path_cost = costToNode[goal];
}</pre>
```

```
return true;
    for (int i = 0; i < numNodes; i++) {
        if (graph[current][i] && !visited[i]) {
            if (costToNode[current] + cost[current][i] < costToNode[i])</pre>
                costToNode[i] = costToNode[current] + cost[current][i];
            }
            if (dfs(graph, cost, i, goal, visited, numNodes,
costToNode)) {
                return true;
            }
        }
    return false;
int main() {
    vector<vector<int>> graph(MAX_NODES, vector<int>(MAX_NODES, 0)); //
Adjacency matrix
    vector<vector<int>> cost(MAX NODES, vector<int>(MAX NODES,
numeric_limits<int>::max())); // Cost matrix
    vector<bool> visited(MAX_NODES, false);
    vector<int> costToNode(MAX_NODES, numeric_limits<int>:::max()); //
Cost to node array
    int numNodes, start, goal;
    numNodes = 6;
    path_cost = 0;
    graph[0][1] = 1;
    graph[0][2] = 1;
    graph[1][3] = 1;
   graph[2][3] = 1;
    graph[3][4] = 1;
   graph[4][5] = 1;
    cost[0][1] = 22;
    cost[0][2] = 24;
    cost[1][3] = 11;
    cost[2][3] = 3;
    cost[3][4] = 15;
    cost[4][5] = 19;
    start = 0;
    goal = 4;
    costToNode[start] = 0;
   if (dfs(graph, cost, start, goal, visited, numNodes, costToNode)) {
        cout << "Path to goal node " << goal << " found." << endl;</pre>
        cout << "Path cost is " << path_cost << endl;</pre>
    } else {
        cout << "No path to goal node " << goal << "." << endl;</pre>
```

```
}
cout << "Non-visited nodes:" << endl;
for (int i = 0; i < numNodes; i++) {
    if (!visited[i]) {
        cout << "Node " << i << endl;
    }
}
return 0;
}</pre>
```

Output:-

```
Visited node: 0
Visited node: 1
Visited node: 3
Visited node: 4
Path to goal node 4 found.
Path cost is 48
Non-visited nodes:
Node 2
Node 5
```

Name: Parth Das Sap: 60004220185 Roll No:- C-111 Al Experiment 3

DFID Implementation

. If the Search resches the depth limit without findingthe good,	Poseth Dad Gees 4220189	(-11)
herry - Depth Flort Iterative Despering (DF10): DF10 is a graph traversal algorithm that combines the space efficiently by Depth-First learth (DFS) with the completeness of Broadth trait search (BFS). It playerms a series of depth limited DFS searches with inversaring depth limits will the goal is found. This approach ensures that ever level by the search tose is explored in a depth-first manner, but with the advantage of eventually exploring all possible depths. Player the search: Professor a depth limited DFS with subsent depth limit: I goal is found dwarp DFS, terminate & ochern the solution. I goal is found dwarp DFS, terminate & ochern the solution. I goal is found dwarp DFS, terminate & ochern the solution. I support the search. Depth limited DFS: Initialize a starting node as visited. White starting node as visited. White starting node as visited. The prode from top of the starting node onto the stark. The rode is the goal seturn the solution. The search workers used and push it onto the stork. I the search workers the deard push it onto the stork.	Al Experiment-3	
of Depth teach first leasth (DFS) with the completeness of Beadth test leasth (DFS) with the completeness of Beadth test leasth (DFS). It playsowns a series of depth limited DFS leasthes with inversing depth limits with the goal is found. This approach ensures that ever level of the leasth love is explosed in a depth first manner, but with the advantage of eventually explosing all possible depth. Proposition: Trutialize: lest initial depth limit to 0: Iterative leasth: Proposed is found during DFS with current depth limit: Y goal is found during DFS, terminate & detern the solution. Y not found, increase the depth limit by I and support the search. Depth limited DFS: Initialize a stack & push the starting node and the stack. mark the starting node as visited. White stack is not empty: Prop node from top by the stack. White stack is not empty: The node is the goal seturn the solution. The node is the goal seturn the solution. The starth resides the down is written current depth. Limit mark tills writted and push it ento the stock.	Aim: To implement OFIO search algorith.	m to seach goal.
of Deadth teast branch (BFS) It playsours a series of depth limited DFS branch (BFS). It playsours a series of depth limited DFS branches with increasing depth limits with the goal is found. This approach ensures that ever level of the search love is explosed in a depth first manner, but with the advantage of eventually explosing all possible depth. Proposition: Trutialize: Set initial depth limit to 0: Iterative search: Proposed the search of the depth limit by I and support the search. I not found, increase the depth limit by I and support the search. I not found, increase the depth limit by I and support the search. I notalize a stack & push the starting node and the stack. Mark the starting node as visited. White stack is not empty: I node is the goal seturn the solution To node from top by the stack If node is the goal seturn the solution To the starch resides the dead push it ento the stock. It the depth limited and push it ento the stock.	Theory: - Depth Flort Iterative Despering (OFID):
of breadth test branch (PS) It performs a series of depth. Limited DFS bearches with increasing depth limits until the goal is found. This approach ensures that ever level of the bearch dose is explaned in a depth first manner, but with the advantage of eventually explaning all possible depths. Proposition: Initialize: Set initial depth limit to 0: Iterative bearch: Proferror a depth limited DFS with current depth limit: I goal is found draine DFS, terminate & othern the solution. I goal is found draine DFS, terminate & othern the solution. I post found, increase the depth limit by 1 and support the search. Depth limited DFS: Initialize a stark & push the starting node onto the stark. mark the starting node as visited. White stark is not empty: Prop node from top of the stork If node is the goal, between the solution Son soch unwinted neighbour of the rode: I the depth limit of push board is resition current depth. I the search reaches the dane of the it onto the stork.	OTIO is a graph traversal algorithm that	combines the space
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I soutialize: Set initial depth limit to 0. Iterative search: Perform a depth limited DES with current depth limit. If goal is hound during DES, terminate & seturn the solution. If not found, increase the depth limit by 1 and support the search. Initialize a stack & push the starting node and the stack. mark the starting node as visited. white starting node as visited. White start is not empty: Pop node from top of the stack If node is the goal, seturn the solution. To soch unvisited neighbour of the node: The depth limit of neighbour is writhin current depth limit, mark talks visited and push it onto the stack.	Algeritam:	and the second
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· If goal as found shring DTS, terminate & deturn the solution Y not found, increase the depth limit by I and Support the search. 3. Depth limited DFS: · Initialize a stack & push the starting node and the stack. mask the starting node as virited. · White stark is not empty: - Pop node from top of the stock - If node is the goal seturn the solution - Too such invisited neighbour of the node: - To the depth limit of neighbour is writtin current depth limit, mark ithe visited and push it onto the stock. If the search reaches the death I wish it onto the stock.		
Support the search. 3. Depth-limited DFS: • Initialize a stack & push the starting node and the stack. mark the starting node as visited. • White starting node as visited. • White starting node as visited. • White starting node as visited. • The prode from top of the stack • If node is the goal seturn the solution • Too soch unvisited neighbour of the node: • If the depth limit of neighbour is writhin current depth. limit mark it is visited and push it onto the stack.	· I and is hound during the town and	sent depth limit.
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· white sork is not empty: -> Pop node from top of the stock -> If node is the goal softern the solution >> Too sort invisited neighbour of the node: -> If the depth limit of neighbour is writhin current down limit, mark ithe visited and push it onto the stock.	· Initialize a stack & push the starting node	and the stack.
The hade is the goal seturn the solution The soch unvisited neighbour of the mode: The depth limit of neighbour is writhin current door. The hearth mark it wisted and push it onto the stack.	The state of the state of the state of	
The search writing and push it onto the stark.	- Do and to empty:	or Name of Street, or other
If the depth limit of neighbour is within current dook. I the search reaches the dank hit onto the stock.	- to make from top of the stock	untiles.
If the depth limit of neighbour is within current dook. I the search recipies the dank hit onto the stock.	> Too soch myself seture the solution	84
. If the search receives the dans I wish it onto the stock.	STATE OF THE PROPERTY OF THE PARTY OF	William Control of the Control of th
. If the Search reaches the dans I will brite the stock.	Coping Alberta Re- april 4 con the	within current again.
	. If the Search reaches the dance I will	ordo the stack.

O DIERWHEIGH -13 Consistences DILD is complete, meaning it will find the goal if it exists because it incrementally increases the depth dimet until all possible redes are explosed. 22 Optimality In an unweighted great, 1510 will find the Shortest parties the goal, as it explores noder level by level in incorporary depth 52 Space Upicional of memory compared to BTS - Although it performs multiple dept limited dearther each individual march were the space completely of DET old) where "I is current depth limit 42 ledundant wares of to plasme Hedundart with accord depth-limited seatch security order at another depths However, the is generally maturizated by its space opposity and completeness. Landadia THE IS A hyperial secret strategic that compines the absorption of els and Bis Is as allas for problems where digth of the Sociation to increase and can be bursted by depth limit on to qualificative finding the solution of one waste white maintaining Jeen appearing similar to Oto a is providily integral in severice with Large to enjoyer stanch spaces where are would to improve that to high manage company but a complet season is still sequent. PHE REAL PROPERTY. me complexity: The time complexity of OFID is O (b'd) where b: branching forter and d depth of shallower goal Pupite repeated nodes it remains comparable Space Complexity: The space Complexity of of Dis old where didepth of enallower goal Driv only stores a stude portroper out rode to current node OFID is complete for finite search tree ensuring that it will and a solution if it exists Ofio is optimal when all edges (00 costs) are wifeen, ie. will pind solution with mor eighth or cost

Code:

