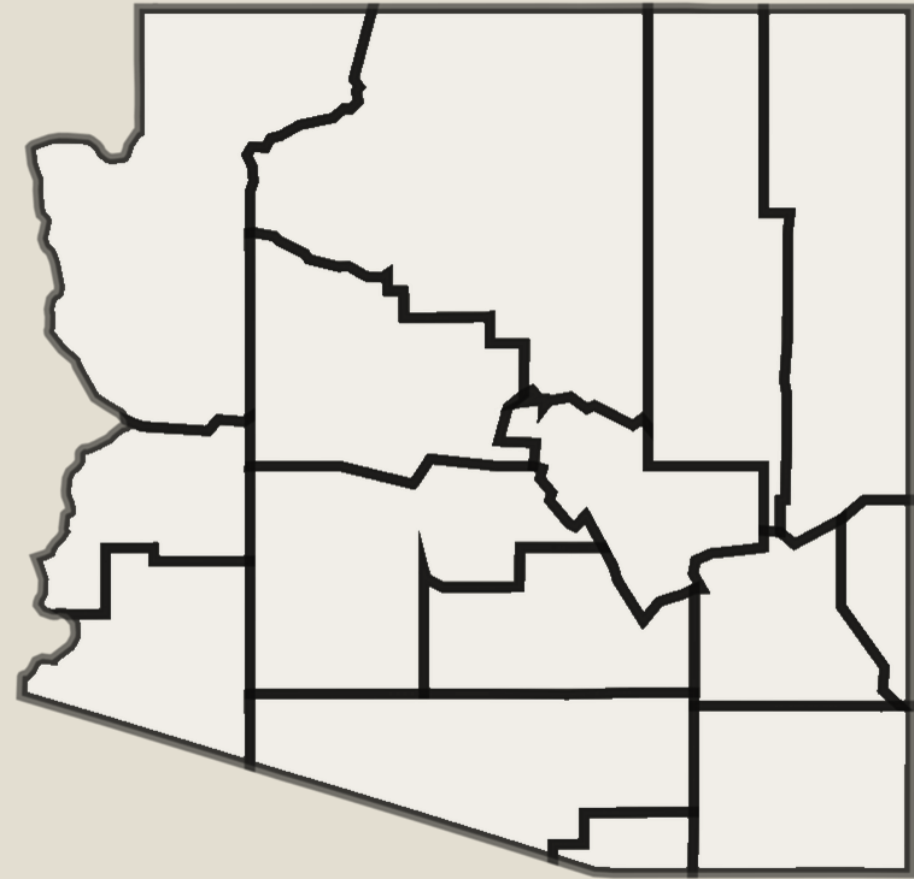


# GRAPH COLORING & SUDOKU

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# Problem 1

- solve the map coloring puzzle
  - given a map with bordered entities (e.g. countries, states, counties) what is the minimum number of colors required to color each entity such that no two adjacent entities have the same color.



# Problem2

- Solve the Sudoku puzzle

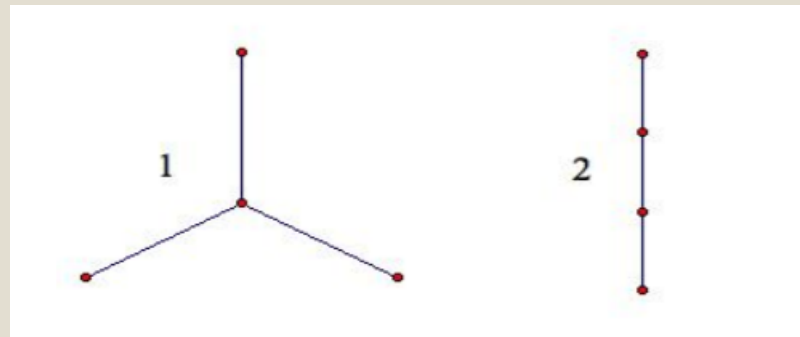
	1	7						
8				5				
2	3		4	1			7	
	8					3		9
		2		7		4		
1		3					6	
	9			3	6		5	7
				2				3
						1	9	

# Some graph theory

- **Definition (Graph)**

A (simple) graph  $G$  consists of a finite or countable vertex set  $V := V(G)$  and an edge set  $E := E(G) \subset \binom{V}{2}$

- For a vertex set  $V$ , we represent edges as  $(v, w)$ ,  $v, w \in V$ . We also write  $v \sim w$  to denote that  $(v, w) \in E$  and say  $v$  and  $w$  are adjacent.
- A very common pictorial representation of graphs is as follows : Vertices are represented as points on plane and edges are lines / curves between the two vertices.



