## **EXPERIMENT: 12**

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#### Part A: Simulating a Deadlock Between Two Transactions

## **Description:**

Given a table StudentEnrollments containing student records, simulate a situation where two concurrent transactions (from different users) try to update overlapping records in different orders, resulting in a deadlock. Demonstrate how such deadlocks are detected and how they can be avoided using proper

transaction ordering.

#### **Input Format:**

- Table **StudentEnrollments** with columns:
  - o student\_id (INT, Primary Key)
  - o student name (VARCHAR(100))
  - o course id (VARCHAR(10))
  - o enrollment\_date (DATE)

#### **Output Format:**

Demonstrate that one transaction will be rolled back automatically by the database to resolve the deadlock.

#### **Constraints:**

- Use two user sessions to run START TRANSACTION simultaneously.
- Ensure the transactions access rows in reverse order to trigger a deadlock.
- Database must support deadlock detection (e.g., MySQL, PostgreSQL).

## **Sample Input:**

#### **StudentEnrollments**

student_id	student_name	course_id	enrollment_date
1	Ashish	CSE101	2024-06-01
2	Smaran	CSE102	2024-06-01
3	Vaibhav	CSE103	2024-06-01

#### **Sample Output:**

Transaction 2 is aborted due to a detected deadlock.

# **Explanation of Output:**

Both transactions try to lock each other's rows in reverse order. This causes a deadlock, and the database automatically rolls back one transaction (usually the one that waited longest) to break the cycle.

#### Query:

DROP TABLE IF EXISTS StudentEnrollments;

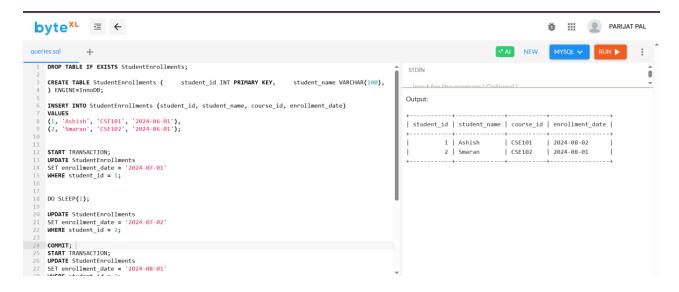
```
CREATE TABLE StudentEnrollments (
student_id INT PRIMARY KEY,
student_name VARCHAR(100),
course_id VARCHAR(10),
enrollment_date DATE
) ENGINE=InnoDB;
```

INSERT INTO StudentEnrollments (student\_id, student\_name, course\_id, enrollment\_date) VALUES

```
(1, 'Ashish', 'CSE101', '2024-06-01'),
(2, 'Smaran', 'CSE102', '2024-06-01');
```

```
START TRANSACTION;
UPDATE StudentEnrollments
SET enrollment_date = '2024-07-01'
WHERE student_id = 1;
DO SLEEP(1);
UPDATE StudentEnrollments
SET enrollment date = '2024-07-02'
WHERE student_id = 2;
COMMIT;
START TRANSACTION;
UPDATE StudentEnrollments
SET enrollment_date = '2024-08-01'
WHERE student_id = 2;
DO SLEEP(1);
UPDATE StudentEnrollments
SET enrollment_date = '2024-08-02'
WHERE student_id = 1;
COMMIT;
SELECT * FROM StudentEnrollments;
```

#### **OUTPUT:**



# **Explanation**:

- Two transactions try to update the same rows in reverse order.
- Each transaction holds a lock the other needs  $\rightarrow$  **deadlock**.
- MySQL detects it and rolls back one transaction automatically to resolve it.

**Key point:** Access rows in the **same order** in all transactions to avoid deadlocks.

## Part B: Applying MVCC to Prevent Conflicts During Concurrent

#### **Reads/Writes Description:**

Use the MVCC (Multiversion Concurrency Control) approach to allow User A to read a record and User B to update the same record concurrently without blocking or conflict. Demonstrate how MVCC provides a consistent snapshot to the reader while allowing the writer to update.

#### **Input Format:**

• Table StudentEnrollments with the same structure.

## **Output Format:**

User A sees the old value during the transaction.
User B successfully updates the row without waiting.

#### **Constraints:**

- Use databases that support MVCC (e.g., PostgreSQL, MySQL InnoDB).
- Avoid SELECT FOR UPDATE; use normal SELECT in repeatable read or snapshot isolation mode.

#### **Sample Input:**

student_id	student_name	course_id	enrollment_date
1	Ashish	CSE101	2024-06-01

## **Sample Output:**

- User A sees: enrollment\_date = 2024-06-01
- User B updates to: 2024-07-10
- User A continues to see the old value in the transaction until commit.

