

→ we define  $(m, n)$  as number of different corresponding bits in a binary search representation of  $x$  and  $y$ .

$$\rightarrow \begin{bmatrix} 1, 3, 5 \\ 2, 7 \end{bmatrix} \rightarrow 8$$

→ Brute force:

- we iterate over the loop  $i \rightarrow n$
- then we iterate from  $i+1$  to  $n$  inner loop
- then compare all the elements  $(i, j)$   
then check if their bits are same or  
not count the different bits.
- then at the end double the count  
then return the count.

```
int count = 0;
for (int i = 0, i < n, i++) {
    int count = 0;
    for (int j = i + 1, j < n, j++) {
        count += __builtin_popcount(i ^ j);
    }
}
```

return  $2 * count$ ;

## Optimal Soln

- Take a array of  $3^l$  size that all the index represent the bits from  $1 \rightarrow 3^l$
- for every element in original array if the bit is set then add 1 to its corresponding index of the array bit.
- after iterating all over the array all the bits that are set at the element in the array are reflecting to the array.
- Then iterate over the freq array for all the non-zero element formula becomes:

$$\text{freq}[i] * (N - f[i]);$$

→  $[1, 3, 5]$   
 for 3  
 $3 * (3-3)$   
 $= 0$

→ then next  
 $0 + 1 + (3-1) + 1 + (3-1)$   
 $= 0 + 2 + 2$   
 $= 4$

→ Double it  $\Rightarrow 4 \times 2 = 8$  (Because we are counting both for  $(A_i, A_j) \& (A_j, A_i)$ ) return the Ans 28.

Time complexity =  $O(N) * O(3^l)$

Space complexity =  $O(3^l)$ .