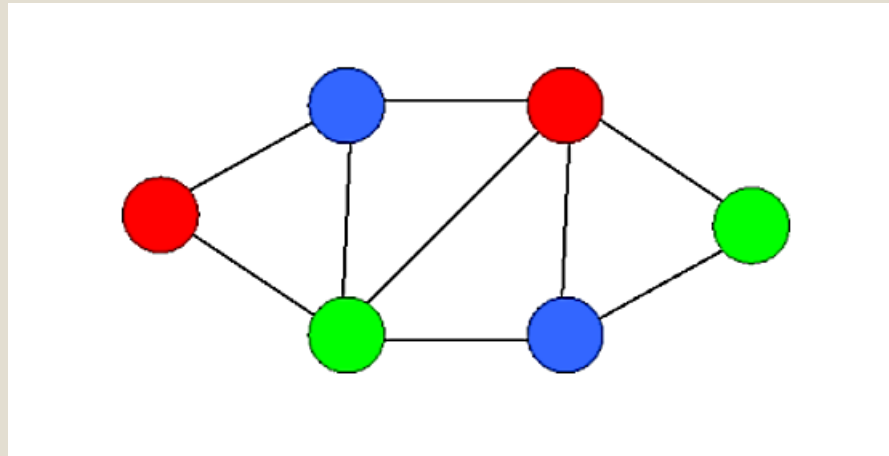
# Graph Coloring

- **Definition (Vertex Coloring of a Graph)**

A proper vertex coloring of a graph  $G$  is an assignment of colors to the vertices of the graph such that no two adjacent vertices are assigned the same color.

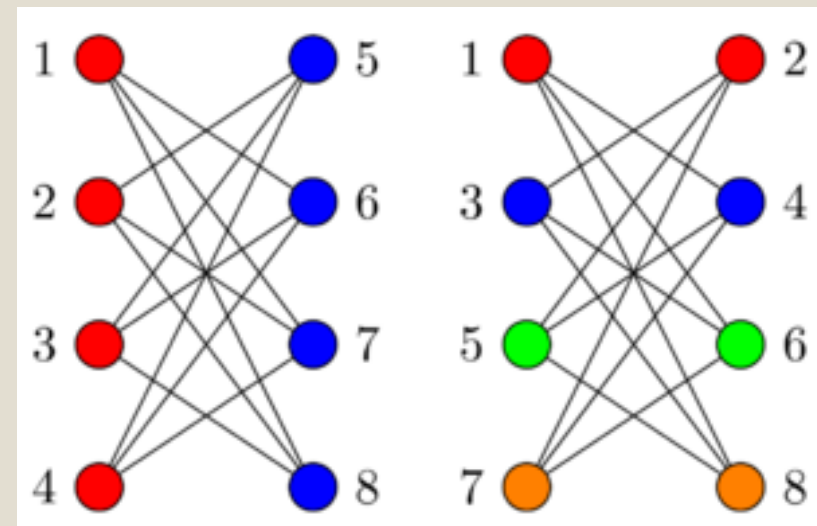
- **Definition (Chromatic Number)**

The minimal number of colors needed to properly vertex color a graph, denoted  $\chi(G)$



# Greedy coloring

- Greedy coloring is a method to color a graph  
Label vertices  $v_1, v_2, \dots, v_n$
- Color each vertex with lowest color available
- There is always an ordering to give the chromatic number, but there are  $n!$  different orderings
- There exist orderings that lead to "bad" colorings
- The number of colors used greatly depends on the labeling of the vertices in a greedy coloring



# Applications of Graph Coloring

- **Assigning radio frequencies**

Draw an edge between two vertices if the corresponding radio stations are too close together so that the frequencies would interfere

- **Zoos and pet stores**

Can't have some animals in the same enclosures, some fish in the same tank, etc.

