

Q

AJ Number \rightarrow Maximum no. of AJ no. present in string \rightarrow return count

✓✓
0, 1 \rightarrow AJ \times

[No.] \rightarrow AJ \checkmark

- 1. 0 and 1 are not AJ numbers.
 - 2. 2, 3, 5, 7, 11, 13, 17, 19, 23, 29 are AJ numbers.
 - Any number not divisible by the numbers in point 2 (Given above) are also AJ numbers.
- $x \% [No.] = 0$

let
 $x \% [No.] = 0 \rightarrow$ AJ \checkmark
 Not divisible

given no.

In 4991, both 499 and 991 are AJ numbers but you can choose either 499 or 991, not both.

Common substring.

Test Cases

for finding maximum

case-1

1 2 3 4 5

let suppose
this is an
AJ

Cnt = 1

case-2

let

1 2 3 4 5
check 5 is an
AJ no.

Cnt = 1

" AJ "

if i counted → then my task is to
remove this no.

can i use this substring from

2 3 4

1 2 3 4 5

We are Just Assuming

NO

1 2 3 4

Cnt = 2

1 2 3

Cnt = 3

1 2

Not
an AJ
Cnt = 3

AJ ✓
Cnt
4

Cnt = 3

return 4

i.e. the maxi. no. of AJ no.

* Actual example for understanding "AJ no."

Ex 81 615

1> $0, 1 \rightarrow AJ \otimes =$ ~~Cnt++~~ ✓

2> $[2, 3, 5, 7, 11, 13, 17, 19, 23, 29] \rightarrow AJ \checkmark$

3> $(x \% [] \{ 1 = 0 \}) \rightarrow AJ \checkmark$ Cnt++ ✓

81 Not an AJ no. AJ cnt = 2 (no change)
in AJ Ent

$(8 \% 2 \{ 0) \rightarrow AJ \times$ false

5 $\rightarrow AJ \text{cnt} = \emptyset 1$

new no. 8161

remove 5 $x \% [\text{no.}]$ remainders
6 % 2 == 0 divisible -

AJcnt = 1 $6, 1 \rightarrow$ are not AJ no.

AJcnt = 1 AJcnt = 2

remove 61 from 8161 $(61 \% 2 \{ 0) \rightarrow \text{Yes} \rightarrow AJ \checkmark$

81 True
81 \% 2 { 0) false;
(81 \% 3 ! { 0) break;

Q. stronger string

$$n = 4$$

(a, b, c, d)

$$\text{ans} = cabd$$

$$n = 6$$

a, b, c, d, e, f

$$\text{ans}$$

e c a b d f
4 2 0 1 3 5

even idx
↓
insert.

