

$(a \vee b \vee c) \rightarrow (a \vee y) \wedge (\bar{y} \vee b \vee c)$: can't convert to 2-SAT using what we did earlier

2-SAT isn't polynomial time

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Ex HAMPATH = $\{ \langle G, s, t \rangle \mid \text{There is a hamiltonian path from } s \text{ to } t \}$

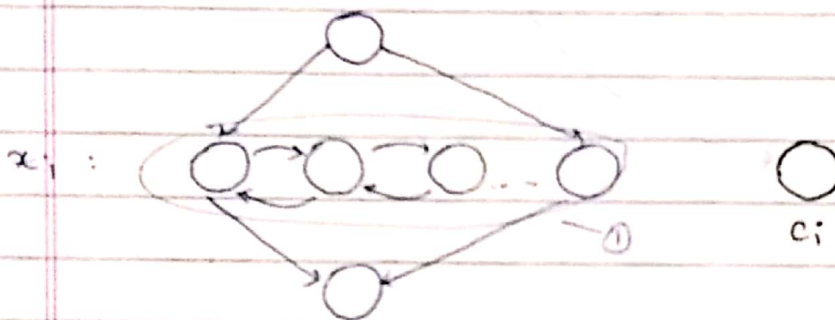
3-SAT \leq_p HAMPATH

literal (x_i, \bar{x}_i)

$$\phi = (a_1 \vee b_1 \vee c_1) \wedge (a_2 \vee b_2 \vee c_2) \dots \wedge (a_k \vee b_k \vee c_k)$$

\downarrow class C_1 \downarrow C_2 \downarrow C_k

For each literal, we'll create a graph here.



For each class, we'll add a node

Now, we'll connect the above components. (x_i and c_i)

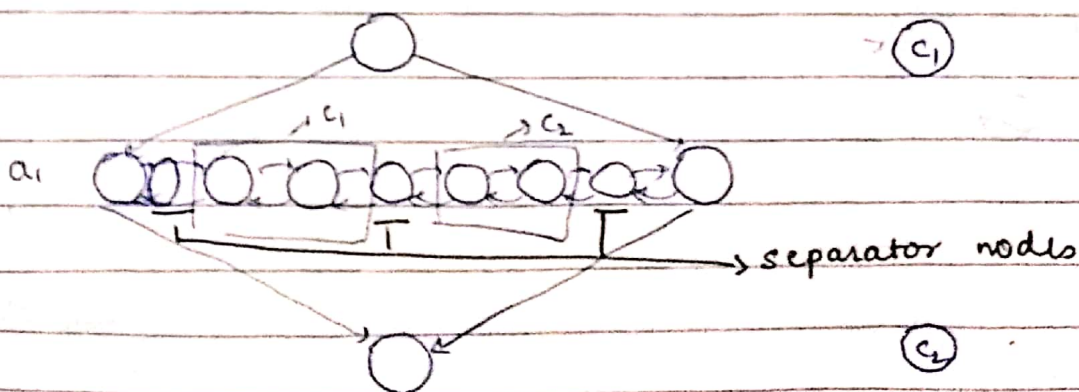
If $x_i \in C_i$, then we'll connect literal x_i to node c_i .

① For each class, we add 2 nodes separated with other classes by a node

$$(a_1 \vee b_1 \vee c_1) \wedge (a_2 \vee b_2 \vee c_2)$$

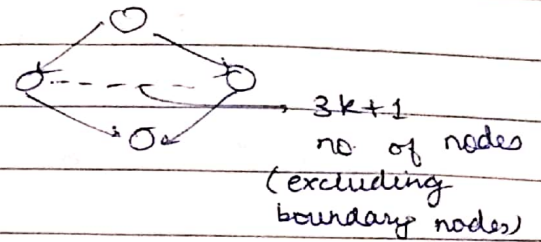
\downarrow C_1 \downarrow C_2

separator node
↓
at boundary also



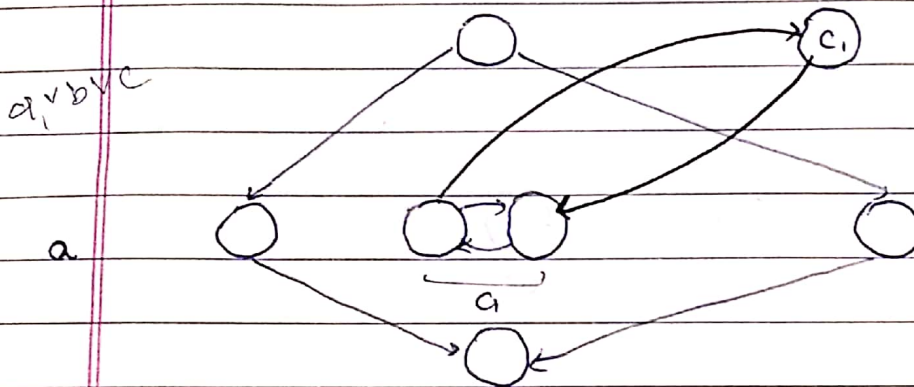
→ Will have similar structure of graph for every literal