```
#include <bits/stdc++.h>
using namespace std;
class Solution
public:
    vector<int> dfsOfGraph(int V, vector<pair<int, int>> adj[], int
goal, int maxDepth)
    {
        vector<int> bfs;
        vector<int> cost(V, INT_MAX);
        vector<bool> vis(V, false);
        vector<int> depth(V, -1);
        queue<tuple<int, int, int>> q;
        q.push({0, 0, 0});
        bool goalFound = false;
        int foundDepth = -1;
        unordered_map<int, vector<int>> nodesAtDepth;
        while (!q.empty())
        {
            int node, currentCost, currentDepth;
            tie(node, currentCost, currentDepth) = q.front();
            q.pop();
            if (currentDepth > maxDepth)
                continue; // Skip nodes beyond maxDepth
            if (vis[node])
                continue;
            vis[node] = true;
            bfs.push_back(node);
            cost[node] = currentCost;
            depth[node] = currentDepth;
            nodesAtDepth[currentDepth].push_back(node);
            if (node == goal)
            {
                goalFound = true;
                foundDepth = currentDepth;
            for (auto it : adj[node])
                int neighbor = it.first;
                int edgeCost = it.second;
                if (!vis[neighbor])
                    q.push({neighbor, currentCost + edgeCost,
currentDepth + 1});
                }
```

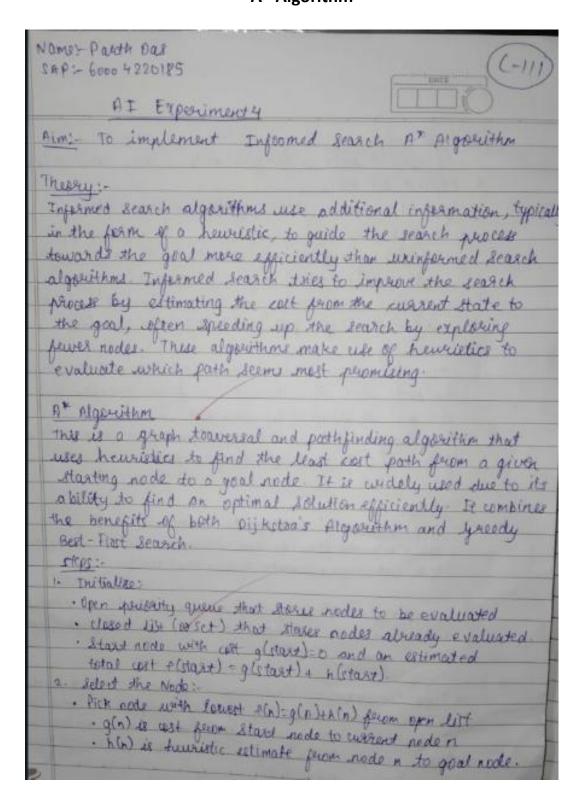
```
for (int d = 0; d <= maxDepth; ++d)</pre>
    cout << "At depth " << d << ":" << endl;</pre>
    set<int> visitedNodes;
    for (int i = 0; i <= d; ++i)
         for (int node : nodesAtDepth[i])
             visitedNodes.insert(node);
    }
    cout << "Visited nodes: ";</pre>
    for (int node : visitedNodes)
         cout << node << " ";</pre>
    cout << endl;</pre>
    cout << "Unvisited nodes: ";</pre>
    vector<int> unvisitedNodes;
    for (int i = 0; i < V; ++i)
    {
         if (visitedNodes.find(i) == visitedNodes.end())
         {
             unvisitedNodes.push_back(i);
    sort(unvisitedNodes.rbegin(), unvisitedNodes.rend());
    for (int node : unvisitedNodes)
    {
         cout << node << " ";</pre>
    cout << endl;</pre>
    cout << "Path cost to goal: ";</pre>
    if (cost[goal] != INT_MAX)
        cout << cost[goal] << endl;</pre>
    }
    else
    {
        cout << "Goal not reachable" << endl;</pre>
    cout << endl;</pre>
if (goalFound)
    cout << "Goal found at depth: " << foundDepth << endl;</pre>
```