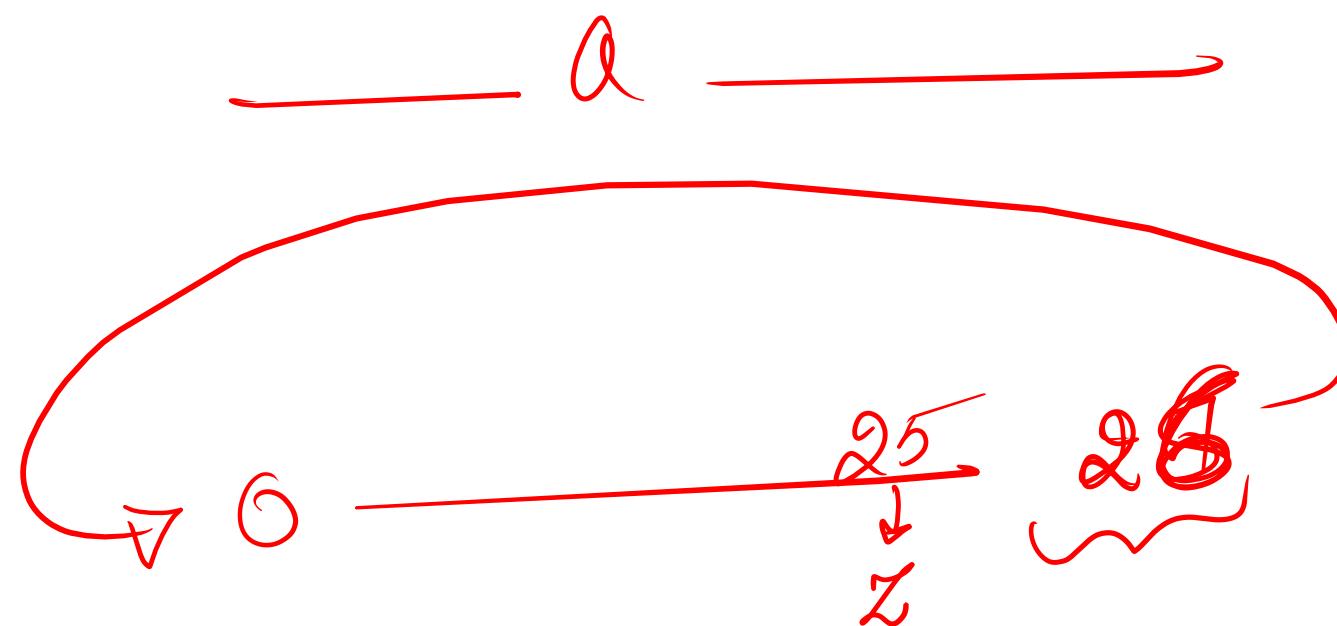
odd idx
↓
append the character

$n = 26$

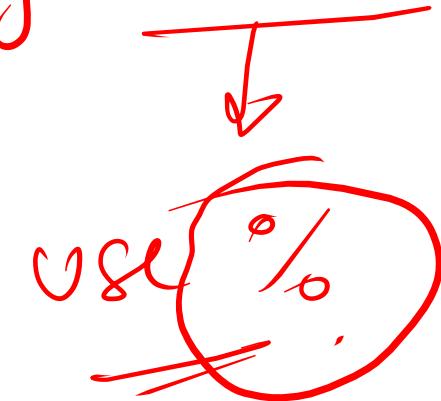


return

$n = 27$



cyclic order



```

static String strangeString(int n) {
    //Write your code here
    StringBuilder ans = new StringBuilder("");
    ans.append('a');
    for(int i = 1 ; i<n ; i++){
        char temp = (char)('a' + (i%26));
        if(i % 2 == 0){ // even idx ---> add the char in the starting
            ans.insert(0,temp);
        }
        else{ // odd ---> add char at the ending
