

# The C - Assembly connection

Systems Software

# Displacement addressing

## ■ Used to access a memory address plus a constant displacement

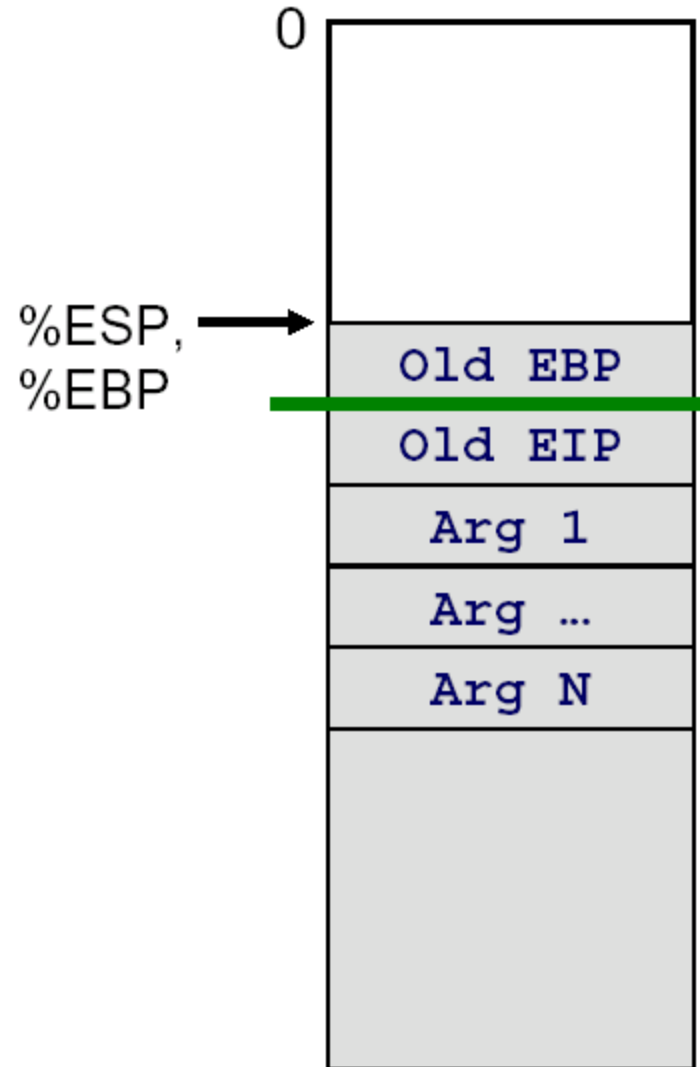
- `movl disp(reg1), reg2`
  - `reg1, reg2` are registers
  - `disp` is a 32-bit constant displacement value
  - Moves the contents of memory at address `(reg1 + disp)` to `reg2`

