

Chapter 4: Threads & Concurrency





Chapter 4: Threads

- Overview
- Multicore Programming
- Concurrency vs. Parallelism
- Amdahl's Law
- Multithreading Models





Objectives

- Identify the basic components of a thread, and contrast threads and processes
- Describe the benefits and challenges of designing multithreaded applications





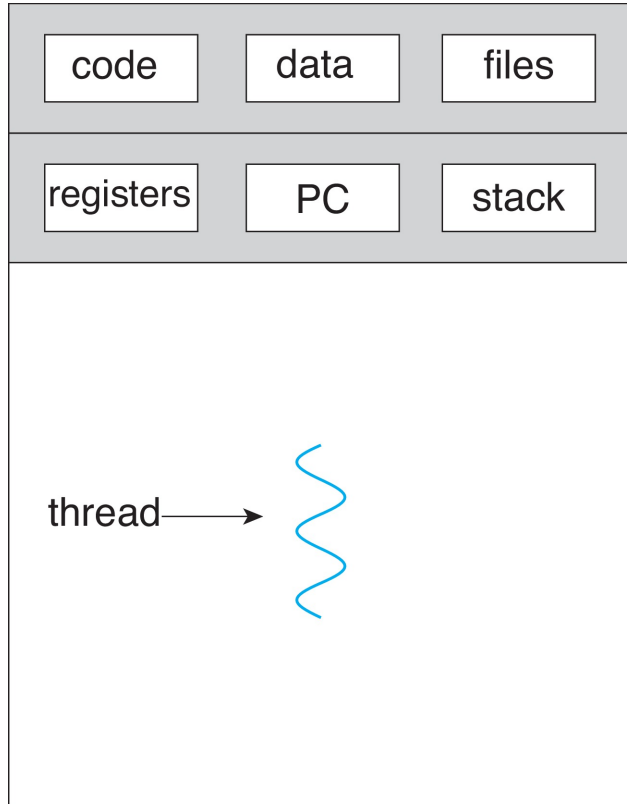
Motivation

- Most modern applications are multithreaded
- Threads run within application
- Multiple tasks with the application can be implemented by separate threads
 - Update display
 - Fetch data
 - Spell checking
 - Answer a network request
- Process creation is heavy-weight while thread creation is light-weight
- Can simplify code, increase efficiency
- Kernels are generally multithreaded

