

How to Win Coding Competitions: Secrets of Champions

Week 4: Sorting and Search Algorithms 2 Lecture 6: Priority queue and binary heap

Maxim Buzdalov Saint Petersburg 2016

- ► Assume you are a cinema critic, and your job is to watch movies once a day
- ► You may process requests in the order of their appearance
- ► Ordinary queue (First In, First Out)

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Priority queue: Interface and implementation

	Vector, Queue, Deque,	Same with sorting
[+]: add an element	O(1)	O(n)
[?]: query the smallest element	O(n)	O(1)
[-]: remove the smallest element	O(n)	O(1)

Priority queue: Interface and implementation

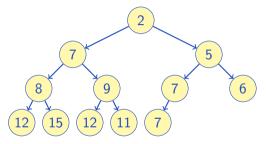
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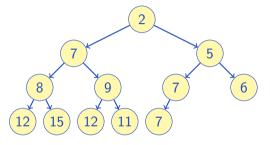
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Is this the best we can do?

Binary heap: a tree-like structure



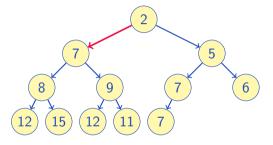
Binary heap: a tree-like structure



Properties of a binary heap:

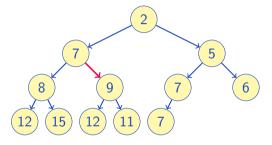
► A node is not less than its parent

Binary heap: a tree-like structure



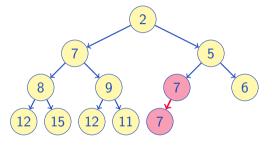
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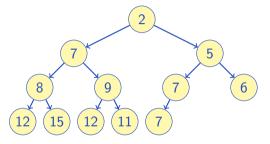
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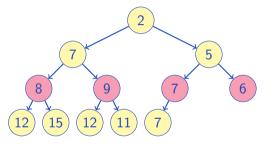


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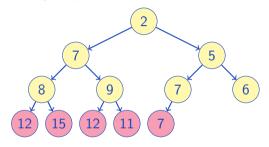
► A node is **not less** than its parent



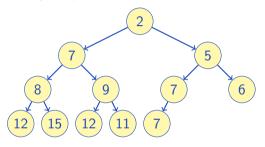
- ► A node is not less than its parent
- ► Every node has at most two children



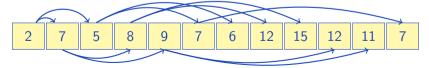
- ► A node is not less than its parent
- ► Every node has at most two children
- ► All levels but the last one are complete

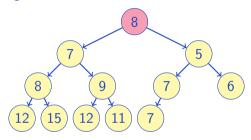


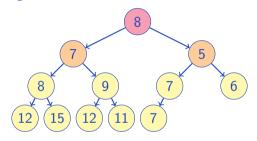
- ► A node is not less than its parent
- ► Every node has at most two children
- ► All levels but the last one are complete
- ► The last level is filled from the left



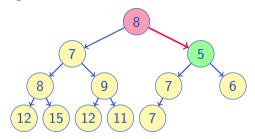
- ► A node is not less than its parent
- ► Every node has at most two children
- All levels but the last one are complete
- ► The last level is filled from the left
- Can be stored in an array or a vector
- ▶ Node at index $i \to \text{parent}$ at index $\lfloor \frac{i}{2} \rfloor$



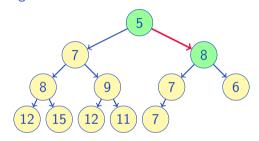




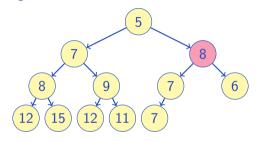
► Look at the children



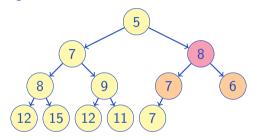
- ► Look at the children
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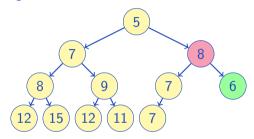
- ► Look at the children
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- Swap with it



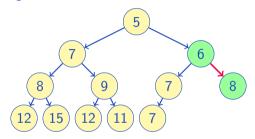
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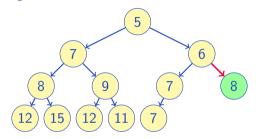
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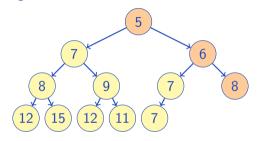
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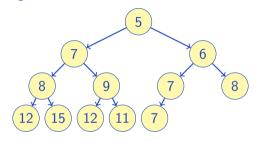
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Complexity of the operation:

 $O(\text{height of a broken node}) = O(\log n)$

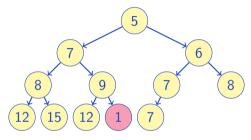


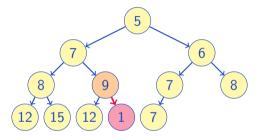
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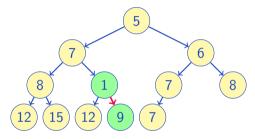
 $O(\text{height of a broken node}) = O(\log n)$

We will call this operation "Sift down"

