

12/3/19

Evolution of Rand. N/w:

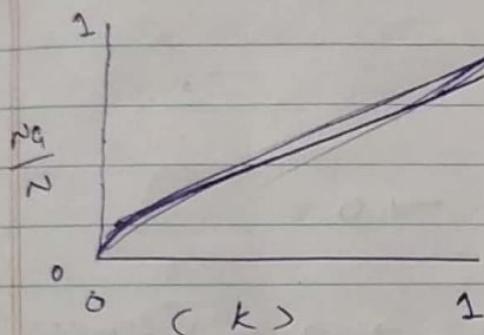
N_G : hub: Giant component: biggest connected component.

How size of N_G varies, when we inc. $\langle k \rangle$ from 0 to $N-1$:

$$\langle k \rangle = (n-1)p$$

if $\langle k \rangle = 0$ then $p = 0$. - all nodes isolated
 extreme cases $N_G = 2$, $\frac{N_G}{N} \rightarrow 0$.

if $\langle k \rangle = n-1$ then $p = 1$ - complete graph
 $N_G = N$, $\frac{N_G}{N} \rightarrow 1$.



linear (expectation of how this ratio varies with $\langle k \rangle$)

But this is not the case!

$\langle k \rangle$ when small \rightarrow lacks large structure, $\frac{N_G}{N} \rightarrow 0$
 ↳ isolated pairs of vertices joined together.
 ↳ no big clusters.

Once, $\langle k \rangle$ inc. critical value $\rightarrow \frac{N_G}{N}$ inc.

$\frac{N_G}{N} \rightarrow$ non linear fashion,
 grows not gradually,
 but substantially.

behaves in diff. way.

$\langle k \rangle \approx$ Giant component \Leftrightarrow each node has avg.
 more than 2 link
 $\langle k \rangle \rightarrow$ does not depend on no. of nodes.
 in the network.

You need to have at least one edge to make that node connected.

\hookrightarrow 1 link is sufficient for a giant comp to emerge.

$$\frac{1}{2} = \langle k \rangle = (n-1) p$$

$$\left[p = \frac{1}{n-1} \right] \rightarrow \text{inversely proportional}$$

~~Giant~~ ^{Huge} no. of nodes \rightarrow less probability \rightarrow we can get a giant component.

Graph

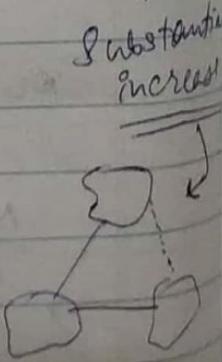
$$\langle k \rangle = (0, 1) \rightarrow \frac{n_g}{n} \rightarrow 0$$

$$\langle k \rangle = (2, 2) \rightarrow \frac{n_g}{n} \rightarrow \text{non-linear fashion}$$

$$\langle k \rangle > 2 \rightarrow \text{almost giant component.}$$

$$\langle k \rangle = 6 \rightarrow \frac{n_g}{n} \rightarrow 1$$

$\langle k \rangle$ beyond 9 \rightarrow giant component.



Subcritical region: $\langle k \rangle < 1$ \rightarrow ratio \rightarrow grey part Date: _____

Critical pt: $\langle k \rangle = 1$ } supercritical region: $\langle k \rangle > 1$ yellow part. ($\langle k \rangle, \log N$)

Connected regime: $\log N \rightarrow$ blue part \rightarrow ratio $\rightarrow 1$

\downarrow every node will be in the giant component.
 $\langle k \rangle \text{ degree} = \log N$ will be sufficient for this

When you inc. $N \uparrow$, N_g inc. slowly ($\log N$)

$\langle k \rangle < 1 \rightarrow$ there might not be any ~~cycles~~ cycle
 they will be corner, like forest
 $(4-5 \text{ nodes})$ many small trees.

Critical pt: $\langle k \rangle = 1 \rightarrow$ form cycles.

- may get a giant comp.
 $N_g = N^{2/3}$

~~two third ratio will be zero~~ $\frac{N_g}{N} = \frac{1}{N^{1/3}} = 0$ when $N \rightarrow \infty$ twothird of the vertices will be in giant comp.

Ex: $\langle k \rangle < 1 \rightarrow \log(7 \times 10^9) = 22.7$.

when $\langle k \rangle = 1 \rightarrow 3 \times 10^6$ - jump of 5 orders.

~~super critical~~ $\langle k \rangle > 1 \therefore (p > \frac{1}{\langle k \rangle})$. substantial jump

almost relevant to real n/w.

Vicinity of critical pt: $\frac{N_g}{N} \sim \langle k \rangle - 1$.

Away from critical pt: $N_g \sim (p - p_c)^n$.

$(p_c = \frac{1}{\langle k \rangle})$ \rightarrow critical pt.

