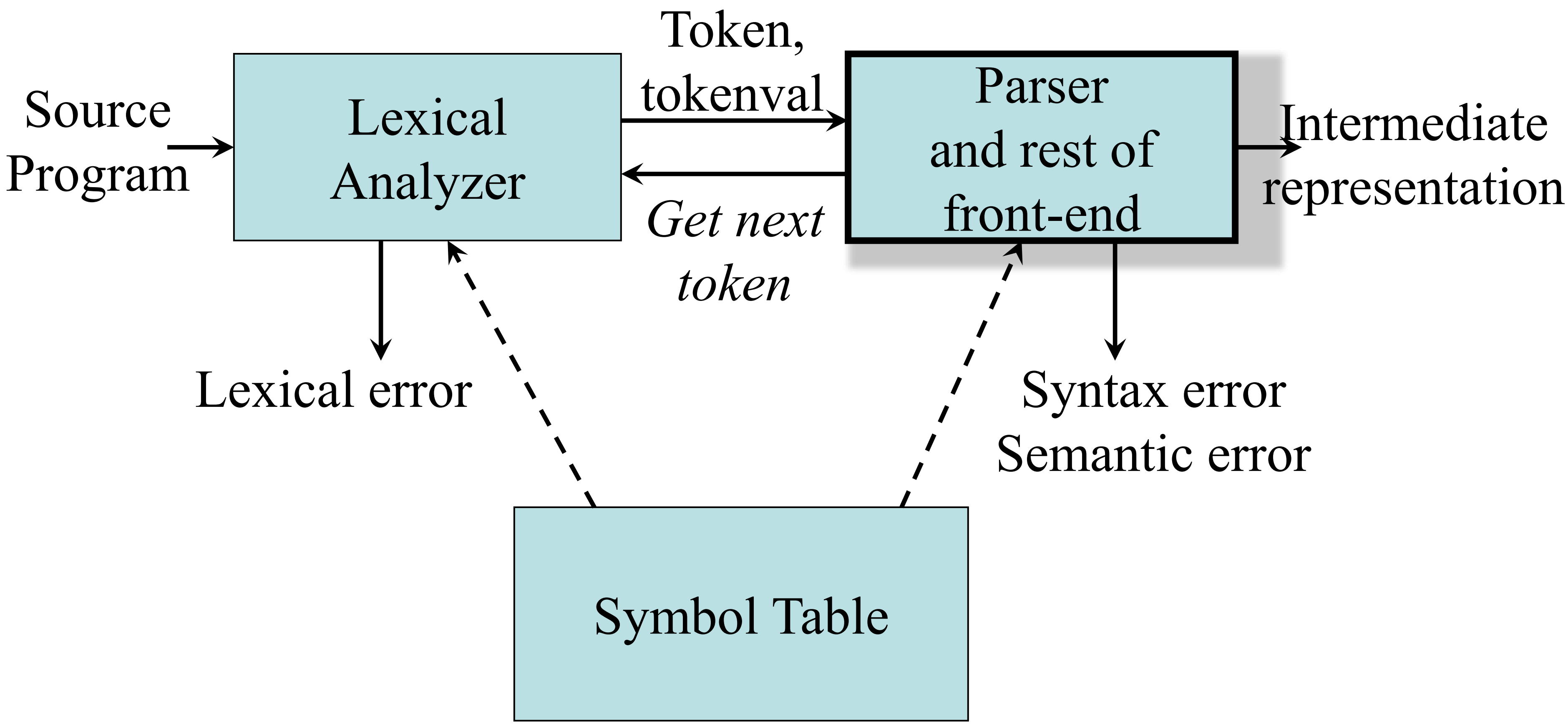
Syntax Analysis

Position of a Parser in the Compiler Model



The Parser

- A parser implements a C-F grammar as a recognizer of strings
- The role of the parser in a compiler is twofold:

1. To check syntax (= string recognizer)
    - And to report syntax errors accurately
2. To invoke semantic actions
    - For static semantics checking, e.g. type checking of expressions, functions, etc.
    - For syntax-directed translation of the source code to an intermediate representation

Syntax-Directed Translation

- One of the major roles of the parser is to produce an *intermediate representation* (IR) of the source program using *syntax-directed translation* methods
- Possible IR output:
  - Abstract syntax trees (ASTs)
  - Control-flow graphs (CFGs) with triples, three-address code, or register transfer list notation
  - WHIRL (SGI Pro64 compiler) has 5 IR levels!

