

Triple sum Equal to Target

idx	0	1	2	3	4	5	6
arr:	1	2	3	4	5	6	7
i=0	↑	↑	↑	↑	↑		
j=1							
K=2							
K=3							
K=4							
fix fix							

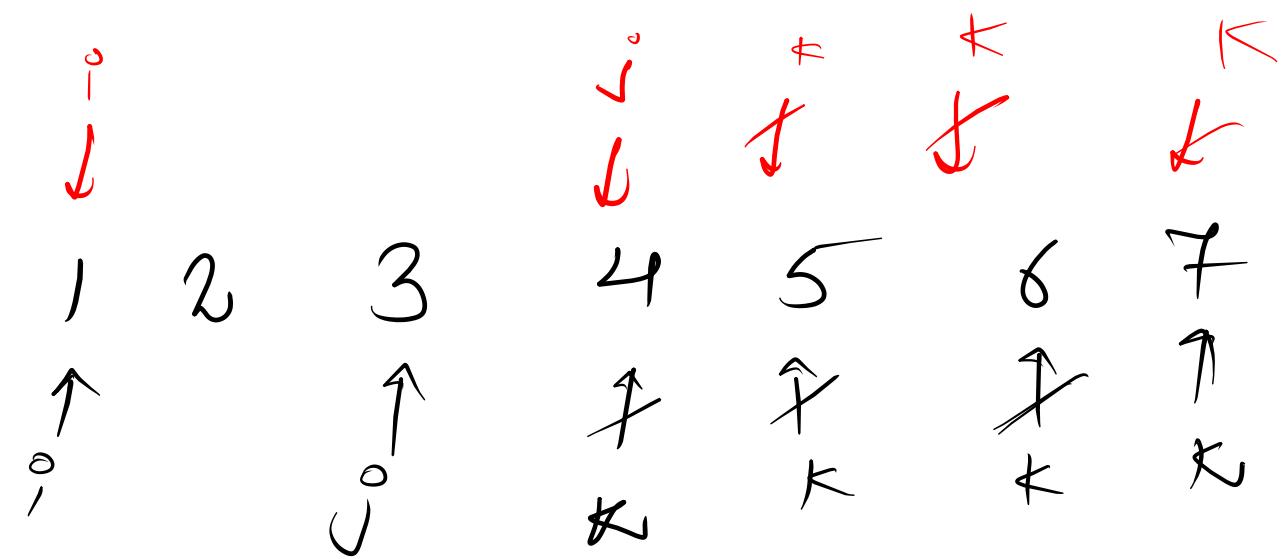
Target = 12

$$[1, 2, 3] \rightarrow 1+2+3 \Rightarrow 6 == 12 \times \text{Cnt} = 0$$

$$[1, 2, 4] \rightarrow 1+2+4 \Rightarrow 7 == 12 \times \text{Cnt} = 0$$

$$[1, 2, 5] \rightarrow 1+2+5 \Rightarrow 8 == 12 \times \text{Cnt} = 0$$

⋮



10^3
 $\downarrow 10^6 \rightarrow$
 $O(N^2)$

10^8
 $O(N^3)$

$[1, 3, 4] \rightarrow \text{sum} = T$

$[1, 3, 5]$ "

$[1, 3, 6]$ "

$[1, 3, 7]$ "

if true

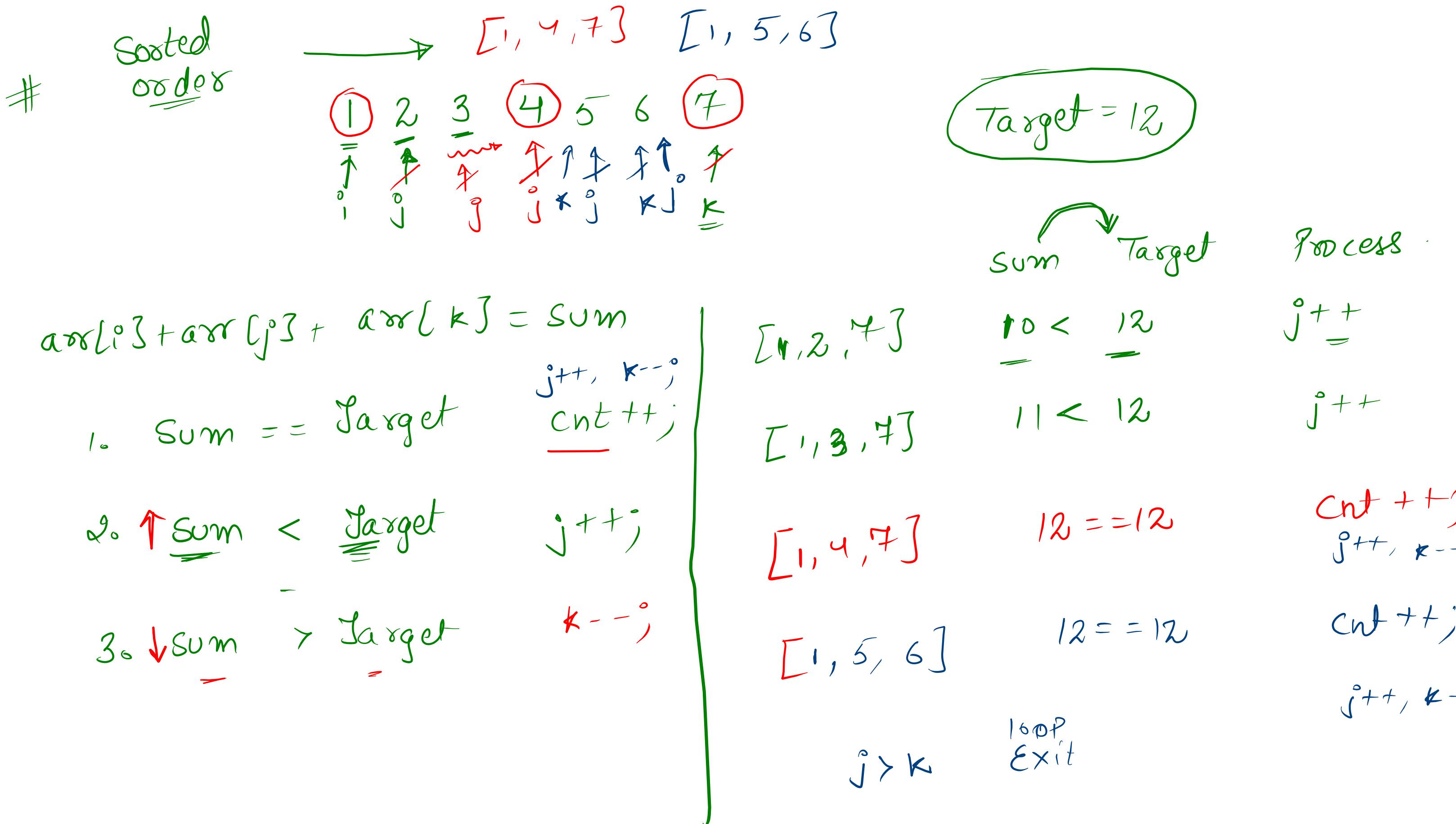
cnt++;