            ans.append(temp);
        }
    }
    return ans.toString();
}

```

$$\begin{array}{l} \text{n=6} \\ \text{n=0} \end{array}$$

$\underline{\text{ans}} = \underline{\text{a}}$

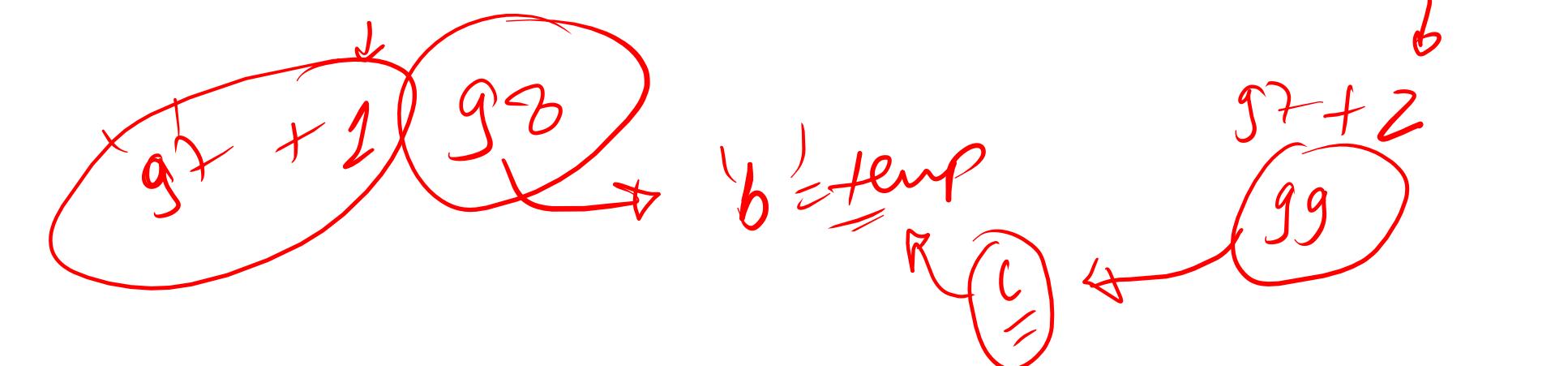
$\underline{\text{ans}} \geq \underline{\text{ab}}$

$$\text{temp} = (\text{char}) ('a' + (i \% 26))$$

1-1.26

$$\begin{array}{l} \text{n=1} \\ \downarrow \\ \text{odd} \end{array}$$

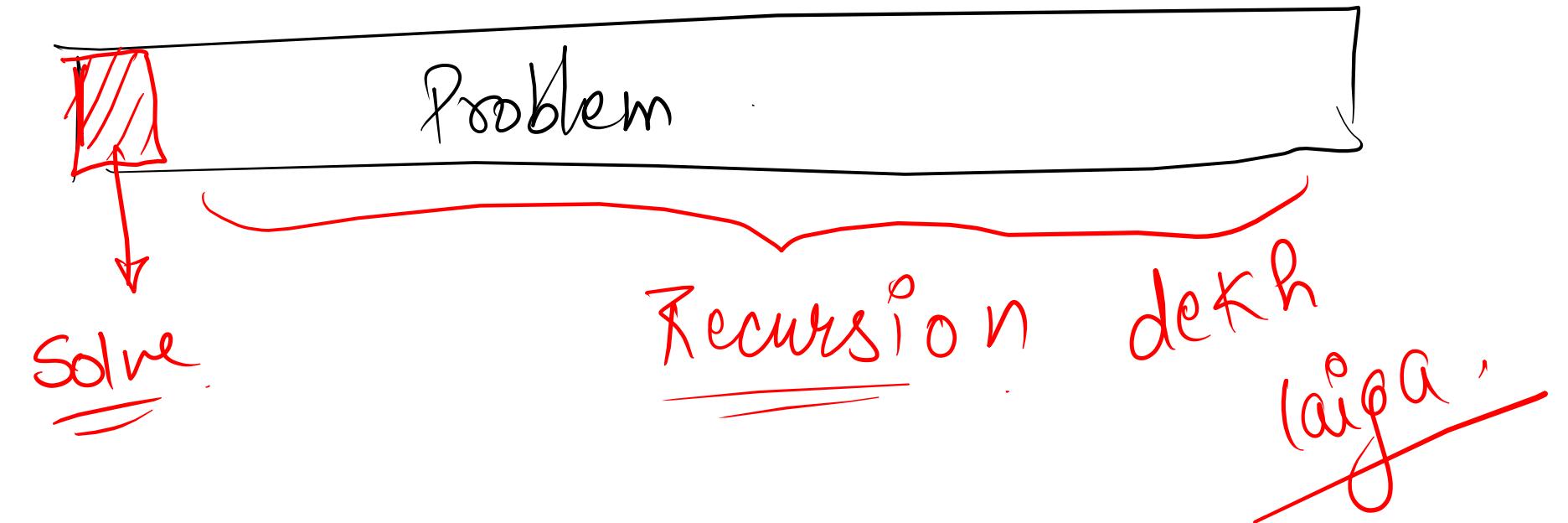
$\underline{\text{ans}} = \underline{\text{g}}$ c a b d f