Error Handling

- A good compiler should assist in identifying and locating errors
  - *Lexical errors*: important, compiler can easily recover and continue
  - *Syntax errors*: most important for compiler, can almost always recover
  - *Static semantic errors*: important, can sometimes recover
  - *Dynamic semantic errors*: hard or impossible to detect at compile time, runtime checks are required
  - *Logical errors*: hard or impossible to detect

Viable-Prefix Property

- The *viable-prefix property* of parsers allows early detection of syntax errors
  - Goal: detection of an error *as soon as possible* without further consuming unnecessary input
  - How: detect an error as soon as the prefix of the input does not match a prefix of any string in the language

Prefix

```
...
for (;)
...
```

Error is detected here

Prefix

```
...
DO 10 I = 1;0
...
```

Error is detected here

Error Recovery Strategies

- *Panic mode*
  - Discard input until a token in a set of designated synchronizing tokens is found
- *Phrase-level recovery*
  - Perform local correction on the input to repair the error
- *Error productions*
  - Augment grammar with productions for erroneous constructs
- *Global correction*
  - Choose a minimal sequence of changes to obtain a global least-cost correction

Grammars (Recap)

- Context-free grammar is a 4-tuple $G = (N, T, P, S)$ where
  - T is a finite set of tokens (*terminal* symbols)
  - N is a finite set of *nonterminals*
  - P is a finite set of *productions* of the form

    $\alpha \to \beta$

    where $\alpha \in (N \cup T)^* \, N \, (N \cup T)^*$ and $\beta \in (N \cup T)^*$
  - $S \in N$ is a designated *start symbol*

Notational Conventions Used

- Terminals
  $a,b,c,\ldots \in T$
  specific terminals: **0**, **1**, **id**, **+**
- Nonterminals
  $A,B,C,\ldots \in N$
  specific nonterminals: *expr*, *term*, *stmt*
- Grammar symbols
  $X,Y,Z \in (N \cup T)$
- Strings of terminals
  $u,v,w,x,y,z \in T^*$
- Strings of grammar symbols
  $\alpha,\beta,\gamma \in (N \cup T)^*$

Derivations (Recap)

- The *one-step derivation* is defined by
  $\alpha\, A\, \beta \Rightarrow \alpha\, \gamma\, \beta$ id+id
  where $A \rightarrow \gamma$ is a production in the grammar
- In addition, we define
  - $\Rightarrow$ is *leftmost* $\Rightarrow_{lm}$ if α does not contain a nonterminal
  - $\Rightarrow$ is *rightmost* $\Rightarrow_{rm}$ if β does not contain a nonterminal
  - Transitive closure $\Rightarrow^*$ (zero or more steps)
  - Positive closure $\Rightarrow^+$ (one or more steps)
- The *language generated by G* is defined by
  $L(G) = \{w \in T^* \mid S \Rightarrow^+ w\}$

Derivation (Example)

Grammar $G = (\{E\}, \{+,*,(,),-,\mathbf{id}\}, P, E)$ with productions $P =$

$$
\begin{aligned}
E \rightarrow\ & E + E \\
|\ & E * E \\
|\ & \mathbf{(}\ E\ \mathbf{)} \\
|\ & \mathbf{-}\ E \\
|\ & \mathbf{id}
\end{aligned}
$$

Example derivations:

$$E \Rightarrow \mathbf{-}\ E \Rightarrow \mathbf{-\ id}$$

$$E \Rightarrow_{rm} E + E \Rightarrow_{rm} E + \mathbf{id} \Rightarrow_{rm} \mathbf{id + id}$$

$$E \Rightarrow^{*} E$$

$$E \Rightarrow^{*} \mathbf{id + id}$$

$E \Rightarrow_{\mathrm{lm}}$ =>E*E=>**id***E=>id*E+E=>**id*id**+E=>**id * id + id**

