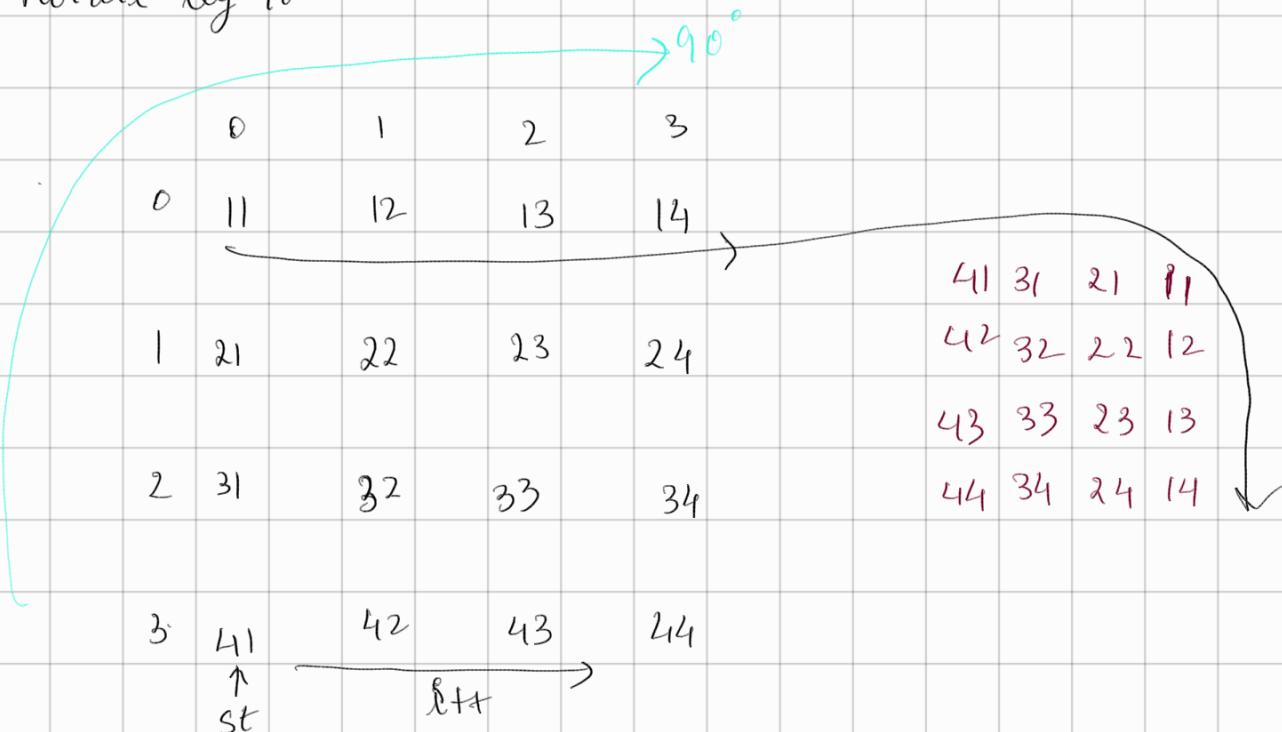


Today's Questions.

- ① Rotate by 90°
- ③ Saddle Point
- ② Ring Rotate(RW)
- ④ Matrix Multiplication

- ① Rotate by 90°



first row left to right → last column top to bottom.

Approach

Start from last row & print to first row
 ↓ bottom to ↓ top

in increasing order columns

```

for (int j=0; j<m; j++) {
    for (int i=n-1; i>=0; i--) {
        System.out.print(mat[i][j] + " ");
    }
    System.out.println();
}
    
```

OR

take Transpose.

$$a[i][j] = a[j][i]$$

Traverse any one of upper triangle or lower triangle matrix & swap

$\text{mat}[i][j]$ & $\text{mat}[j][i]$.

Dont swap complete matrix, only any one triangle diagonal remain same & others are swapped

after transpose, reverse the rows or swap columns in the matrix

① Transpose Code:

```
for (int i=0; i<n; i++) {  
    for (int j=i+1; j<n; j++) {  
        Swap (mat, i, j); // swaps mat[i][j] & mat[j][i]  
    }  
}
```

② Swap Code:

```
int temp = mat[i][j];  
mat[i][j] = mat[j][i];  
mat[j][i] = temp;
```

③ Reverse Rows code:

```
for (int i=0; i<n; i++) {  
    for (int j=0; j<n/2; j++) {  
        int temp = mat[i][j];  
        mat[i][j] = mat[i][n-1-j];  
        mat[i][n-1-j] = temp;  
    }  
}
```

for ①, ② & ③ make functions to make code modular.