```
else
        {
            cout << "Goal not found within depth limit" << endl;</pre>
        return bfs;
   void printGraph(int V, vector<pair<int, int>> adj[])
    {
        cout << "Graph representation:" << endl;</pre>
        for (int i = 0; i < V; i++)
            cout << "Node " << i << ": ";</pre>
            for (auto it : adj[i])
                cout << "(" << it.first << ", " << it.second << ") ";</pre>
            cout << endl;</pre>
        }
void addEdge(vector<pair<int, int>> adj[], int u, int v, int cost)
    adj[u].push_back({v, cost});
    adj[v].push_back({u, cost});
int main()
    int V = 14;
   vector<pair<int, int>> adj[V];
   addEdge(adj, 0, 1, 2);
   addEdge(adj, 0, 2, 3);
   addEdge(adj, 1, 3, 4);
   addEdge(adj, 1, 4, 7);
   addEdge(adj, 2, 5, 8);
    addEdge(adj, 3, 7, 1);
   addEdge(adj, 2, 6, 2);
    addEdge(adj, 4, 9, 2);
   addEdge(adj, 4, 10, 3);
    addEdge(adj, 5, 11, 5);
    addEdge(adj, 11, 8, 5);
    addEdge(adj, 6, 12, 10);
    addEdge(adj, 9, 13, 1);
    Solution obj;
   obj.printGraph(V, adj);
   int goal = 13;
   int maxDepth = 4;
   vector<int> ans = obj.dfsOfGraph(V, adj, goal, maxDepth);
   return 0;}
```

Output:-

```
Graph representation:
Node 0: (1, 2) (2, 3)
Node 1: (0, 2) (3, 4) (4, 7)
Node 2: (0, 3) (5, 8) (6, 2)
Node 3: (1, 4) (7, 1)
Node 4: (1, 7) (9, 2) (10, 3)
Node 5: (2, 8) (11, 5)
Node 6: (2, 2) (12, 10)
Node 7: (3, 1)
Node 8: (11, 5)
Node 9: (4, 2) (13, 1)
Node 10: (4, 3)
Node 11: (5, 5) (8, 5)
Node 12: (6, 10)
Node 13: (9, 1)
At depth 0:
Visited nodes: 0
Unvisited nodes: 13 12 11 10 9 8 7 6 5 4 3 2 1
Path cost to goal: 12
At depth 0:
Visited nodes: 0
Unvisited nodes: 13 12 11 10 9 8 7 6 5 4 3 2 1
Path cost to goal: 12
At depth 1:
Visited nodes: 0 1 2
Unvisited nodes: 13 12 11 10 9 8 7 6 5 4 3
Path cost to goal: 12
At depth 2:
Visited nodes: 0 1 2 3 4 5 6
Unvisited nodes: 13 12 11 10 9 8 7
Path cost to goal: 12
At depth 3:
Visited nodes: 0 1 2 3 4 5 6 7 9 10 11 12
Unvisited nodes: 13 8
Path cost to goal: 12
At depth 4:
Visited nodes: 0 1 2 3 4 5 6 7 8 9 10 11 12 13
Unvisited nodes:
Path cost to goal: 12
Goal found at depth: 4
```

Name:Parth Das Sap: 60004220185 Roll No: C-111 Al Experiment 4 A* Algorithm



3. Goal Test · If current nock is the goal node, seturn solution (path & cost) 4. Expand the Node: · Generate the neighbors of the current node · For each neighbor: · Calculate of (neighbox) = of arrient) + cost (current neighbox) · update f(neighboo) = q(neighboo) + h (neighbox) · If this neighbor has not been evaluated before (00 if it effects a sheaper path), add it to the open list. 5 Repeat -. more the current node to the closed list and continue selecting the node with the smallest f(n) from the open list until the goal is reached as spen list is empty. · Once the goal node is swached, backtoack through the parents of the nodes to constauct the final parts. Tentary: · templeteness - + is complete meaning it will find a solution is one exists, as long as the heuristic hin) is admissible Idea not overestimate the actual cost. · Optimality - At is optimal of the housistic as admissible and consistent. It quarantees the shortest park (reast cure) Fine complexity: The same complexity of no depends on the boarching factor b, depth of solution d, and the hewastice accuracy in weest last, it may explore a large portion of the graph, leading to expenential time

```
· Space complexity: At stores all explained nodes in memory
       so its space complainty is also O(h'd). Memory requirement
        are typically a limiting factor of A"
       Houristic Function:
       . Adminible heuristic: The heuristic h(n) is admissible if it
       never overestimates the cost to reach the goal. For eg.
      straight line distance is an admissible he writer for portafinding
       peroblems in grands in graphs.
       · sonsistent huristic - A heuristic is consistent if , fer every
       node a and its neighbor n', the estimated cost of searking
       the goal frame one no greater than the cost of swaching
       the neighbor i plus the estimated cost of scharling
        the gal from a
                  h (n) < cost (n,n') + h (n')
       Limitorious:
        POLERY.
       memory Usage: At sequises storing all explored nodes,
       what can result in high enemony usage for large
       Conclusion :-
       while At is complete and quaranteer optimal solution, its time
      and space complexity son become a limiting factor, especially in
      large search spaces, due to its requirement to store all
     explored nodes Dogite this, A* Gramains on of the most
      efficient search algorithms for a variety of posts finding and
1919 geoph soaversal publims when the heurestic function
                                 FOR EDUCATIONAL USE
```

Code:-

```
#include<bits/stdc++.h>

using namespace std;
// Node structure to store the node, gCost, hCost, and fCost
struct Node {
    string name;
    int gCost, hCost, fCost;
    Node(string name_, int gCost_, int hCost_)
        : name(name_), gCost(gCost_), hCost(hCost_) {
        fCost = gCost + hCost;
    }
    // Comparator for priority queue (min-heap based on fCost)
    bool operator<(const Node& other) const {</pre>
```