Chomsky Hierarchy: Language Classification

- A grammar G is said to be
  - *Regular* if it is *right linear* where each production is of the form
    $A \rightarrow w\, B$ or $A \rightarrow w$
    or *left linear* where each production is of the form
    $A \rightarrow B\, w$ or $A \rightarrow w$
  - *Context free* if each production is of the form
    $A \rightarrow \alpha$
    where $A \in N$ and $\alpha \in (N \cup T)^*$
  - *Context sensitive* if each production is of the form
    $\alpha\, A\, \beta \rightarrow \alpha\, \gamma\, \beta$
    where $A \in N$, $\alpha,\gamma,\beta \in (N \cup T)^*$, $|\gamma| > 0$
  - *Unrestricted*

Chomsky Hierarchy

$\mathbb{L}(\textit{regular}) \subset \mathbb{L}(\textit{context free}) \subset \mathbb{L}(\textit{context sensitive}) \subset \mathbb{L}(\textit{unrestricted})$

Where $\mathbb{L}(T) = \{ L(G) \mid G \text{ is of type } T \}$
That is: the set of all languages
generated by grammars G of type T

Examples:

Every *finite language* is regular! (construct a FSA for strings in $L(G)$)

$L_1 = \{ \mathbf{a}^n\mathbf{b}^n \mid n \geq 1 \}$ is context free

$L_2 = \{ \mathbf{a}^n\mathbf{b}^n\mathbf{c}^n \mid n \geq 1 \}$ is context sensitive

Parsing

- *Universal* (any C-F grammar)
  - Cocke-Younger-Kasimi
  - Earley
- *Top-down* (C-F grammar with restrictions)
  - Recursive descent (predictive parsing)
  - LL (Left-to-right, Leftmost derivation) methods
- *Bottom-up* (C-F grammar with restrictions)
  - Operator precedence parsing
  - LR (Left-to-right, Rightmost derivation) methods
    - SLR, canonical LR, LALR

Top-Down Parsing

- LL methods (Left-to-right, Leftmost derivation) and recursive-descent parsing

Grammar:

$E \rightarrow T + T$

$T \rightarrow \textbf{(}\ E\ \textbf{)}$

$T \rightarrow \textbf{-}\ E$

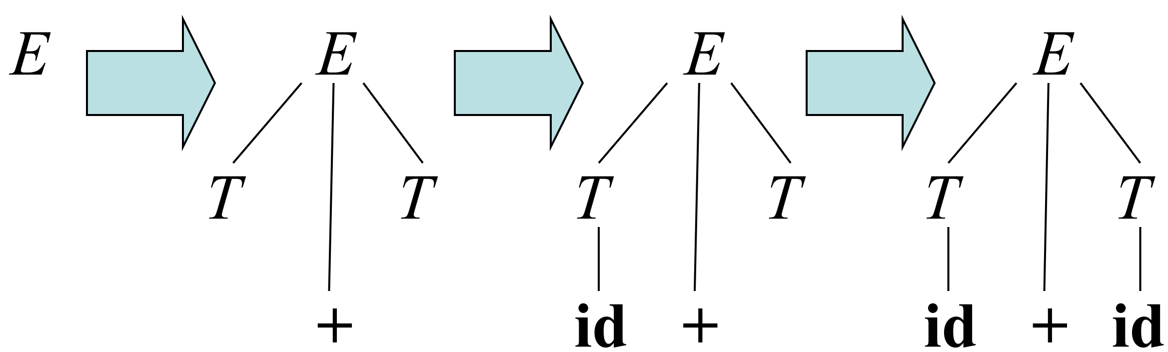
$T \rightarrow \textbf{id}$

Leftmost derivation:

$E \Rightarrow_{lm} T + T$

$\Rightarrow_{lm} \textbf{id} + T$

$\Rightarrow_{lm} \textbf{id} + \textbf{id}$



Left Recursion (Recap)

- Productions of the form

$$\begin{aligned} A \rightarrow\ & A\ \alpha \\ \mid\ & \beta \\ \mid\ & \gamma \end{aligned}$$

  are left recursive
- When one of the productions in a grammar is left recursive then a predictive parser loops forever on certain inputs