Time complexity

$O(N^3)$

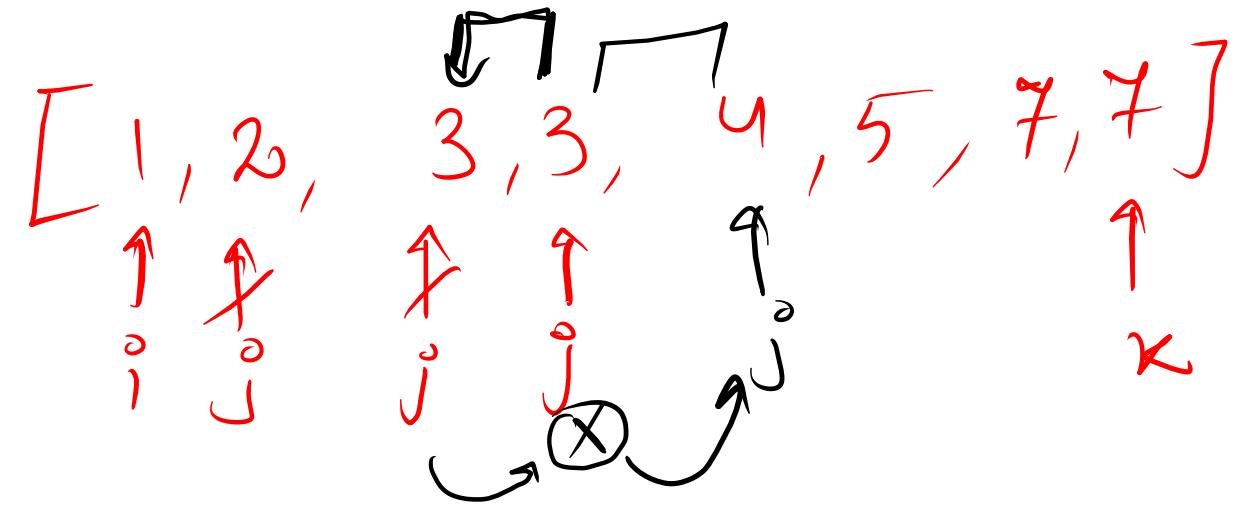
TLE

optimise
 $\hookrightarrow \approx O(N^2)$



$\text{prev} = \text{current}$

let assume



$[1, 2, 7]$

$10 \leftarrow 11 \quad j++$

$[1, 3, 7]$

$11 = 11 \quad \text{cnt}++ \quad j \uparrow$

$\hookrightarrow [1, 3, 7]$

$11 = = \quad \text{cnt}++$

we are increasing cut for

the same subset.

