
Assembly Language

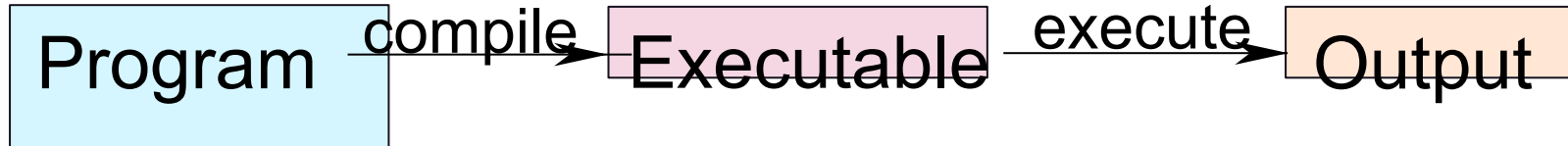
What Can a Computer Understand ?

- * Computer can clearly **NOT** understand instructions of the form
 - * Compute the determinant of a matrix
 - * Find the shortest path between Mumbai and Delhi

The Language of Instructions

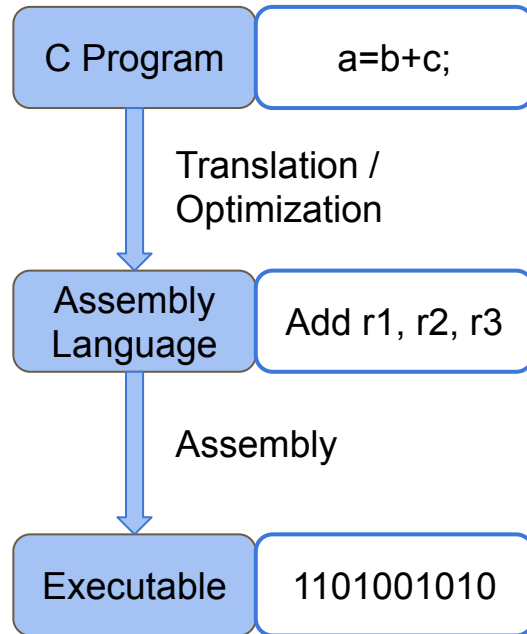
- * Humans can understand complicated sentences
- * Computers can understand
 - * Very simple instructions
 - * The semantics of all the instructions supported by a processor is known as its ***instruction set architecture*** (ISA). This includes the semantics of the instructions themselves, along with their operands, and interfaces with peripheral devices.

How to Instruct a Computer ?



- * Write a program in a high level language – C, C++, Java
- * **Compile** it into a format that the computer understands
- * Execute the program

Compilation and Assembly



Compilation

Compilation and Assembly

- Write a simple program, and compile it using gcc
- The resultant *a.out* is in machine language, and in binary form
- Try `gcc -S test.c -o test.s` . *test.s* is the assembly program
- To get the binary from the assembly program, run `gcc test.s`

Popular Instruction Set Architectures (ISAs)

x86	Laptops, desktops, servers
ARM	Mobiles, Raspberry Pi
SPARC v8	Leon3
PowerPC	IBM machines
RISC-V	IITM's Shakti processors
SimpleRISC	CS301
ToyRISC	CS311

Features of an ISA

- * **Complete**

- * It should be able to implement all the programs that users may write.

- * **Concise**

- * The instruction set should have a limited size. Typically an ISA contains 32-1000 instructions.

Features of an ISA – II

* Generic

- * Instructions should not be too specialized, e.g. add14 (adds a number with 14) is too specialized

* Simple

- * Should not be very complicated.

Designing an ISA

- * Important questions that need to be answered :
 - * How many instructions should we have ?
 - * What should they do ?
 - * How complicated should they be ?

RISC vs CISC

A ***reduced instruction set computer*** (RISC) implements simple instructions that have a simple and regular structure. The number of instructions is typically a small number (64 to 128). Examples: ARM, IBM PowerPC, HP PA-RISC

A ***complex instruction set computer*** (CISC) implements complex instructions that are highly irregular, take multiple operands, and implement complex functionalities. Secondly, the number of instructions is large (typically 500+). Examples: Intel x86, VAX

Let us now design an ISA ...