A General Systematic Left Recursion Elimination Method

Input: Grammar G with no cycles or ε-productions

Arrange the nonterminals in some order $A_1, A_2, \ldots, A_n$

for $i = 1, \ldots, n$ **do**

    for $j = 1, \ldots, i\text{-}1$ **do**

        replace each

            $A_i \rightarrow A_j\,\gamma$

        with

            $A_i \rightarrow \delta_1\,\gamma \mid \delta_2\,\gamma \mid \ldots \mid \delta_k\,\gamma$

        where

            $A_j \rightarrow \delta_1 \mid \delta_2 \mid \ldots \mid \delta_k$

    enddo

    eliminate the *immediate left recursion* in A_i

enddo

Immediate Left-Recursion Elimination

Rewrite every left-recursive production

$A \rightarrow A\ \alpha$ A$\rightarrow$S α | γ

$\quad | A\ \delta$ S$\rightarrow$A β

$\quad | \beta$

$\quad | \gamma$ S $\rightarrow$S$\alpha\beta$ | $\gamma\beta$

into a right-recursive production:

$A \rightarrow \beta\ A_R$ $S\rightarrow\gamma\beta$S'

$\quad | \gamma\ A_R$

$A_R \rightarrow \alpha\ A_R$ $S'\rightarrow\alpha\beta$S'| ε

$\quad | \delta\ A_R$

$\quad | \varepsilon$

Example Left Recursion Elim.

$$\left.\begin{aligned} A &\to B\ C \mid \mathbf{a} \\ B &\to C\ A \mid A\ \mathbf{b} \\ C &\to A\ B \mid C\ C \mid \mathbf{a} \end{aligned}\right\} \text{Choose arrangement: } A, B, C$$

$i = 1$: nothing to do

$i = 2, j = 1$: $B \to C\ A \mid \underline{A}\ \mathbf{b}$

$\Rightarrow$ $B \to C\ A \mid \underline{B\ C}\ \mathbf{b} \mid \underline{\mathbf{a}}\ \mathbf{b}$

$\Rightarrow_{\text{(imm)}}$ $B \to C\ A\ B_R \mid \mathbf{a}\ \mathbf{b}\ B_R$

$B_R \to C\ \mathbf{b}\ B_R \mid \varepsilon$

$i = 3, j = 1$: $C \to \underline{A}\ B \mid C\ C \mid \mathbf{a}$

$\Rightarrow$ $C \to \underline{B\ C}\ B \mid \underline{\mathbf{a}}\ B \mid C\ C \mid \mathbf{a}$

$i = 3, j = 2$: $C \to \underline{B}\ C\ B \mid \mathbf{a}\ B \mid C\ C \mid \mathbf{a}$

$\Rightarrow$ $C \to \underline{C\ A\ B_R}\ C\ B \mid \underline{\mathbf{a}\ \mathbf{b}\ B_R}\ C\ B \mid \mathbf{a}\ B \mid C\ C \mid \mathbf{a}$

$\Rightarrow_{\text{(imm)}}$ $C \to \mathbf{a}\ \mathbf{b}\ B_R\ C\ B\ C_R \mid \mathbf{a}\ B\ C_R \mid \mathbf{a}\ C_R$

$C_R \to A\ B_R\ C\ B\ C_R \mid C\ C_R \mid \varepsilon$

Left Factoring

- When a nonterminal has two or more productions whose right-hand sides start with the same grammar symbols, the grammar is not LL(1) and cannot be used for predictive parsing
- Replace productions

$$A \to \alpha\ \beta_1 \mid \alpha\ \beta_2 \mid \ldots \mid \alpha\ \beta_n \mid \gamma$$

with

$$A \to \alpha\ A_R \mid \gamma$$
$$A_R \to \beta_1 \mid \beta_2 \mid \ldots \mid \beta_n$$

Predictive Parsing

- Eliminate left recursion from grammar
- Left factor the grammar
- Compute FIRST and FOLLOW
- Two variants:
  - Recursive (recursive-descent parsing)
  - Non-recursive (table-driven parsing)

Recursive-Descent Parsing (Recap)

- Grammar must be LL(1)
- Every nonterminal has one (recursive) procedure responsible for parsing the nonterminal's syntactic category of input tokens
- When a nonterminal has multiple productions, each production is implemented in a branch of a selection statement based on input look-ahead information

Using FIRST and FOLLOW in a Recursive-Descent Parser

$expr \rightarrow term\ rest$
$rest \rightarrow +\ term\ rest$
$\quad\quad\ \mid -\ term\ rest$
$\quad\quad\ \mid \varepsilon$
$term \rightarrow \mathbf{id}$

procedure *rest*();
begin
  if *lookahead* in FIRST(+ *term rest*) **then**
    match('+'); *term*(); *rest*()
  else if *lookahead* in FIRST(- *term rest*) **then**
    match('-'); *term*(); *rest*()
  else if *lookahead* in FOLLOW(*rest*) **then**
    return
  else error**()**
end;

where FIRST(+ *term rest*) = { + }
FIRST(**-** *term rest*) = { **-** }
FOLLOW(*rest*) = { **$** }

