

# Definitions, Laws, and Axioms

Tanisha Gupta, Aidan Nguyen, Jennifer Jiang, Parth Parikh

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This document contains the official list of definitions, laws, and axioms that students are permitted to reference and use for CS 2050 assignments and examinations.

## 1 Definitions

1. **Negation ( $\neg$ ):** the negation of  $p$  is the statement “not  $p$ .”
2. **Conjunction ( $\wedge$ ):** the conjunction of  $p$  and  $q$  is the statement “ $p$  and  $q$ .” For the statement to be true, both  $p$  and  $q$  must be true.
3. **Disjunction ( $\vee$ ):** the disjunction of  $p$  and  $q$  is the statement “ $p$  or  $q$ .” For the statement to be true, at least one of  $p$  or  $q$  must be true.
4. **Exclusive-or ( $\oplus$ ):** the exclusive-or of  $p$  and  $q$  is the statement “ $p$  XOR  $q$ .” For the statement to be true, exactly one of  $p$  and  $q$  must be true.
5. **Conditional statement ( $\rightarrow$ ):** “If  $p$ , then  $q$ .” Written  $p \rightarrow q$ . This is false when  $p$  is true and  $q$  is false, and true in every other instance.
6. **Converse:**  $q \rightarrow p$ .
7. **Inverse:**  $\neg p \rightarrow \neg q$ .
8. **Contrapositive:**  $\neg q \rightarrow \neg p$  (logically equivalent to the given conditional).
9. **Bi-conditional statement ( $\leftrightarrow$ ):** “ $p$  if and only if  $q$ .” One cannot exist without the other. Statement is true when both statements are true or both statements are false.  $q$  is sufficient for  $p$  and  $p$  is sufficient for  $q$ .
10. **Definition of Biconditional:**
$$p \leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$$
11. **Tautology:** a compound proposition that is always true. Examples:  $p \rightarrow p, p \vee \neg p$ .
12. **Contradiction:** a compound proposition that is always false. Example:  $\neg p \wedge p$ .
13. **Contingency:** a compound proposition that is neither a tautology nor a contradiction. Most statements are contingencies. Example:  $p \wedge q$ .
14. **Logical Equivalence:** Two compound propositions  $p$  and  $q$  are logically equivalent if  $p \leftrightarrow q$  is a tautology. Notation:  $p \equiv q$ .
15. **Predicate (Propositional Function):** A predicate is a statement involving one or more variables that becomes either true or false when values are assigned to the variables. Written like  $P(x)$ .
16. **Universal Quantifier ( $\forall$ ):** The universal quantifier  $\forall$  is read as “for all” or “for every.” The statement  $\forall x P(x)$  means that the predicate  $P(x)$  is true for every element  $x$  in the specified domain.

17. **Existential Quantifier ( $\exists$ ):** The existential quantifier  $\exists$  is read as “there exists.” The statement  $\exists xP(x)$  means that there is at least one element  $x$  in the domain for which  $P(x)$  is true.
18. **Unique Existential Quantifier ( $\exists!$ ):** The unique existential quantifier  $\exists!$  is read as “there exists exactly one.” The statement  $\exists!xP(x)$  means that there is exactly one element  $x$  in the domain for which the predicate  $P(x)$  is true.
19. **Sets of Numbers:**
- $\mathbb{N}$ : Natural numbers  $\{0, 1, 2, 3, \dots\}$
  - $\mathbb{Z}$ : Integers  $\{\dots, -2, -1, 0, 1, 2, \dots\}$
  - $\mathbb{Q}$ : Rational numbers  $\{\frac{a}{b} \mid a, b \in \mathbb{Z}, b \neq 0\}$
  - $\mathbb{R}$ : Real numbers (all points on the number line)
  - $\mathbb{C}$ : Complex numbers  $\{a + bi \mid a, b \in \mathbb{R}, i^2 = -1\}$
  - $\mathbb{I}$ : Irrational numbers (reals not in  $\mathbb{Q}$ , e.g.,  $\pi, \sqrt{2}$ )
  - $\mathbb{Z}^+$ : Positive integers  $\{1, 2, 3, \dots\}$
  - $\mathbb{Z}^-$ : Negative integers  $\{-1, -2, -3, \dots\}$
20. **Nested Quantifier:** When one quantifier is within the scope of another. Ex:  $\forall x \exists y (x - 2y = 5)$
21. **Argument:** Is a collection of premises and a conclusion. An argument is valid if and only if the truth of the premises guarantees the conclusion. Or, in the language of the class (where  $t$  is the conclusion and  $p_1, p_2, \dots, p_k$  are premises):

$$(p_1 \wedge p_2 \wedge \dots p_n) \rightarrow t$$

## 2 Laws & Axioms

- **De Morgan’s Laws:**

1.  $\neg(p \wedge q) \equiv \neg p \vee \neg q$
2.  $\neg(p \vee q) \equiv \neg p \wedge \neg q$

- **Conditional Disjunction Equivalence (CDE):**  $p \rightarrow q \equiv \neg p \vee q$ .

Table 1: **Only equivalences present on this table will be accepted.**

<i>Equivalence</i>	<i>Name</i>
$p \rightarrow q \equiv \neg p \vee q$	Conditional Disjunction Equivalence
$p \rightarrow q \equiv \neg q \rightarrow \neg p$	Contrapositive Law
$p \leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$	Definition of Biconditional
$\neg(\neg p) \equiv p$	Double Negation Law
$p \wedge \mathbf{T} \equiv p$ $p \vee \mathbf{F} \equiv p$	Identity Laws
$p \vee \mathbf{T} \equiv \mathbf{T}$ $p \wedge \mathbf{F} \equiv \mathbf{F}$	Domination Laws
$p \vee p \equiv p$ $p \wedge p \equiv p$	Idempotent Laws
$p \vee q \equiv q \vee p$ $p \wedge q \equiv q \wedge p$	Commutative Laws
$(p \vee q) \vee r \equiv p \vee (q \vee r)$ $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$	Associative Laws
$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$ $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$	Distributive Laws
$\neg(p \wedge q) \equiv \neg p \vee \neg q$ $\neg(p \vee q) \equiv \neg p \wedge \neg q$	De Morgan's Laws
$p \vee (p \wedge q) \equiv p$ $p \wedge (p \vee q) \equiv p$	Absorption Laws
$p \vee \neg p \equiv \mathbf{T}$ $p \wedge \neg p \equiv \mathbf{F}$	Negation Laws

• **De Morgan's Laws for Quantifiers:**

1.  $\neg(\forall x P(x)) \equiv \exists x \neg P(x)$
2.  $\neg(\exists x P(x)) \equiv \forall x \neg P(x)$

Table 2: Only equivalences and rules of inference present on this table will be accepted.

<i>Rule of Inference</i>	<i>Name</i>
$p$ $p \rightarrow q$ <hr/> $\therefore q$	Modus Ponens
$\neg q$ $p \rightarrow q$ <hr/> $\therefore \neg p$	Modus Tollens
$p$ <hr/> $\therefore p \vee q$	Addition
$p \wedge q$ <hr/> $\therefore p$	Simplification
$p$ $q$ <hr/> $\therefore p \wedge q$	Conjunction
$p \rightarrow q$ $q \rightarrow r$ <hr/> $\therefore p \rightarrow r$	Hypothetical Syllogism
$p \vee q$ $\neg p$ <hr/> $\therefore q$	Disjunctive Syllogism
$p \vee q$ $\neg p \vee r$ <hr/> $\therefore q \vee r$	Resolution