$\langle k \rangle > \log N$ — connected region

$$\underline{N \approx N}$$

Every node is connected

True, Avg. degree ~~depends on N~~ depends on N.

$$\frac{\log N}{N} \rightarrow 0$$

Complete graph when $\rightarrow \langle k \rangle = n - 1$.



$$P = N^2$$

$$z = (\begin{smallmatrix} 1 \\ 0 \end{smallmatrix}, 0)$$

↓ ↓
 $P=0, P=1$

$\rightarrow 0 \rightarrow$ isolated

$-S_2, -S_1, -S_4 \rightarrow$ trees

$-1 \rightarrow$ cycles.

further \rightarrow more complete graphs

Int'l

\rightarrow
empirical
obs.

In real netw, even when $\langle k \rangle < \log N$,
we get connected n/w.

In rand. n/w when $\langle k \rangle > \log N$, then
we get a connected n/w.

~~no proof
or logic~~

Internet \leftarrow fully connected giant component,
 $\langle k \rangle \sim \log N$

~~substantially
less~~

$\langle k \rangle = 1 \rightarrow$ independent of N . apsara
but for $\langle k \rangle = \log N \rightarrow$ dependent on N Date: _____

Erdos Renyi showed only when $\langle k \rangle > \log N$
then n/w is connected (giant comp.) , in
real n/w , we see difference.

real n/w are Supercritical: (yellow region)

$$X \rightarrow \underline{\langle k \rangle}$$

depends on N

when $N \uparrow$, $\log N$ depends
accordingly

\Rightarrow Why to study rand. n/w :

to show that rand $n/w \neq$ real n/w .

In 2000, Barabasi showed that

for
real
 n/w

Base pt. to refer to \rightarrow rand. n/w .

\downarrow not same

\rightarrow it is scale free n/w .

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Small World Phenomenon:

Neural world, Brain research - comparison b/w
an abnormal brain and an average brain.

Comparison of rand. n/w and real n/w w.r.t.
different phenomena:

① Degree Distribution.

② Small World Phenomenon.

Erdos \rightarrow like a star, more than 500
co-authors, that's why erdos no.
 \hookrightarrow lived a nomadic life.

→ Graphs in around & intermediates we are able to reach most of the targets.

→ WWW 2008, MS messenger.

- If tree like str. with $\langle k \rangle$ friends / node
 - ↳ triadic nature. → small no. of loops.
- In real n/w, triadic closure reduces loops

Diameter → max. length n/w should have:

$\langle k \rangle$ — one node is connected to k no. of nodes (on avg.)
 with 2 loops: $\langle k \rangle^2$ possibilities. → no. of nodes
 with d loops: $\langle k \rangle^d$ nodes. on level 2
avg. (approx)

$$N(d) = 1 + \langle k \rangle + \dots + \langle k \rangle^d = \frac{\langle k \rangle^{d+1} - 1}{\langle k \rangle - 1}$$

$$N(d_{\max}) = N$$

$$\langle k \rangle^{d_{\max}} \approx N$$

↑
no. of nodes it
is covering.
(all nodes
covered: N)

$$d_{\max} \approx \frac{\log N}{\log \langle k \rangle}$$

it holds good in rand
n/w, but not in real
network.

if k fixed, vary N , $d_{\max} \propto \log N$

$$d \approx \frac{\log N}{\log \langle k \rangle}$$

→ average length in the
graph will be smaller to it.
(Not diameter, avg. length)

fb \rightarrow someone has 50000 friends, some have 1, 2 friends.

\hookrightarrow not seen this behaviour in Rand. n/w.

apsara

average length: considers all avg. length these extreme cases will give big diameters: affected by outliers

Diameter is distorted due to some external factors also, so here, we talk about avg. length and not diameter.

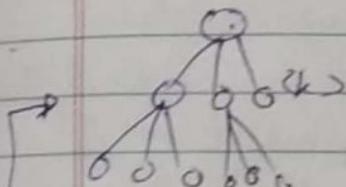
(d)

Borabadi

Table 3.2

$\langle d \rangle \rightarrow$ avg. length b/w any 2 vertices.

$$\langle d \rangle \approx \ln N / \ln k \quad \begin{matrix} \text{random} \\ \text{(no. of loops)} \end{matrix}$$



\rightarrow it is perfectly balanced, but it is not so in real graph.

\hookrightarrow Rand. n/w

that's why d_{\max} of Rand. n/w $\approx \langle d \rangle$ in real.

In real life, friends of a nodes also ~~might~~ be friends themselves.

\hookrightarrow Clustering coeff.: measure of transitivity.

Every link has prob. 'p'

total no. of links b/w neighbours: $\frac{k_i(k_i-1)}{2}$

$$\langle L_i \rangle = p \cdot \frac{k_i(k_i-1)}{2}, \quad p = \frac{\langle k \rangle}{N}$$

$$C_i = \frac{2 \langle L_i \rangle}{k_i(k_i-1)} = p = \frac{\langle k \rangle}{N}$$

$\frac{L}{N} \propto \alpha^2$

expected
clustering
coeff.

\rightarrow not following.

\rightarrow expected $\left(\frac{1}{n}\right)$ behavior.

no. of nodes-

\rightarrow Clustering Coeff
is independent
of n , in
real Network

Watts - Strogatz Model

Small World (b) → connecting to a node
outside their group.

gives Almost a regular lattice: close to real n/w.

