

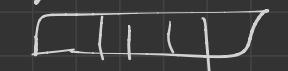
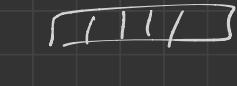

Recursion

{ Day-6 }

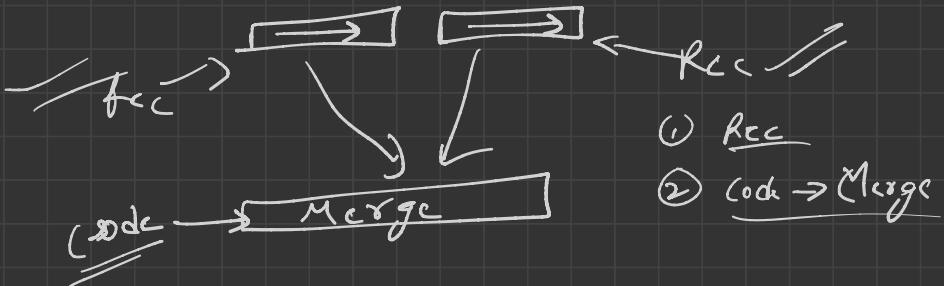
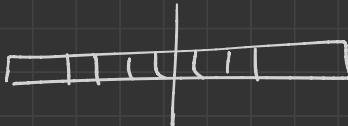
→ Quick Sort

i/p → array

o/p →



Merge Sort :-



→ Quick sort

0	1	2	3	4	5
3	5	1	8	2	4

1st step

element →

right place



partition

1	5	3	8	2	4
---	---	---	---	---	---

1	1	2	3	1	8	1	5	1	4
---	---	---	---	---	---	---	---	---	---



a			()	\
---	--	--	---	---	---

position pr dealo → right place

↓

<a

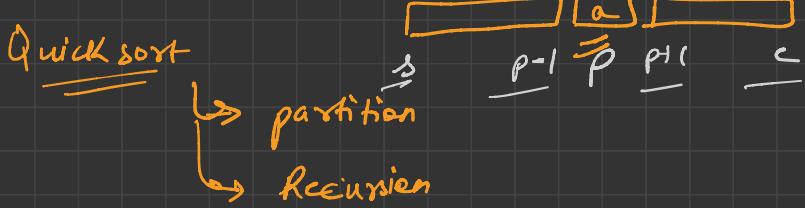
>a

partition

		(a)		
--	--	---	---	---	--	--

<a

>a



```

void quickSort ( int arr[], int s , int e )
{
    // base case
    if ( s >= e )
        return;

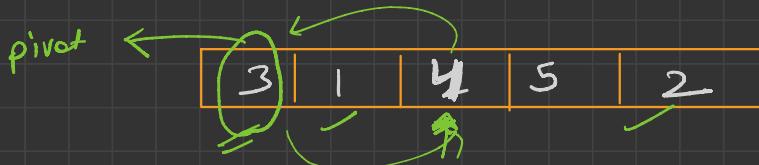
    // (1) partition
    P = partition ( arr, s, e );

    // Recur
    quickSort ( arr, s, P-1 );
    quickSort ( arr, p+1, e );
}