- **Resource allocation**

Create a vertex for each task and connect two vertices when the corresponding tasks require the same resource

A proper coloring ensures no two tasks that require the same resource are done at the same time

Chromatic number would give the optimal way to do the tasks simultaneously

- **Sudoku puzzle solvers**

# Assignment Integer Linear Program

- $x_{vi} = 1$  if and only if vertex  $v$  is assigned color  $i$
- $w_i = 1$  if at least one vertex is assigned color  $i$

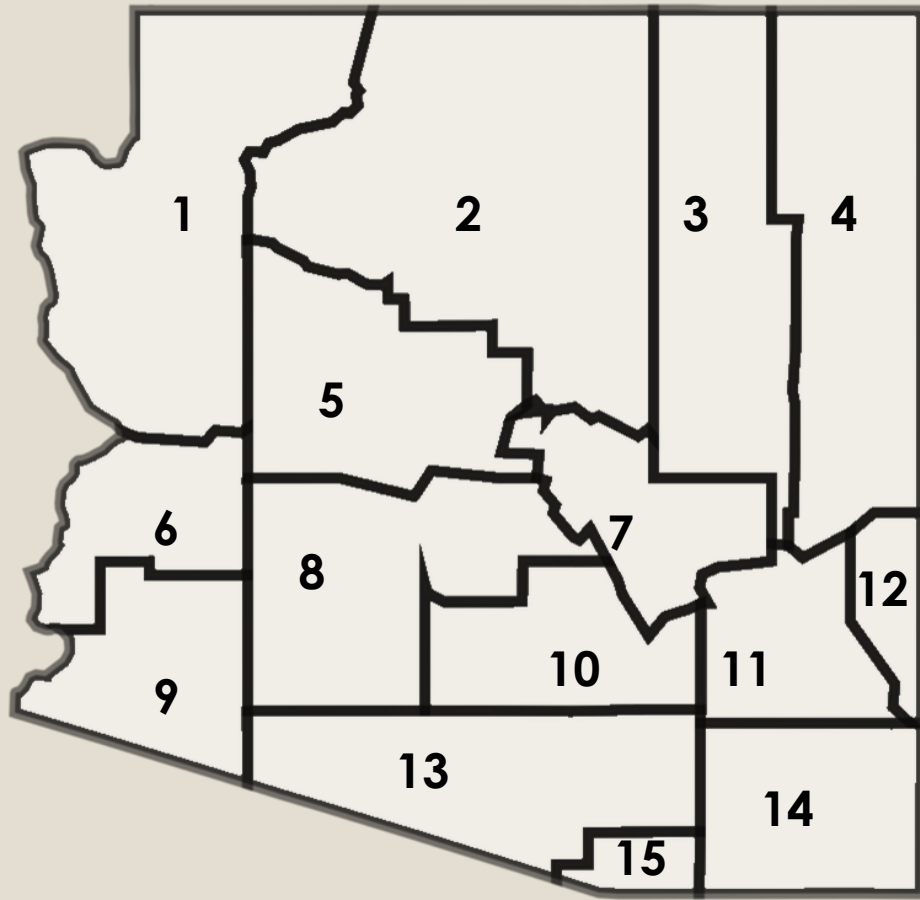
$$\begin{aligned} (\text{ASS-S}) \quad & \min \sum_{1 \leq i \leq H} w_i \\ \text{s.t.} \quad & \sum_{i=1}^H x_{vi} = 1 \quad \forall v \in V \\ & x_{ui} + x_{vi} \leq w_i \quad \forall (u, v) \in E, i = 1, \dots, H \\ & x_{vi}, w_i \in \{0, 1\} \quad \forall v \in V, i = 1, \dots, H \end{aligned}$$