No: 19/3/19 In random n/w a super critical region occurs when $k > 1$ & $k < m, n$, but in real network the isolated nodes in super critical region is in giant component - it differentiates both.

& Small World Phenomenon

(The graph will be very sparse)
but it is connected and ~~it~~ takes for any two nodes that will be connected through 5-6 hops.

Small World Phenomenon:

'has'
weak links & strong links

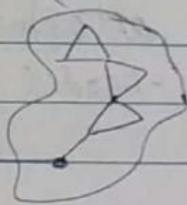
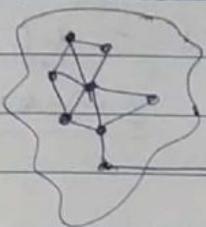
The real n/w are not random n/w but it has small world phenomenon. Can you proof it has small world (to model it some thing which has small world phenomenon).

Watt Strogatz Model :

The main reason real is not random - we are not having the transition (treads) in rand. n/w.

Homophily (close friends)

- weak ties (global nature) as to connect to outer n/w
- Homophily creates triangles - (local in nature).



→ A grid will solve the homophily but not capture weak tie.

* regularity is homophily
randomness will capture the weak ties.

→ till then grid also it is not equivalent to real n/w.

Milgram - decentralized search - every port where person wants to send optimizes a certain strategy.

This random (weak ties) not capture by grid model
Two random nodes will be connected with a prob.

$$\frac{1}{d(v,w)^2} \quad g = 0, \rightarrow g = 1 \quad (\text{same as grid})$$

g small you are connected to your friends (not random friends, you are connected together).

$$d^2 \cdot \frac{1}{d^2} = 1 \quad : g = 0 \text{ to } \infty \quad g = 1.5 \text{ to } 2$$

Blogging - Medium

→ Distance in graph:

(i) geographical dis.

(ii) closest friend

(iii) 2nd level of friend.

Model which is generative in nature - closer to real world phenomena

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WWW Network : testbed for many models

→ Preferential attachment Model - defined by Barabasi and Albert (BA Model)

→ Scale free network :

www → say - in development of n/w theory.

started with it as a rand. n/w, but then they observed it is not. - testbed for most n/w. main diff. b/w rand. n/w and www n/w.

Graph: green: poisson straight \rightarrow real graph for www: directed graph kin, leant.

→ Pareto distbn : power law distbn.

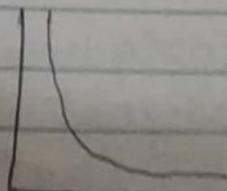
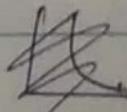
$$p_k = k^{-\gamma}$$

$$\log p_k \sim -\gamma \log k$$

$$\gamma_{in} = 2.1 \\ \gamma_{out} = 2.45$$

we have taken log hence the straight line with slope $-\gamma$

if we don't take log then power law distbn will be like:





- Can't avoid it when
soc. ~~aspects~~ like
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Pareto principle - (80/20 rule) (management)

↳ 80% of wealth is with 20% ppl
in ~~fact~~

- Path b/w any 2 nodes in scale free - may not have many hops.
- More no. of nodes will have less degrees - ~~reduces~~
An poison $\rightarrow \langle k \rangle$ will have most no. of nodes
↳ not observed in real n/w.

Q1 www is scale free n/w. - PL Pietron.
Many others n/w are also scale free, how they converge to a sim. scale free n/w?

- Erdos Renyi $\rightarrow G(n, p)$
 $\frac{1}{n}$ fixed. (nodes: fixed)
but no. of nodes \nearrow
in www \rightarrow nodes \nearrow from 2 to trillions -
or citations \rightarrow " "

The nodes having high degree: rich and rich becoming rich.

- Growth of the n/w : Erdos and Renyi did Preferential attachment. not look at

* BA Model

→ no nodes, each node has at least 2 degree

Growth: add a new node with m links

preferential attachment: prob. of association

$$(A \text{ node } v_i \text{ with degree } k_i) \rightarrow \frac{k_i}{\sum_j k_j}$$

if $k_i \uparrow$ then more prob.
↳ rich becomes richer

Drawbacks: - not explain the initial graph to start with

- if nodes have to be added 2 by 2 or all at one time?

↓
refined by LCA:
Given $G_i(t-1)$ → build $G_i(t)$ - by adding node $v_i(t)$

t : iteration → adding $v_i(t)$ node

v_i → already in the graph:

to join v_i and v_k prob:

$$\left\{ \begin{array}{l} \frac{k_i}{2t-1} \\ \frac{1}{2t-1} \end{array} \right. \quad \begin{array}{l} \text{if } i \leq t-1 \\ \text{if } i = t \end{array}$$

connecting
node to a node
already in the
graph

very less. (self loop)

C: clustering coeff.

apsara

Simulated Graph : BA model

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Avg. Path length is almost the same.
Small World Phenomenon ✓
but Clustering coeff is different.