$j < k \& \&$
if ($\text{arr}[j] == \text{arr}[j-1]$) $j++$

$j < k \& \&$
if ($\text{arr}[k] == \text{arr}[k+1]$) $k--$

$[1, 2, 2, 2, 2, 2, 3, \dots, 7, 8, 8]$

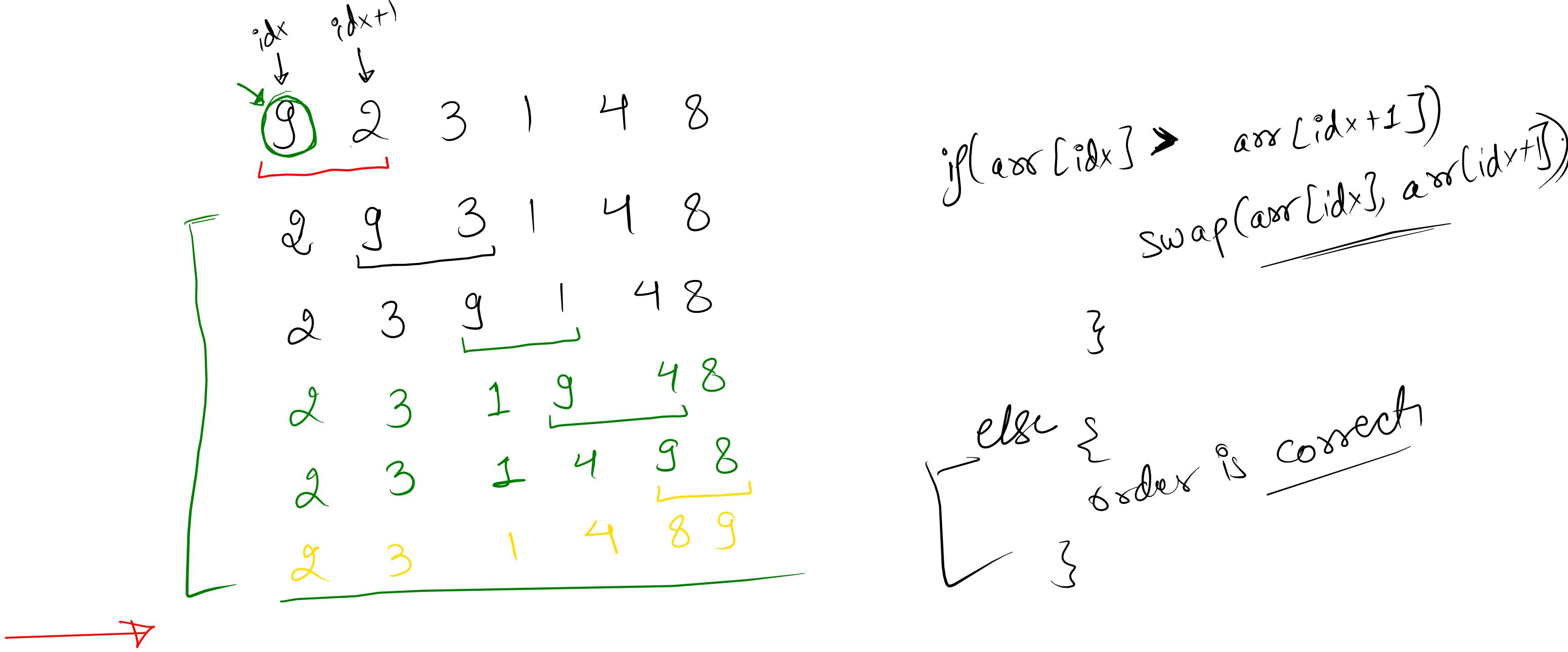
while

Bubble Sort → Push max Element at last by adjacent swapping

→ sorting

I/P: 8 2 3 1 4 8

SORT → O/P: 1 2 3 4 8 9



2 3 1 4 8 9

2 3 1 4 8 9

2 1 3 4 8 9

2 1 3 4 8 9

2 1 3 4 8 9

2 1 3 4 8 9

2 1 3 4 8 9

1 2 3 4 8 9

1 2 3 4 8 9

1 2 3 4 8 9

1 2 3 4 8 9

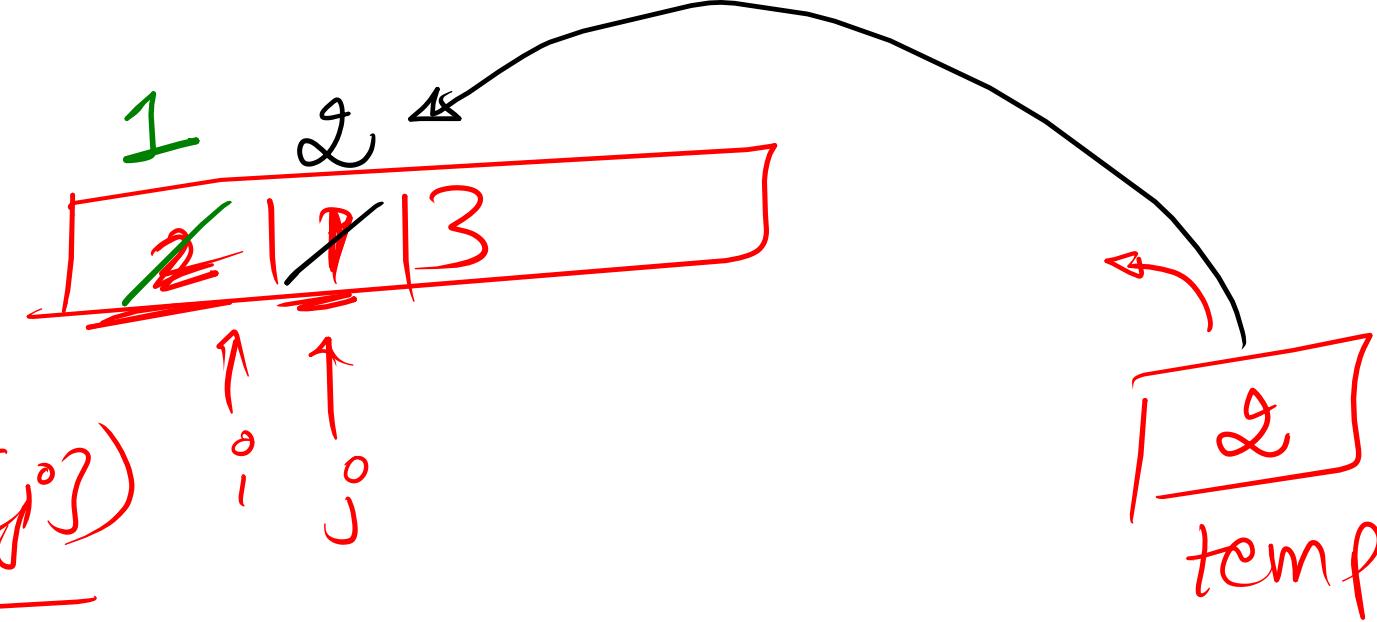
1 2 3 4 8 9

Sorting done.

*

wrong order &

swap(arr[i], arr[j])



Internal working
of Swap

Java ↓
int temp = arr[i]
arr[i] = arr[j]
arr[j] = temp

Selection Sort

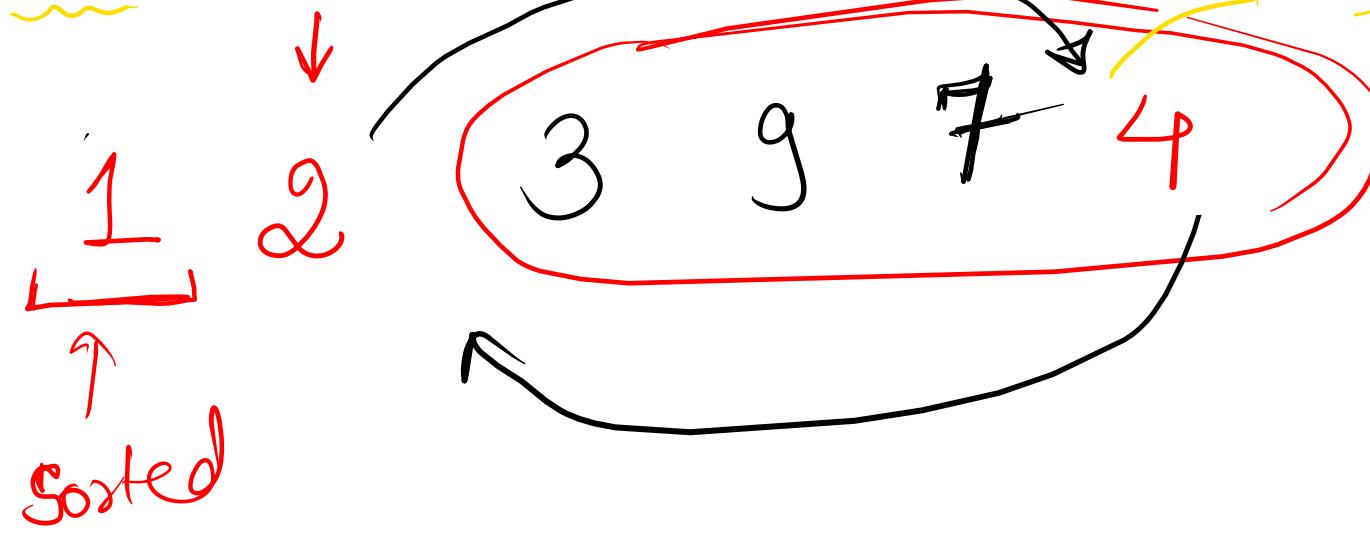
mini check?

{select mini from the remaining arr and swap it}



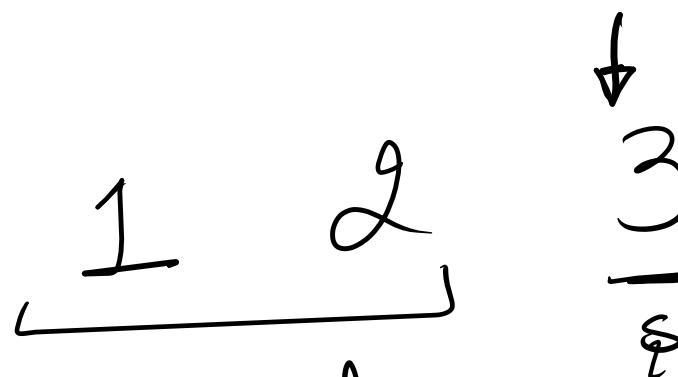
$$4 > 1$$

Sorted
(contain mini)



