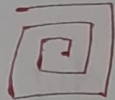


# 2-D arrays

Creation-

```
int matrix[][ ] = new int [size1][size2];
```

## Spiral Matrix



amazon, google, apple, oracle

eg:

1	2	3	4
5	6	7	8
9	10	11	12
13	14	15	16

output:

1 2 3 4 8 12 16 15 14 13

9 5 6 7 11 10

## Approach -

Iteration: →		1 <sup>st</sup>	2 <sup>nd</sup>	matlab
start Row	= 0	1	++	
start Col	= 0	1	++	
End Row	= n-1 (3)	2	--	
end Col	= n-1 (3)	2	--	

// Code in LAPTOP

```
ps void printSpiral (int matrix[ ][ ]) {
```

```
    int sRow = 0 , sCol = 0;
```

```
    int eRow = eCol = matrix.length - 1;
```

```
    while (sRow <= eRow && sCol <= eCol) {
```

```
        // top
```

```
        for (int j = sCol; j <= eCol; j++) {
```

```
            System.out.print (matrix[sRow][j] + " ");
```

```
        }
```

```
        // right
```

```
        for (int i = sRow + 1; i <= eRow; i++) {
```

```
            System.out.print (matrix[i][eCol] + " ");
```

```
        }
```

```
        // bottom
```

```
        for (int j = eCol - 1; j >= sCol; j--) {
```

```
            System.out.print (matrix[eRow][j] + " ");
```

```
            if (sRow == eRow) break;
```

```
        }
```

```
        // left
```

```
        for (int i = eRow - 1; i >= sRow + 1; i--) {
```

```
            if (sCol == eCol) break;
```

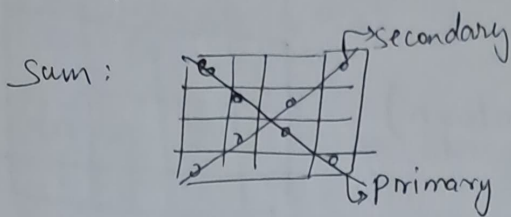
```
            System.out.print (matrix[i][sCol] + " ");
```

```
        } sCol++; sRow++; eRow--; eCol--;
```

# Diagonal Sum

Concept :

$O(n)$



Sum of ele of primary array =  
condition  $\Rightarrow i = j$

$$\text{Sum} += \text{arr}[i][j] \quad \text{where } i = j$$

So we can replace  $j$

$$\text{Sum} += \text{arr}[i][i]$$

Sum of ele of secondary array =

condition: If  $(i + j = n - 1)$

$$\text{Then } j = n - 1 - i$$

Therefore,

$$\text{sum} += \text{arr}[i][n - 1 - i]$$

Code -  
ps ~~is~~ int diagonalSum(int matrix[][]) {

for (int i = 0; i < matrix.length; i++) {

sum += matrix[i][i];

if (i != matrix.length - 1 - i)

sum += matrix[i][matrix.length - 1 - i];

}

return sum;

}

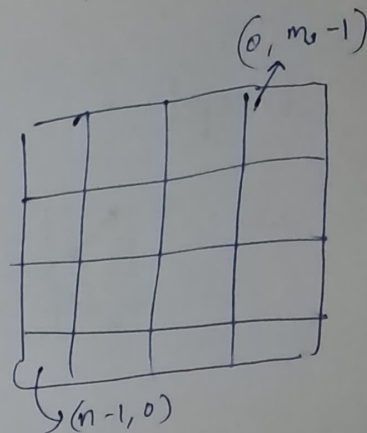


# Search in SORTED MATRIX

① Brute force  $O(n^2)$   
Search everywhere

② Row wise / Column wise  $O(n \log n)$   
Search in each R/C separately

③ Staircase Approach  $O(m+n)$



$(n-1, 0)$

key < cell value

TOP

key > cell value

RIGHT

$(0, m-1)$

key < cell value

LEFT

key > cell value

BOTTOM

CODE - ps ~~void~~ <sup>boolean</sup> staircaseSearch (int matrix[][], int key)

```
{
    int row = 0, col = matrix.length - 1;
    while (row < matrix.length & col >= 0) {
        if (matrix[row][col] == key) {
            System.out.print(" — ");
            return true;
        }
        else if (key < matrix[row][col])
            col--;
        else
            row++;
    }
    System.out.print("key not found");
}
```