BA model is able to capture small world phenomenon by growth and preferential attachment but it is not able to produce C° → triadic closure (the triangles it is not able to produce).

Fitness value of each node.

BA → rich becomes rich → earlier alta-vista search engine was very popular but then when Google came → it turned upside & down.
↳ rich becoming rich not true any more
↳ new entrant was rich now.
Fitness value of BA went down.

prob. of associating a new node with already existing nw → need to consider both fitness values.
↳ Barabasi - Barabasi Model

⇒ Information Diffusion

⇒ Community Detection.

→ Through BA, we are able to produce Scale Free Graph.

Page No.: _____

BA follows approach: rich becomes richer.
Eg: Windows → becoming popular and popular
by this model who ever comes later on will not
become popular at all.

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Bianconi Barabasi Model

y_i = fitness value
 k_i = degree

assoc. one num. value with each of the nodes: FITNESS

Weighted avg: $\frac{\sum y_i k_i}{\sum y_j k_j}$: Preferential Attachment

Eg: NOKIA v. popular earlier → but after few years it vanished:

(external factors)

fitness takes care of exogenous / external
Degree " " endogenous / factors.
internal factors.

~~Zafarani~~

Information Diffusion in Social Networks

Oreo → reached a lot of people without spending

Online forestom → a neg. message of a corporate becoming viral

info. spreading through physical / non-physical interactions.

Sociology, Epidemiology: like spread of disease
ethnography: ~~culture~~ culture study

Modelling the Communication

When info spreads : it has a source; senders ①

↳ receives ② : who receives diffused info.

in Diff. sets Act of all recv. >> set of all senders.

③ medium: to pass the info - personal communication, or social media sites.

Types:

① ~~#~~ Herd : global dependence - based on other's ppl decision - they will take their decision what your friend or nbg. is doing, not based on proper rational thinking.

② Info. Cascades : Just observe their nbg, in prev. case we were looking globally, at all ppl, here we just look at our immediate nbg's.

↳ Both are kind of similar.

How many ppl adopted new technology?

↳ network is not visible here, we only know rate of adoptability of the new technology that is Diffusion of Innovation.

↳ They want to diffuse the innovation at very fast rate to become popular.

④ Epidemics - similar to ③, pathogen-transmitting from one individual to another ind.

In diffusion of innov. → you make a decision to buy a smart phone, but in epidemics' you aren't making a decision, pathogen if is in your body then it will do whatever it wants.

Explicit : Now is observable, in
Implicit : .. . not"

→ Eg. of Herd Behavior: Auction - bidding ant. of all the others (global)
(reputation and expertise also comes into picture.)
But. of trust is also there in bidding.

→ If you want to make a decision then we go with a known devil rather than unknown devil → herd behavior → wherever crowd goes → its evolutionary impact → not completely irrational

Eg: → Solomon Asch's Expt., - elevator video.

→ Herd behavior → common factor - global pressure (public)

more like a complete graph.
- how to be some decision made, in seq. orders.

HERO BEHAVIOUR: (4 conditions)

Decisions are not mindless → it is based on what you observe.

↳ no message passing is possible: public info.

→ Woman → 3 marbles - red or blue

2 red, 1 blue or 2 blue, 1 red

decision → publicly visible.

if the first two decisions are BB -- then every next dec. will be B only. $\rightarrow \text{prob: } \frac{1}{2} > \frac{1}{2} = \frac{1}{3}$

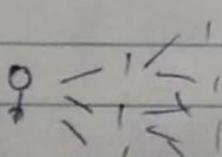
Bayesian Model: to study herd behaviors.

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new slide

Herd intervention:

- ↳ To stop it, you need to intervene by passing some unknown information. - rumors.
- A child shouting: emperor doesn't have any clothes was an intervention.

 → a person starts rumors, then it spreads. Graph - may not be complete.

info. cascades →
b/w friends.
(immediate)

Diff. b/w herding and cascading).

Hotmail - free mail service - "Get your free email at Hotmail" → Advertising through email itself!
↳ It became VIRAL.

$\approx/\approx \rightarrow$ directed graph. \rightarrow if undirected then assume there are 2 edges - \leftrightarrow

Decisions - binary (node)

Active node: adopted the behaviour
Inactive: has not got the info.

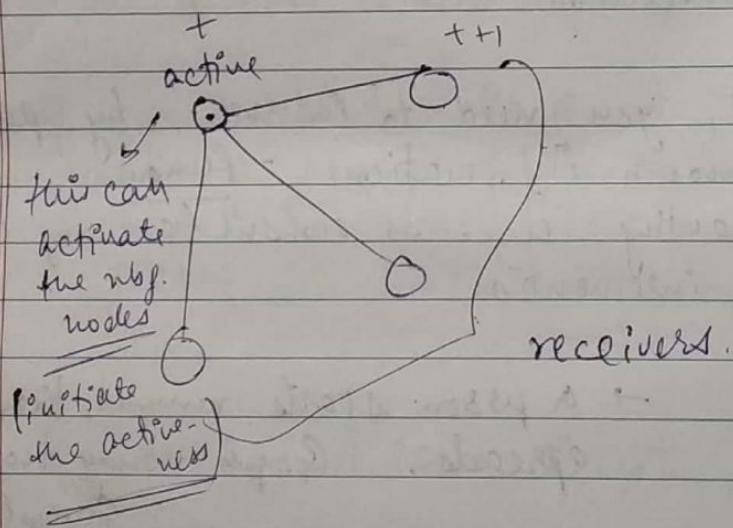
Can activate the wbf. nodes -

Activation \rightarrow inactive \rightarrow active, not vice versa

How can you make it other way also true? - Ppl are studying this

\rightarrow ICM: sender centric \rightarrow focus on senders rather than receivers.

