#### ECE408/CS483/CSE408 Fall 2022

Applied Parallel Programming

# Lecture 7: Convolution, Constant Memory and Constant Caching

### Course Reminders

- Lab updates
  - We are grading Lab 2 now, expect it to be graded soon
  - Lab 3 is due this Friday
- Take a note of Midterm 1 date
  - Details to be provided soon

## Objective

- To learn convolution, an important parallel computation pattern
  - Widely used in signal, image and video processing
  - Foundational to stencil computation used in many science and engineering applications
  - Critical component of Neural Networks and Deep Learning
- Important GPU technique
  - Taking advantage of cache memories

## **Convolution Mathematics**

$$f(x) * g(x) = \int_{-\infty}^{\infty} f(\tau) \cdot g(x - \tau) d\tau$$

$$f[x] * g[x] = \sum_{k=-\infty}^{\infty} f[k] \cdot g[x-k]$$

# **Convolution Applications**

- A popular operation that is used in various forms in signal processing, digital recording, image processing, video processing, computer vision, and machine learning.
- Convolution is often performed as a filter that transforms the input signal (audio, video, etc) in some context-aware way.
  - Some filters smooth out the signal values so that one can see the bigpicture trend
  - Or Gaussian filters to blur images, backgrounds

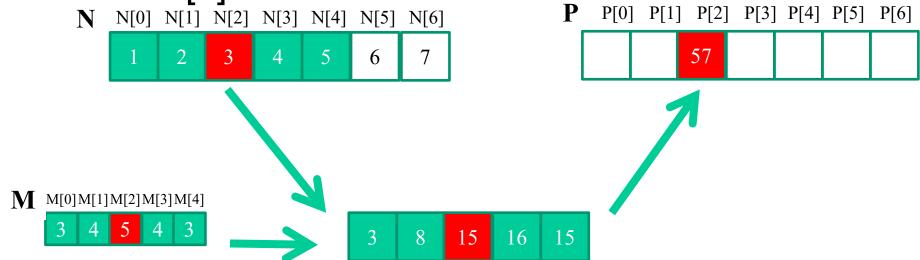
# **Convolution Computation**

- An array operation where each output data element is a weighted sum of a collection of neighboring input elements
- The weights used in the weighted sum calculation are defined by an input mask array, commonly referred to as the convolution kernel
  - We will refer to these mask arrays as convolution masks or convolution filters to avoid confusion.
  - The same convolution mask is typically used for all elements of the array.

# 1D Convolution Example

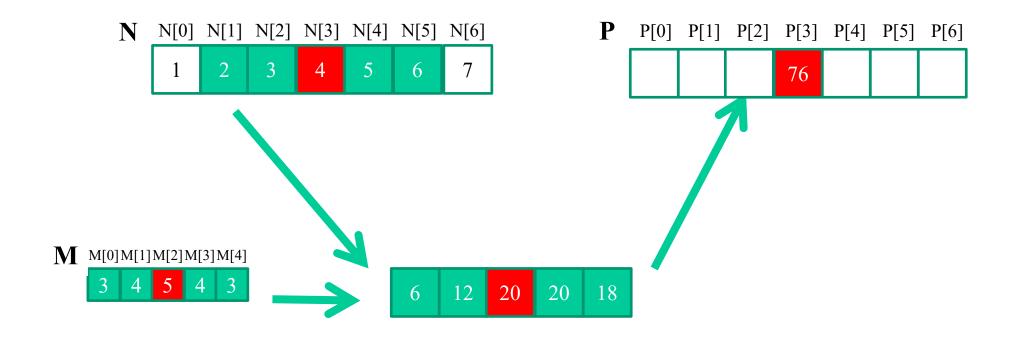
- Commonly used for audio processing
  - MASK\_WIDTH is usually an odd number of elements for symmetry (5 in this example)
  - MASK\_RADIUS is the number of elements used in convolution on each side of the pixel being calculated (2 in this example).