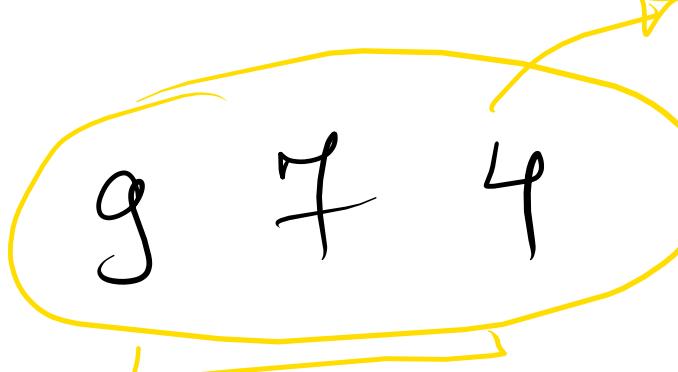
$$4 > 2$$

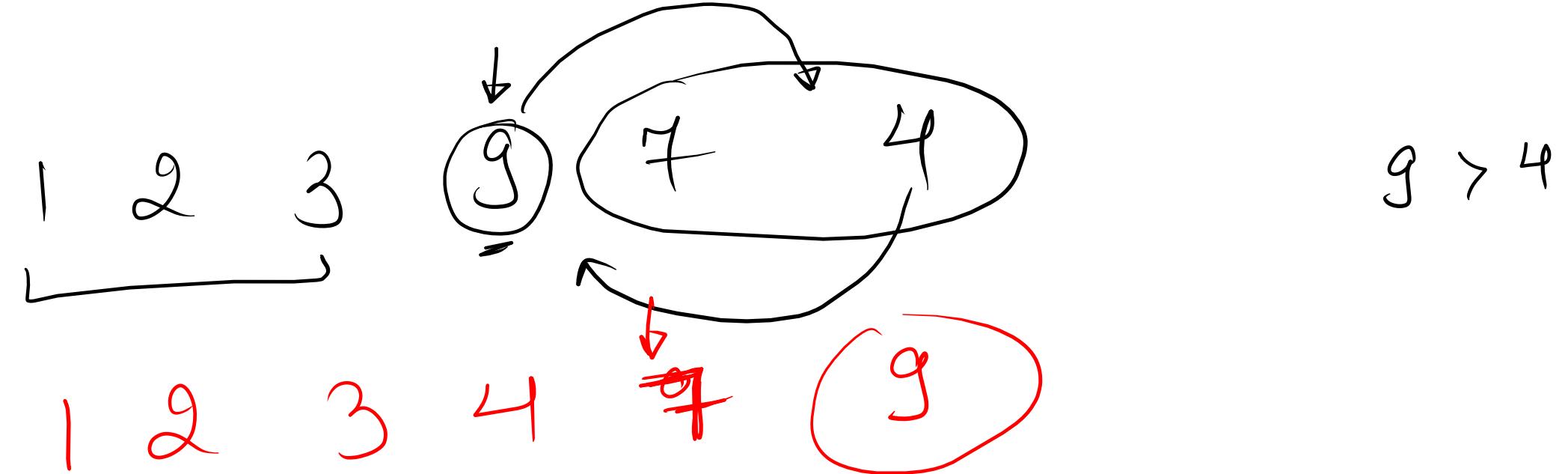
Sorted



Sorted

3
4





Maximum Occurrence

```

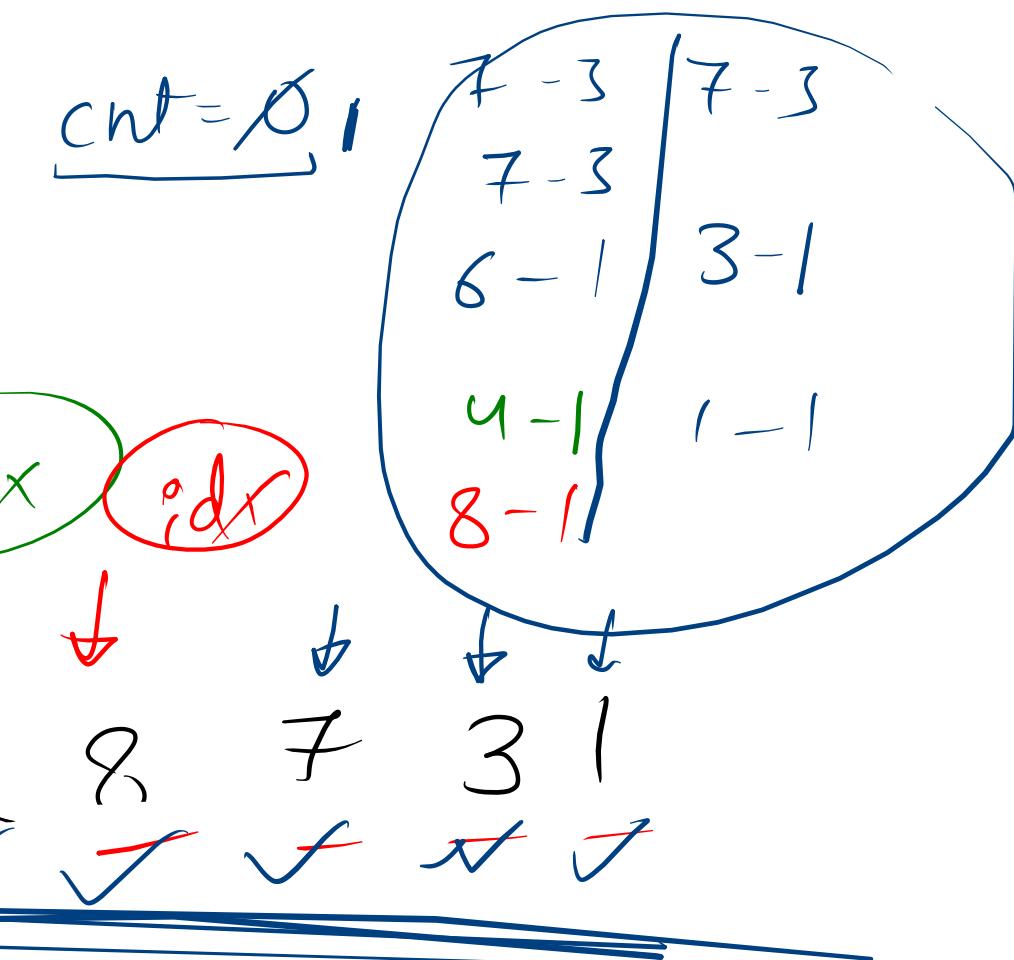
if(arr[idx] == arr[idx+x]) cnt++;
if(cnt > maxcnt) {
    maxcnt = cnt;
    res = arr[idx];
}