④ Swapping columns.

```
int leftCol = 0, rightCol = arr[0].length - 1;
while (leftCol < rightCol) {
    for (int i = 0; i < arr.length; i++) {
        swap(arr, i, leftCol, rightCol);
        leftCol++;
        rightCol--;
    }
}
```



⑤ Overloading swap function for swapping column

```
void swap(int[][] arr, int i, int leftCol, int rightCol) {
    int temp = arr[i][leftCol];
    arr[i][leftCol] = arr[i][rightCol];
    arr[i][rightCol] = temp.
}
```

* Function with same name but different number or types of arguments are called overloaded functions/methods.

② Matrix multiplication.

$A_{n_1 \times m_1} \times B_{n_2 \times m_2}$ is only possible if $m_1 = n_2$

$$\begin{array}{c}
 A_{n_1 \times m_1} \quad \times \quad B_{n_2 \times m_2} \\
 \begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} \\ a_{21} & a_{22} & a_{23} & a_{24} \\ a_{31} & a_{32} & a_{33} & a_{34} \end{bmatrix}_{3 \times 4} \quad \times \quad \begin{bmatrix} b_{11} & b_{12} \\ b_{21} & b_{22} \\ b_{31} & b_{32} \\ b_{41} & b_{42} \end{bmatrix}_{4 \times 2}
 \end{array} = C_{n_1 \times m_2}$$

$$\begin{bmatrix} c_{11} & c_{12} \\ c_{21} & c_{22} \\ c_{31} & c_{32} \end{bmatrix}_{3 \times 2}$$

$$\begin{aligned}
 C_{11} &= a_{11} b_{11} + a_{12} b_{21} + a_{13} b_{31} + a_{14} b_{41} \\
 C_{12} &= a_{11} b_{12} + a_{12} b_{22} + a_{13} b_{32} + a_{14} b_{42} \\
 C_{21} &= a_{21} b_{11} + a_{22} b_{21} + a_{23} b_{31} + a_{24} b_{41} \\
 C_{22} &= a_{21} b_{12} + a_{22} b_{22} + a_{23} b_{32} + a_{24} b_{42} \\
 C_{31} &= a_{31} b_{11} + a_{32} b_{21} + a_{33} b_{31} + a_{34} b_{41} \\
 C_{32} &= a_{31} b_{12} + a_{32} b_{22} + a_{33} b_{32} + a_{34} b_{42}
 \end{aligned}
 \quad \left. \begin{aligned}
 C_{ij} &= \sum a_{ik} \times b_{kj}
 \end{aligned} \right\}$$

① if ($m_1 \neq n_2$) → Print ("Invalid input"); return;

approach ①:

make a new matrix → $\text{ans}[n_1][m_2]$
 and loop over for all variables in ans
 calculate them & place in ans .

CODE

```

int[][] ans = new int[n1][m2];
for (int i=0; i<n1; i++) {
    for (int j=0; j<m2; j++) {
        ans[i][j] = 0;
        for (int k=0; k<m1; k++) {
            ans[i][j] += mat1[i][k] * mat2[k][j];
        }
    }
}
    
```

// print output

```

for (int i=0; i<n1; i++) {
    for (int j=0; j<m2; j++) {
        System.out.print(ans[i][j] + " ");
    }
    System.out.println();
}
    
```

Formulae for answer matrix C^o

$$C_{ij}^o = \sum a_{ik}^o \times b_{kj}^o$$

③ Saddle Point.

Given square matrix, output saddle point

Saddle point \rightarrow largest in column, smallest in row
all elements are unique

Eg:

34	42	30	59
1x	1x	1v 2x	1x
66	85	44	97
1x	1x	1v 2x	1x
42	16	25	36
1x	1v 2x	1x	1x
15	21	05	50
x	1x	1v 2x	1x

Saddle Point

(Minimum in row & Maximum in column)

① Zero saddle point possible?

Yes Eg: Replace 66 with 22 in above matrix

[1][0]

② more than one saddle point

No, Proof below

Assume 2 saddle point s_1, s_2
mat

	s_1		s_3	
	(x_1, c_1)			

$$s_3 > s_1 \text{ & } s_3 < s_2 \quad \textcircled{1}$$

		(or 2)
s_4		s_2

$$s_4 > s_2 \quad \underline{s_4 < s_1} \quad \textcircled{2}$$

∴ $s_1 < s_3 \quad \& \quad s_2 > s_3 \quad , \quad s_2 < s_4 \quad \& \quad s_1 > s_4.$

$$\begin{array}{l} s_1 < s_3 < s_2 \\ \downarrow \qquad \qquad \qquad \downarrow \\ s_1 < s_2 \quad \& \quad s_2 < s_1 \end{array}$$

Cannot be possible because elements are all distinct.

Therefore, there will always be atmost 1 saddle point.

Approach

find minVal rowwise for each row

check if minVal is maxVal for it's column.

if saddle found the print & return.

* Math Needed:

- * ** ① Permutations and Combinations.
- * ② Sequence and Series. → AP, GP, AGP.
- * ③ Number Theory.
 - Prime numbers
 - GCD & LCM
 - Factorial

*
④ Logarithm.

⑤ Sets, Relations and Functions.