## ■ Example: `movl 8(%ebp), %edx`

- Moves value at memory address `%ebp + 8` to `%edx`

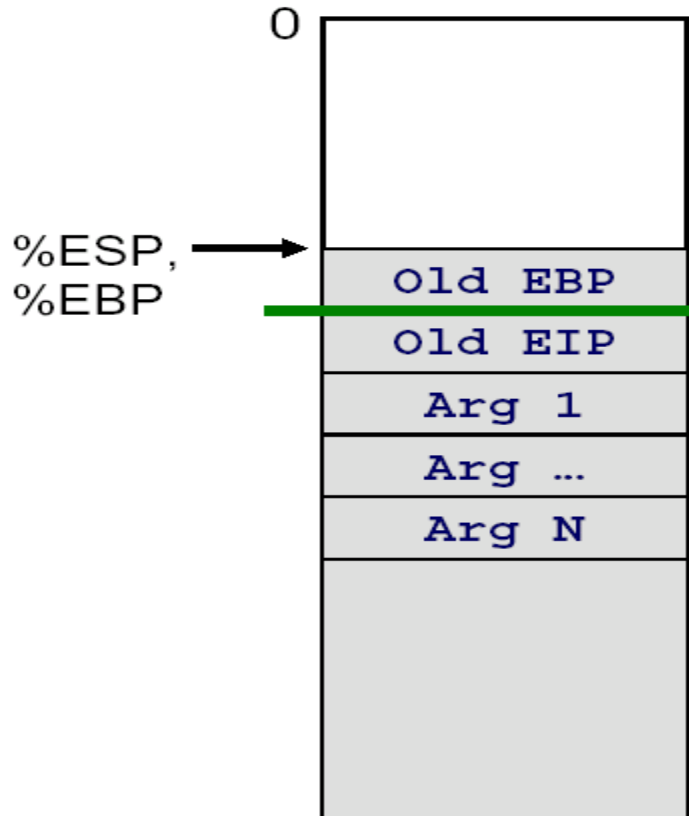
# Base Pointer: EBP

- **Save caller's base pointer**
  - `pushl %ebp`
- **Set value of EBP for callee's use**
  - Current value of stack pointer used as callee's base pointer
  - `movl %esp, %ebp`



# Using displacement addressing

```
/* sum first 3 elements of  
array a and store result in  
a[3]. */  
void sumup(int *a) {  
    int temp;  
    temp = a[0] + a[1] + a[2];  
    a[3] = temp;  
}
```



sumup:

```
pushl %ebp  
movl  %esp,%ebp  
movl  8(%ebp), %edx
```

} Set  
Up

```
movl  4(%edx), %eax  
addl  (%edx), %eax  
addl  8(%edx), %eax  
movl  %eax, 12(%edx)
```

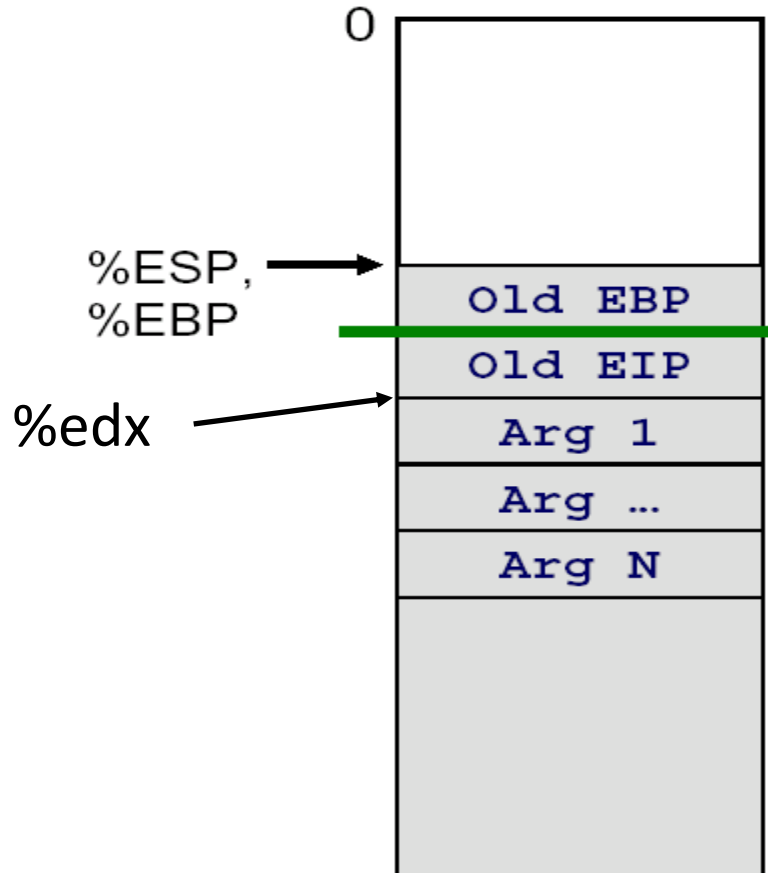
} Body

```
popl %ebp  
ret
```

} Finish

# Using displacement addressing

```
void sumup(int *a) {  
    int tmp;  
    temp = a[0] + a[1] + a[2];  
    a[3] = tmp;  
}
```



```
sumup:  
    pushl %ebp  
    movl  %esp, %ebp  
    movl  8(%ebp), %edx
```