- If we've k -classes within the 3-SAT formula, no. of vertices in b/w : $3k+1$ (nodes)



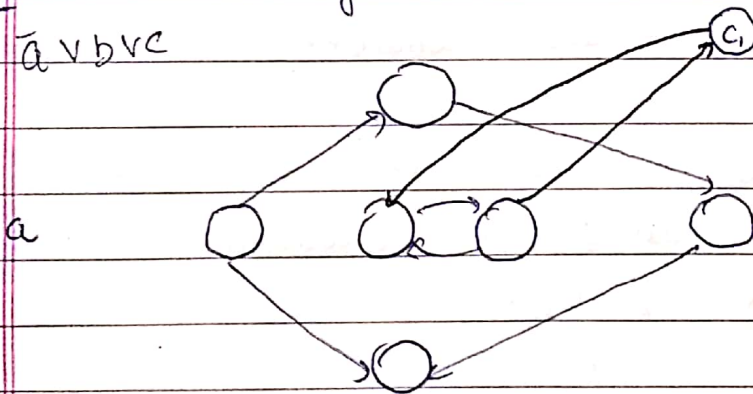
- How to connect C_i with literal?

For a particular class C_i , we've nodes inside literal graphs



- If $\bar{a} \in C_i$, we'll join it like:

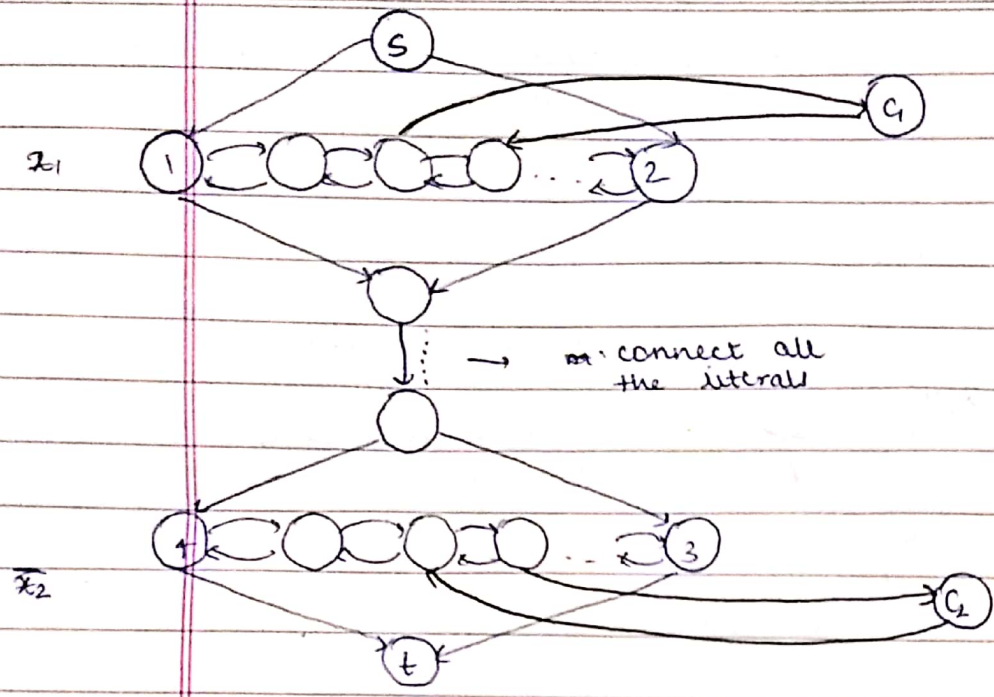
(reverse the arrow)



- If $x_i \notin C_j$: don't make any edge.

②

let $x_1 \in C_1$
 $x_2 \in C_2$



we've to get HAMPATH from this graph.

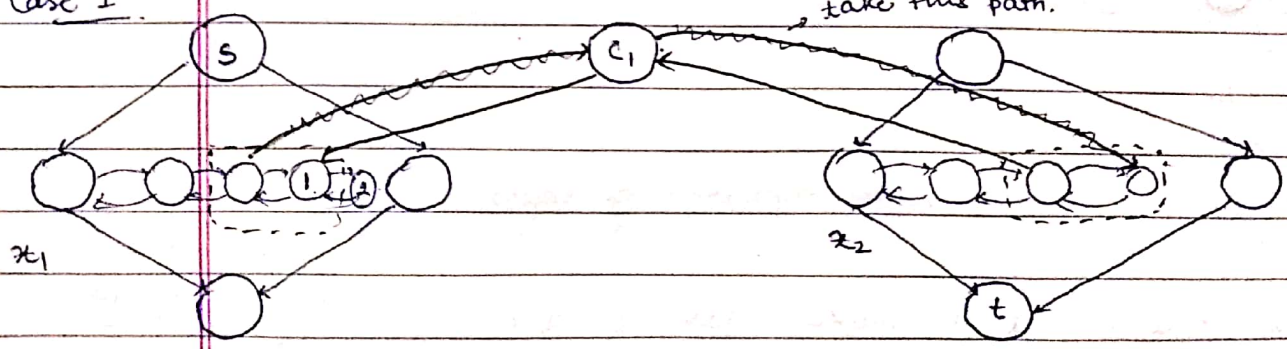
- ↓
- if $x_1 \in C_1$: go from s to 1 move inside, visit C_1 , move to C_2 & move down further. (left → right)
 - $\bar{x}_2 \notin C_2$: go from s to 3 ~ ~ ~ C_2 , move to 4 & move down further. (right → left)