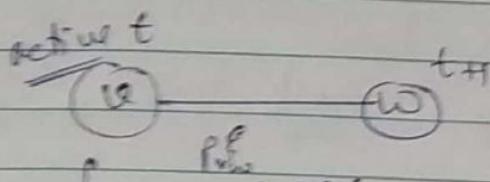
every nodes get 2 chance to activate it wbf.



Temporal \rightarrow t : 1 node active, at $t+1$

another node (3) activates

max. no. of nodes that you can reach which has a particular info.?



$P(\text{node } v \text{ activating } w)$

assumption: has only 2 chance to activate its info.

If a person is not retweeting at that time, it won't retweet that after 2 days. So, this assumption works!

Header: ICM: On an avg. what are max. no. of nodes it can initiate in a set of nodes.

- Context*: news agency \leftarrow pass info., select nodes that will help you to cascade that info.
- ppl who ~~are~~ have good reach among public
 - your reach wants to go to the max. ppl
 - marketing strategy

What if we started with $(v_2) \rightarrow ?$

Q. Set of initiators to reach max. no. of ppl with less no. of steps?

Maximizing the spread of cascades.

S: initial set

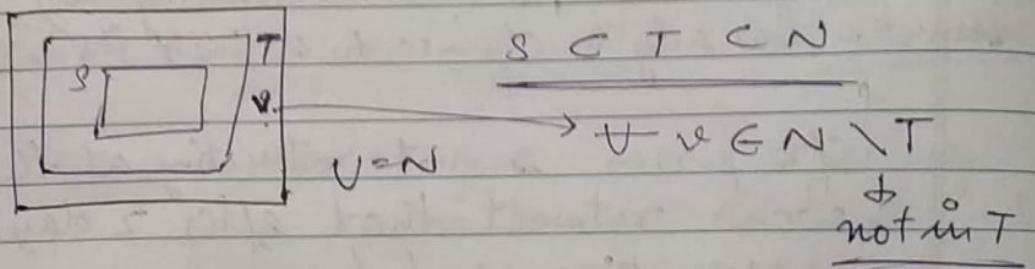
f(S): Spread

\hookrightarrow maximization problem.

$$f(s) = \{v_{1,2,4,5,6}\} \text{apsara}$$

$f(s) \rightarrow$ at any time it is going to be
(non-avg)

Submodulars:
func.



$f(s) \rightarrow$ set of all nodes that can
reach when start with s .

$$f(s+v) - f(s) \geq f(T+v) - f(T)$$

$$f(s) < f(t)$$

→ impact that is created on adding v to S
is more than adding v to T.

Greedy set $\rightarrow S$: Add e to S to max: $f(S \cup \{e\}) - f(S)$.
 ↳ 63% of opt^{imal} soln. we will be able to achieve.

Date: _____

directive: suspended from nodes
 ① Herd
 ② Info. cascading } also were observable - in public
 ③ Info. effect } - public observable in herd but not in cascading.
 ④ Date: _____

Diffusion of Innovations: (not only info. passing)

N/bs are not observable, pub. observable.
 Unlike ① & ②
why and how info spreads.

Innov. - Observable: Apple builds a new phone.

Advantage: Adv. from your competitor.

Compatible: cannot change society to completely (connec. b/w old ad new) new.

Triability: limited basis (good)

Not Complex: should be easy to use.

Segregated the adopters!

more information
 weak tie (global input) strong tie (close friend)

Adoption: external as well as internal influence
 (mass) (personal) more influence
 exogenous factors endogenous factors not information

①

Innovators: people who innovate (min. no. of rel)

Early Adopters: Fans of the product: immediately purchase it.