} Set Up

```
    movl  4(%edx), %eax  
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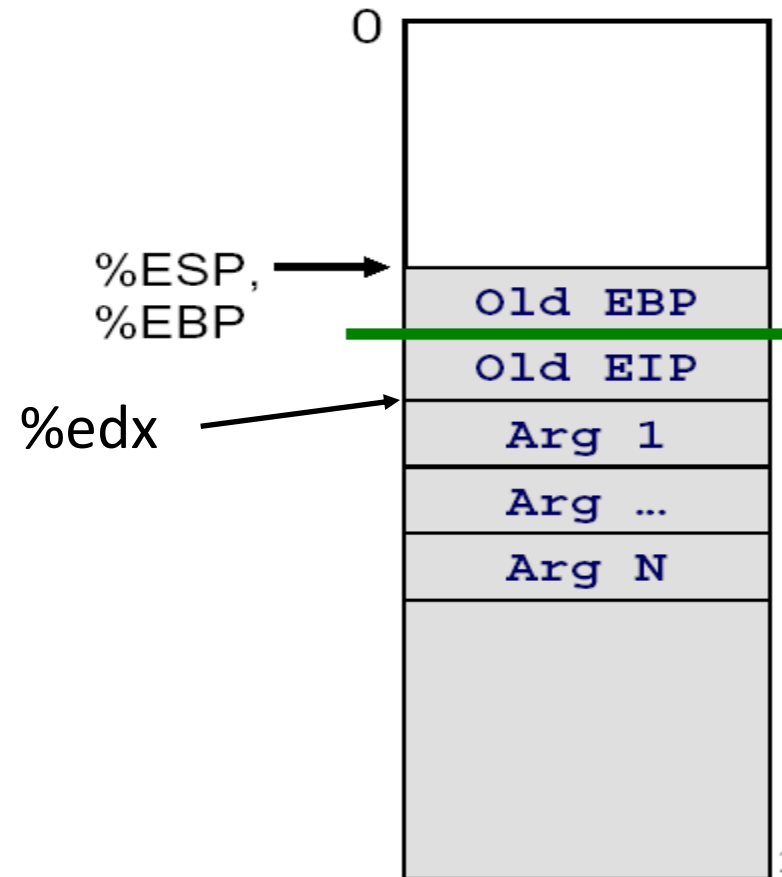
} Body

```
    popl %ebp  
    ret
```

} Finish

# Understanding sumup

```
/* sum first 3 elements
 * of array a and store result
 * in a[3].
 */
void sumup(int *a) {
    int tmp;
    temp = a[0] + a[1] + a[2];
    a[3] = temp;
}
```



Value of a (address of first array element) in %edx

```
movl 4(%edx), %eax    ; %eax = a[1]
addl (%edx), %eax     ; %eax += a[0]
addl 8(%edx), %eax     ; %eax += a[2]
movl %eax, 12(%edx)   ; a[3] = %eax
```

# Indexed addressing

## ■ A very general purpose addressing mode

- **disp(reg1, reg2, scale)** accesses memory address:  $\text{reg1} + (\text{reg2} * \text{scale}) + \text{disp}$ 
  - disp: constant displacement (32 bits)
  - reg1: Base register: Any of the 8 general-purpose registers
  - reg2: Index register: Any register except for %esp
  - scale: Must be 1, 2, 4, or 8

## ■ Example: `movl 10(%eax, %ebx, 4), %ecx`

- Accesses memory at address  $\%eax + (\%ebx * 4) + 10$
- Stores result in %ecx

# Address Computation Examples

<b>%edx</b>	<b>0xf000</b>
<b>%ecx</b>	<b>0x100</b>

Expression	Computation	Address
<b>0x8 (%edx)</b>		
<b>(%edx, %ecx)</b>		
<b>(%edx, %ecx, 4)</b>		
<b>0x80 (, %edx, 2)</b>		



# Address Computation Examples

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-------------	---------------

<b>%ecx</b>	<b>0x100</b>
-------------	--------------

Expression	Computation	Address
<b>0x8 (%edx)</b>	<b>0xf000 + 0x8</b>	<b>0xf008</b>
<b>(%edx, %ecx)</b>		
<b>(%edx, %ecx, 4)</b>		
<b>0x80 (, %edx, 2)</b>		

# Address Computation Examples

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<code>0x8(%edx)</code>	<code>0xf000 + 0x8</code>	<code>0xf008</code>
<code>(%edx,%ecx)</code>	<code>0xf000 + 0x100</code>	<code>0xf100</code>
<code>(%edx,%ecx,4)</code>		
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<code>(%edx,%ecx,4)</code>	<code>0xf000 + 4*0x100</code>	<code>0xf400</code>
<code>0x80(,%edx,2)</code>		

# Address Computation Examples

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<code>(%edx,%ecx,4)</code>	<code>0xf000 + 4*0x100</code>	<code>0xf400</code>
<code>0x80(,%edx,2)</code>	<code>2*0xf000 + 0x80</code>	<code>0x1e080</code>

# Computing addresses: `leal` instruction

- Sometimes want to save the result of an address computation for later use
  - Rather than reading or writing memory at the computed address
- The `leal` instruction is used for this: “load effective address”
- Example: `leal 4(%eax, %ebx, 8), %ecx`
  - Computes the address of the “source” argument, stores in the destination
  - Does not access the memory address itself: use `movl` for that!
- Can also be used to compute arithmetic expressions!
  - Anything of the form  $x + y * k$ , where  $k = 1, 2, 4, \text{ or } 8$

# A more complex example

## ■ What does this code do?