```

$$Cnt = \emptyset \vee 3$$

$$cnt = 0$$

$F - \emptyset \vee 3$



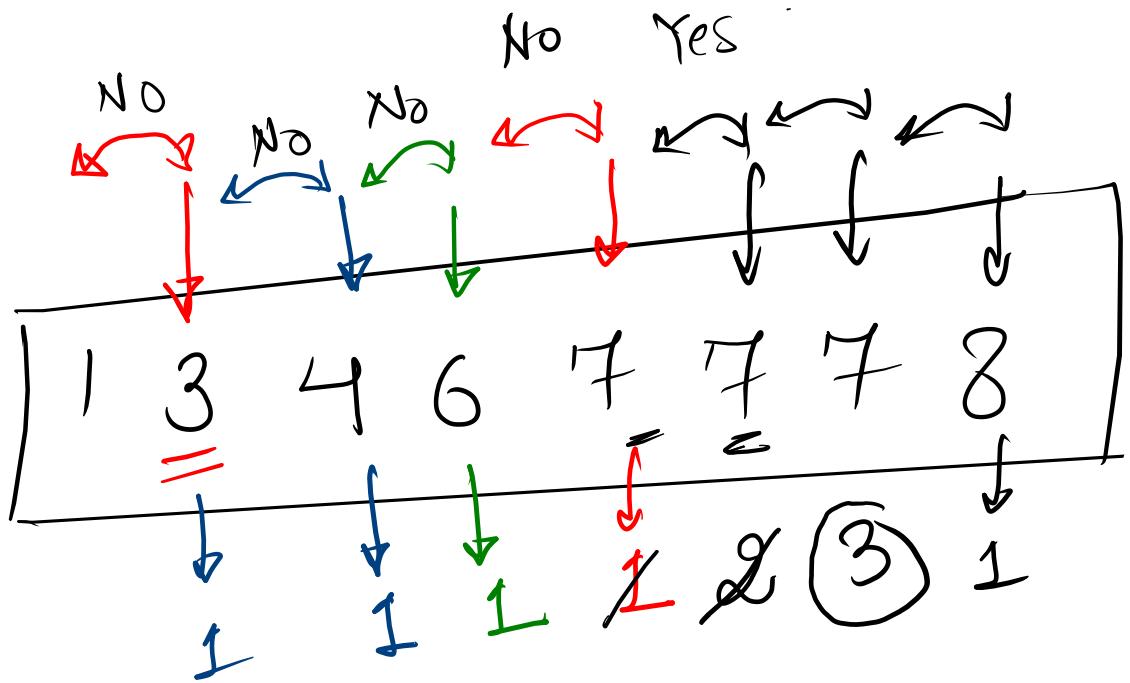
$\# 10^6 \rightarrow$ Sorting

77648731

Sorting

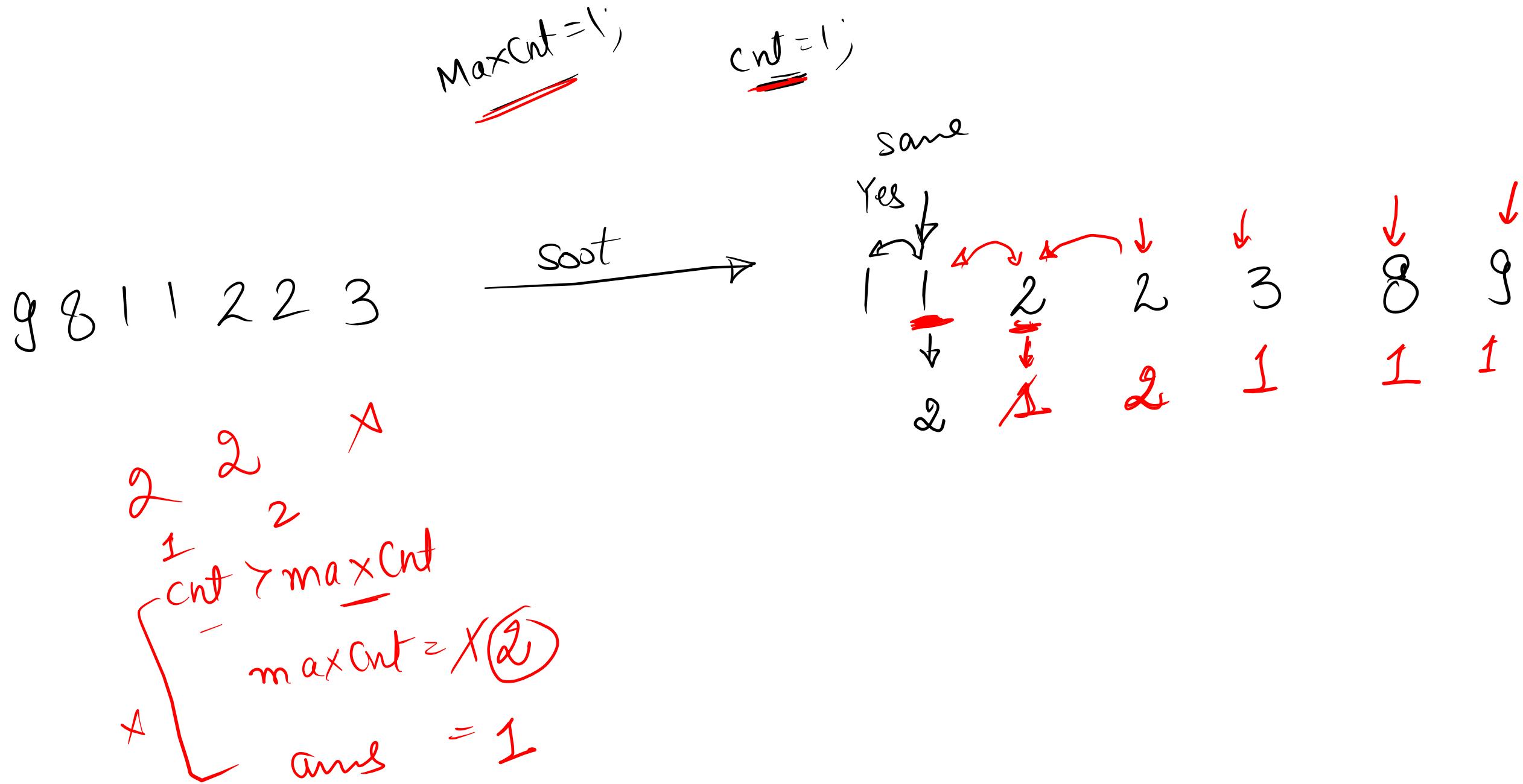
Cnt:

if (cnt > maxCnt) {
 maxCnt = cnt;
 answer = arr[i];
}



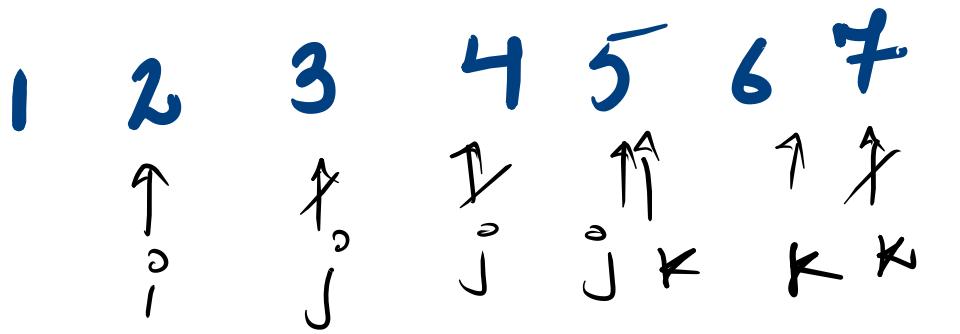
if (arr[i] == arr[i-1]) {
 cnt++;
}
}
else {
 cnt = 1;
}
}

ans we $x = 7$



$[2, 3, 7]$

while ($j < k$)



$[2, 3, 7]$

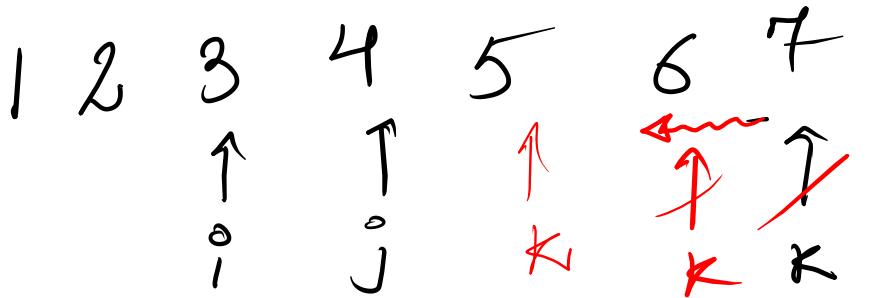
$i = 12$ $\frac{\text{cnt}++}{\text{cnt}++};$