Recursion

Defination: A function calling itself till the base case is not true :-

divide & conquer



*

factorial

$$\underline{5!} = \underline{5} \times \underline{4} \times \underline{3} \times \underline{2} \times \underline{1}$$
$$= \underline{\underline{120}}$$

factorial (5)



5 *

fact (4)

Recursion

(1)

factorial

$n=5$

~~f(2)
f(1)
f(0)~~

~~f(1)~~

~~f(2)~~

~~f(3)~~

~~f(4)~~

~~f(5)~~

~~call stack~~

$5 \times \underline{\text{fact}(4)}$

Recursion Tree

$4 \times \underline{\text{fact}(3)}$

$3 \times \underline{\text{fact}(2)}$

$2 \times \underline{\text{fact}(1)}$

$0 \times \underline{\text{fact}(-1)}$

$1 \times \underline{\text{fact}(-1)}$

$5 \times 4 \times 3 \times 2 \times 1$

if you didn't add
base case.
Then the
function call
itself at ∞



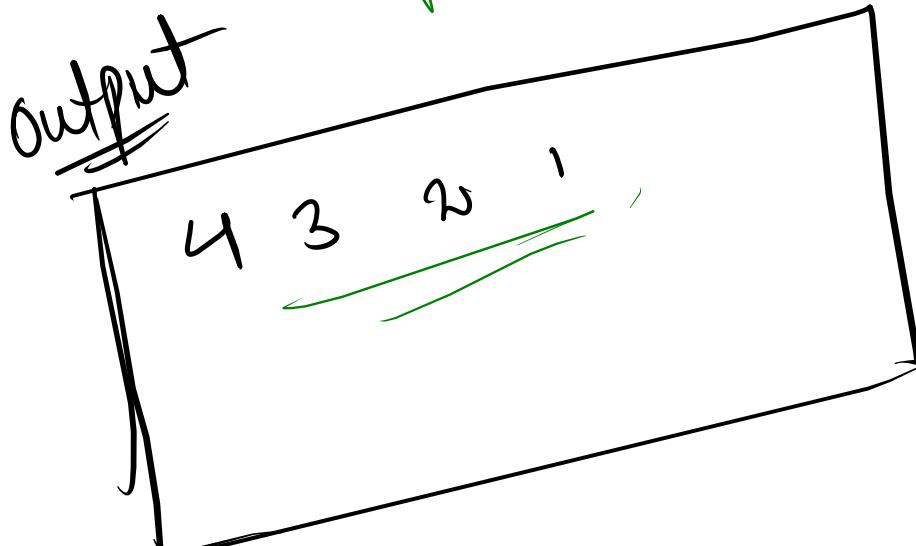
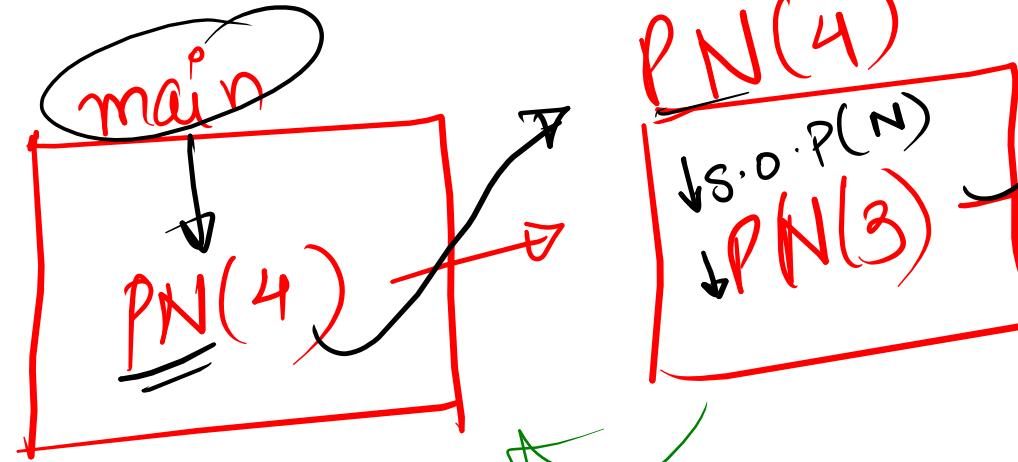
stack over flow

To avoid stack overflow
i have to add base
case.

