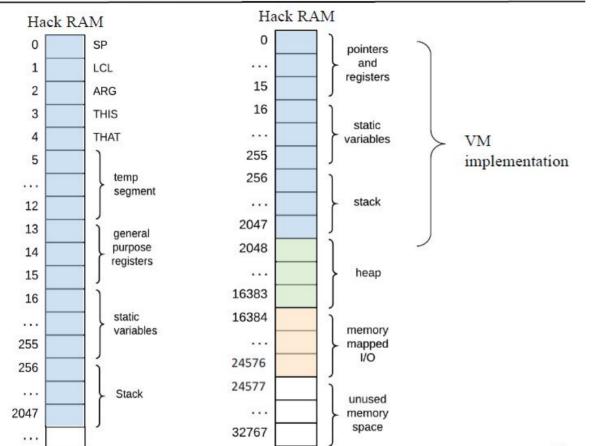
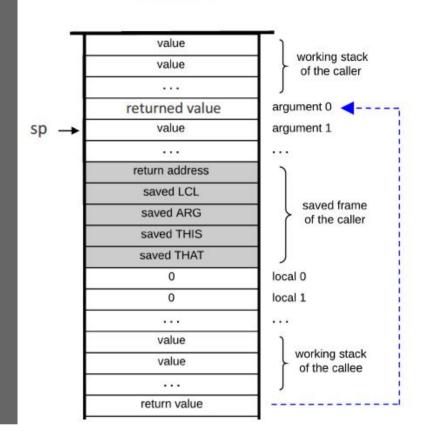
Standard mapping of the VM on the Hack platform



return



Name	Location	Usage	
SP	RAM[0]	Stack Pointer: the memory address just following the memory address containing the topmost stack value	
LCL	RAM[1]	Base address of the local segment	
ARG	RAM[2]	Base address of the argument segment	
THIS	RAM[3]	Base address of the this segment	
THAT	RAM[4]	Base address of the that segment	
TEMP	RAM[5-12]	Holds the temp segment	
R13 R14 R15	RAM[13-15]	If the assembly code generated by the VM translator needs variables, it can use these registers.	

VM command	Assembly (pseudo) code, generated by the VM translator		
call f nArgs (calls a function f, informing that nArgs arguments were pushed to the stack before the call)	push returnAddress push LCL push ARG push THIS push THAT ARG = SP-5-nArgs LCL = SP goto f (returnAddress)	// generates a label and pushes it to the stack // saves LCL of the caller // saves ARG of the caller // saves THIS of the caller // saves THAT of the caller // repositions ARG // repositions LCL // transfers control to the callee // injects the return address label into the code	
function f nVars (declares a function f, informing that the function has nVars local variables)	(f) repeat nVars times: push 0	// injects a function entry label into the code // nVars = number of local variables // initializes the local variables to 0	
return (terminates the current function and returns control to the caller)	<pre>frame = LCL retAddr = *(frame-5) *ARG = pop() SP = ARG+1 THAT = *(frame-1) THIS = *(frame-2) ARG = *(frame-3) LCL = *(frame-4) goto retAddr</pre>	// frame is a temporary variable // puts the return address in a temporary variable // repositions the return value for the caller // repositions SP for the caller // restores THAT for the caller // restores THIS for the caller // restores ARG for the caller // restores LCL for the caller // go to the return address	

Figure 8.5 Implementation of the function commands of the VM language. All the actions described on the right are realized by generated Hack assembly instructions.

Special symbols in translated VM programs

Symbol	Usage	
SP	This predefined symbol points to the memory address within the host RAM just following the address containing the topmost stack value.	
LCL, ARG, THIS, THAT	These predefined symbols point, respectively, to the base addresses within the host RAM of the virtual segments local, argument, this, and that of the currently running VM function.	
R13-R15	These predefined symbols can be used for any purpose.	
Xxx.i symbols	Each static variable <i>i</i> in file Xxx.vm is translated into the assembly symbol Xxx. <i>j</i> ., where <i>j</i> is incremented each time a new static variable is encountered in the file Xxx.vm. In the subsequent assembly process, these symbolic variables will be allocated to the RAM by the Hack assembler.	
functionName\$label	Let foo be a function within a VM file Xxx. Each label bar command within foo should generate and insert into the assembly code stream a symbol Xxx.foo\$bar.	
	When translating goto bar and if-goto bar commands (within foo) into assembly, the full label specification Xxx.foo\$bar must be used instead of bar.	
functionName	Each function foo command within a VM file Xxx should generate and insert into the assembly code stream a symbol Xxx.foo that labels the entry point to the function's code. In the subsequent assembly process, the assembler will translate this symbol into the physical memory address where the function code starts.	
functionName\$ret.i	Let foo be a function within a VM file Xxx.	
	Within foo, each function call command should generate and insert into the assembly code stream a symbol Xxx.foo\$ret.i, where i is a running integer (one such symbol should be generated for each call command within foo).	
	This symbol serves as the return address to the calling function. In the subsequent assembly process, the assembler will translate this symbol into the physical memory address of the command immediately after the function call command.	