# Code(AMPL)

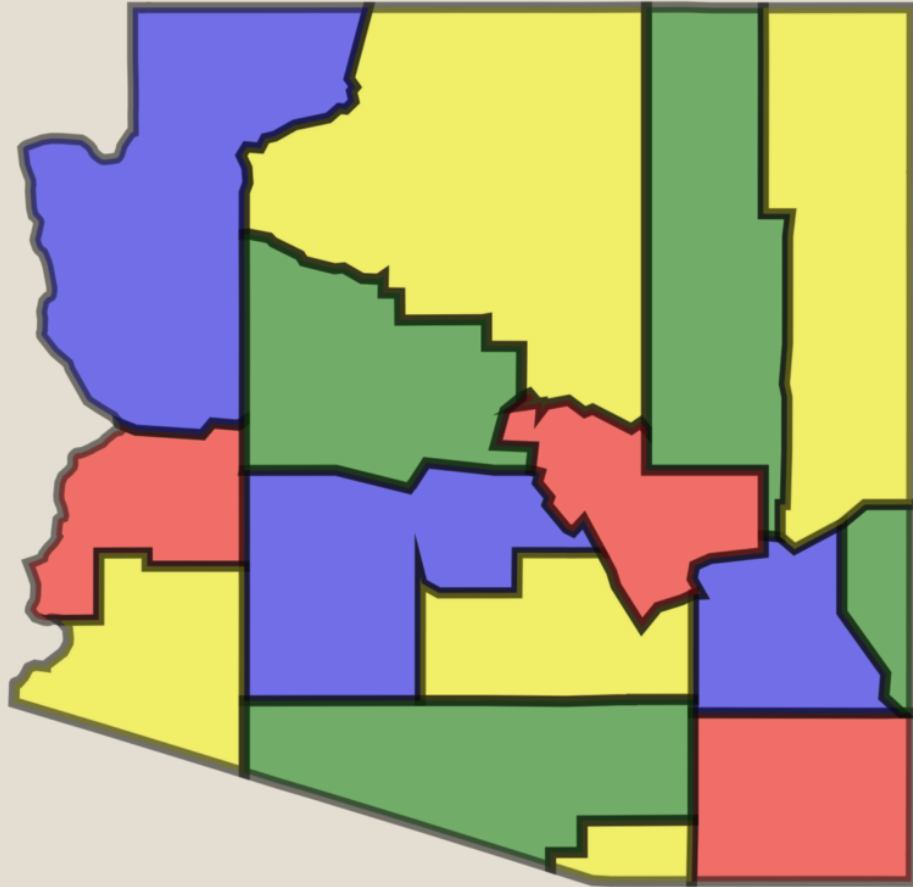
```
graph_coloring.mod ✖  
  
param H;  
set V := 1..H;  
set E within (V cross V);  
  
var w{1..H} binary;  
var x{V,1..H} binary;  
  
#objective function that yields the chromatic number of a graph  
minimize Colors: sum{i in 1..H} w[i];  
  
#Ensures each vertex is assigned a color  
subject to Assigned {i in V}:  
sum{j in 1..H} x[i,j]=1;  
  
#Ensures no two adjacent vertices are assigned the same color  
subject to Edges {(i,j) in E, k in 1..H}:  
x[i,k] + x[j,k] <= w[k];
```

# Problem1



```
graph_coloring.dat
param H := 15;
set E :=
(1,2) (1,5) (1,6)
(2,3) (2,5) (2,7)
(3,4) (3,7) (3,11)
(4,12) (5,6) (5,7)
(5,8) (6,8) (6,9)
(7,8) (7,10) (7,11)
(8,9) (8,10) (8,13)
(9,13) (10,11) (10,13)
(11,12) (11,14) (13,14)
(13,15) (14,15);
```

# Solution



```
x [*,*]  
: 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 :=  
1 1 0 0 0 0 0 0 0 0 0 0 0 0 0  
2 0 1 0 0 0 0 0 0 0 0 0 0 0 0  
3 1 0 0 0 0 0 0 0 0 0 0 0 0 0  
4 0 1 0 0 0 0 0 0 0 0 0 0 0 0  
5 0 0 1 0 0 0 0 0 0 0 0 0 0 0  
6 0 1 0 0 0 0 0 0 0 0 0 0 0 0  
7 0 0 0 1 0 0 0 0 0 0 0 0 0 0  
8 1 0 0 0 0 0 0 0 0 0 0 0 0 0  
9 0 0 1 0 0 0 0 0 0 0 0 0 0 0  
10 0 1 0 0 0 0 0 0 0 0 0 0 0 0  
11 0 0 1 0 0 0 0 0 0 0 0 0 0 0  
12 1 0 0 0 0 0 0 0 0 0 0 0 0 0  
13 0 0 0 1 0 0 0 0 0 0 0 0 0 0  
14 1 0 0 0 0 0 0 0 0 0 0 0 0 0  
15 0 1 0 0 0 0 0 0 0 0 0 0 0 0  
;  
Colors = 4
```

# Remarks

- **Four color theorem:**

- given any separation of a plane into contiguous regions, producing a figure called a *map*, no more than four colors are required to color the regions of the map so that no two adjacent regions have the same color.
- The four color theorem was proved in 1976 by [Kenneth Appel](#) and [Wolfgang Haken](#) after many false proofs and counterexamples