problem. C_1 might be connected to other diamond structure also. we have to rule this possibility out in order to get hamiltonian path.

↓ contradiction

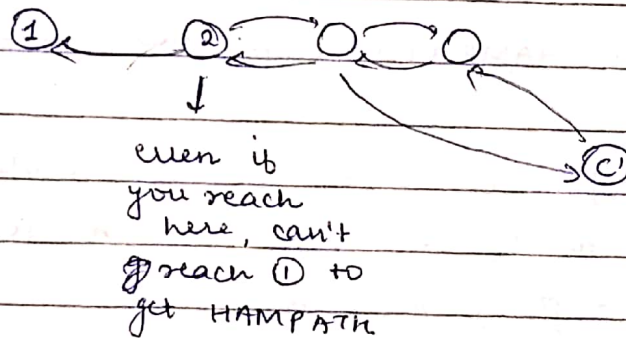
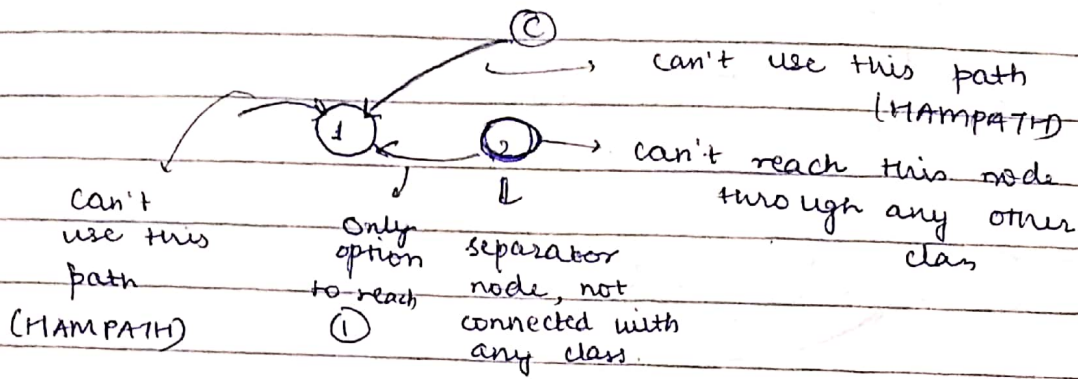
let it move to other diamond structure.

Case I

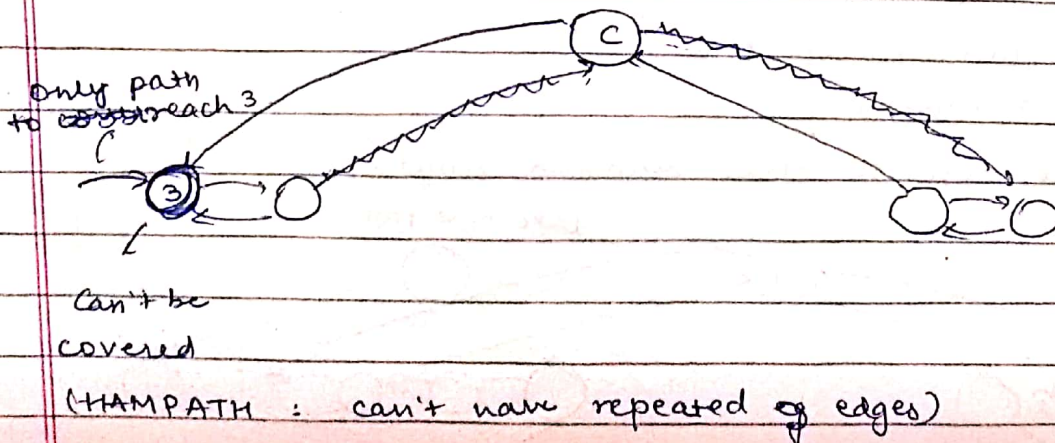


Problem here: we won't be able to cover certain vertices in both the diamond structures (either one of ① or ②)

How can we reach 1.?



Case II: Arrows to C are opposite. (complement case)
Apply same reasoning



if \mathcal{K}_1 is true: it'll make class C true.

If more than 1 literal: make any one true.

Each class is at least connected with one of diamond structures