```
return fCost > other.fCost;
   }
};
// A* algorithm implementation
void AStar(unordered map<string, vector<pair<string, int>>> graph,
unordered_map<string, int> heuristics, string start, string goal) {
   // Priority queue for open nodes (min-heap)
   priority_queue<Node> openList;
    // Cost map to store gCost of nodes (cost from start)
    unordered_map<string, int> gCostMap;
   // Parent map to track the path (backtracking from goal to start)
    unordered_map<string, string> parentMap;
   // Set gCost of all nodes to infinity initially
   for (const auto& node : graph) {
        gCostMap[node.first] = numeric limits<int>::max();
   // Closed set to track visited nodes
    unordered_map<string, bool> closedSet;
    openList.push(Node(start, 0, heuristics[start]));
    gCostMap[start] = 0;
    parentMap[start] = ""; // Start node has no parent
    int iteration = 0;
   while (!openList.empty()) {
        iteration++;
        Node current = openList.top();
        openList.pop();
        cout << "Iteration " << iteration << ": Processing node " <<</pre>
current.name << " with fCost: " << current.fCost << " (gCost: " <<
current.gCost << ", hCost: " << current.hCost << ")\n";
        if (current.name == goal) {
            cout << "Goal '" << goal << "' reached with total cost: "</pre>
<< current.gCost << endl;</pre>
            stack<string> path;
            string node = goal;
            while (node != "") {
                path.push(node);
                node = parentMap[node];
            cout << "Path: ";</pre>
            while (!path.empty()) {
                cout << path.top();</pre>
                path.pop();
                if (!path.empty()) cout << " -> ";
            cout << endl;</pre>
            return;
```

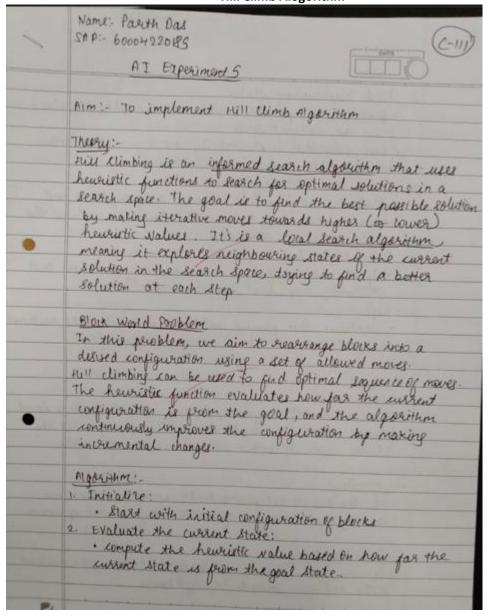
```
closedSet[current.name] = true;
        for (const auto& neighbor : graph[current.name]) {
            string neighborName = neighbor.first;
            int pathCost = neighbor.second;
            if (closedSet[neighborName]) continue;
            int newGCost = current.gCost + pathCost;
            if (newGCost < gCostMap[neighborName]) {</pre>
                gCostMap[neighborName] = newGCost;
                int hCost = heuristics[neighborName];
                // Push the neighbor to the priority queue
                openList.push(Node(neighborName, newGCost, hCost));
                parentMap[neighborName] = current.name;
                cout << " Adding neighbor " << neighborName << " with</pre>
gCost: " << newGCost << ", hCost: " << hCost << ", fCost: " << newGCost
 hCost << endl;
        }
        cout << " Priority Queue: ";</pre>
        priority queue < Node > tempQueue = openList;
        while (!tempQueue.empty()) {
            Node node = tempQueue.top();
            tempQueue.pop();
            cout << node.name << "(fCost: " << node.fCost << ") ";</pre>
        }
        cout << endl << endl;</pre>
    cout << "Goal '" << goal << "' cannot be reached.\n";</pre>
int main() {
    unordered_map<string, int> heuristics = {
        {"A", 10}, {"B", 8}, {"C", 5}, {"D", 7}, {"E", 3},
        {"F", 6}, {"G", 5}, {"H", 3}, {"I", 1}, {"J", 0}
    };
   unordered_map<string, vector<pair<string, int>>> graph = {
        {"A", {{"B", 6}, {"F", 3}}},
        {"B", {{"A", 6}, {"C", 3}, {"D", 2}}},
        {"C", {{"B", 3}, {"D", 1}, {"E", 5}}},
        {"D", {{"B", 2}, {"C", 1}, {"E", 8}}},
        {"E", {{"C", 5}, {"D", 8}, {"I", 5}, {"J", 5}}},
        {"F", {{"A", 3}, {"H", 7}, {"G", 1}}},
        {"G", {{"F", 1}, {"I", 3}}},
        {"H", {{"F", 7}, {"I", 2}}},
        {"I", {{"E", 5}, {"G", 3}, {"H", 2}, {"J", 3}}},
        {"J", {{"E", 5}, {"I", 3}}}
    };
    string start = "A";
    string goal = "J";
   AStar(graph, heuristics, start, goal);
```

```
return 0;
}
```

Output:-

```
Iteration 1: Processing node A with fCost: 10 (gCost: 0, hCost: 10)
  Adding neighbor B with gCost: 6, hCost: 8, fCost: 14
  Adding neighbor F with gCost: 3, hCost: 6, fCost: 9
  Priority Queue: F(fCost: 9) B(fCost: 14)
Iteration 2: Processing node F with fCost: 9 (gCost: 3, hCost: 6)
  Adding neighbor H with gCost: 10, hCost: 3, fCost: 13
  Adding neighbor G with gCost: 4, hCost: 5, fCost: 9
  Priority Queue: G(fCost: 9) H(fCost: 13) B(fCost: 14)
Iteration 3: Processing node G with fCost: 9 (gCost: 4, hCost: 5)
  Adding neighbor I with gCost: 7, hCost: 1, fCost: 8
  Priority Queue: I(fCost: 8) H(fCost: 13) B(fCost: 14)
Iteration 4: Processing node I with fCost: 8 (gCost: 7, hCost: 1)
  Adding neighbor E with gCost: 12, hCost: 3, fCost: 15
  Adding neighbor H with gCost: 9, hCost: 3, fCost: 12
  Adding neighbor J with gCost: 10, hCost: 0, fCost: 10
  Priority Queue: J(fCost: 10) H(fCost: 12) H(fCost: 13) B(fCost: 14) E(fCost: 15)
Iteration 5: Processing node J with fCost: 10 (gCost: 10, hCost: 0)
Goal 'J' reached with total cost: 10
                                                                             (X) The lan
Path: A \rightarrow F \rightarrow G \rightarrow I \rightarrow J
```

Name: Parth Das SAP:- 60004220185 Roll No: C-111 AI EXperiment 5 Hill Climb Alogorithm



3. Find Neighbouring States: · yenerate all pessible valid moves (neighbour states) from the current state. · Each neighbour is breated by moving a block either onto the table or onto another block. 4. Charge the Best Neighbour: · Evaluate the heuristic value of each neighbour and choose the one that reduces the heuristic (beings the state closes to the goal) 5. Move to the Best Neighbour. · If the best reighbour is an improvement, more to it otherwise, stop (local optimum reached). 6. Repeat · Continue this process until the current state matches the goal state or a local optimum is reached. Features: · completenes:-Hill climb is not complete, it may get stuck in a local optimum where no neighbouring states improve the sola solution, even if better solution exists also where · Optimality:-Hill dembing is net guaranteed to be optimal because it may find local optima ratherthan global optimum · space complexity:-O(1) as the algorithm only needs to store the current State and the immediate neighbourse

```
3. Find Neighbouring States:
   · generate all pessible valid moves (neighbour states)
   from the current state.
   · Each neighbour is breated by moving a block either
   onto the table or onto another block.
4. Charge the Best Neighbour:
   · Evaluate the heuristic value of each neighbour and
   choose the one that reduces the heuristic (beings the
    state closes to the goal)
 5. Move to the Best Neighbour:
    · If the best neighbour is an improvement, more to it.
     otherwise, stop (local optimum reached).
 6. Repeat
    · Continue this process until the current state matches
    the goal state or a local optimum is reached
  Features:
  · completenes:
  Hill ilimb is not complete, it may get stuck in a local
  optimum where no neighbouring states improve the solo
   solution even if better solution exists also where
  · Optimality:-
   till dembing is not guaranteed to be optimal because
    it may find local optima ratherthan global optimum
   · space complexity:
   O(1) as the algorithm only needs to store the current
    State and the immediate neighbourse
```

Code:-