① Point no. $[1, n]$

$$\underline{n = 4}$$

O/P $\Rightarrow 4, 3, 2, 1$



Base case

\dots
 \dots
 \dots

$PN(0)$

```
if(n <= 0)
    return;
```

otherwise
function
calling

$PN(2)$

$\downarrow S.O.P(2)$
 $PN(1)$

$PN(1)$
 $\downarrow S.O.P(1)$
 $PN(0)$

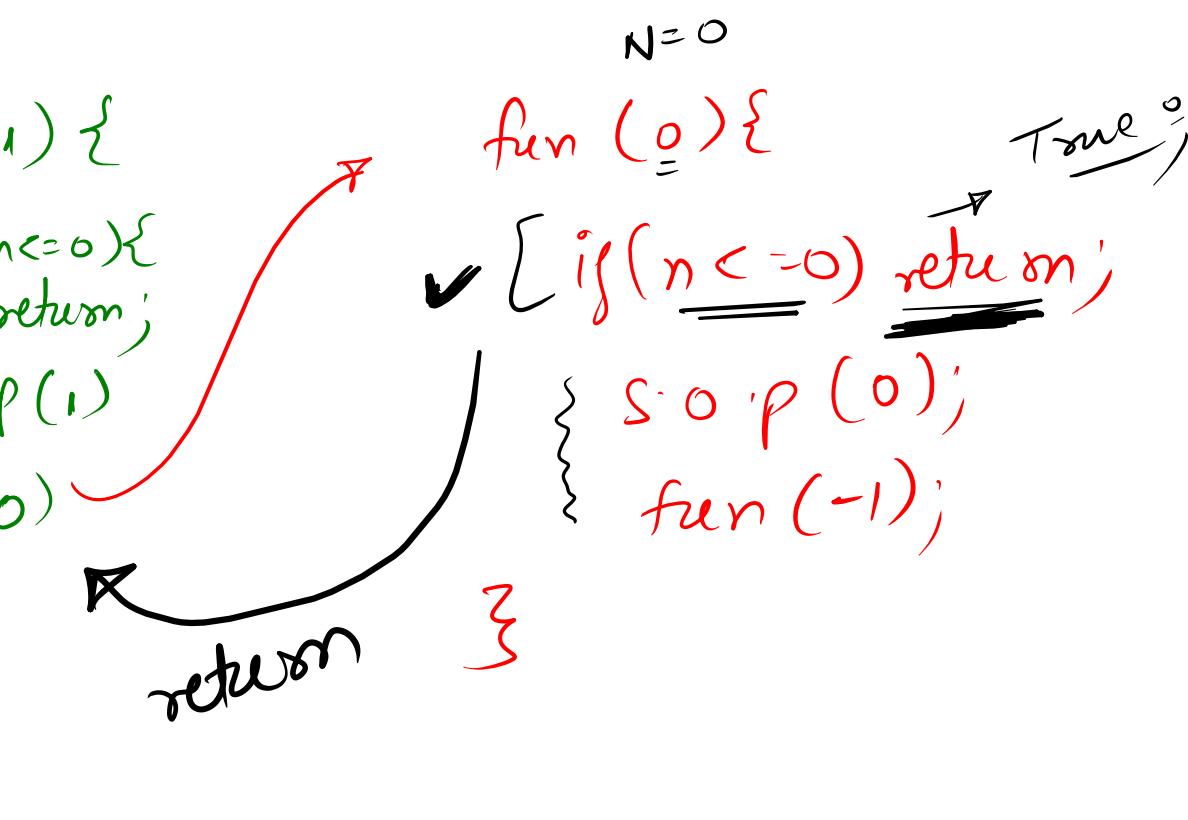
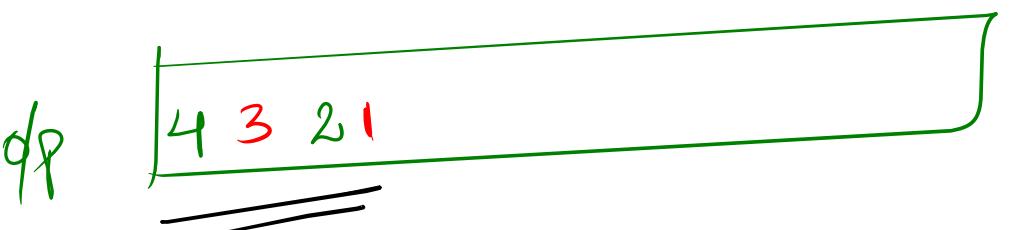
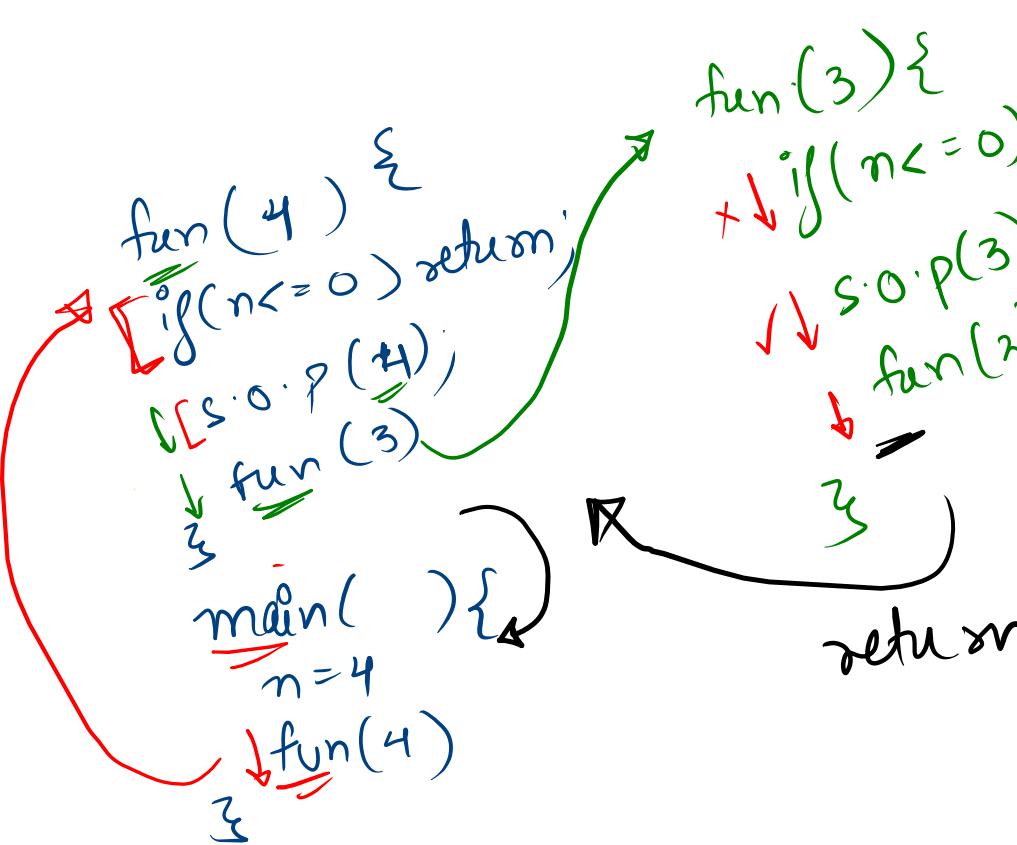
$PN(3)$
 $\downarrow S.O.P(3)$
 $PN(2)$