#### **Explanation of Output:**

MVCC ensures User A reads a consistent snapshot taken at the start of the transaction, unaffected by concurrent updates. This enables non-blocking concurrency.

#### **Query:**

**DROP TABLE IF EXISTS StudentEnrollments;** 

```
CREATE TABLE StudentEnrollments (
student_id INT PRIMARY KEY,
student_name VARCHAR(100),
course_id VARCHAR(10),
enrollment_date DATE
) ENGINE=InnoDB;
```

**INSERT INTO StudentEnrollments (student\_id, student\_name, course\_id, enrollment\_date)** 

**VALUES** 

```
(1, 'Ashish', 'CSE101', '2024-06-01');
```

START TRANSACTION;

SELECT student\_id, student\_name, course\_id, enrollment\_date

FROM StudentEnrollments

```
WHERE student_id = 1;
```

#### **UPDATE StudentEnrollments**

**SET enrollment\_date = '2024-07-10'** 

WHERE student\_id = 1;

#### **COMMIT**;

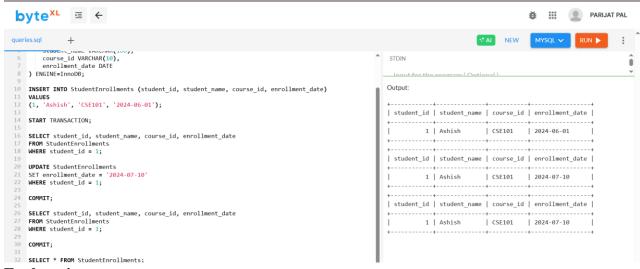
SELECT student\_id, student\_name, course\_id, enrollment\_date

FROM StudentEnrollments

WHERE student\_id = 1;

# **COMMIT**;

## **SELECT \* FROM StudentEnrollments;OUTPUT:**



#### **Explanation:**

- User A starts a transaction and reads a row.
- User B updates the same row and commits.
- User A still sees the old value until they commit.

## Part C: Comparing Behavior With and Without MVCC in High-

#### **Concurrency Description:**

Evaluate how MVCC vs. traditional locking behaves when multiple users access the same row for read and write. Use SELECT FOR UPDATE to demonstrate blocking in a non-MVCC system and contrast that with MVCC-based reads and updates.

#### **Input Format:**

Same StudentEnrollments table and data.

### **Output Format:**

#### Two scenarios:

- With Locking: Readers are blocked until the writer commits.
- With MVCC: Readers get consistent data without blocking.

#### **Constraints:**

- MVCC-supported database (e.g., PostgreSQL).
- Use different isolation levels or query techniques to simulate both cases.

## **Sample Input:**

student_id	student_name	course_id	enrollment_date
1	Ashish	CSE101	2024-06-01

# **Sample Output:**

- Without MVCC: Reader blocks until writer commits.
- With MVCC: Reader sees 2024-06-01 even while the writer updates to 2024-07-10.

## **Explanation of Output:**

- · Traditional locking causes blocking and delays.
- MVCC enables concurrent operations with no blocking, ensuring performance and consistency.

# **Query:**

DO SLEEP(2);

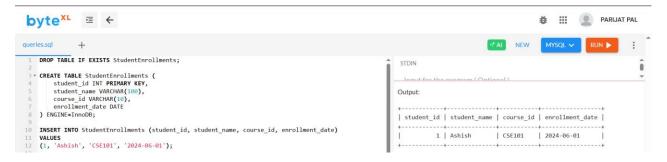
DROP TABLE IF EXISTS StudentEnrollments;

```
CREATE TABLE StudentEnrollments (
student_id INT PRIMARY KEY,
student_name VARCHAR(100),
course_id VARCHAR(10),
enrollment_date DATE
) ENGINE=InnoDB;

INSERT INTO StudentEnrollments (student_id, student_name, course_id, enrollment_date)
VALUES
(1, 'Ashish', 'CSE101', '2024-06-01');
START TRANSACTION;
SELECT * FROM StudentEnrollments WHERE student_id = 1 FOR UPDATE;
```

```
UPDATE StudentEnrollments
SET enrollment_date = '2024-07-10'
WHERE student_id = 1;
COMMIT;
START TRANSACTION;
SET TRANSACTION ISOLATION LEVEL REPEATABLE READ;
SELECT * FROM StudentEnrollments WHERE student_id = 1;
UPDATE StudentEnrollments
SET enrollment_date = '2024-08-15'
WHERE student_id = 1;
COMMIT;
SELECT * FROM StudentEnrollments WHERE student_id = 1;
COMMIT;
SELECT * FROM StudentEnrollments;
```

# **Output:**



# **Explanation:**

- With SELECT FOR UPDATE (locking): Readers block if a writer has locked the row.
- With normal SELECT in REPEATABLE READ (MVCC): Readers see a consistent snapshot while writers update concurrently.