```
from copy import deepcopy

def get_position_index(final):
    fs = final[0]
    p = dict()
    for i in range(len(fs)):
        if i == 0:
            p[fs[i]] = None
        else:
            p[fs[i]] = fs[i - 1]
    return p

def get_next_states(s, visited):
    states = []
    for i in range(len(s)):
        if len(s[i]) == 0:
```

```
continue
        for j in range(len(s)):
            if i == j:
                continue
            ns = deepcopy(s)
            ns[j].append(ns[i].pop())
            if ns not in visited:
                states.append(ns)
    return states
def get_h(s, f):
    c = 0
    for x in s:
        for i in range(len(x)):
            if i == 0:
                c += 1 if f[x[i]] is None else -1
            else:
                c += 1 \text{ if } f[x[i]] == x[i - 1] \text{ else } -1
    return c
# Count of maximum platforms to be used
# Taking user input in a way that works for Python 3
initial_input = input("Enter the initial state (e.g., 'A B C'):
").split()
final_input = input("Enter the final state (e.g., 'A B C'): ").split()
initial = (initial_input, [], [])
final = (final_input, [], [])
fi = get_position_index(final)
visited = []
not visited = []
current = deepcopy(initial)
while get_h(current, fi) != get_h(final, fi):
    moves += 1
    visited.append(current)
    # Generate possible next states
    ns = get_next_states(current, visited)
    ns_h = []
    # Calculate heuristic value for each state
    for s in ns:
        ns_h.append((get_h(s, fi), s))
    # Sort by the heuristic value
    ns_h = sorted(ns_h, key=lambda x: x[0])
    try:
        # Get the best state to explore
        c, current = ns_h.pop()
        # Add the remaining to not visited for future use
        not_visited = not_visited + ns_h
    except IndexError:
```

```
# If there are no new next states, backtrack
    print('No new unique states, backtracking...')
    not_visited = sorted(not_visited, key=lambda x: x[0])
    c, current = not_visited.pop()
    print("Current heuristic value: " + str(c), current)
print('Found solution in ' + str(moves) + ' moves')
```

Output:-

```
Enter the initial state (e.g., 'A B C'): A B C D
Enter the final state (e.g., 'A B C'): B D C A

Current heuristic value: -4 (['A', 'B', 'C'], [], ['D'])

Current heuristic value: -2 (['A', 'B'], [], ['D', 'C'])

Current heuristic value: 0 (['A'], ['B'], ['D', 'C'])

Current heuristic value: 2 ([], ['B'], ['D', 'C', 'A'])

Current heuristic value: 2 (['B'], [], ['D', 'C', 'A'])

Current heuristic value: 0 (['B'], ['A'], ['D', 'C'])

Current heuristic value: 0 (['B', 'A'], ['C'], ['D'])

Current heuristic value: 0 (['B'], ['C', 'A'], ['D'])

Current heuristic value: 2 (['B', 'D'], ['C', 'A'], [])

Current heuristic value: 2 (['B', 'D'], ['C'], ['A'])

Current heuristic value: 2 (['B', 'D'], ['C'], ['A'])

Current heuristic value: 4 (['B', 'D', 'C', 'A'], [], [])

Found solution in 13 moves
```

Name: Parth Das Sap: 60004220185

Batch: C-22 Roll no: C-111

Subject : Artificial Intelligence

EXPERIMENT NO. 6

Topic: Implementation of Genetic Algorithm

THEORY:

	Parth Das (C	-111)
-	6000 4220 MS C2-2	
	AI Experiment-6	
	Aim: To implement genetic Algorithm	-
Markey.	Theory:-	
	A yenetic Algorithm (CAA) is a type of comb. Optimizat algorithm imported by the process of natural selection biological evalution. It is widely used to solve optimization	r in miration
6	and search phoblems, pasticularly when the solution wast and complex. Get belong to the family of evolution algorithms, which generate solutions to optimization puring techniques inspired by natural selection.	nary
	tey Concepts:	
	to be the problem and is often knowled as chromosome (stain array of variables)	solution gos
	2 Filmer Furntion'-	
0	on a given task, whether it's optimization or classifican	
	current population to approduce.	
-2	The individual with higher fitness scores are more likely to selected for the next generation (just like "survival or the	be fitter!
	· Reulette wheel Selection: Poobability of selection is proportion · Townsment selection: A set of individuals is shown none and the best among them is ellected.	al to filme dently,

4. Course over (Recombination):	
-> The conscious operation takes	2 parent individuals and compines
parts of their chromosomes:	to produce offering new solutions.
-> The idea is that by combini	ng sucressful toats from two parents,
you can releate offering in	eith even better traite.
8 Autations	
an individual.	handom changes in the encomprome of
-> Mutation ensures divibility in	the repulation and helps awoid local
optime by exploring new a	the population and helps avoid local was of the solution space. If
6. Terminator:	
-> The also supports the process.	of selection, crossover and mutation
	the termination condition is met.
-> This could be a fixed number	of generations, mooning a sime limit
or achieving a contain fitness -	hrepholo.
optimality: Not good anteed, o	pproximates near optimal solutions.
completeness: Not complete, as	by not find a solution in finite time. &
Time Completity: 0 (A * P * + f (n	1) when Cregenerations, & Population sizes
f(n): fixma	es fundion cost
space comploxity. O(PAN) w	shere for populations ire no chapmany langua.
Conclusion: -	The state of the s
In the content of A I, year	the Algorithms are useful for problems o
where search apole is dust an	deposity diderstand, and they offer a the
However the dentity could	utions therough an evolutionary approved
in conjunction with other	plinization techniques on
gall real wested At applicant	ul turing and are typically used principalisms and principalisms and principalisms and principalisms.
	DUCATIONAL USE

CODE:

import random

Parameters

 $POPULATION_SIZE = 100$

GENOME_LENGTH = 10 # Number of bits (for each dimension)

 $MUTATION_RATE = 0.1$

GENERATIONS = 100

Function to convert binary to decimal within a specific range def binary_to_decimal(binary):

```
return int(binary, 2) / (2 ** GENOME LENGTH - 1) * 5.12 - 2.56 # Scale to [-2.56, 2.56]
# Sphere function
def sphere(genome):
  x = [binary to decimal(genome[i:i + GENOME LENGTH])] for i in range(0, len(genome),
GENOME LENGTH)]
  return sum(xi^{**}2 \text{ for } xi \text{ in } x)
# Generate initial population
def generate population(size):
  return [".join(random.choice('01') for _ in range(GENOME_LENGTH * 2)) for _ in range(size)] #
Double genome length for two dimensions
# Selection: Tournament selection
def selection(population):
  tournament = random.sample(population, 5)
  return min(tournament, key=sphere) # Minimize the Sphere function
# Crossover: Two-point crossover
def two point crossover(parent1, parent2):
  point1 = random.randint(1, len(parent1) - 2)
  point2 = random.randint(point1 + 1, len(parent1) - 1)
  child1 = parent1[:point1] + parent2[point1:point2] + parent1[point2:]
  child2 = parent2[:point1] + parent1[point1:point2] + parent2[point2:]
  return child1, child2
# Mutation: Flip a bit
def mutate(genome):
  return ".join(bit if random.random() > MUTATION RATE else random.choice('01') for bit in
genome)
# Main genetic algorithm
```

def genetic_algorithm():

population = generate population(POPULATION SIZE)