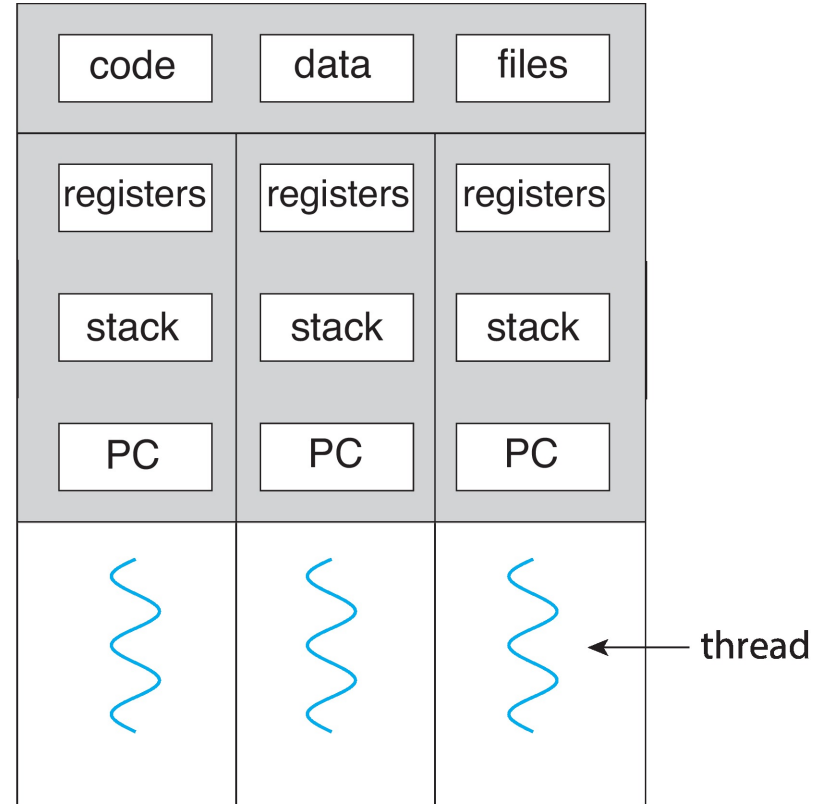




Single and Multithreaded Processes



single-threaded process

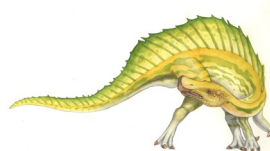
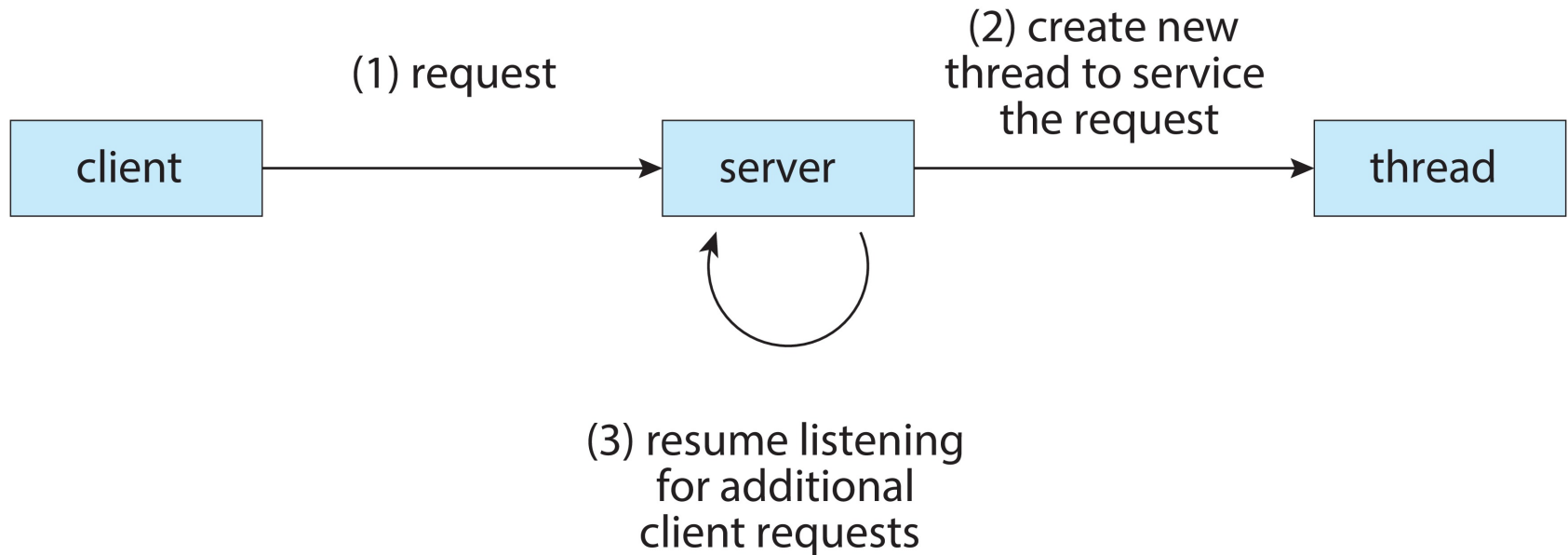


multithreaded process





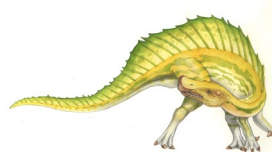
Multithreaded Server Architecture





Benefits

- ❑ **Responsiveness** – may allow continued execution if part of process is blocked, especially important for user interfaces
- ❑ **Resource Sharing** – threads share resources of process, easier than shared memory or message passing
- ❑ **Economy** – cheaper than process creation, thread switching lower overhead than context switching
- ❑ **Scalability** – process can take advantage of multicore architectures





Multicore Programming

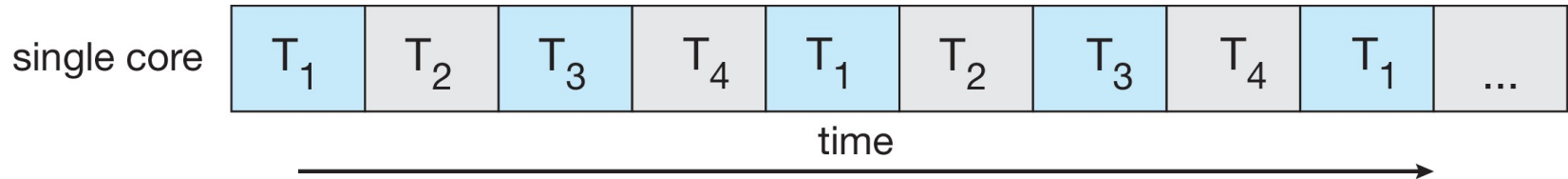
- **Multicore** or **multiprocessor** systems putting pressure on programmers, challenges include:
 - **Dividing activities**
 - **Balance**
 - **Data splitting**
 - **Data dependency**
 - **Testing and debugging**
- *Parallelism* implies a system can perform more than one task simultaneously
- *Concurrency* supports more than one task making progress
 - Single processor / core, scheduler providing concurrency



