Arrows for Parallel Computations

MARTIN BRAUN

University Bayreuth, 95440 Bayreuth, Germany

Glasgow Un 2 8QQ, Scotland OLEG LO ΕV

euth, 9544 University euth, Germany

Abs

Arrows formOL: are? a general interface for nputation and pose therefore as an alternative to monads for API design. We express parallelist using this concept. This is a new way to represent parallel computation. We define an Arrows-b ed interface for parallelism and implement it using multiple parallel Haskells. OL: Benefits: In this manner we are able to bridge across various parallel Haskells with a common interface.

benefit of being portable across flavours of parallel This new way of writing parallel program, Haskells.OL: Wdh? Each parallel comput an arrow, they can be composed and transformed as such. We thus introduce some syntactic sug to provide parallelism-aware arrow combinators.

We also define several parallel skeletons with our framework. Benchmarks shows that our framework does not induce too much overhead performance-wise.

Contents

Contents

1 Introduction

OL: todo, reuse 5.5, "Impact" at the end and more

blablabla arrows, parallel, haskell.

Contribution HIT HERE REALLY STRONG

Structure The remaining text is structures as follows. Section ?? briefly introduces known parallel Haskell flavours and gives an overview of Arrows to the reader (Sec. ??). Section ?? discusses related work. Section ?? defines Parallel Arrows and presents a basic interface. Section ?? defines Futures for Parallel Arrows, this concept enables better communication. Section ?? presents some basic algorithmic skeletons (parallel map with and without load

main

balancing, map-reduce) in our newly defined dialect. More advanced ones are showcased in Section ?? (pipe, ring, torus). Section ?? shows the benchmark results. Section ?? discusses future work and concludes.

2 Background

2.1 Show on to parallel Haskells

There are already several w rograms in Haskell. As we will base our s to write pai lel Haskells parallel arrows on existing vill now give a short introduction to the ones we use as backends in baper.

In its purest form, parallel computation (actions) can be looked at as the execution of some functions a -> b in parallel, as a igure ?? symbolically shows:

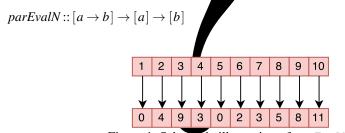


Figure 1: Schematic illustration of parEvalN.

Before we go into detail on how we can use this idea of parallelism for parallel Arrows, as a short introduction to parallelism in Haskell we will now implement parEvalN with several different parallel Haskells.

2.1.1 Multicore Haskell

Multicore Haskell (??) is way to do parallel processing found in standard GHC. It ships with parallel evaluation strategies (??) for several types which can be applied with using :: a -> Strategy a -> a. For parEvalN this means that we can just apply the list of functions [a -> b] to the list of inputs [a] by zipping them with the application operator \$. We then evaluate this lazy list [b] according to a Strategy [b] with the using :: a -> Strategy a -> a operator. We construct this strategy with parList :: Strategy a -> Strategy [a] and rdeepseq :: NFData a => Strategy a where the latter is a strategy which evalutes to normal form. To ensure that programs that use parEvalN have the correct evaluation order, we annotate the computation with pseq :: a -> b -> b which forces the compiler to not reorder multiple parEvalN computations. This is particularly necessary in circular communication topologies like in the torus or ring skeleton that we will see in chapter ?? which resulted in deadlock scenarios when executed without pseq during testing for this paper.

Multicore Haskell on Hackage is available under https://hackage.haskell.org/package/ parallel-3.2.1.0, compiler support is integrated in the stock GHC.