S-
stepped

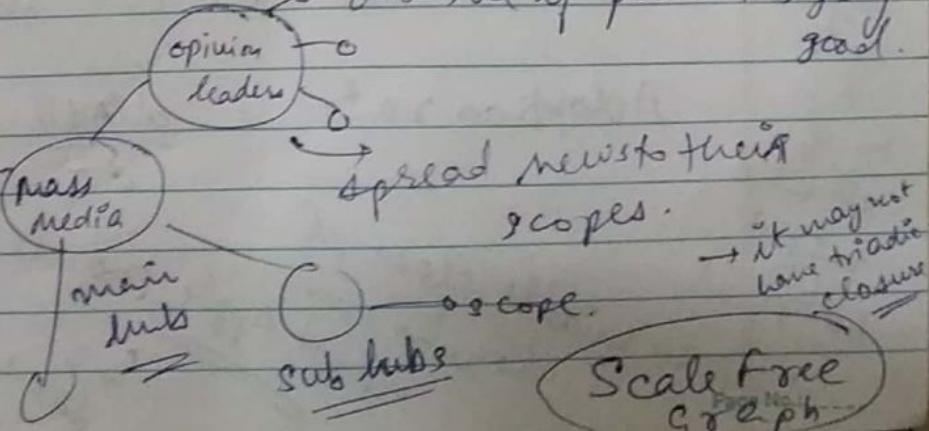
Early Majority: based on reviews - if positive

Late ..: a little late

Laggards: wait for 2-3 years, then ~~wait~~ to observe if product is going good.

② Kat 2

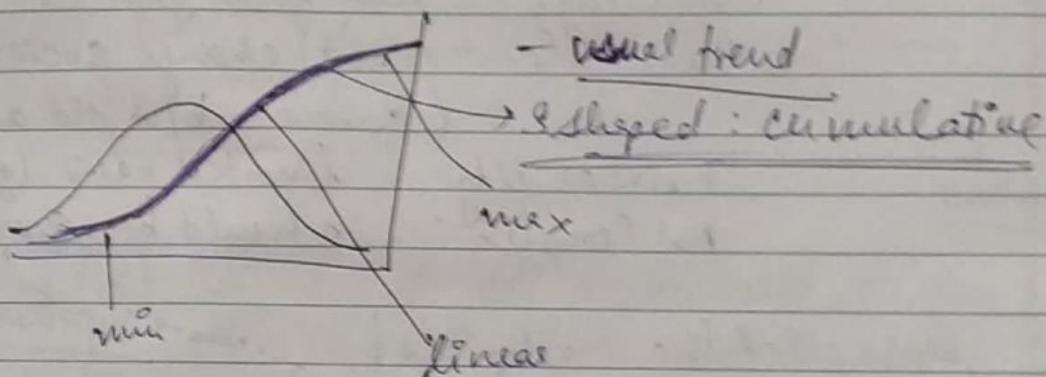
Two step



- ③ Logistics : Awareness : a new product is there
 Interest : show some interest
 Evaluation : reviews, analysis.
 Trial : if friend got it, then decide.
 Adoption : after checking.

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Marketing people → benefit from Early Adopters.



→ Modeling of Diffusion of innovation :

differential eq → to find rate of growth .

$A(t)$ → ppl : adopted the product, t : time

$i(t)$ → coeff. of diffusion : innov. rate - measure
 ↳ we can define it. (Company - good brand value → high $i(t)$) (Apple, Google)

define P → till time t → no. of ^{potential} adopters. (^{potential} adopter)

$$\text{Adoption rate} = \frac{d(A(t))}{dt} = i(t) [P - A(t)]$$

+
 no. of ppl who might adopt it
 ↳ upper middle class ppl.
 ↳ who can afford Apple products.
 ↳ Samsung might add lower class ppl also.

$i(t) \rightarrow$ how innov. the product is.
 rate affects the remaining ppl who might
 adopt the product.

$i(t) :$

External Influence Model

Internal

Mixed

$$\begin{aligned} i(t) &= \alpha \xrightarrow{\text{Innov. factor}} \\ i(t) &= \beta A(t) \xrightarrow{\text{imitation factor}} \\ i(t) &= \alpha + \beta(A(t)). \end{aligned}$$

based on product reviews - (websites) - α

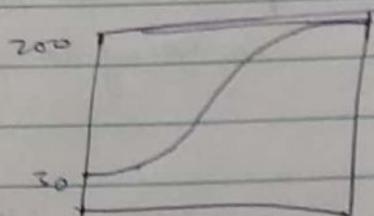
(not under any peer pressure or friends influence)

ppl who will not see any external factors, and
 rely on their friends-influence (Herd Behaviour)
 ↳ not much of rational. - $\beta(A(t))$

both: $\alpha + \beta(A(t))$

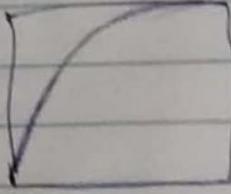
→ Internal Influence Model - Pure imitation

starting with 30 → then reaches max. - flat



→ Mixed →

most of the cases



→ Intervention : if a faulty product - address it

Reduce interest - don't advertise.

Epidemics : - not only spreading of disease
→ viral attack, worm attack
on computers

Some cascading effects without n/w observable

→ Segment the phenomenon - classify what types among these
can be combination

→ Pathogen → virus transmitted among PCs. also
→ retweeting of tweets.

→ Population - can be electronic devices also

Implicit n/w - sim. to innov. diffusion model

↓
person is not going to decide,
whether you want to carry
pathogen or not.

↓
person who is adopting
is going to decide

→ Contact n/w : connection - medium : air

→ Fully Mixed

Since

basic Epidemic Models :

→ SIR, SIRS, SIS.