- * Single Instruction ISA

- * sbn – subtract and branch if negative

```
1: sbn a, b, 3    // a = a - b; if a < 0, jump to 3
2: sbn c, d, 3
3: sbn e, f, 1
...
```

Let us now design an ISA ...

- * Add ($a + b$) (assume $\text{temp} = 0$)

```
1: sbn temp, b, 2  
2: sbn a, temp, exit
```

Single Instruction ISA - II

*

Initialization:

one = 1

index = 10

sum = 0

1: sbn temp, temp, 2

2: sbn temp, index, 3

3: sbn sum, temp, 4

4: sbn index, one, exit

5: sbn temp, temp, 6

6: sbn temp, one, 1

exit

Single Instruction ISA - II

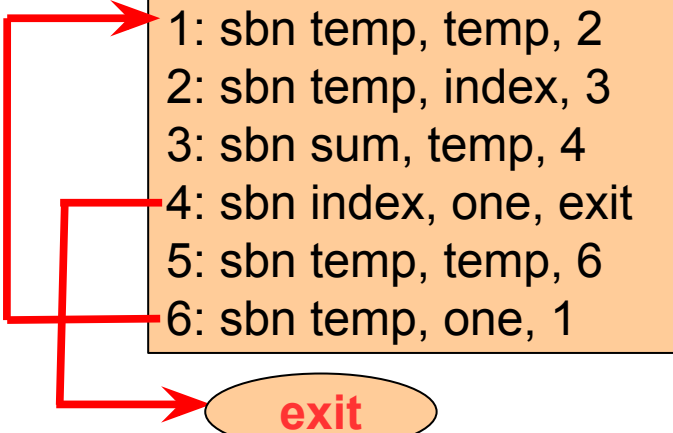
*

Initialization:

one = 1

index = 10

sum = 0



```
graph TD; L1[1: sbn temp, temp, 2] --> L2[2: sbn temp, index, 3]; L2 --> L3[3: sbn sum, temp, 4]; L3 --> L4[4: sbn index, one, exit]; L4 --> L5[5: sbn temp, temp, 6]; L5 --> L6[6: sbn temp, one, 1]; L6 --> Exit([exit]); Exit --> L1;
```

1: sbn temp, temp, 2	// temp = 0
2: sbn temp, index, 3	// temp = -1 * index
3: sbn sum, temp, 4	// sum += index
4: sbn index, one, exit	// index -= 1
5: sbn temp, temp, 6	// temp = 0
6: sbn temp, one, 1	// (0 - 1 < 0), hence goto 1

Single Instruction ISA - II

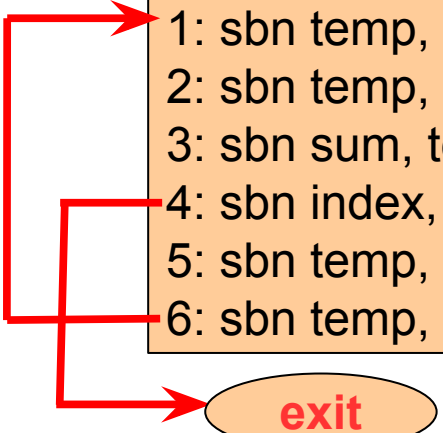
* Add the numbers – 1 ... 10

Initialization:

one = 1

index = 10

sum = 0



```
graph TD; 1[1: sbn temp, temp, 2] --> 2[2: sbn temp, index, 3]; 2 --> 3[3: sbn sum, temp, 4]; 3 --> 4[4: sbn index, one, exit]; 4 --> 5[5: sbn temp, temp, 6]; 5 --> 6[6: sbn temp, one, 1]; 6 --> exit([exit]);
```

1: sbn temp, temp, 2	// temp = 0
2: sbn temp, index, 3	// temp = -1 * index
3: sbn sum, temp, 4	// sum += index
4: sbn index, one, exit	// index -= 1
5: sbn temp, temp, 6	// temp = 0
6: sbn temp, one, 1	// (0 - 1 < 0), hence goto 1