# Sudoku

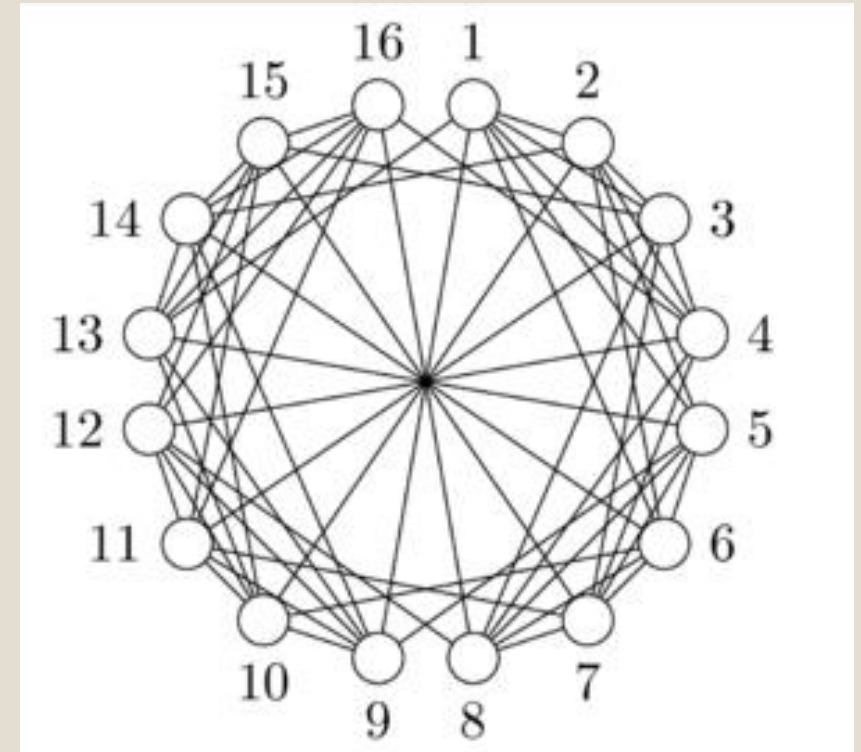
- The word Sudoku is a Japanese abbreviation for the phrase, “*suji wa dokushin ni kagiru*” which translates as the “the digits must remain single.”
- The puzzle consists of a 9×9 grid in which some of the entries of the grid have a number from 1 to 9.
- Filling the table with the numbers must follow these rules:
  - Numbers in rows are not repeated
  - Numbers in columns are not repeated
  - Numbers in 3 × 3 blocks are not repeated
  - Order of the numbers when filling is not important

8	3	4	6	7	1	9	2	5
1	2	5	8	3	9	6	4	7
7	9	6	5	2	4	3	1	8
9	5	7	3	1	8	4	6	2
2	4	1	9	5	6	8	7	3
3	6	8	2	4	7	5	9	1
6	8	2	7	9	3	1	5	4
5	1	9	4	8	2	7	3	6
4	7	3	1	6	5	2	8	9

# Sudoku as a Graph Coloring Problem

1	2	3	4
5	6	7	8
9	10	11	12
13	14	15	16

- We consider a 4x4 Sudoku puzzle.
- Label grid squares in Sudoku puzzle and construct corresponding graph by having two vertices be adjacent if:
  - they are in the same row
  - they are in the same column
  - they are in the same 2x2 subgrid
- A proper coloring of the graph translates to a solution to the Sudoku puzzle.



# Binary integer linear program (BILP) for general $n \times n$ puzzles

- In general, any  $n \times n$  game can be created, where  $n = m^2$  and  $m$  is any positive integer.
- In the BILP, the objective function is kept arbitrary. However, the constraints are important to follow to rules in order to find a feasible solution.

Define:

$$x_{ijk} = \begin{cases} 1, & \text{if element } (i, j) \text{ of the } n \times n \text{ Sudoku matrix contains the integer } k \\ 0, & \text{otherwise.} \end{cases}$$

$$\begin{aligned} \min \quad & \mathbf{0}^T \mathbf{x} \\ \text{s.t.} \quad & \sum_{i=1}^n x_{ijk} = 1, \quad j=1:n, k=1:n \quad (\text{only one } k \text{ in each column}) \end{aligned} \quad (1)$$

$$\sum_{j=1}^n x_{ijk} = 1, \quad i=1:n, k=1:n \quad (\text{only one } k \text{ in each row}) \quad (2)$$

$$\sum_{j=mq-m+1}^{mq} \sum_{i=mp-m+1}^{mp} x_{ijk} = 1, \quad k=1:n, p=1:m, q=1:m \quad (\text{only one } k \text{ in each submatrix}) \quad (3)$$