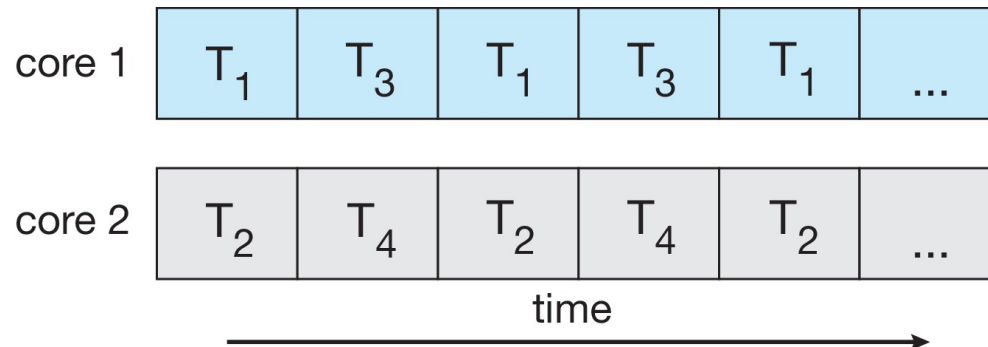


Concurrency vs. Parallelism

□ Concurrent execution on single-core system:



□ Parallelism on a multi-core system:





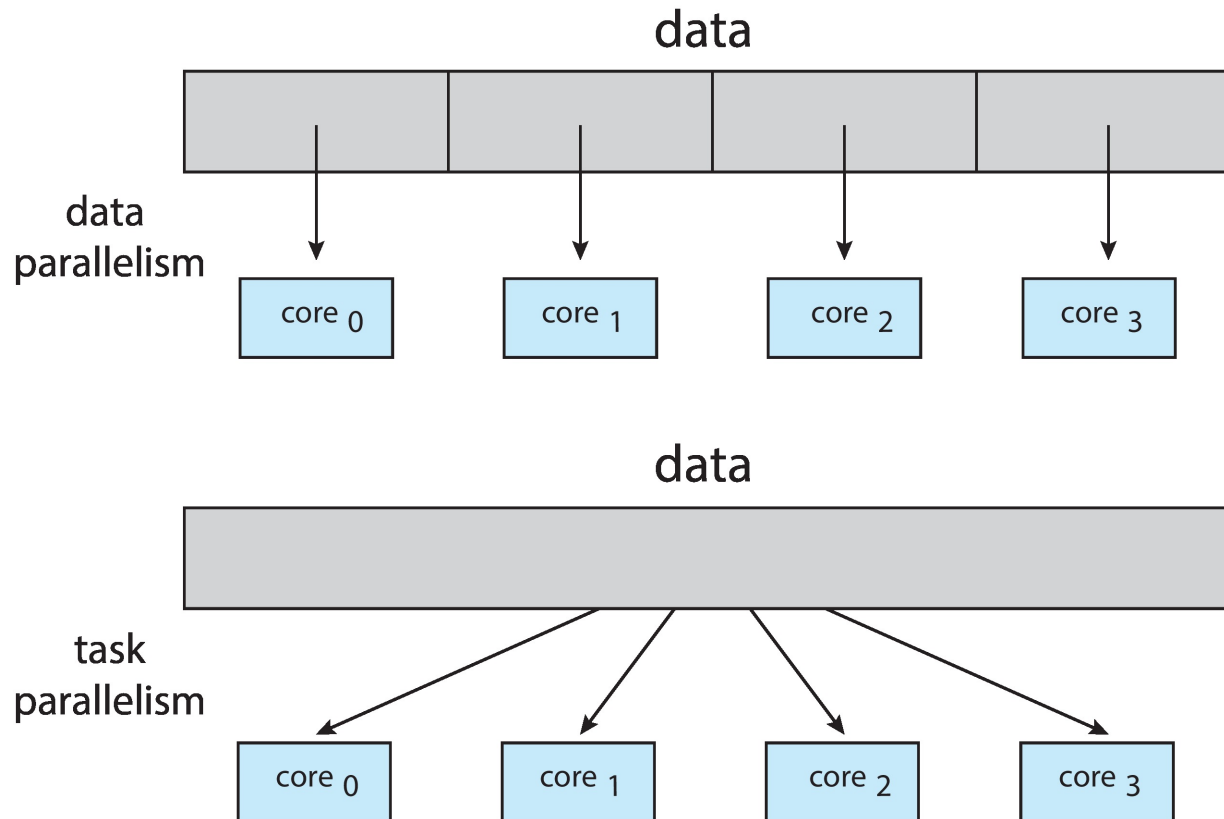
Multicore Programming

- Types of parallelism
 - **Data parallelism** – distributes subsets of the same data across multiple cores, same operation on each
 - **Task parallelism** – distributing threads across cores, each thread performing unique operation





Data and Task Parallelism





Amdahl's Law

- Identifies performance gains from adding additional cores to an application that has both serial and parallel components
- S is serial portion
- N processing cores

$$speedup \leq \frac{1}{S + \frac{(1-S)}{N}}$$

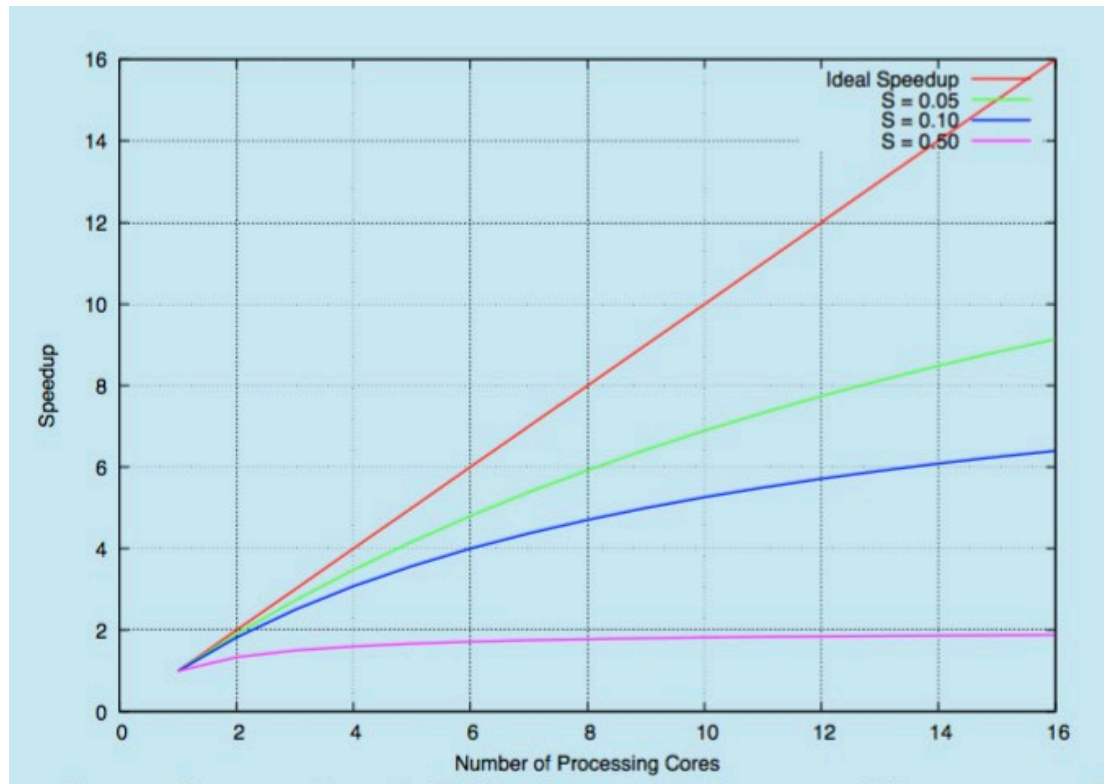
- That is, if application is 75% parallel / 25% serial, moving from 1 to 2 cores results in speedup of 1.6 times
- As N approaches infinity, speedup approaches $1 / S$

Serial portion of an application has disproportionate effect on performance gained by adding additional cores