$PN(4)$
 $\downarrow S.O.P(4)$
 $PN(3)$

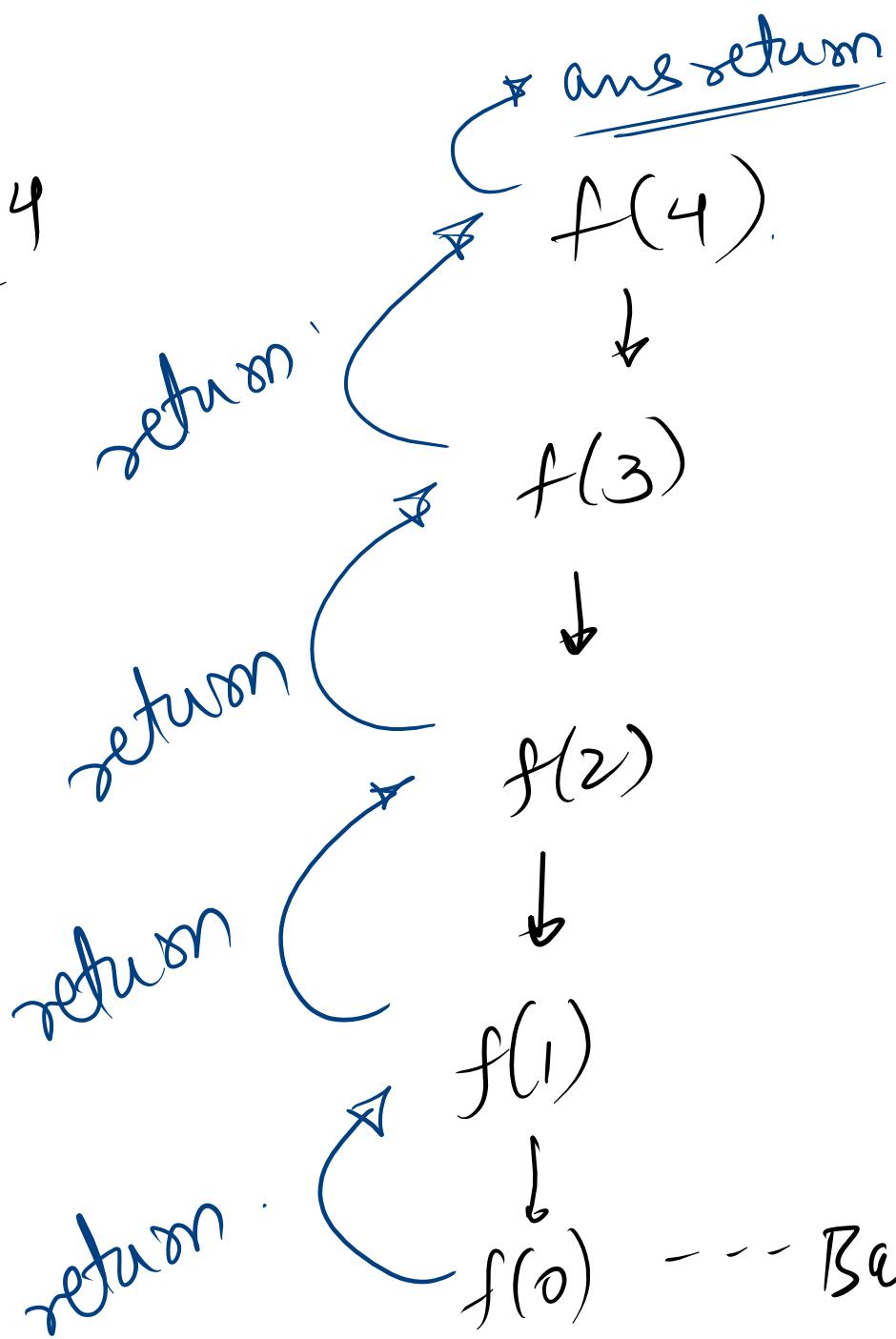
$\underline{PN(4)}$

② Point No $[1, N]$

$N = 4$



$n=4$



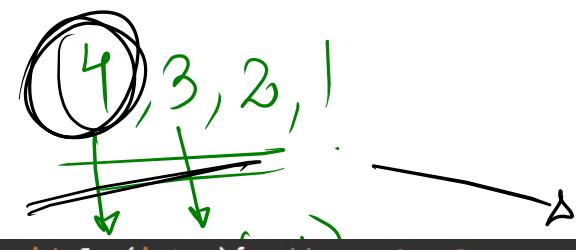
Recursion Tree

* Time complexity: $O(N)$ $\sim n$

function calling
 $n \rightarrow (n+1)$

* Space complexity: $O(N)$



* Point 

```

public static void fun(int n){ //n == 4 3
    if(n <= 0) return ;
    System.out.print(n+ " ");
    fun(n-1); // fun(3) fun(2)
}

