

Arrays

Definition of an Array:

List of elements of the same type placed in a contiguous memory location.

Creation of an Array:

```
dataType arrayName[] = new dataType [size];
```

Length of array:

```
length = array-name.length;
```

Passing array as argument:

```
func declaration : ps void fn-name (int array-name[])  
                    {  
                    }  
                    }
```

```
fn call : fn-name (array-name);
```

Linear search

```
public static Array{
```

```
    public static int linearSearch (int a[], int key)  
    {
```

```
        for (int i=0; i<a.length; i++) {
```

```
            if (a[i] == key)
```

```
                return i;
```

```
        }
```

```
    }  
    return -1;
```

```
}
```

Time complexity : $O(n)$

Binary Search

- Array should be sorted.

Code -

```
public static int binary_search (int A[], int key)
{
    int start = 0, end = A.length - 1;
    while (start <= end) {
        int mid = (start + end) / 2;
        if (A[mid] == key) {
            return mid;
        }
        if (A[mid] < key) {
            start = mid + 1;
        }
        else {
            end = mid - 1;
        }
    }
    return -1;
}
```

Time Complexity : $O(\log_2 n)$

Reverse an array:

```
public static void reverse (int A[]) {
```

```
    int start = 0, end = A.length - 1;
```

```
    while (first start < end) {
```

```
        int temp = A[start];
```

```
        A[start] = A[end];
```

```
        A[end] = temp;
```

```
        for start ++;
```

```
        end ++; end --;
```

```
    }
```

```
}
```

Pairs in an array:

```
public static void pairs (int A[]) {
```

```
    for (int i = 0; i < A.length; i++) {
```

```
        int curr = A[i];
```

```
        for (int j = i + 1; j < A.length; j++) {
```

```
            System.out.print ("(" + curr + ", " + A[j] + ") ");
```

```
        }
```

```
        System.out.println();
```

```
    }
```

```
}
```

Time complexity: $O(n^2)$

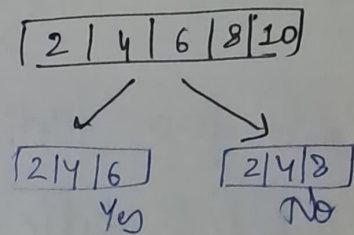
$$\text{Total pairs} = \frac{n(n-1)}{2}$$

2	4	6	8	10
---	---	---	---	----

Pairs: (2,4) (2,6)
 (2,8) (2,10)
 (4,6) (4,8) (4,10)
 (6,8) (6,10)
 (8,10)

Print SubArrays

a continuous part of array



Total no of subarrays: $\frac{n(n+1)}{2}$

code-

```
public static void printSubArrays (int A[]) {  
    //start for (int i=0; i < A.length; i++) {  
        int start=i;  
        for (int j=0; j < A.length; j++) {  
            //end int end=j;  
            for (int k=start; k <= end; k++) {  
                //print System.out.print(A[k] + " ");  
            }  
            System.out.println();  
        }  
        System.out.println();  
    }  
}
```

Sum of Subarrays

Brute force Method

```
public static void maxSubarray (int A[]) {
```

```
    int currSum = 0;
```

```
    int maxSum = Integer.MIN_VALUE;
```

```
    for (int i = 0; i < A.length; i++) {
```

```
        for (int j = i; j < A.length; j++) {
```

```
            currSum = 0;
```

```
            for (int k = i; k <= j; k++) {
```

```
                currSum += A[k];
```

```
            }
```

```
            if (maxSum < currSum)
```

```
                maxSum = currSum;
```

```
        }
```

```
    }
```

```
}
```

Time complexity: $O(n^3)$

Very bad

Sum of Subarrays (II)

20/9/23

Prefix Method

In this we have a new array: Prefix array
That stores sum of previous indexes at that index

1	-2	6	-1	3
---	----	---	----	---

$-2 + 6 + -1 = 3$

1	-1	5	4	7
---	----	---	---	---

Prefix array



$$\text{prefix}[\text{end}] - \text{prefix}[\text{start} - 1]$$

$$= 4 - 1 = \underline{\underline{3}}$$

```
public static void prefixSum  
(int A[]) {
```

```
    int currSum = 0;
```

```
    int maxSum = Integer.MIN_VALUE;
```

```
    int prefix[] = new int[A.length];  
    prefix[0] = A[0];
```

```
    for (int i = 1; i < A.length; i++) {  
        prefix[i] = prefix[i-1] + A[i];  
    } // calculate prefix array
```

```
    for (int i = 0; i < A.length; i++) {
```

```
        for (int j = i; j < A.length; j++) {
```

```
            currSum = 0;
```

```
            currSum = start = 0;
```

```
            currSum = i = 0 ? prefix[j] : prefix[j] - prefix[i-1];
```

```
            if (maxSum < currSum)
```

```
            { maxSum = currSum; }
```

```
        }  
    }
```

Time
Complexity
: $O(n^2)$

Max Subarray Sum (III)

KADANE'S Algorithm

* Do consider -ve sum as Zero in prefix method

eg:

-2, -3, 4, -1, -2, 1, 5, -3

Csum	0	0	4	3	1	2	7	4	1
Msum	0	0	4	4	4	4	7	4	1

Code -

```
public static void kadanes(int A[]) {  
    int ms = Integer.MIN_VALUE;  
    int cs = 0;  
    for (int i = 0; i < A.length; i++) {  
        cs = cs + A[i];  
        if (cs < 0) {  
            cs = 0;  
        }  
        ms = Math.max(cs, ms);  
    }  
}
```

Time complexity: $O(n)$

Trapping Rainwater

Points -

- ★ min no of bars $> \underline{\underline{2}}$
- ★ Asc/Des \rightarrow no water is trapped

$$\star \quad \boxed{\text{trapped water} = (\text{water level} - \text{bar height}) \times \text{width}}$$

$$\hookrightarrow \boxed{\text{Water level} = \min(\text{max left}, \text{max Right})}$$

This is done with help of
Auxiliary arrays
(helping arrays)

1st) Generate 2 arrays maxleft & max Right, these will store max left & maxRight of that index respectively.

2nd) loop -

$$\text{waterLevel} = \min(\text{max left}, \text{max Right})$$

$$\text{trapped water} = \text{water level} - \text{bar height}$$

PTO for CODE

code -

```
public static voidint trapping (int A[]) {  
    int n = A.length;  
    int maxleft[] = new int [n];  
    maxleft[0] = A[0];  
    for (int i=0 ; i<n ; i++) {  
        maxleft[i] = Math.max(A[i], maxleft[i-1]);  
    }  
    int maxRight[] = new int [n];  
    maxRight[n-1] = A[n-1];  
    for (int i=n-1 ; i>=0 ; i--) {  
        maxRight[i] = Math.max(A[i], maxRight[i+1]);  
    }  
    int trap-water = 0;  
    for (int i=0 ; i<n ; i++) {  
        int water-level = Math.min(maxleft[i], maxRight[i]);  
        trapwater += water-level - A[i];  
    }  
    return trap-water;  
}
```

Time Complexity: $O(n)$