Amdahl's Law





Multithreading Models

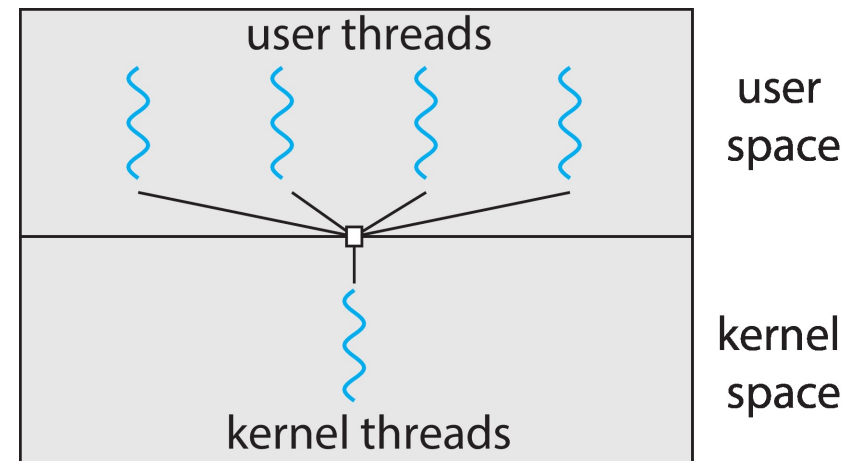
- Many-to-One
- One-to-One
- Many-to-Many





Many-to-One

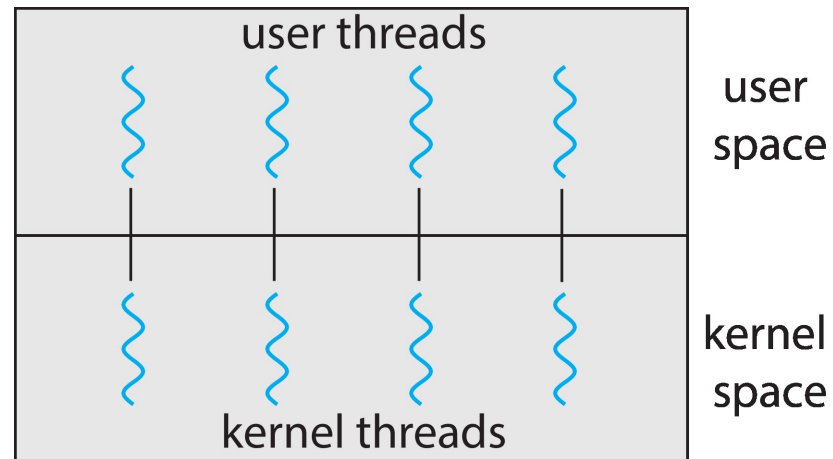
- ❑ Many user-level threads mapped to single kernel thread
- ❑ One thread blocking causes all to block
- ❑ Multiple threads may not run in parallel on multicore system because only one may be in kernel at a time
- ❑ Few systems currently use this model
- ❑ Examples:
 - ❑ **Solaris Green Threads**
 - ❑ **GNU Portable Threads**





One-to-One

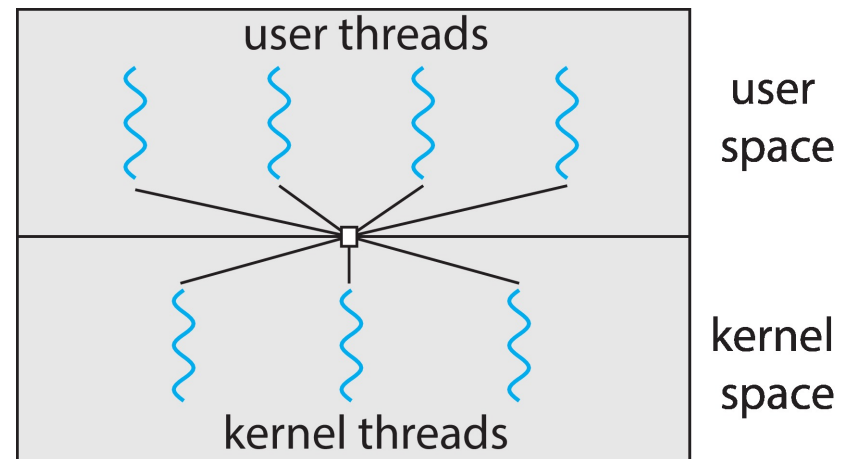
- Each user-level thread maps to kernel thread
- Creating a user-level thread creates a kernel thread
- More concurrency than many-to-one
- Number of threads per process sometimes restricted due to overhead
- Examples
 - Windows
 - Linux





Many-to-Many Model

- ❑ Allows many user level threads to be mapped to many kernel threads
- ❑ Allows the operating system to create a sufficient number of kernel threads
- ❑ Not very common



End of Chapter 4