```
for generation in range(GENERATIONS):
    best fitness = min(sphere(genome) for genome in population)
    print(f"Generation {generation}: Best fitness = {best fitness}")
    new population = []
    while len(new population) < POPULATION SIZE:
      parent1 = selection(population)
      parent2 = selection(population)
      child1, child2 = two point crossover(parent1, parent2)
      new population.extend([mutate(child1), mutate(child2)])
    population = new population
  # Best solution
  best genome = min(population, key=sphere)
  print(f''Best genome: {best genome}, Best fitness: {sphere(best genome)}'')
if name == " main ":
  genetic algorithm()
OUTPUT:
Generation 0: Best fitness = 0.18342056837412085
Generation 1: Best fitness = 0.10326408919389689
Generation 2: Best fitness = 0.005974162588901006
Generation 3: Best fitness = 0.0024172188252786
Generation 4: Best fitness = 0.00031311124679774926
Generation 5: Best fitness = 1.2524449871909348e-05
Generation 6: Best fitness = 0.00021291564782247003
Generation 7: Best fitness = 6.262224935954675e-05
```

Generation 8: Best fitness = 6.262224935954675e-05 Generation 9: Best fitness = 1.2524449871909348e-05 Generation 10: Best fitness = 1.2524449871909348e-05 Generation 11: Best fitness = 1.2524449871909348e-05 Generation 12: Best fitness = 1.2524449871909348e-05 Generation 13: Best fitness = 1.2524449871909348e-05 Name: Parth Das Sap: 60004220185 Roll No:- C-111

AI Experiment 7:- Perceptron Learning

Parth Das 60004220125 C-22	(-11)
AI Assignment - 7	
	,
Ain :- To implement Preception Learning	
The Perceptoen is one of the simplest types notworks used for binary classification of first introduced by Frank Rosenblatt in 195	s and ferm the
learning algorithm is an iterative process weights of a linear classifier	is to supdate the
· Inputs: Feature values (+9 x, 70, 10, whights: largerients associated with the search of we, we, was	orputs 3 the track see
· Bias: the additional parameter (b) that boundary to be shifted. · Activation Function - A step function that (binary output) based on whether the will write exceeds a shreshold.	ecopute either los-1
And the same of th	
where output y is either 100-1 depend	ligh sleepath bould
1 Instintuation: Start with small reador and hims b.	
and apply the activation function to class	late weighted sum by the sample
FOREDUCATIONAL	

3. Update Weights: · If the prediction matches the actual label, do nothing · To the prediction is inspired, adjust the weights and bigg w: wi + 2 x (y- y) *2; b = b+ 2 * (y- 5) where a is the learning hate y is actual label, is provided label 4. Repeat the process for a specified number of iterations or until convergence. Fratures: -> completeness = only complete if dataset is linearly separable. In opportunity, it can jud a solution if there raists a linear boundary that separates the down into two classes -> optimality: This alforithm deen't grananter an opinal solution in terms of minimizing dasification execuse. -> Time complexity: Otton & d) where To no of iterations n-number of toaining samples d'number of features per sample. -> Space complexity: o(n +d) if weights o(d) and Bias(o(1)) with input data O(n+d) Conclusion: This algorithm is foundation in ML particularly for binary lines classification problems. It efficiently finds a decision boundary when the data is linearly separable, though it does not 2110124 quarantee optimality in terms of maximizing the margin.

Code:-

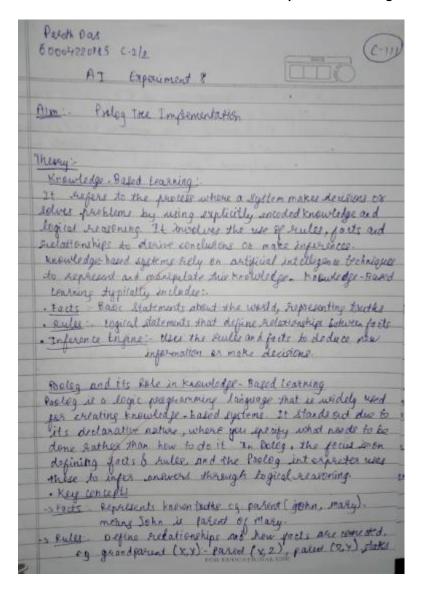
```
import numpy as np
import pandas as pd
class PerceptronNetwork:
   def __init__(self, input_size):
        self.weights = np.random.randn(input_size)
        self.bias = np.random.rand(1)
    def predict(self, x):
        weighted_sum = np.dot(self.weights, x) + self.bias
        return 1 if weighted_sum >= 0 else 0
    def train(self, X, y, learning_rate, epochs):
        results = []
        for epoch in range(epochs):
            for i in range(len(X)):
                prediction = self.predict(X[i])
                error = y[i] - prediction
                self.weights += (learning_rate * error * X[i])
                self.bias += (learning_rate * error)
                # Collecting results for this epoch
                results.append({
                    "Sample": i,
                    "Input": X[i],
                    "Weights": self.weights.copy(),
                    "Epoch": epoch + 1,
                    "Bias": self.bias[0],
                    "Prediction": prediction,
                    "Error": error,
                    "Learning Rate": learning rate
                })
        # Convert results to a DataFrame for better visualization
        df_results = pd.DataFrame(results)
        print(df_results)
# Create a perceptron network and train it
perceptron_network = PerceptronNetwork(input_size=2)
X = np.array([[0, 0], [0, 1], [1, 0], [1, 1]])
y = np.array([0, 1, 1, 1])
# Training with 1 epoch for initial results
perceptron_network.train(X, y, 0.1, 1)
# Predictions after initial training
for x in X:
    prediction = perceptron_network.predict(x)
    print(f"Input: {x}, Perceptron Prediction: {prediction}")
# Further training for 10 epochs
perceptron_network.train(X, y, 0.1, 10)
# Final predictions after training
for x in X:
   prediction = perceptron_network.predict(x)
   print(f"Input: {x}, Perceptron Prediction: {prediction}")