B→B or T | T

T→T and F | F

F→not B | (B) | true | false | id

(id and true)

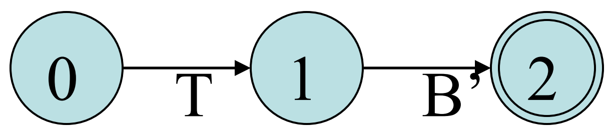
B→TB'

B'→or TB'| ε

T→FT'

T'→and FT'| ε

F→not B | (B) | true | false | id



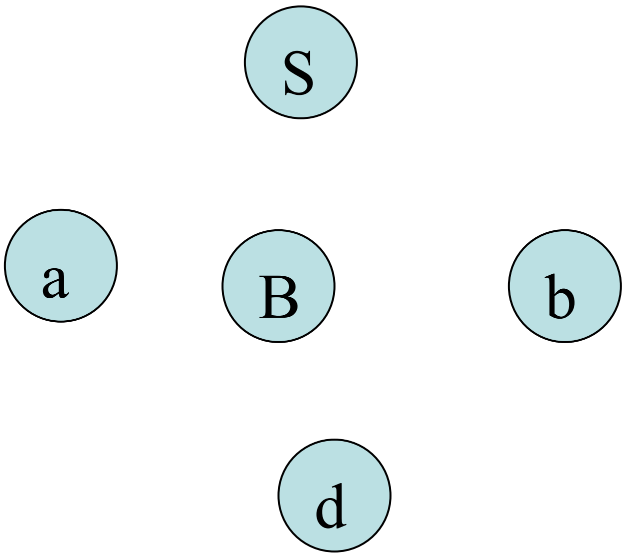
S→aBb

B→cd|d

```
S(){
Match a;
B();
Match b;
}
```

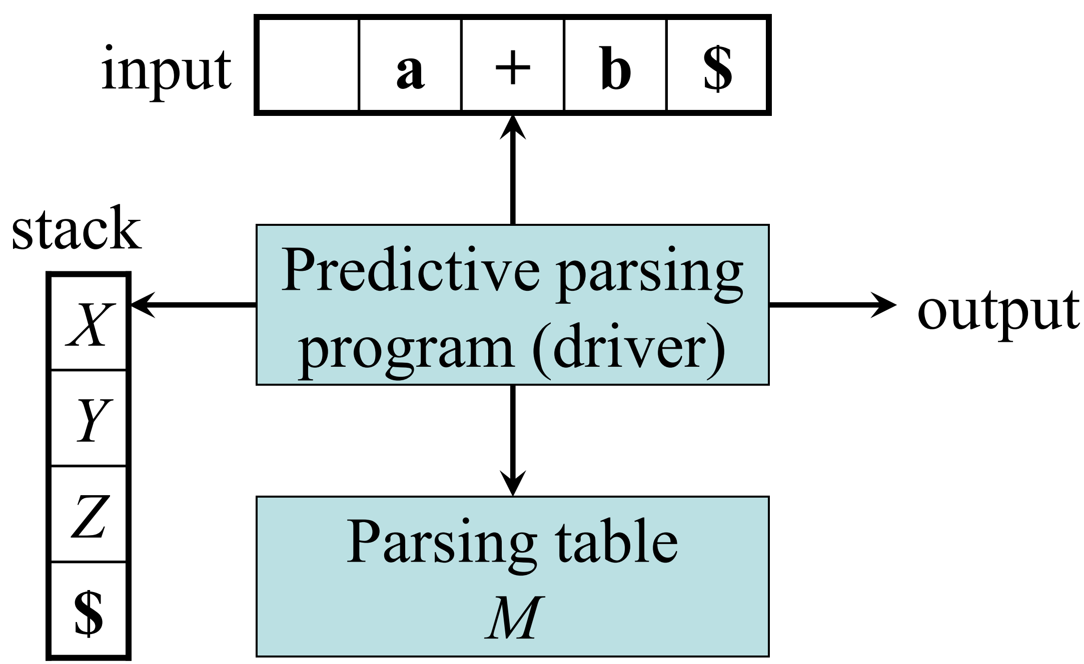
```
B()
(match c && match d) || match d;
```

adb



Non-Recursive Predictive Parsing: Table-Driven Parsing

- Given an LL(1) grammar $G = (N, T, P, S)$ construct a table $M[A,a]$ for $A \in N$, $a \in T$ and use a *driver program* with a *stack*