→ ppl types:

- susceptible

- infected

- recovered

} based on this we have defined
4 models

① SIR → no recovery, like AIDS.

no. of susceptible individuals → normalized = $\frac{S(t)}{N}$
 .. . infected .. . → $I(t)/N$

β : Contact Probability.

if $\beta=1$, everyone contacts everyone.

if $\beta=0$ → no connections.

$N = S(t) + I(t)$. — only two sets of
ppl.

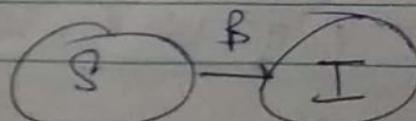
Model: Infected ~~person~~ meet βN people on avg.
and infect βS of them.

$$\frac{\text{rate of change}}{\text{dt}} = -\beta IS, \quad \frac{dI}{dt} = \beta IS$$

No. of susceptibles decreasing → they are infected now.
No. of infected will increase.

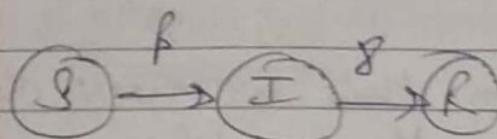
by same value,
infected ppl
will increase.

rate of infection: $I(t)$ ✓



(2) SIR model

we can be cured, and then further we will not be infected.



once you get recovered, you will hold the immune.

but immune can be time dependent.

SIRS

for sometime
(years)

$$\frac{dS}{dt} = -\beta IS$$

$$\frac{dI}{dt} = \beta IS - \gamma I$$

$$\frac{dR}{dt} = \gamma I$$

→ recovery rate.

(3) SIS model

no recovery - you become but become susceptible again.

(4) SIRS model

Intervention: - Immunisation

bulgur wheat
rat plague
goat

Mad cow disease → anyone can sell anything to anyone - animals affected

remove weak ties → Animals were killed → intervene, soil where animals were living

✓ Evolution

* Community Detection / Analysis / Evaluation:

group of ind. - having common interests or similar characteristics.