```

Output:-

Samp	ole Input		Weights	Epoch	Bias F	rediction	Error	Learning Rate
0	0 [0, 0]	[0.2104908952127543.	-0.32273565879466526]		0.586074	1	-1	0.1
1	1 [0, 1]		-0.32273565879466526]		0.586074	ī	0	0.1
2	2 [1, 0]		-0.32273565879466526]		0.586074	1	0	0.1
3	3 [1, 1]		-0.32273565879466526]		0.586074	1	ē	0.1
		eptron Prediction: 1						
		eptron Prediction: 1						
		eptron Prediction: 1						
		eptron Prediction: 1						
	ple Input		Weights	Epoch	Bias	Prediction	Error	Learning Rate
0	0 [0, 0]		3, -0.32273565879466526			1		0.1
1	1 [0, 1]		3, -0.32273565879466526		0.486074	1	0	
2	2 [1, 0]		3, -0.32273565879466526		0.486074	1	0	
3	3 [1, 1]		3, -0.32273565879466526		0.486074	1	0	0.1
4	0 [0, 0]	0.210490895212754	3, -0.32273565879466526	2	0.386074	1	-1	0.1
5	1 [0, 1]	[0.210490895212754]	3, -0.32273565879466526	2	0.386074	1	0	0.1
6	2 [1, 0]	[0.210490895212754]	3, -0.32273565879466526	2	0.386074	1	0	0.1
7	3 [1, 1]	[0.210490895212754]	3, -0.32273565879466526	2	0.386074	1	0	0.1
8	0 [0, 0]	[0.210490895212754]	3, -0.32273565879466526	3	0.286074	1	-1	0.1
9	1 [0, 1]	[0.210490895212754]	3, -0.22273565879466525	3	0.386074	e	1	0.1
10	2 [1, 0]	[0.210490895212754	, -0.22273565879466525	3	0.386074	1	0	0.1
11	3 [1, 1]	[0.210490895212754	3, -0.22273565879466525	3	0.386074	1	0	0.1
12	0 [0, 0]	[0.210490895212754	3, -0.22273565879466525	4	0.286074	1	-1	0.1
13	1 [0, 1]	[0.210490895212754	3, -0.22273565879466525	4	0.286074	1	0	0.1
14	2 [1, 0]	[0.210490895212754	3, -0.22273565879466525	4	0.286074	1	0	0.1
15	3 [1, 1]	[0.210490895212754	3, -0.22273565879466525	4	0.286074	1	0	0.1
16	0 [0, 0]	[0.2104908952127543	, -0.22273565879466525]	5	0.186074	1	-1	0.1
17	1 [0, 1]	[0.2104908952127543	, -0.12273565879466525]	5	0.286074	0	1	0.1
18	2 [1, 0]	[0.2104908952127543	, -0.12273565879466525]	5	0.286074	1	0	0.1
19	3 [1, 1]		, -0.12273565879466525]	5		1	0	0.1
20	0 [0, 0]	[0.2104908952127543	, -0.12273565879466525]	6	0.186074	1	-1	0.1
21	1 [0, 1]		, -0.12273565879466525]	6		1	0	0.1
22	2 [1, 0]		, -0.12273565879466525]	6		1	0	0.1
23	3 [1, 1]		, -0.12273565879466525]	6		1	0	0.1
24	0 [0, 0]		, -0.12273565879466525]	7		1	-1	0.1
25	1 [0, 1]		-0.022735658794665242]		0.186074	0	1	0.1
26	2 [1, 0]	TABLE 8 FOR CONTRACTOR (\$1.00)	-0.022735658794665242]	7		1	0	0.1
27	3 [1, 1]		-0.022735658794665242]	7		1	0	0.1
28	0 [0, 0]		-0.022735658794665242]	8		1	-1	0.1
29	1 [0, 1]		-0.022735658794665242]	8		1	0	0.1
30	2 [1, 0]		-0.022735658794665242]	8		1	0	0.1
31	3 [1, 1]		-0.022735658794665242]	8		1	0	0.1
32	0 [0, 0]		-0.022735658794665242]		-0.013926	1	-1	0.1
33	1 [0, 1]		3, 0.07726434120533476]	9		0	1	0.1
34	2 [1, 0]		3, 0.07726434120533476]	9	0.086074	1	0	0.1
35	3 [1, 1]		3, 0.07726434120533476]	9	0.086074	1	0	0.1
36	0 [0, 0]		3, 0.07726434120533476]		-0.013926	1	-1	0.1
37	1 [0, 1]		3, 0.07726434120533476]		-0.013926	1	0	0.1
38	2 [1, 0]		3, 0.07726434120533476]		-0.013926	1	0	0.1
39	3 [1, 1]		3, 0.07726434120533476]	10	-0.013926	1	0	0.1
	-	eptron Prediction: 0						
		eptron Prediction: 1						
		eptron Prediction: 1						
input:	[1 1], Perc	eptron Prediction: 1						

Name: Parth Das SAP:- 60004220185 Roll No: C-111

AI EXperiment 8 Prolog Tree



that X to a quandparent of Y of X is parent of z and 2 is garent of Y. - Duries: - Allow users to ask questions based on the facts and sules and Isolog will attempt to provide answers e.g. 2 - grandparent (john, mary). · Posty Take (search Tau) A Prolog Tree is a vibualization of thelogic that Polog was to solve queries. It suggestents the Step by- step process by which Pooling wies to satisfy a guery by to aversing possible paths of foils and scular. This seasch persons is dypically conducted using backtoocking. How Pooley Take Wooks: · Outry Initiation: when a query is posed begins to search for a solution starting from the scot of the tree. 2 Groal matching - It matches the goal (query) against the facts and rule defined in the knowledge base. 3. Boanching - The stree boanches but at each step as Prolog trees different rule or facts that night satisfy the inverse goal. 4. Backtracking - If a boarch leads to a dead and I failure to satury the goal). Poolog backtracks to the previous level and tries another branch. Solution Path - when Psolog finds a path that satisfies the goal, it returns the solution.

Papperotice. optimelity: Not optimal by default uses dopth first search with backtoacking finds the flost solution without necessarily being the best one. completeness: complete if the seasch space is finite and peoperly structured for get stuck in infinite loops Time complexity: Exponential in well case O(b) Highly sometime to size of the search b boarding foctors d: dopin of search tree. Space complexity: Linear in depth of search trove old). more memory afficient than RTS optimization. To impliance optimality in Poelog, you can add continuit or use stockeries like best post search or houridice, which are not part of standard "Poolog but can be implemented. Conclusion: while Pooling is a privilegal for knowledge-based systems due to its logical reasoning capabilities, its default search method (depth piset with backtocking) may not always se optimal or efficient. Understanding the limitations in terms of optimality, completeness, and the computational lest hops an designing more efficient lugic programs or choosing alternatives search stradegies for specific published domains

Relationships:-