FIRST (Revisited)

- FIRST(α) = { *the set of terminals that begin all strings derived from* α }

FIRST(a) = $\{a\}$ if $a \in T$

FIRST(ε) = $\{\varepsilon\}$

FIRST(A) = $\cup_{A\to\alpha}$ FIRST(α) for $A\to\alpha \in P$

FIRST($X_1X_2\ldots X_k$) =

if for all $j = 1, \ldots, i\text{-}1 : \varepsilon \in$ FIRST(X_j) **then**
add non-ε in FIRST(X_i) to FIRST($X_1X_2\ldots X_k$)

if for all $j = 1, \ldots, k : \varepsilon \in$ FIRST(X_j) **then**
add ε to FIRST($X_1X_2\ldots X_k$)

$E \rightarrow E+T|T$

$T \rightarrow T*F|F$

$F \rightarrow (E)|id$

$E \rightarrow TE'$

$E' \rightarrow +TE' \mid \varepsilon$

$T \rightarrow FT'$

$T' \rightarrow *FT' \mid \varepsilon$

$F \rightarrow (E) \mid id$

First(E')=+, ε

First(T')=*, ε

First(E)=First(T)=First(F)=(,id

$S \rightarrow ABCDE$

$A \rightarrow Ba | \varepsilon$ $\quad A \rightarrow a | \varepsilon$ $\quad$ (coz $B \rightarrow \varepsilon$)

$B \rightarrow b | \varepsilon$

$C \rightarrow c$

$D \rightarrow d | \varepsilon$

$E \rightarrow De | \varepsilon$ $\quad E \rightarrow e | \varepsilon$ $\quad$ (coz $D \rightarrow \varepsilon$)

First (S) {a,b,ε,c}

First (A) {a,b,ε}

First (B) {b,ε}

First (C) {c}

First (D) {d,ε}

First (E) {d,ε,e}

FOLLOW

- FOLLOW(A) = { *the set of terminals that can immediately follow nonterminal A* }

FOLLOW(A) =

2 **for** all ($B \to \alpha\ A\ \beta$) $\in P$ **do**
FOLLOW(A) = FIRST(β) except {ε}

3 **for** all ($B \to A\ \beta$) $\in P$ and $\beta \to \varepsilon$ **do**
FOLLOW(A) = FOLLOW(B)

4 **for** all ($B \to \alpha\ A$) $\in P$ **do**
FOLLOW(A) = FOLLOW(B)

1 **if** A is the start symbol S **then**
add **$** to FOLLOW($A$)

E→TE’

E’→+TE’| ε

T→FT’

T’→*FT’| ε

F→(E)|id

Rule	E	E’	T	T’	F
1	$				
2	)		+		*
3			$,)		+,),$
4		$,)		+,),$	

Find FIRST and FOLLOW

Example

S→Bb|Cc {a,b,c,ε}{$}

B→aB|ε {a,ε} {b}

C→cC|ε {c,ε}{c}

Example

S→aABb {a}{$}

A→c|ε {c,ε}{d,$,b}

B→d|ε {d,ε}{b}

Example

S→ACB|CbB|Ba

A→da|BC

B→g|ε

C→h|ε

Example

S→aBDh

B→cC

C→bC|ε

D→EF

E→g|ε

Constructing an LL(1) Predictive Parsing Table

```
for each production A → α do
    for each a ∈ FIRST(α) do
        add A → α to M[A,a]
    enddo
    if ε ∈ FIRST(α) then
        for each b ∈ FOLLOW(A) do
            add A → α to M[A,b]
        enddo
    endif
enddo
Mark each undefined entry in M error
```

Example Table

$$E \to T\, E_R$$
$$E_R \to +\, T\, E_R \mid \varepsilon$$
$$T \to F\, T_R$$
$$T_R \to *\, F\, T_R \mid \varepsilon$$
$$F \to (\, E\,) \mid \mathbf{id}$$