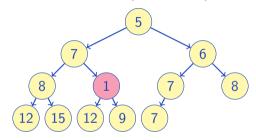




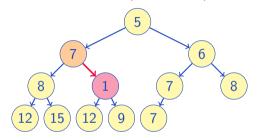
► Look at the parent



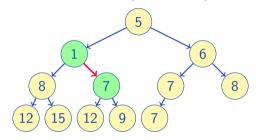
- ► Look at the parent
- ▶ If the parent is greater, swap with it



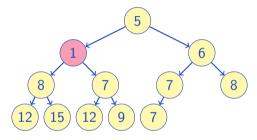
- ► Look at the parent
- ▶ If the parent is greater, swap with it
- ▶ If still too small, continue



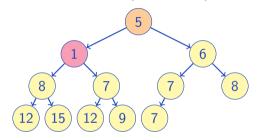
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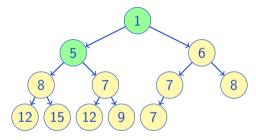
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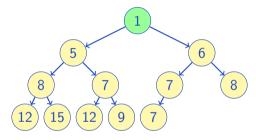
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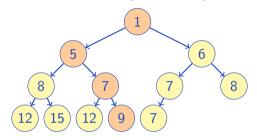
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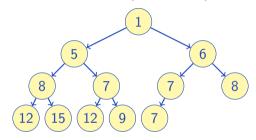
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Complexity of the operation:

 $O(\text{depth of the broken node}) = O(\log n)$



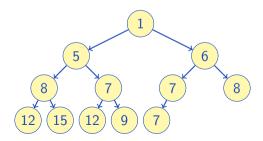
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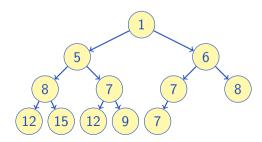
 $O(\text{depth of the broken node}) = O(\log n)$

We will call this operation "Sift up"

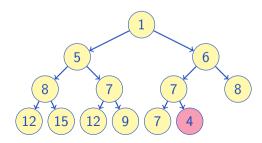
- ► [+]: add an element
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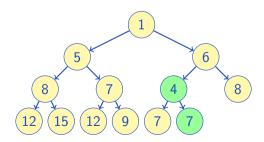
- \blacktriangleright [+]: add an element \rightarrow add the new node, then sift it up
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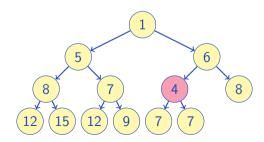
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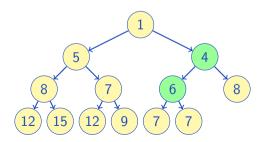
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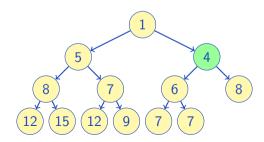
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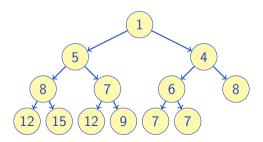
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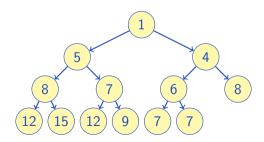
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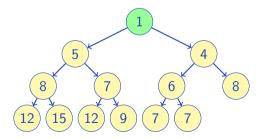
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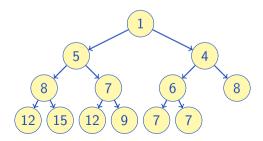
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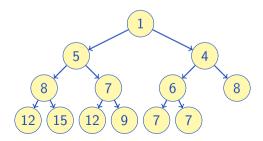
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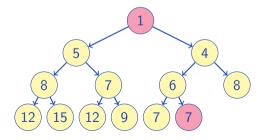
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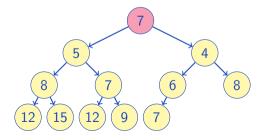
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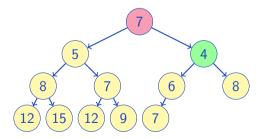
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 - → move the last node to the root, then sift it down



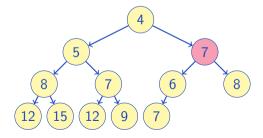
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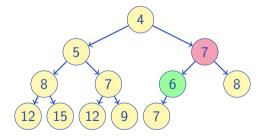
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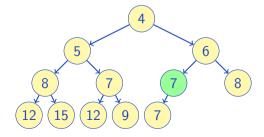
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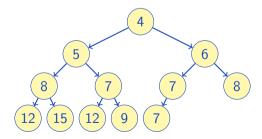
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 - \rightarrow move the last node to the root, then sift it down $\rightarrow O(\log n)$



Priority queue: Implementations and their performance

	Vector, Queue, Deque,	Same with sorting	Binary heap
[+]	O(1)	O(n)	$O(\log n)$
[?]	O(n)	O(1)	O(1)
[-]	O(n)	O(1)	$O(\log n)$

- ▶ Java
 - java.util.PriorityQueue<T>
- ► C++
 - ► Warning: C++ assumes MAXIMUM value at root
 - everything is the same, but min \leftrightarrow max, $\leq \leftrightarrow \geq$, etc.
 - ► High-level data structure from <queue>
 - ▶ std::priority_queue<T>
 - ► Low-level operations from <algorithm>
 - ▶ std::make_heap
 - ▶ std::push_heap
 - ▶ std::pop_heap
 - ▶ std::sort_heap
- Python
 - ► The heapq module