```
# Address of array in %ecx
movl 8(%ecx), %eax
movl 4(%ecx), %edx
leal (%edx, %eax, 2), %eax
movl (%ecx), %edx
leal (%edx, %eax, 2), %eax
movl %eax, 12(%ecx)
```

# A more complex example

## ■ What does this code do?

```
# Address of array[] in %ecx
movl 8(%ecx), %eax          ; eax = array[2]
movl 4(%ecx), %edx          ; edx = array[1]
leal (%edx, %eax, 2), %eax  ; eax = array[1] + (2*array[2])
movl (%ecx), %edx           ; edx = array[0]
leal (%edx, %eax, 2), %eax  ; eax = array[0] + (2 * eax)
movl %eax, 12(%ecx)         ; array[3] = eax
```

## ■ Original source code

```
void dosomething(int *array) {
    int tmp;
    tmp = array[0] + (2 * array[1]) + (4 * array[2]);
    array[3] = tmp;
}
```

# Structures

```
struct rec {  
    int i;  
    int a[3];  
    int *p;  
};
```



# Structures

```
struct rec {  
    int i;  
    int a[3];  
    int *p;  
};
```

## Memory Layout



## ■ Characteristics

- Contiguously-allocated region of memory
- Refer to members within structure by names
- Members may be of different types

# Structures

## ■ Accessing Structure Member

- Given an instance of the struct, we can use the `.` operator:

- `struct rec r1; r1.i = val;`

- What if we have a *pointer* to a struct: `struct rec *r = &r1;`

```
struct rec {  
    int i;  
    int a[3];  
    int *p;  
};
```

# Structures

## ■ Accessing Structure Member

- Given an instance of the struct, we can use the `.` operator, just like Java:
  - `struct rec r1; r1.i = val;`
- What if we have a *pointer* to a struct: `struct rec *r = &r1;`
  - Using `*` and `.` operators: `(*r).i = val;`
  - Or, use `->` operator for short: `r->i = val;`
- Pointer indicates first byte of structure; access members with offsets

```
struct rec {  
    int i;  
    int a[3];  
    int *p;  
};
```

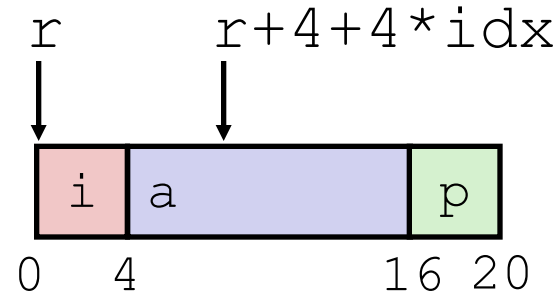
```
void  
set_i(struct rec *r,  
      int val)  
{  
    r->i = val;  
}
```

## IA32 Assembly

```
# %eax = val  
# %edx = r  
movl %eax, (%edx)    # Mem[r] = val
```

# Generating Pointer to Structure Member

```
struct rec {  
    int i;  
    int a[3];  
    int *p;  
};
```



## ■ Generating Pointer to Array Element

- Offset of each structure member determined at compile time