$A \to \alpha$	FIRST(α)	FOLLOW(A)
$E \to T\, E_R$	**(id**	**$)**
$E_R \to +\, T\, E_R$	+	**$)**
$E_R \to \varepsilon$	ε	
$T \to F\, T_R$	**(id**	**+ $)**
$T_R \to *\, F\, T_R$	*	**+ $)**
$T_R \to \varepsilon$	ε	
$F \to (\, E\,)$	**(**	***** **+ $)**
$F \to \mathbf{id}$	**id**	***** **+ $)**

	id	+	*	**(**	**)**	**$**
E	$E \to T\, E_R$			$E \to T\, E_R$		
E_R		$E_R \to +\, T\, E_R$			$E_R \to \varepsilon$	$E_R \to \varepsilon$
T	$T \to F\, T_R$			$T \to F\, T_R$		
T_R		$T_R \to \varepsilon$	$T_R \to *\, F\, T_R$		$T_R \to \varepsilon$	$T_R \to \varepsilon$
F	$F \to \mathbf{id}$			$F \to (\, E\,)$		

LL(1) Grammars are Unambiguous

Ambiguous grammar

$S \rightarrow \mathbf{i}\ E\ \mathbf{t}\ S\ S_R \mid \mathbf{a}$

$S_R \rightarrow \mathbf{e}\ S \mid \varepsilon$

$E \rightarrow \mathbf{b}$

$A \rightarrow \alpha$	FIRST(α)	FOLLOW(A)
$S \rightarrow \mathbf{i}\ E\ \mathbf{t}\ S\ S_R$	**i**	**e $**
$S \rightarrow \mathbf{a}$	**a**	
$S_R \rightarrow \mathbf{e}\ S$	**e**	**e $**
$S_R \rightarrow \varepsilon$	ε	
$E \rightarrow \mathbf{b}$	**b**	**t**

Error: duplicate table entry

	a	**b**	**e**	**i**	**t**	**$**
S	$S \rightarrow \mathbf{a}$			$S \rightarrow \mathbf{i}\ E\ \mathbf{t}\ S\ S_R$		
S_R			$S_R \rightarrow \varepsilon$ $S_R \rightarrow \mathbf{e}\ S$			$S_R \rightarrow \varepsilon$
E		$E \rightarrow \mathbf{b}$				

Predictive Parsing Program (Driver)

push(**$**)
push(S)
a := *lookahead*
repeat
    X := pop()
    if X is a terminal or X = **$** **then**
        match(X) // *moves to next token and a := lookahead*
    else if $M[X,a] = X \to Y_1 Y_2 \ldots Y_k$ **then**
        push(Y_k, $Y_{k\text{-}1}$, …, Y_2, Y_1) // *such that* Y_1 *is on top*
        ... invoke actions and/or produce IR output …
    else   error()
    endif
until X = **$**

Example Table-Driven Parsing

$$E \rightarrow T\ E_R$$
$$E_R \rightarrow +\ T\ E_R \mid \varepsilon$$
$$T \rightarrow F\ T_R$$
$$T_R \rightarrow *\ F\ T_R \mid \varepsilon$$
$$F \rightarrow (\ E\ ) \mid \textbf{id}$$