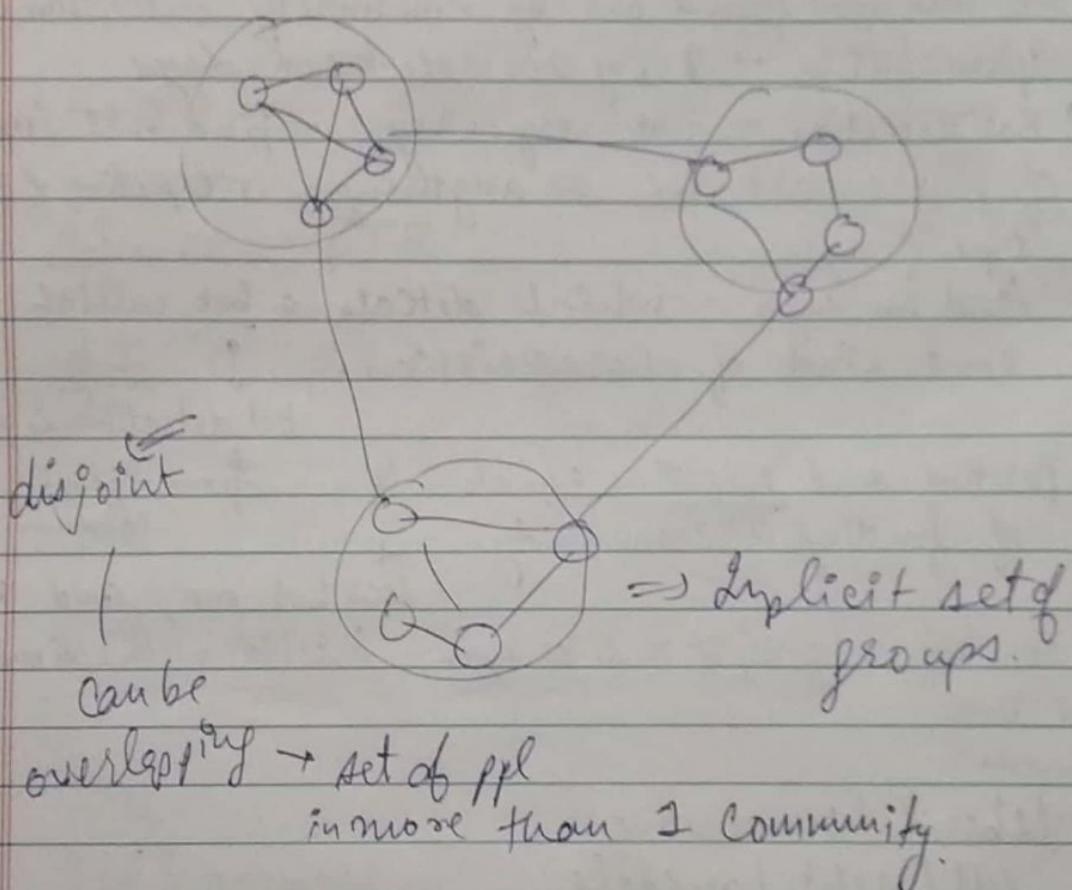
economical, social,
political, personal

- late 1980s → Dutch sys. - 30 million - sending spam mails. (billion)
 - ↳ this was found out by community detection algo.
 - Spam filters → very accurate these days.
- Bot detection - not very easy. - find bots irrespective of geographical loc. or anything, irrespective of OS, type of sys.
 - find an algo - which detects a bot which has some kind of characteristics.
 - bot detection is diff from people detection (community)
- Positive and Negative aspects of finding a community.
 - this bot one, and sys. group of ind. or sys. infected with virus.
- liberal / conservative left / right / middle.
- group atmosphere study : noise ↓ individual : noise ↑
- Recomm. sys.
- Explicit Community / Group v/s Implicit

↑ ppl who know they are forming this group, know about the members
↳ Google group, Yahoo group.

↓ finding bots, need to find characteristics.

In Social media, we will find many Implicit groups - confining with those who have similar idea only 'left right', and blocking others with diff. idea - kind of bubbles!



Amazon - recomm sys. - 1 group: action - other group: thriller, suspense \longleftrightarrow oriented these movies.

can be
overlap as well

→ Protein interaction only happens when they are similar in nature.

→ www comm. : ppl having interest in same topic

- How can you evaluate comm? - good/bad/neutral
values it is having
- Clustering method in Graphs
- Strong connec : internal in a comm.
Weak " : external ties.
- Partition the graph into various components: comm.
Ideal decomposition - completely disjoint comm.

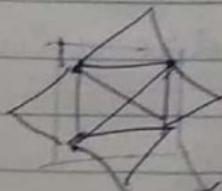
Zachary's Karate Club - from 2000, citation of this paper grew exponentially.
more: in past 2 years, if they communicated more than 10 times then they have a link b/w them.
Tussle ~~is~~ between trainer and admin.
Ppl ~~who had~~ more interaction with node 1 moved with him and others " node 34.
↳ It is implicit at beginning but it ended as explicit.

- Clustering v/s Community.
- member based comm. | group based.

1. nodes with similar char.

↳ Eg: all the nodes having k degrees - clique
↳ such graph

which is completely connected



8 nodes

clique over 4 nodes.

→ finding a clique — mostly used in SNA
 ↳ NP-hard problem

maybe we can relax it a bit,
 when $n \rightarrow$ very huge, then
 if one node is connected to $n-1$ or $n-2$... $n-k$ nodes then

weakly k-clique (k -cliques) subgraph which ↳ also fine.
 k -clubs are almost a clique → T.C.: quicker.
 k -clans because clique is NP-hard

→ similarity measures

Q) → (maximum clique, maximal clique)
 ↳ one way of finding communities.

→ maximum, maximal (N , a/b) → a div. b
 $\cancel{a/b}$

minimum = 2 → $a \neq 1$ divides everything
 maximum = ∞

($N = 813$, a/b)

Minimum

minimal → prime no.s. (all)

(2, 3, 5, 7, ...)

not comparing
these else with
other else. and

they are not minimum
because they can't be
compared.

can't find lesser else. that is
minimum All these are minimal

- minimum ele is a minimal ele.
- but minimal is not minimum
- only 1 minimum but minimal ≥ 1

apsara

v_n - starting vertex. - Brute force technique
 - it will find maximal clique but may not maximum because starting vertex is arbitrary

2^n - Complexity. - if $n \gg$ huge then high complexity
 Pruning \rightarrow clique \rightarrow nodes should have degree of $k-1$ or more.
 prune all nodes with degree $< k-1$
 To reduce no. of nodes we need to go through.

Even after pruning, complexity is undesirable

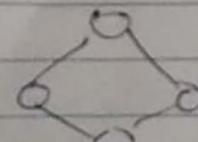
\rightarrow k -plex : set of vertices which have $d_v \geq |V|-k$
 previously $\rightarrow d \geq k-1$ but now \rightarrow
~~too nodes~~ When $k \uparrow$
 $(|V|-k) \downarrow$

$$6-1 = 5 \quad , \quad \underline{d_v \geq 5}$$

$\frac{|V|}{k} \downarrow$
 $\downarrow 1\text{-plex}$

$$2\text{-plex} \rightarrow \underline{\underline{d_v \geq 6-2 = 4}}$$

\Rightarrow k -core - degree at least k .



\Rightarrow k -shell - nodes part of k -core but not of $(k+1)$ core. \rightarrow relaxed clique

It is 2-plex as well as 2-core.

~~national open
networked
conversations~~

→ connected

→ 0-core : same graph, at least 0

0-shell :

degree

↳ isolated pt.

1-core : nodes with atleast

graph itself

1 edge → then maximal

connected.

→ k -cores of complete graph : — same

complete graph.

n nodes → $\underline{\underline{k \geq n}}$

16/4/19

k -plex → also no

holes

1-plex → complete graph.

(IV) Using clique as a seed of a core:

1st CPM: clique as a seed (smaller clique)