```
int *find_a  
(struct rec *r, int idx)  
{  
    return &r->a[idx];  
}
```

```
# %ecx = idx  
# %edx = r  
leal 0(,%ecx,4),%eax    # 4*idx  
leal 4(%eax,%edx),%eax  # r+4*idx+4
```

# Practice question

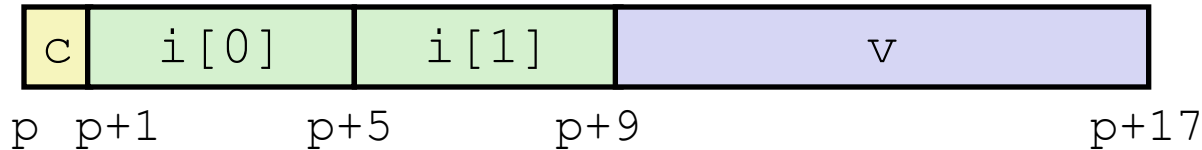
```
struct prob {  
    int *p;  
    struct {  
        int x;  
        int y;  
    } s;  
    struct prob *next;  
} ;
```

```
void sp_init(struct prob *sp)  
{  
    sp->s.x = _____;  
    sp->p   = _____;  
    sp->next = _____;  
}
```

```
1    movl 8(%ebp), %eax  
2    movl 8(%eax), %edx  
3    movl %edx, 4(%eax)  
4    leal 4(%eax), %edx  
5    movl %edx, (%eax)  
6    movl %eax, 12(%eax)
```

# Structures & Alignment

## ■ Unaligned Data

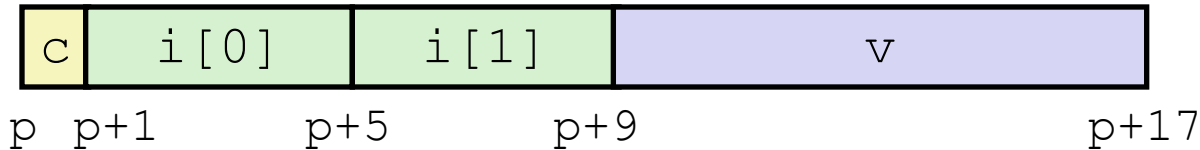


```
struct S1 {  
    char c;  
    int i[2];  
    double v;  
} *p;
```

- How would it look like if data items were *aligned* (address multiple of type size)?

# Structures & Alignment

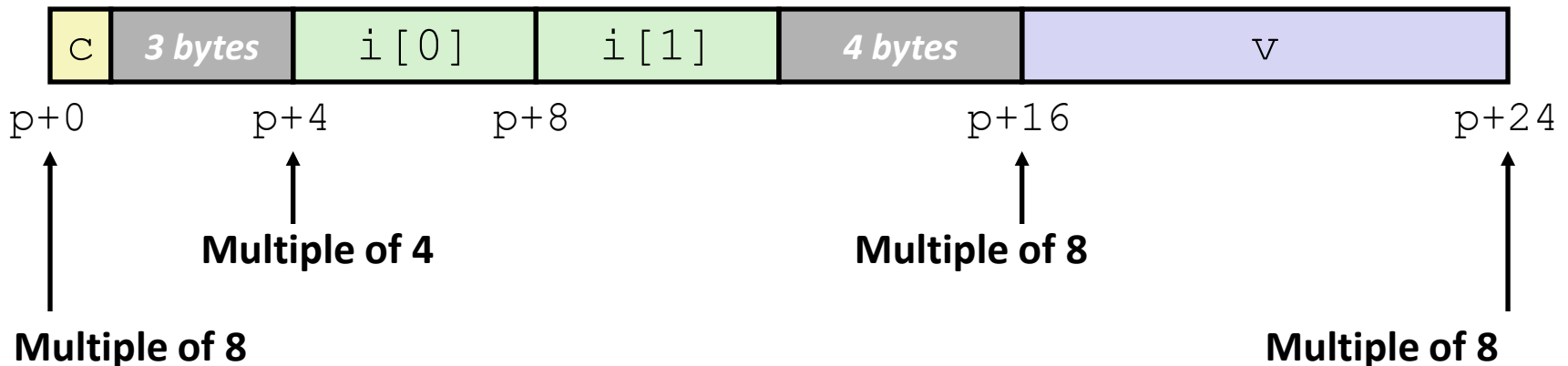
## ■ Unaligned Data