$i = 12$ $\frac{j++, k--;}{\text{cnt}++};$

$[2, 4, 6]$

$[2, 5, 5]$ $j \geq k$ exit

$$\begin{aligned}
 i^0 &= \\
 j^0 &= i^0 + 1 \\
 k &= n - 1
 \end{aligned}$$



$[3, 4, 7]$

$\downarrow 14 > 12$

$[3, 4, 6]$

$13 > 12$

$[3, 4, 5]$

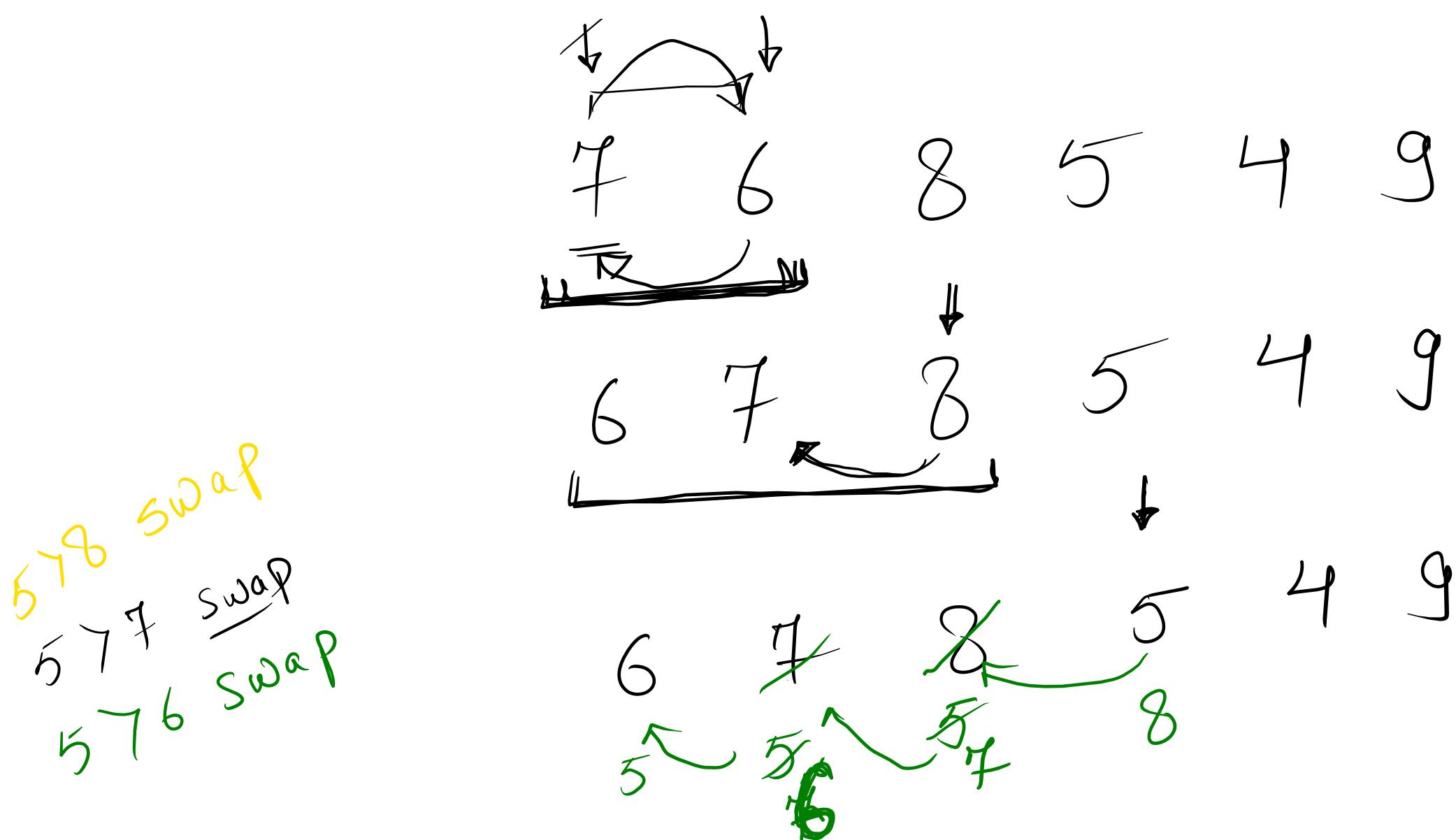
$12 == 12$

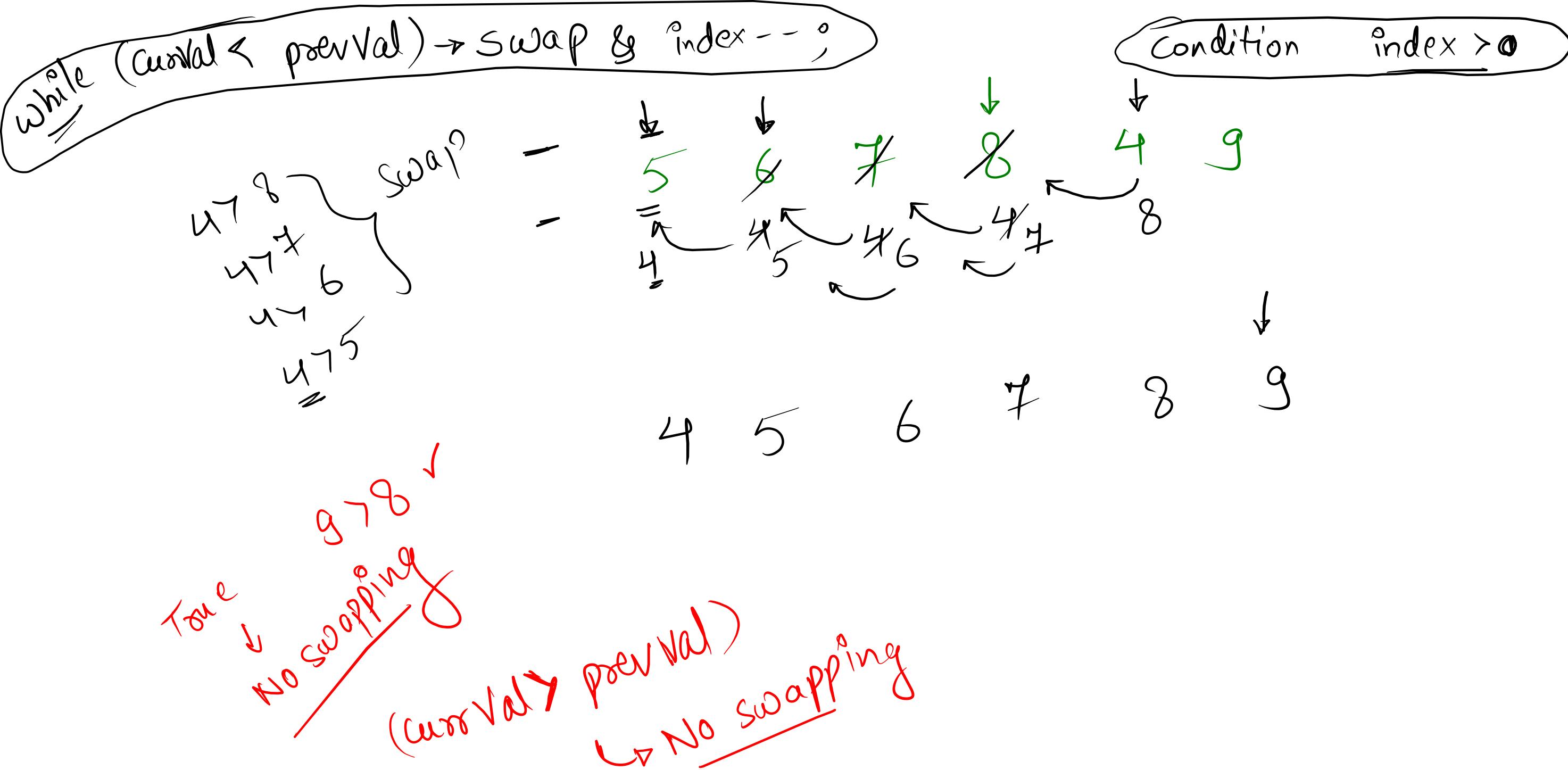
sum > Target $k--;$

$k--;$

$cnt++;$

Insertion Sort





Example - 3

~~if($a[i] > a[i-1]$)~~

~~swap~~

$12 > 11 \rightarrow$ True
 $\cancel{12 > 11} \rightarrow$ No swapping

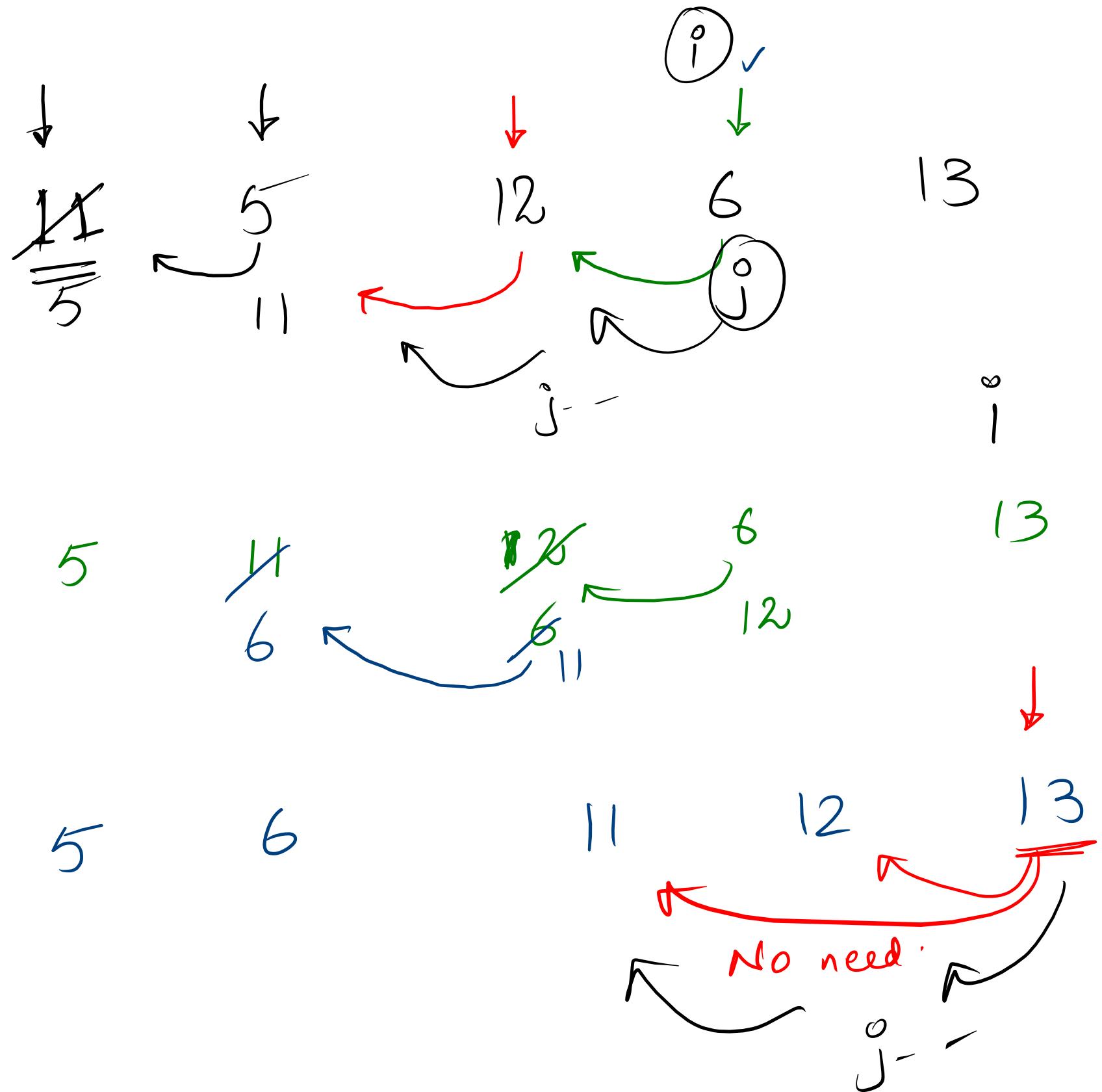
~~~~6 > 12~~ swap~~

~~~~6 > 11~~ swap~~