$$\sum_{k=1}^n x_{ijk} = 1 \quad i=1:n, j=1:n \quad (\text{every position in matrix must be filled}) \quad (4)$$

$$x_{ijk} = 1 \quad \forall (i, j, k) \in G \quad (\text{given elements } G \text{ in matrix are set "on"}) \quad (5)$$

$$x_{ijk} \in \{0, 1\} \quad (6)$$



# Code(AMPL)

```
sudoku.mod

set N := 1..9 ;
set G within {N cross N cross N};
var x {( i,j,k ) in {N cross N cross N}} binary ;

minimize nothing : x[1 ,1 ,1];

subject to Columns {j in N, k in N}:
    sum{i in N}x[i,j,k]=1;

subject to Rows {i in N, k in N}:
    sum{j in N}x[i,j,k]=1;

subject to Squares{k in N, p in 1..3 ,q in 1..3}:
    sum{j in (3*p -2) ..(3*p) } sum{i in (3*q -2) ..(3*q) }x[i,j,k]=1;

subject to all_filled{i in N, j in N}:
    sum{k in N}x[i,j,k]=1;

subject to known {(i,j,k) in G }:x[i,j,k]=1;
```

# Problem2

	1	7						
8				5				
2	3		4	1			7	
	8					3		9
		2		7		4		
1		3					6	
	9			3	6		5	7
				2				3
						1	9	

```
sudoku.dat
set G:=
(1,2,1) (1,3,7) (2,1,8)
(2,5,5) (3,1,2) (3,2,3)
(3,4,4) (3,5,1) (3,8,7)
(4,2,8) (4,7,3) (4,9,9)
(5,3,2) (5,5,7) (5,7,4)
(6,1,1) (6,3,3) (6,8,6)
(7,2,9) (7,5,3) (7,6,6)
(7,8,5) (7,9,7) (8,5,2)
(8,9,3) (9,7,1) (9,8,9);
```

# Solution

5	1	7	6	8	3	9	2	4
8	4	9	2	5	7	6	3	1
2	3	6	4	1	9	5	7	8
7	8	4	5	6	2	3	1	9
9	6	2	3	7	1	4	8	5
1	5	3	8	9	4	7	6	2
4	9	8	1	3	6	2	5	7
6	7	1	9	2	5	8	4	3
3	2	5	7	4	8	1	9	6

```
x [1,*,*]
: 1 2 3 4 5 6 7 8 9 :=
1 0 0 0 0 1 0 0 0 0
2 1 0 0 0 0 0 0 0 0
3 0 0 0 0 0 0 1 0 0
4 0 0 0 0 0 1 0 0 0
5 0 0 0 0 0 0 0 1 0
6 0 0 1 0 0 0 0 0 0
7 0 0 0 0 0 0 0 0 1
8 0 1 0 0 0 0 0 0 0
9 0 0 0 1 0 0 0 0 0
```

```
[2,*,*]
: 1 2 3 4 5 6 7 8 9 :=
1 0 0 0 0 0 0 1 0
2 0 0 0 1 0 0 0 0 0
3 0 0 0 0 0 0 0 0 1
4 0 1 0 0 0 0 0 0 0
5 0 0 0 0 1 0 0 0 0
6 0 0 0 0 0 0 1 0 0
7 0 0 0 0 0 1 0 0 0
8 0 0 1 0 0 0 0 0 0
9 1 0 0 0 0 0 0 0 0
```

```
[3,*,*]
: 1 2 3 4 5 6 7 8 9 :=
1 0 1 0 0 0 0 0 0
2 0 0 1 0 0 0 0 0 0
3 0 0 0 0 0 1 0 0 0
4 0 0 0 1 0 0 0 0 0
5 1 0 0 0 0 0 0 0 0
6 0 0 0 0 0 0 0 0 1
7 0 0 0 0 1 0 0 0 0
8 0 0 0 0 0 0 1 0 0
9 0 0 0 0 0 0 0 1 0
```

# Bibliography:

- [Colour Maps using Integer Programming by Mohamed Leila](#)
- [Wikipedia-Four color theorem](#)
- [An Integer Programming Model for the Sudoku Problem Andrew C. Bartlett Amy N. Langville March 18, 2006](#)
- [An Integer Linear Programming Approach to Graph Coloring by Megan Duff and Rebecca Robinson](#)
- [Solving sudoku as an Integer Programming problem by Olszowy Wiktor](#)
- [Wikipedia- AMPL \(A Mathematical Programming Language\)](#)
- [MathWorks](#)

Thank You