```

}

→ Partition



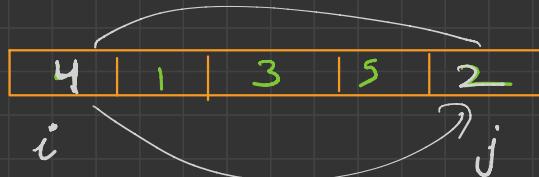
↳ take a pivot

↳ count all elements < pivot →

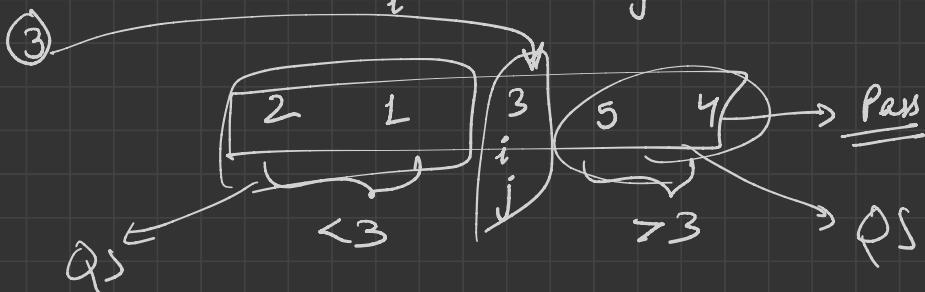
↳ pivot → $\Rightarrow + \text{count}$

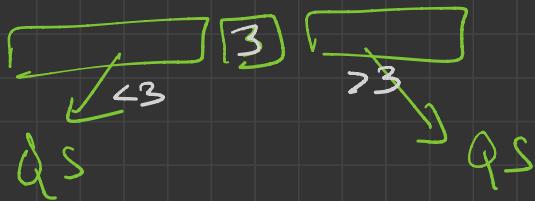
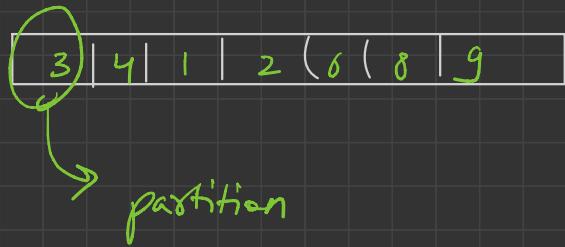


↳ $\langle a | a | >a$ → condn

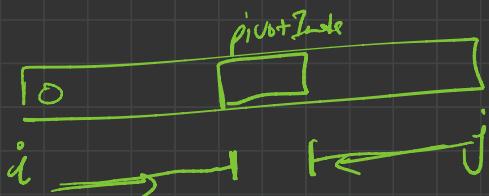
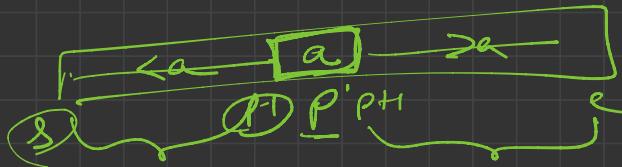


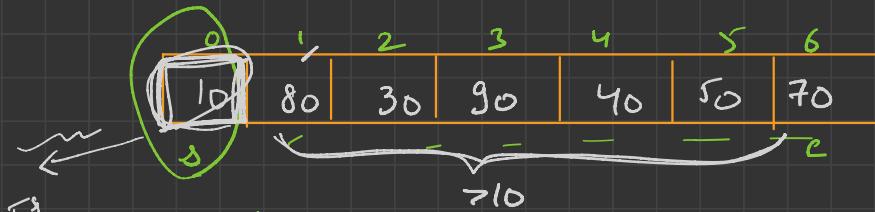
2 1 3 5 4
 $i =$ j





SORTED





$$\text{pivot} = 10$$

$$\text{cnt} = 0$$

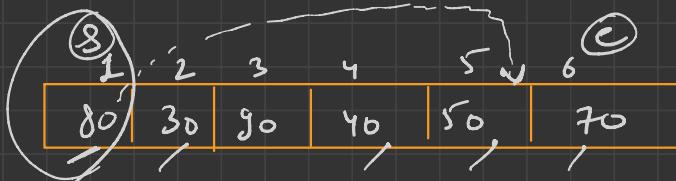
$$\Rightarrow \text{pivotIndex} = i + \text{cnt}$$

$$= 0 + 0 = 0$$

$$i = 0, j = 1$$

$$0 < 0$$

return 0;