Calculation of P[2]:



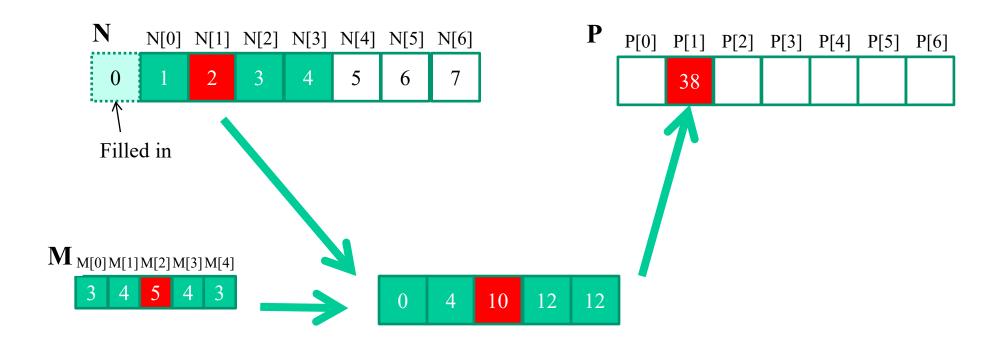
## 1D Convolution Example

Calculation of P[3]



## 1D Convolution Boundaries

- Calculation of output elements near the boundaries of the input array need to deal with "ghost" elements
  - Different policies (0, replicates of boundary values, etc.)



## A 1D Convolution Kernel with Boundary Handling

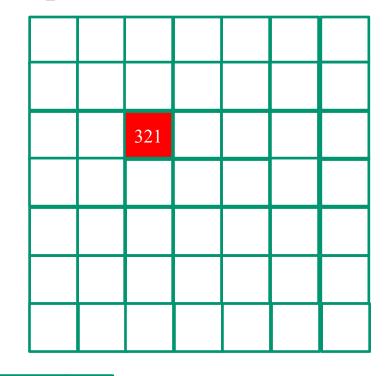
This kernel forces all elements outside the valid range to 0

```
global void
convolution 1D basic kernel(float *N, float *M, float *P, int Mask Width, int Width)
  int i = blockIdx.x*blockDim.x + threadIdx.x;
  float Pvalue = 0;
  int N start point = i - (Mask Width/2);
  for (int j = 0; j < Mask Width; <math>j++) {
    if (((N \text{ start point} + j) >= 0) \&\& ((N \text{ start point} + j) < Width)) {}
      Pvalue += N[N start point + j]*M[j];
  P[i] = Pvalue;
```

# 2D Convolution

N

<b>1</b>						
1	2	3	4	5	6	7
2	3	4	5	6	7	8
3	4	5	6	7	8	9
4	5	6	7	8	5	6
5	6	7	8	5	6	7
6	7	8	9	0	1	2
7	8	9	0	1	2	3



M

IVE						
1	2	3	2	1		
2	3	4	3	2		
3	4	5	4	3		
2	3	4	3	2		
1	2	3	2	1		

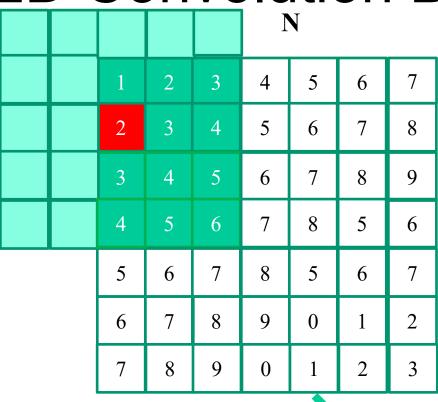


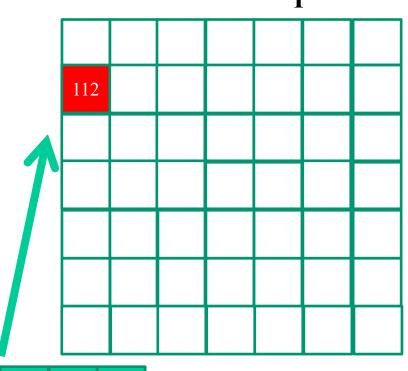


1	4	9	8	5
4	9	16	15	12
9	16	25	24	21
8	15	24	21	16
5	12	21	16	5

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# 2D Convolution Boundary Condition