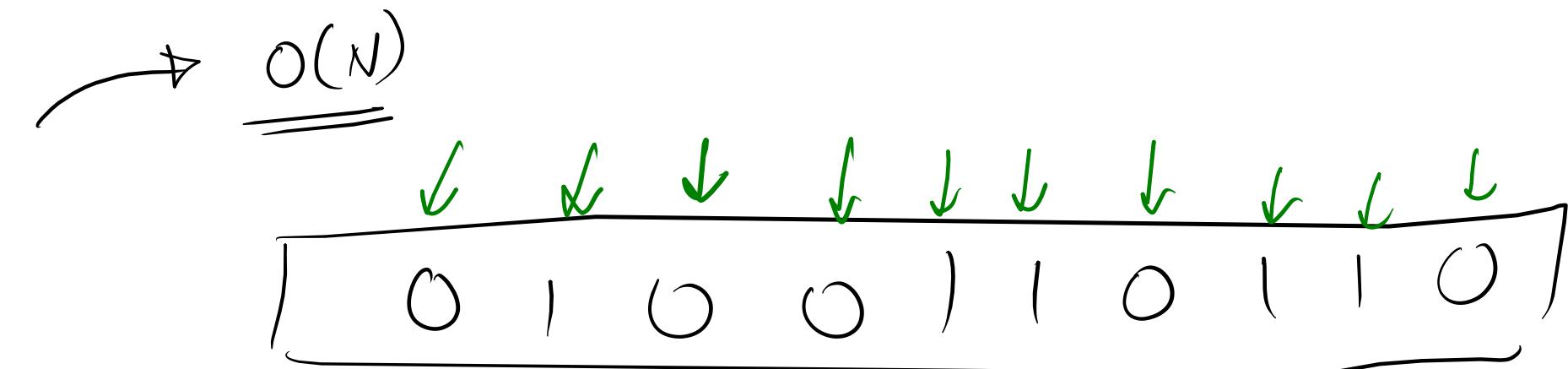
false

True

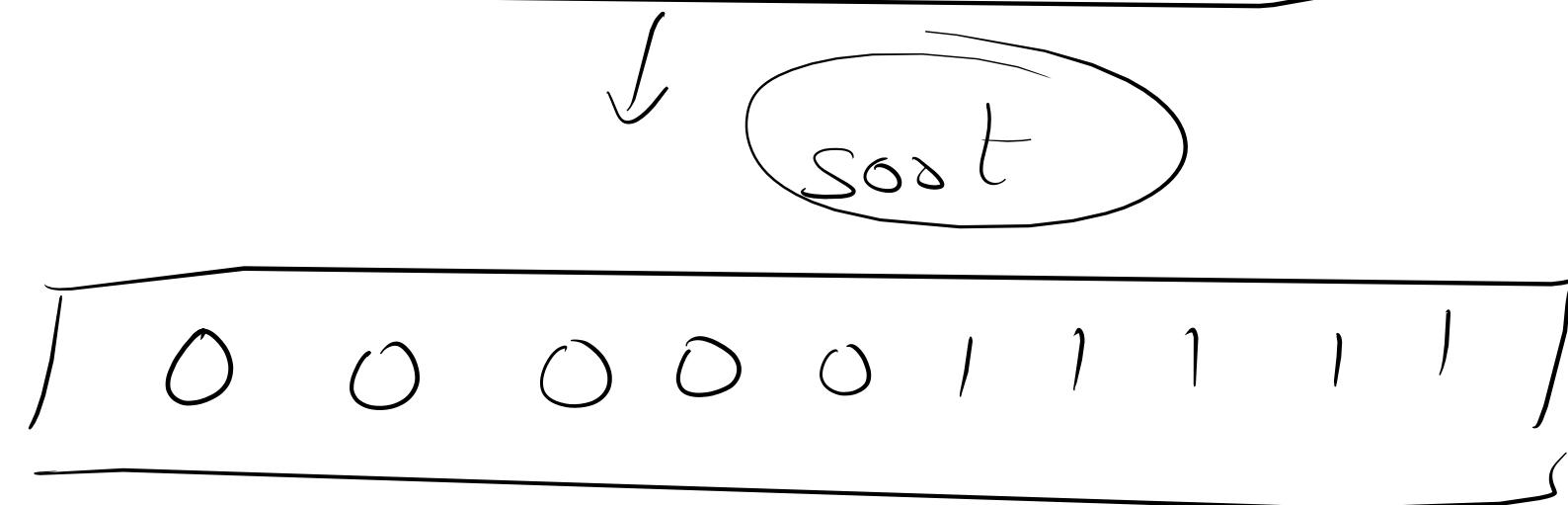
$6 > 5$



0-1 Sorting



BS } $\frac{TLE}{\underline{\underline{O(N^2)}}$
SS }
IS



Approach - 1 : Arrays.Sort(arr) ✓

Approach 2 : result = [1 1 1 1 1 1 1 - 1 1 1]

$$\text{cntZero} = 0 \cancel{1} \cancel{2} \cancel{3} \cancel{4} \cancel{5}$$

$$\text{cntOne} = \cancel{0} \cancel{1} \cancel{2} \cancel{3} \cancel{4} \cancel{5}$$

↓

while (cntZero > 0) {

 result[i] = 0;

 i++; cntZero--;

}

while (cntOne > 0) {

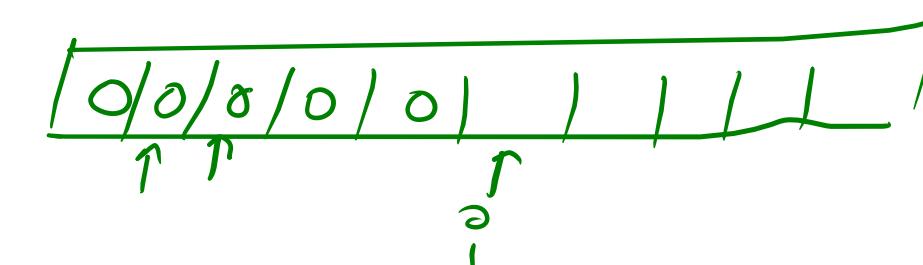
 result[i] = 1;

 i++;
 cntOne--;