```

$$\text{fun}(n) = n, \text{fun}(n-1)$$

Base case 

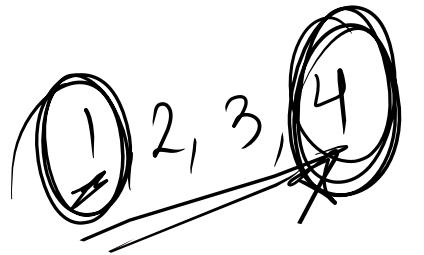
Point(4) ✓
function call ✓
(n-1) ↗

Point(3) ✓
func.all ✓
(n-1) ↗

Point(2) ✓
fu. call(n-1) ↗

PN(1) ✓
f call(n-1) ↗

P1

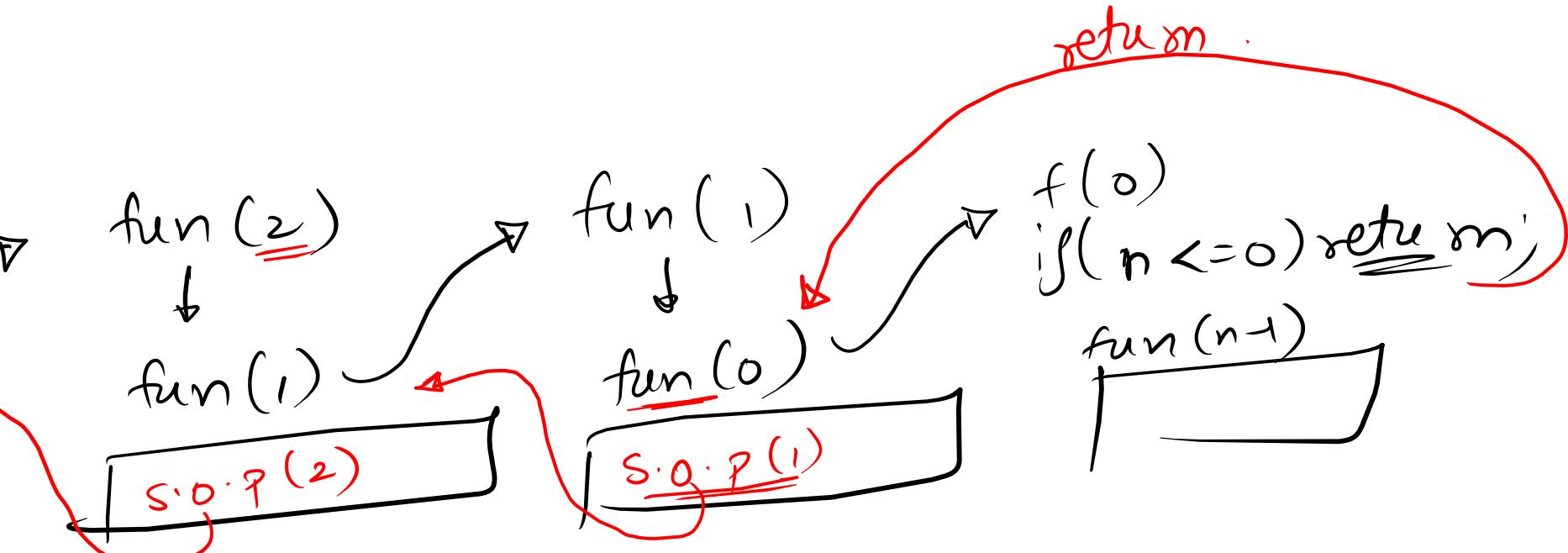
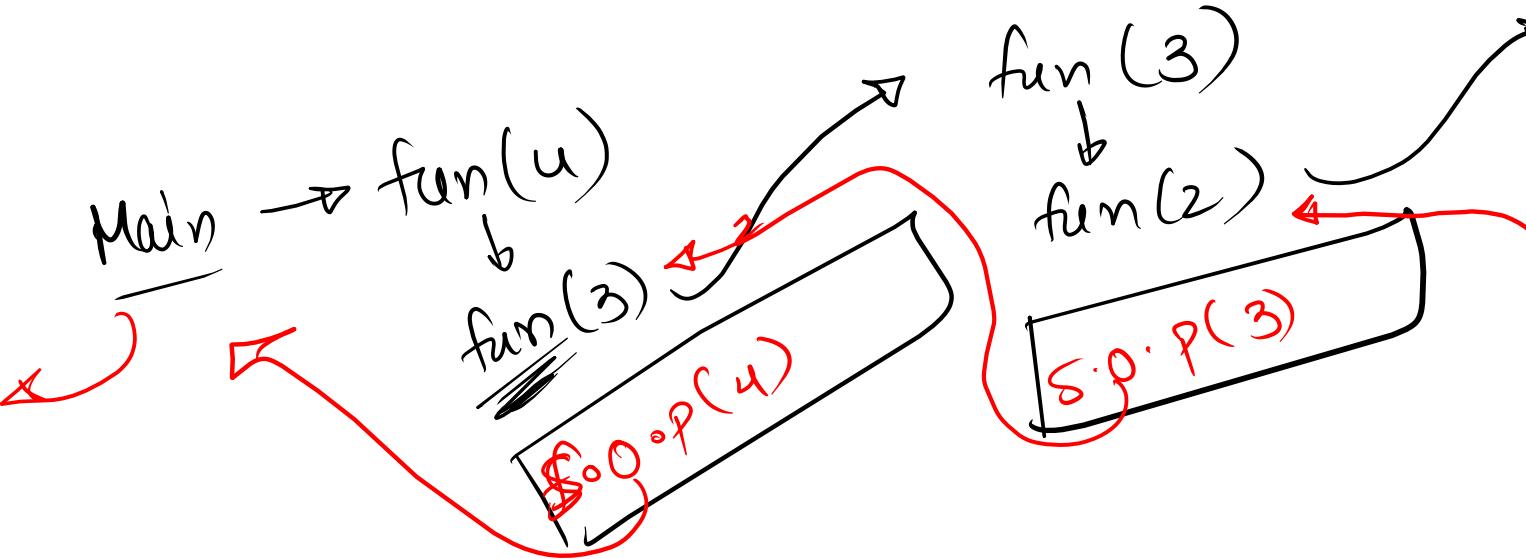
* Point 

```

public static void fun2(int n){
    if(n <= 0) return ;
    fun2(n-1); //fun(3) fun(2) fun(1) fun(0)
    System.out.print(n+ " ");
}

