Sorting

Bubble Sort: large dement come to end of array by swapping with adjacent element

for (int twon = 0; twon (and length; i++) {

for (int j=0; j< Andergth-1-twon; j++) {

if (arr [j] > arr [j+1]) {

int temp = arr [j]

arr [j+1] = arr [j+1]

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Best & Worst -line complexity: O(n2)

Selection Sort: select the smallest & put at in the front.

for (int 1=0; 1< orr. length-1; 1++) { int min Pos = i; for (int j=i+1; j(arr. length; j++) { if (arr [minfos] > arr [j]) } int temp = arz [minPos]; arr [minfos] = arr [i]; arr [i] = temp;

Best & worst case: O(n2)

Insertion sort: Pick an element (from ansorted)

and insert it in right position (cortigue)

is me peeche ke saare elements se compare karle hain.

for (int i = 1; i< oor, length; i++) }

int curr = arr[i];

prev = i-1;

while (prev>=0 && orr [prev] > curr) {
our [prev +1] = orr [prev];
prev -- .

our [prev+1] = curr;

Time complexity: O(n2)

Merge Sort Divide & Conquer 5 | 2 | 8 8

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Base cases of Recursion (Merge sort)-1) Si > Ci (invalid) Si = Ci (single element) Steps for recursion--j) divide 2) Merge Sort (left) 3 call for (right) Jinner fr -3) Merge (temp array) PS void extender (and 25' any 25) ! ( [ mailine to

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Merge Sont Gode
  PS void main (String args[]) ?
       int arr = {6,3,9,5,2,839
   merge Sort (arr, 0, arr, length-1);

? PrintArr (arr);
 PS void merge Sort (int arr [], intsi, int ei) {
Here int mid= Sit (Bi-Si)/2
                                  //si-ei/2_
   merge Sort (arr, o, mid); // left
    merge Sort (arr, mid+1, arr. leagth-1)://right
     merge (arr, si, mid, li); // merging them
   Ef (si >= ei) }
     geturn;
     void print Arr (ant arr []) {
        for (int i=0; i <= arr-length; i++)
```

PS void merge ("int arr [7, "int Si, int mid, intei) { unt demp[]= new unt [ei-si+1]; unt i = Si; 11 for left part unt g = mid+1; Il for right part int k = 0; 11 for temp array while (i <= mid 22 j <= ei) { if (arr [i] < arr [j]) { 11 compairing. temp [K] = arr [i];

i++; }
else } & putting them in temp array temp [k] = orr [j]; 11 remaining left part while (i <= mid) { temp[k++] = arr [i++]; Il remaining right while (j <= ei) { q femp [Ktt] = arr [j+t]; for (K=0, P=Si; k <temp-length; K++, P++) } g arr CiJ= temp[R]; 11 copy them to original array=

Quick Sort Prot & Partition

Steps-

1) Pivot (last element)

2) Partition (parts)

3) Call quick sort (left)

Base case-

\* Si>= ei;

Code -

ps void main (string args[]) {
int orr = { 6, 7, 1, 4, 5, 9 }
quick Sort (arr, 0, orr.length-1);
print Array (arr);

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Ps void quick Sort (int arr[], intsi, intei) {

if (si>=ei) {

3 return;

int pIdx = partition (arr, si, ei); quick Sort (arr, si, pIdx-1); quick Sort (arr, pIdx+1, ei);

11 left 11 rigat

```
PS int partition (int arr [], int si, int ei) {
        unt pivot = arr [ei];
        int i = si-1; // to make place for els
Smaller than pivot
       for (int j=0;j< ei;j++) {
    if (arr Cj] <= pivot) {
        i++;
             int temp = orr [j];
             our [j] = our [i];
           ear [i] = ar [j];
        int temp = pivot;
 wrong pivot = arr [i]; arr [ei] = arr [i];
        orr li] = temp;
        return i;
```

Scarch in Rotated Sorted Array Input: sorted, notated array with distinct nos
It is rotated at a pulot point. Final index of green dement [3 | 5 | 6 | 0 | 1 | 2 | Oitput: 3 for this we will use Modified binary Search Approach: target=0

[3 [5 [6]] [1] [2]

Si mid ei case 1: mid on L! [our [si] = mid] \_ case a: 4 feft (si = far (= mid) case b: rock 11 right (else) case 2: mid on L2 [arr [mid] (= arr[ei]] case c: l2 right [mid (= tar (= ei) cored: 12 left (clx

, aut tar, int si, int eil PS int search (int 1/ Jase case if (si) { 3 return -1; Int mid= si+ (ei-si)/2 if (arr [mid] == tar) { || case found return mid; if (arr [si] (= arr [mid]) // mid on [1 LI left if (our [si] (= tar & tar (= our (mid]) }

geturn search (our, tor, si, mid-1); LI rightelse {
return search (arr, tar, mid+1, ei):
2 else { if (arr [mid] (= tar 4d tar (= arr [ei]) 12 right if { return search (arr, tar, mid+1, ei); } 12 left else ? roun search (an, tar, si, mid-1); }

Counting Sort o(n + largest) \* first we count the freq of nos from 0 to 9 and where the freq in count array. \* The we place (sort) nos in main array acc to count array.

Code - array | 0 | 2 | 0 | 1 | 3 | 5 |

PS void counting Sort (int arr []) { int largest = Integer. MIN-VALUE; I for oize for (int i=0; i < arr. length; i++) } largest = Math.max (largest, over [i]): A count int count [] = new int [largest + ]); for (int i=0; ixarr.length; i++) { Count [arr[i]] ++; // plate that no for (int i=0; ix count. length; i++) { nohile (count [j] >0) } 11 sorting arr GJ = i; 3 3 3 cout[i]--;