}

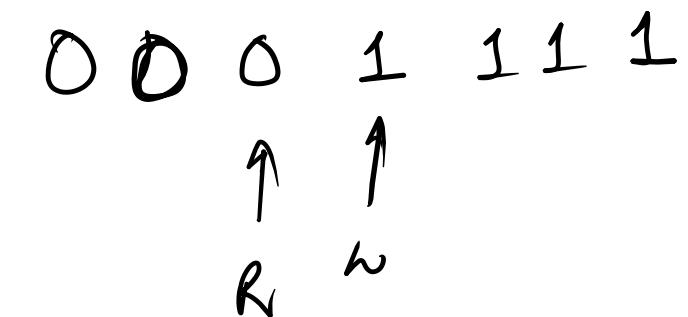
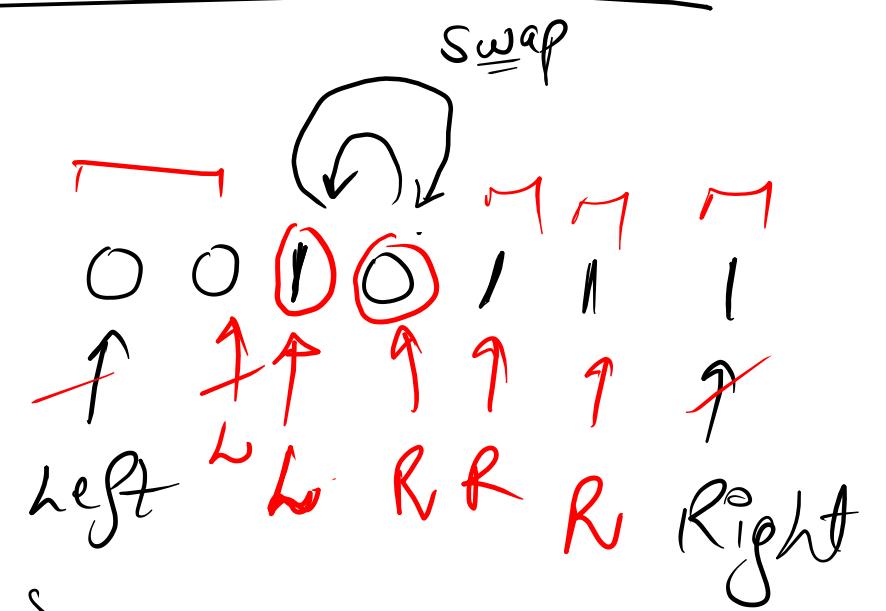
return result;

cntZero = 5 & 3 ... yo
cntOne = 5



Best / optimal Approach

~~while(left < Right)~~

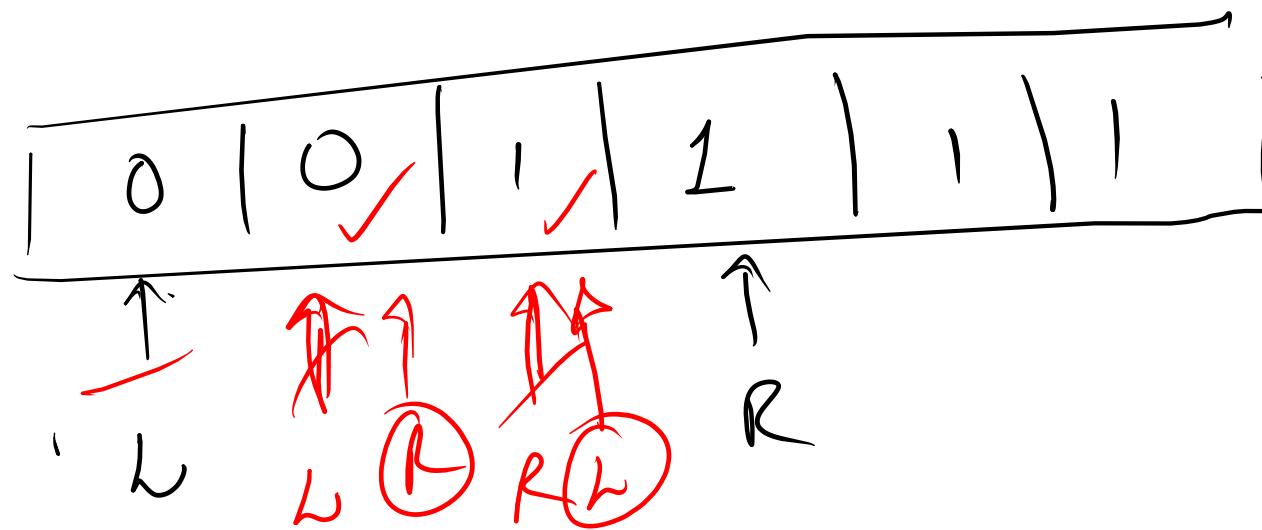
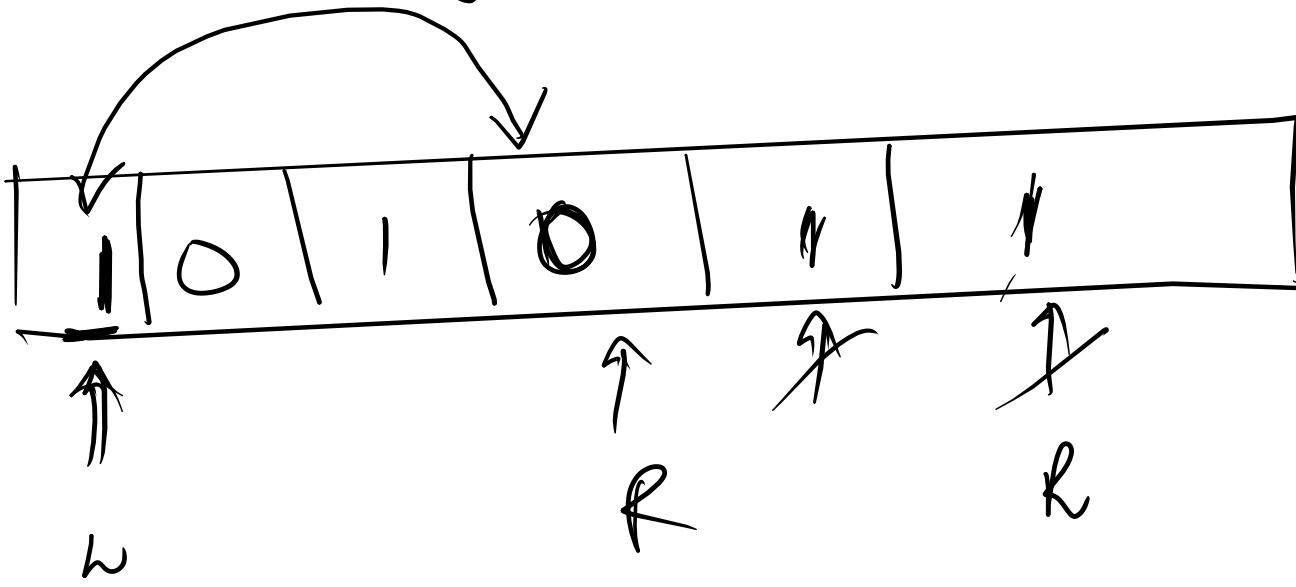


Left pointer $\Rightarrow 0$ point
Right pointer $\Rightarrow 1$ point

2-pointer

↓
if conflict
↓
swap

Swapping done



while
(($L < R$) false
 exit .

Sorting ①

BC → Sort ① ②

→ Beautiful array

→ 2sum & 3sum

→ H/w

* IW = 7 day -
↳ 4 day Test

(8:30 pm)

* Camera → ✓

* Practice