Single Instruction ISA - III

* Find whether a number is positive

- The given number is at address 'number'
- If num is positive, write '0' to address 'result'; if negative, write '-1'
- Initialization: 'one' = 1

Single Instruction ISA - III

* Find whether a number is positive

- The given number is at address 'number'
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- Initialization: 'one' = 1

```
1: sbn result, result, 2      // result = 0
2: sbn temp, temp, 3          // temp = 0
3: sbn temp, number, 5        // if positive, then branch
4: sbn result, one, 5          // if negative, result = -1
5: exit
```

Multiple Instruction ISA

- * Arithmetic and Logical Instructions

- * add, subtract, multiply, divide, or, and, not

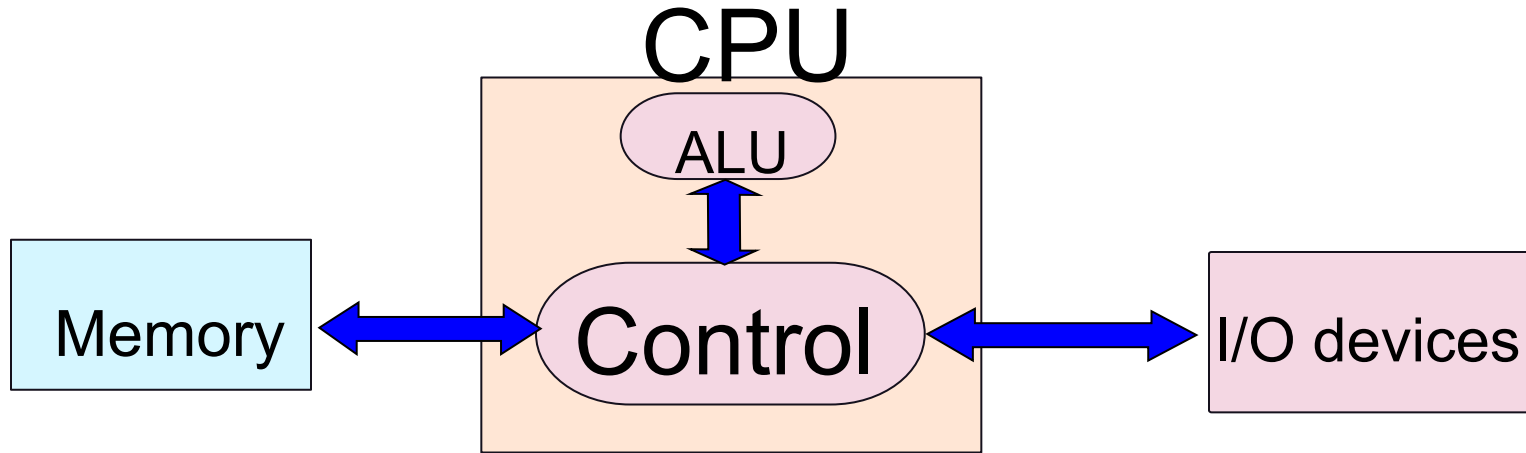
- * Move instructions

- * Transfer values between memory locations

- * Branch instructions

- * Move to a new program location, based on the values of some memory locations

Von-Neumann Architecture



Problems with Harvard/ Von-Neumann Architectures

- * The memory is assumed to be one large array of bytes



General Rule: Larger is a structure, slower it is

- * **Solution:**

- * Have a small array of named locations (**registers**) that can be used by instructions



Insight: Accesses exhibit locality (tend to use the same variables frequently in the same window of time)

Uses of Registers

- * A CPU (Processor) contains set of registers (16-64)
- * These are named storage locations.
- * Typically values are loaded from memory to registers.
- * Arithmetic/logical instructions use registers as input operands
- * Finally, data is stored back into their memory locations.

Example of a Program in Machine Language with Registers

```
1: r1 = mem[b] // load b
2: r2 = mem[c] // load c
3: r3 = r1 + r2 // add b and c
4: mem[a] = r3 // save the result
```

- * r1, r2, and r3, are registers
- * mem → array of bytes representing memory

Machine with Registers