```

* PN 1, 2, 3, 4

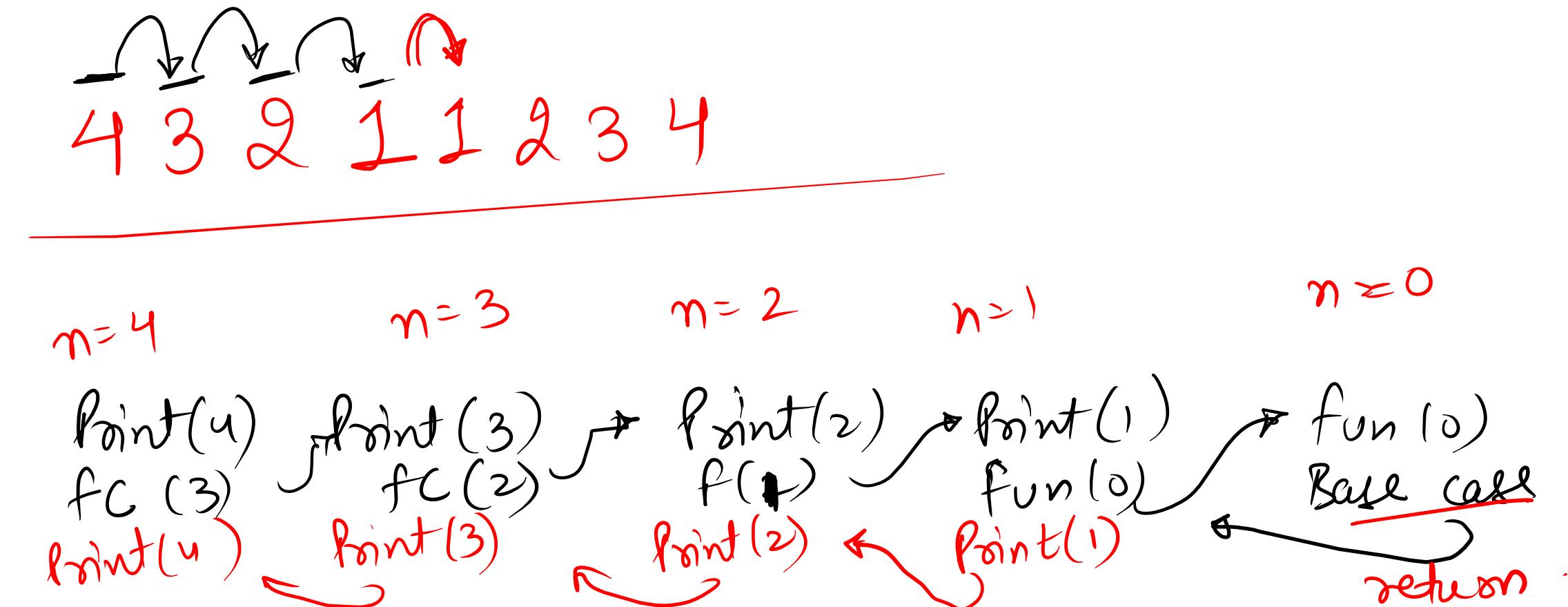


1, 2, 3, 4,

~~point~~

$$n=4$$

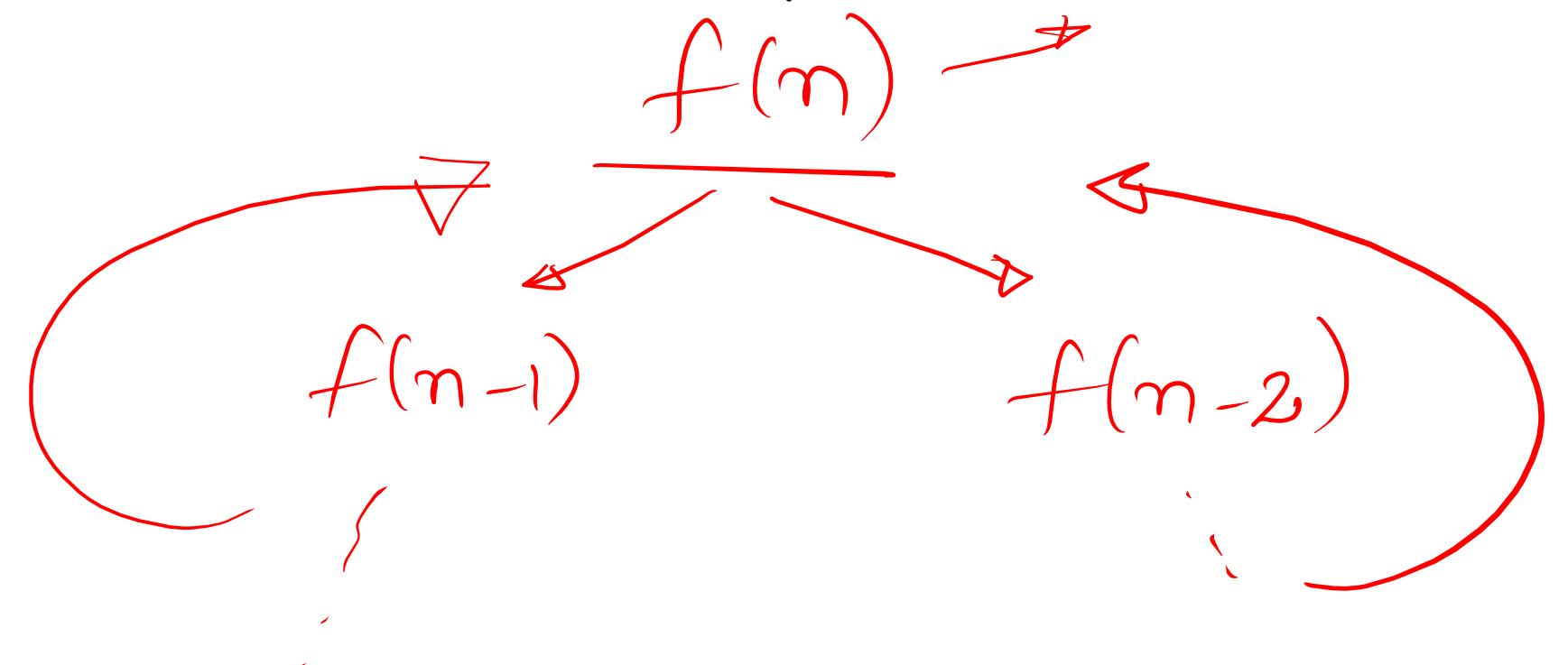
✓ ↴ print()
✓ ↴ func. call()
✓ ↴ print()



* Fibonacci \Rightarrow

1st term 2nd term 3rd term 4th
0, 1, 1 = , 2
 $t_{n-2} + t_{n-1} = t_n$

$t_n = \text{sum}$
 $t_{n-1} + t_{n-2}$



//Base

if ($n = 0$) return 0;

* int $a = \boxed{\text{fun}(n-1)}$

1st term

int $b = \boxed{\text{fun}(n-2)}$

if ($n == 1$) return 1;

2nd term

return $a+b$;

1 0