```
struct S1 {  
    char c;  
    int i[2];  
    double v;  
} *p;
```

## ■ Aligned Data

- Primitive data type requires **K** bytes
- Address must be multiple of **K**



# Alignment Principles

## ■ Aligned Data

- Primitive data type requires K bytes
- Address must be multiple of K

## ■ Aligned data is required on some machines; it is *advised* on IA32

- Treated differently by IA32 Linux, x86-64 Linux, and Windows!

## ■ What is the motivation for alignment?



# Alignment Principles

## ■ Motivation for Aligning Data

- Physical memory is accessed by aligned chunks of 4 or 8 bytes (system-dependent)
  - Inefficient to load or store datum that crosses quad word boundaries

## ■ Compiler

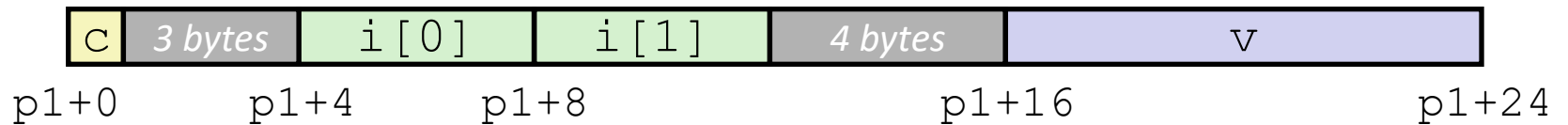
- Inserts padding in structure to ensure correct alignment of fields
- **sizeof()** should be used to get true size of structs

# Specific Cases of Alignment (IA32)

- **1 byte: char, ...**
  - no restrictions on address
- **2 bytes: short, ...**
  - lowest 1 bit of address must be  $0_2$
- **4 bytes: int, float, char \*, ...**
  - lowest 2 bits of address must be  $00_2$
- **8 bytes: double, ...**
  - Windows (and most other OSs & instruction sets): lowest 3 bits  $000_2$
  - Linux: lowest 2 bits of address must be  $00_2$ 
    - i.e., treated the same as a 4-byte primitive data type

# Saving Space

```
struct S1 {  
    char c;  
    int i[2];  
    double v;  
} *p1;
```



# Saving Space

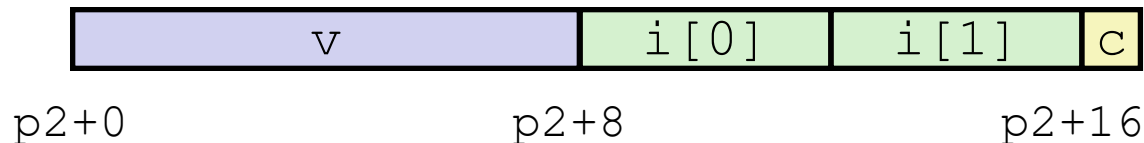
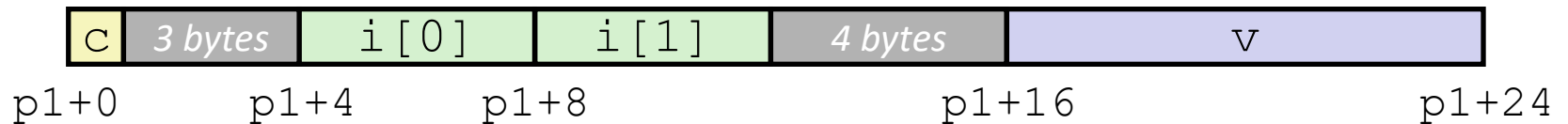
## ■ Put large data types first:

```
struct S1 {  
    char c;  
    int i[2];  
    double v;  
} *p1;
```



```
struct S2 {  
    double v;  
    int i[2];  
    char c;  
} *p2;
```

## ■ Effect (both have K=8)



Unfortunately, doesn't satisfy requirement that struct's *total size* is a multiple of K

# Arrays of Structures

- Satisfy alignment requirement for every element
- How would accessing an element work?

```
struct S2 {  
    double v;  
    int i[2];  
    char c;  
} a[10];
```