```
% Male members
male(john).
male(mike).
male(alex).
male(tom).
male(jake).
male(eric).
% Female members
female(susan).
female(linda).
female(kate).
female(anna).
female(sophia).
% Defining parent relationships
parent(john, mike). % John is the parent of Mike
parent(john, linda). % John is the parent of Linda
parent(susan, mike). % Susan is the parent of Mike
parent(susan, linda). % Susan is the parent of Linda
                       % Mike is the parent of Alex
parent(mike, alex).
                       % Mike is the parent of Kate
% Linda is the parent of Tom
parent(mike, kate).
parent(linda, tom).
parent(linda, sophia). % Linda is the parent of Sophia
                       % Alex is the parent of Jake
parent(alex, jake).
parent(kate, anna).
                       % Kate is the parent of Anna
```

```
Rules:-
```

```
% Grandfather and Grandmother
grandfather(X, Y) :- male(X), parent(X, Z), parent(Z, Y).
grandmother(X, Y) :- female(X), parent(X, Z), parent(Z, Y).
% Father and Mother
father(X, Y) :- male(X), parent(X, Y).
mother(X, Y) :- female(X), parent(X, Y).
% Son and Daughter
son(X, Y) :- male(X), parent(Y, X).
daughter(X, Y) :- female(X), parent(Y, X).
% Uncle and Aunt
uncle(X, Y) :- male(X), parent(Z, Y), sibling(X, Z).
aunt(X, Y) :- female(X), parent(Z, Y), sibling(X, Z).
% Sibling (brother or sister)
sibling(X, Y) :- parent(Z, X), parent(Z, Y), X = Y.
% Cousin (cousin brother or cousin sister)
cousin(X, Y) := parent(Z, X), parent(W, Y), sibling(Z, W).
```

Queries:-

1. Who is John's grandson?

```
grandfather(john, X), male(X)

X = alex

X = tom
```

2. Who are the siblings of Kate?



3. Who is the cousin brother of Sophia?

```
cousin(sophia, X),male(X).
X = alex
```

4. Who is the cousin sister of Sophia?

```
cousin(sophia, X),female(X).

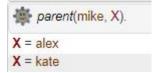
X = kate
```

5. Who is the uncle of Tom?

```
wncle(X, tom).
X = mike
```

6. Who is Anna's grandfather?

7. Who are the children of Mike?

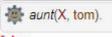


8. Is Jake a son of Alex?



true

9. Who is the aunt of Tom?



false

10. Is Linda the grandmother of Anna?



Names: SapId:

PARTH DAS 60004220185

Subject: AI Batch: C2-2

EXPERIMENT NO. 9

Topic: Game Design

GAME TITLE:

PacMan Lite

IDEA/DESCRIPTION:

The PacMan Lite game is a simplified version of the classic PacMan game. The player controls a PacMan character and navigates through a maze, collecting pellets while avoiding ghosts. The objective is to collect all the pellets in the maze before the player's lives are depleted.

The core elements of the game include:

<u>Maze Generation</u>: The game features a pre-defined maze layout, with walls, pellets, and spawn locations for the PacMan character and ghosts.

<u>Player Control</u>: The player can control the PacMan character using the arrow keys, moving it in four directions (up, down, left, right) to navigate the maze.

<u>Pellet Control:</u> The player must collect all the pellets scattered throughout the maze to complete the level.

Ghost Avoidance: The player must avoid colliding with the ghosts, which will cause the player to lose a life. The ghosts move autonomously using a simple pathfinding algorithm.

RULES:

- 1. The player starts with a set number of lives (e.g., 3).
- 2. The player can move the PacMan character in four directions (up, down, left, right) using the arrow keys.
- 3. The player must collect all the pellets in the maze to complete the level.
- 4. If the PacMan character collides with a ghost, the player loses a life.
- 5. If the player runs out of lives, the game is over.
- 6. The player can restart the game by clicking the "Play Again" button.

TASK ENVIRONMENT:

Property	Value
Observable	Partially
Deterministic	Deterministic
Episodic	Sequential
Static	Dynamic
Discrete	Discrete
Agents	Single

Observable: The environment is **partially observable** because the player can only see the immediate area around the PacMan character, and the positions of the ghosts are only revealed when they are in the player's field of view.

Deterministic:: The environment is **deterministic** because the maze layout, pellet positions, and ghost movements are predetermined and do not change during the game.

Episodic: The task is **sequential** because the player's actions (moving the PacMan character) directly affect the game state and the outcome of the current episode (level).

Static: The environment is **dynamic** during the game, as the PacMan character and ghosts move and the pellets are collected, changing the state of the game.

Discrete:: The environment is **discrete**, as the maze is divided into a grid of cells, and the PacMan character and ghosts move between these cells.

Agents: There is a **single agent** (the player controlling the PacMan character) in the game

Algorithms

- 1. <u>Dijkstra's Algorithm:-</u>: This finds the shortest path from each enemy to the closest resource when the enemy's health is low. By minimizing the distance, enemies can reach resources quickly to "heal."
- 2. <u>Heuristic Function</u>: The heuristic function (Manhattan distance) provides a baseline for calculating distances.

CONCLUSION

The PacMan Lite game presents a partially observable, deterministic, and sequential environment where a single agent (the player) navigates a discrete grid-based maze, collecting pellets while avoiding autonomous ghosts. The game showcases the implementation of pathfinding algorithms for the ghosts and collision detection between game entities. The simplified nature of the game allows for an interactive and accessible experience in understanding classic arcade-style game mechanics.

Name: Parth Das Sap: 60004220185

Batch: C-22 Roll no: C-111

Subject : Artificial Intelligence

Experiment 10:- Case Study of an Al Application

Sr N o	Paper Title	Publicatio n Name (Year)	Algorithm/Techniq ue Used	Dataset (If Any)	Results Mentioned in Paper	Observatio ns
1	DetectGPT: Zero-Shot Machine- Generated Text Detection using Probability Curvature	arXiv (2023)	Probability Curvature	None specifi ed	DetectGPT distinguish es human from AI text with significant accuracy	This paper emphasize s transparen cy in Al content generation detection and has potential application s in education and journalism
2	HuggingGP T: Solving Al Tasks with ChatGPT and its Friends in Hugging Face	arXiv (2023)	Integration of ChatGPT with Hugging Face models	None specifi ed	Shows improved efficiency in performing Al tasks by leveraging existing models	Demonstrat es the versatility of combining language models with other Al tools for diverse tasks

3	Tree of Thoughts: Deliberate Problem Solving with Large Language Models	arXiv (2023)	Large Language Models with reasoning	None specifi ed	LLMs, when combined with deliberate problem-solving strategies, enhance performance in reasoning tasks	Highlights the importance of structured thought processes in Al-driven problem solving
4	Toolformer: Language Models Can Teach Themselve s to Use Tools	arXiv (2023)	Toolformer framework	None specifi ed	Language models demonstrat ed self-supervised learning to use tools independen tly	This research pushes the boundaries of self- learning models, enabling them to autonomou sly acquire and use external tools
5	A Watermark for Large Language Models	arXiv (2023)	Watermarking for LLM-generated content	None specifi ed	Proposed a technique to trace and verify Algenerated content	Vital for Al content authenticity , particularly in legal, educational , and journalistic sectors

6	Is ChatGPT a General- Purpose Natural Language Processing Task Solver?	arXiv (2023)	ChatGPT performance on diverse NLP tasks	None specifi ed	ChatGPT performs well across many NLP tasks but struggles with domain-specific tasks	Shows that while ChatGPT is versatile, domain-specific tuning is necessary for high precision tasks
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