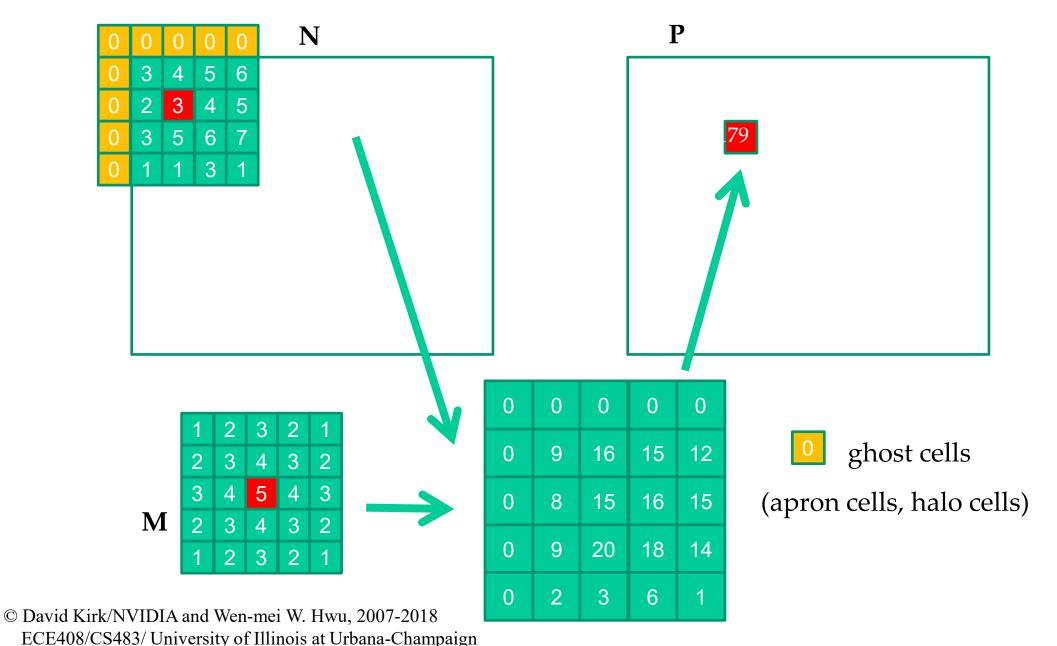


M					
1	2	3	2	1	
2	3	4	3	2	
3	4	5	4	3	
2	3	4	3	2	
1	2	3	2	1	

0	0	0	0	0
0	0	4	6	6
0	0	10	12	12
0	0	12	12	10
0	0	12	10	6

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## 2D Convolution – Ghost Cells



## What does this kernel accomplish?

$$\mathbf{M} = \frac{1}{273} \times \begin{bmatrix} 1 & 4 & 7 & 4 & 1 \\ 4 & 16 & 26 & 16 & 4 \\ 7 & 26 & 41 & 26 & 7 \\ 4 & 16 & 26 & 16 & 4 \\ 1 & 4 & 7 & 4 & 1 \end{bmatrix}$$

Hint: Assume input N is an image

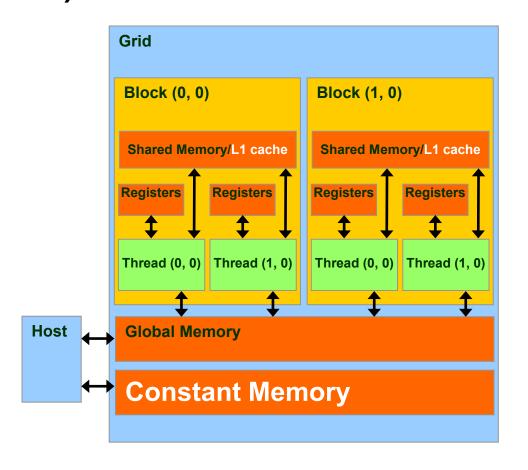
### Access Pattern for M

- Elements of M are called mask (kernel, filter) coefficients (weights)
  - Calculation of all output P elements needs M
  - M is not changed during grid execution
- Bonus M elements are accessed in the same order when calculating all P elements

M is a good candidate for Constant Memory

# Programmer View of CUDA Memories (Review)

- Each thread can:
  - Read/write per-thread registers (~1 cycle)
  - Read/write per-block
     shared memory (~5 cycles)
  - Read/write per-grid
     global memory (~500 cycles)
  - Read/only per-grid constant memory (~5 cycles with caching)



## Memory Hierarchies

- Review: If we had to go to global memory to access data all the time, the execution speed of GPUs would be limited by the global memory bandwidth
  - We saw the use of shared memory in tiled matrix multiplication to reduce this limitation

Another important solution: Caches

## Cache

- A cache is an "array" of cache lines
  - A cache line can usually hold data from several consecutive memory addresses
- When data is requested from the global memory, an entire cache line that includes the data being accessed is loaded into the cache, in an attempt to reduce global memory requests
  - The data in the cache is a "copy" of the original data in global memory
  - Additional hardware is used to remember the addresses of the data in the cache line

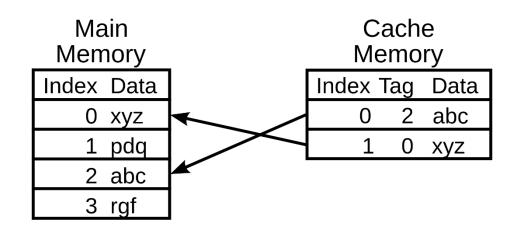
# Caches Store Lines of Memory

#### Recall: memory bursts

- contain around 1024 bits (128B) from
- consecutive (linear) addresses.
- Let's call a single burst a line.