Stack	Input	Production applied
$\$\underline{E}$	$\underline{\textbf{id}}\textbf{+id*id\$}$	$E \rightarrow T\ E_R$
$\$E_R\underline{T}$	$\underline{\textbf{id}}\textbf{+id*id\$}$	$T \rightarrow F\ T_R$
$\$E_RT_R\underline{F}$	$\underline{\textbf{id}}\textbf{+id*id\$}$	$F \rightarrow \textbf{id}$
$\$E_RT_R\underline{\textbf{id}}$	$\underline{\textbf{id}}\textbf{+id*id\$}$	**advance the lookahead**
$\$E_R\underline{T_R}$	$\underline{\textbf{+}}\textbf{id*id\$}$	$T_R \rightarrow \varepsilon$
$\$\underline{E_R}$	$\underline{\textbf{+}}\textbf{id*id\$}$	$E_R \rightarrow +\ T\ E_R$
$\$E_RT\underline{+}$	$\underline{\textbf{+}}\textbf{id*id\$}$	**advance the lookahead**
$\$E_R\underline{T}$	$\underline{\textbf{id}}\textbf{*id\$}$	$T \rightarrow F\ T_R$
$\$E_RT_R\underline{F}$	$\underline{\textbf{id}}\textbf{*id\$}$	$F \rightarrow \textbf{id}$
$\$E_RT_R\underline{\textbf{id}}$	$\underline{\textbf{id}}\textbf{*id\$}$	**advance the lookahead**
$\$E_R\underline{T_R}$	$\underline{\textbf{*}}\textbf{id\$}$	$T_R \rightarrow *\ F\ T_R$
$\$E_RT_RF\underline{*}$	$\underline{\textbf{*}}\textbf{id\$}$	**advance the lookahead**
$\$E_RT_R\underline{F}$	$\underline{\textbf{id}}\textbf{\$}$	$F \rightarrow \textbf{id}$
$\$E_RT_R\underline{\textbf{id}}$	$\underline{\textbf{id}}\textbf{\$}$	**advance the lookahead**
$\$E_R\underline{T_R}$	$\underline{\textbf{\$}}$	$T_R \rightarrow \varepsilon$
$\$\underline{E_R}$	$\underline{\textbf{\$}}$	$E_R \rightarrow \varepsilon$
$\underline{\$}$	$\underline{\textbf{\$}}$	Accept

Panic Mode Recovery

Add synchronizing actions to undefined entries based on FOLLOW

Pro:	Can be automated
Cons:	Error messages are needed

FOLLOW(E) = {) $ }
FOLLOW(E_R) = {) $ }
FOLLOW(T) = { +) $ }
FOLLOW(T_R) = { +) $ }
FOLLOW(F) = { + *) $ }

	id	+	*	(	)	$
E	$E \to T\ E_R$			$E \to T\ E_R$	***synch***	***synch***
E_R		$E_R \to +\ T\ E_R$			$E_R \to \varepsilon$	$E_R \to \varepsilon$
T	$T \to F\ T_R$	***synch***		$T \to F\ T_R$	***synch***	***synch***
T_R		$T_R \to \varepsilon$	$T_R \to *\ F\ T_R$		$T_R \to \varepsilon$	$T_R \to \varepsilon$
F	$F \to$ **id**	***synch***	***synch***	$F \to (E)$	***synch***	***synch***

synch: the driver pops current nonterminal A and skips input till synch token or skips input until one of FIRST(A) is found

Phrase-Level Recovery

Change input stream by inserting missing tokens
For example: **id id** is changed into **id * id**

Pro:	Can be fully automated
Cons:	Recovery not always intuitive

Can then continue here

	id	**+**	*****	**(**	**)**	**$**
E	$E \to T\ E_R$			$E \to T\ E_R$	***synch***	***synch***
E_R		$E_R \to +\ T\ E_R$			$E_R \to \varepsilon$	$E_R \to \varepsilon$
T	$T \to F\ T_R$	***synch***		$T \to F\ T_R$	***synch***	***synch***
T_R	***insert ****	$T_R \to \varepsilon$	$T_R \to *\ F\ T_R$		$T_R \to \varepsilon$	$T_R \to \varepsilon$
F	$F \to$ **id**	***synch***	***synch***	$F \to (\ E\ )$	***synch***	***synch***

insert *: driver inserts missing * and retries the production

Error Productions

$E \to T\ E_R$
$E_R \to +\ T\ E_R \mid \varepsilon$
$T \to F\ T_R$
$T_R \to *\ F\ T_R \mid \varepsilon$
$F \to (\ E\ ) \mid \mathbf{id}$

Add "*error production*":

$$T_R \to F\ T_R$$

to ignore missing *, e.g.: **id id**

Pro:	Powerful recovery method
Cons:	Manual addition of productions

	id	+	*	(	)	$
E	$E \to T\ E_R$			$E \to T\ E_R$	***synch***	***synch***
E_R		$E_R \to +\ T\ E_R$			$E_R \to \varepsilon$	$E_R \to \varepsilon$
T	$T \to F\ T_R$	***synch***		$T \to F\ T_R$	***synch***	***synch***
T_R	$T_R \to F\ T_R$	$T_R \to \varepsilon$	$T_R \to *\ F\ T_R$		$T_R \to \varepsilon$	$T_R \to \varepsilon$
F	$F \to \mathbf{id}$	***synch***	***synch***	$F \to (\ E\ )$	***synch***	***synch***