$$\text{pivot} = 80$$

$$\text{cnt} = 0$$

$$\Rightarrow \text{pivotIndex} = i + \text{cnt}$$

$$= 1 + 0 = 1$$

after swap

\Rightarrow



$i = 2$

$j = 5$

$i < j$

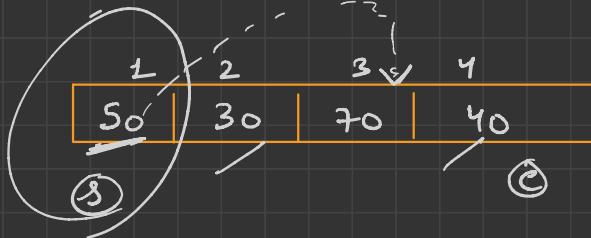
$j > 80$

< 80

sort

right

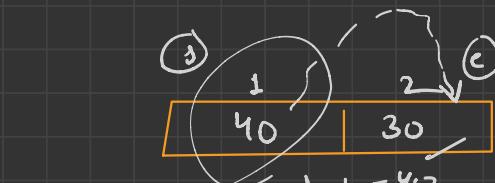
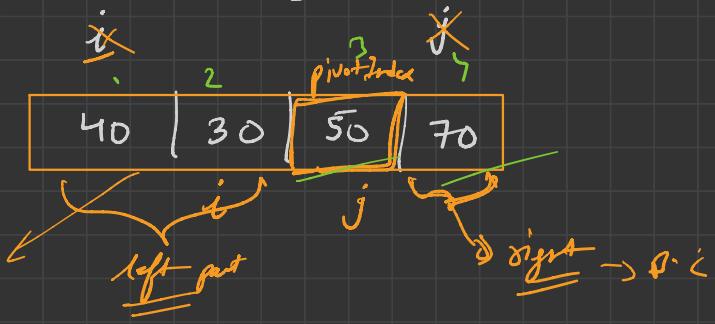
first part



$$\begin{cases}
 \text{pivot} = 50 \\
 \text{cnt} = 0 \\
 \rightarrow 2
 \end{cases}$$

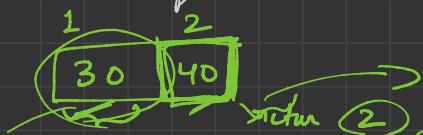
$$\begin{aligned}
 \text{pivotIndex} &= s + \text{cnt} \\
 &= 1 + 2 \\
 &= 3
 \end{aligned}$$

after swap \rightarrow



$$\begin{cases}
 \text{pivot} = 40 \\
 \text{cnt} = 1 \\
 \rightarrow \text{pivotIndex} = s + \text{cnt} = 1 + 1 = 2
 \end{cases}$$

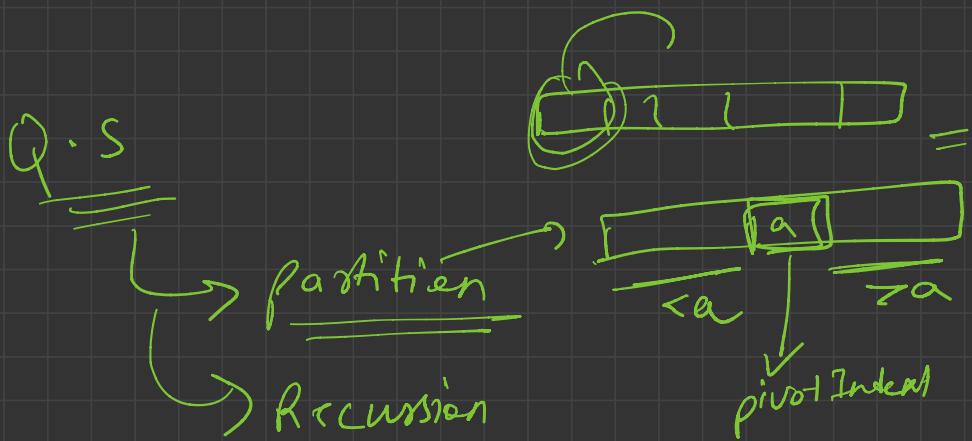
after swap



0	1	2	3	4	5	6
10	30	40	50	70	80	90

/ \ / \ / \ /

End result



$\rightarrow H/w \rightarrow$ in-place ?

\rightarrow stable \rightarrow ?

$\frac{S.C}{T.C} O(n)$

$O(n \log n)$

Worst case $\rightarrow O(n^2)$

Q.S

(T.C)

Avg Case - $O(n \log n)$

But Case - $O(n \log n)$

Worst Case - $O(n^2)$

$\frac{S.C}{\downarrow}$
 $O(n)$

$n \rightarrow$ size of array

H.W

in-place \rightarrow ✓ ✗

stable \rightarrow ✓ ✗

KHVD ^{code}
Kruska with DRP
RUN