#### What's a cache?

- An array of cache lines (and tags).
- Memory read produces a line,
- cache stores a copy of the line, and
- tag records line's memory address.



# Caches and Locality

- Some definitions:
  - Spatial locality: when the data elements stored in consecutive memory locations are accessed consecutively
  - Temporal locality: when the same data element is accessed multiple times in short period of time
- Both spatial locality and temporal locality improve the performance of caches

## Memory Accesses Show Locality

#### An executing program

- loads and store data from memory.
- Consider sequence of addresses accessed.

**Sequence** usually **shows** two types of **locality**:

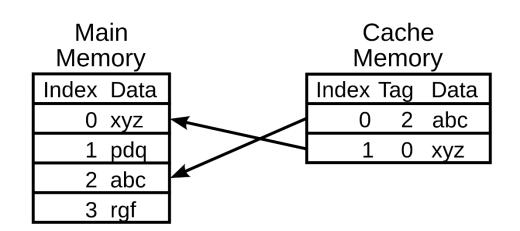
- spatial: accessing X implies
   accessing X+1 (and X+2, and so forth) soon
- temporal: accessing X implies
   accessing X again soon
   (Caches improve performance for both types.)

# Caches Can't Hold Everything

Caches are smaller than memory.

#### When cache is full,

- must make room for new line,
- usually by discarding least recently used line.



# Shared Memory vs. Cache

- Shared memory in CUDA is another type of temporary storage used to relieve main memory contention.
  - In terms of distance from the SMs, shared memory is similar to L1 cache.

- Unlike cache, shared memory does not necessarily hold a copy of data that is also in main memory
  - Shared mem requires explicit data transfer instructions into locations in the shared mem, whereas cache doesn't

# Shared Memory vs. Cache - Cont'd

- Caches vs. shared memory
  - Both on chip\*, with similar performance
  - (As of Volta generation, both using the same physical resources, allocated dynamically!)

#### What's the difference?

- Programmer controls shared memory contents (called a scratchpad)
- Microarchitecture automatically determines contents of cache.

\*Static RAM, not DRAM, by the way—see ECE120/CS233.

## Constant Cache in GPUs

- Modification to cached data needs to be (eventually) reflected back to the original data in global memory
  - Requires logic to track the modified status, etc.

- Constant cache is a special cache for constant data that will not be modified during kernel execution by a grid
  - Data declared in the constant memory not modified during kernel execution.
  - Constant cache can be accessed with higher throughput than L1 cache for some common patterns

## GPU Has Constant and L1 Caches

#### To support writes (modification of lines),

- changes must be copied back to memory, and
- cache must track modification status.
- L1 cache in GPU (for global memory accesses) supports writes.

#### Cache for constant / texture memory

- Special case: lines are read-only
- Enables higher-throughput access than L1 for common GPU kernel access patterns.

## How to Use Constant Memory

Host code is similar to previous versions, but...

Allocate device memory for M (the mask)

- outside of all functions
- using \_\_constant\_\_
   (tells GPU that caching is safe).

For copying to device memory, use

- cudaMemcpyToSymbol(dest, src, size, offset = 0, kind = cudaMemcpyHostToDevice)
- with destination defined as above.

## Host Code Example

```
(MASK WIDTH is the size of the mask.)
// global variable, outside any kernel/function
  constant float Mc[MASK WIDTH] [MASK WIDTH];
// Initialize Mask
float Mask[MASK WIDTH] [MASK WIDTH]
for(unsigned int i = 0; i < MASK WIDTH * MASK WIDTH; i++) {</pre>
    Mask[i] = (rand() / (float)RAND MAX);
    if(rand() % 2) Mask[i] = - Mask[i]
cudaMemcpyToSymbol(Mc, Mask, MASK WIDTH*MASK WIDTH*sizeof(float));
ConvolutionKernel<<<dimGrid, dimBlock>>>(Nd, Pd);
```

## ANY MORE QUESTIONS? READ CHAPTER 7