

## DB2 12 — The ultimate enterprise database for business-critical transactions and analytics

DB2 12 for z/OS
Technical Overview
Part 1

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- The presentation is available for download from the resources list
- The webcasts will be available on replay
- You can submit questions by typing into the questions area of your webcast control
  panel at least 15 minutes before the end of the webcast.
- Questions will be answered as time permits
- Any questions not answered due to time constraints will be answered after the webcast and added to the "Resources list"



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## **Objectives**

- Introduce and discuss the new features of DB2 12 for z/OS
- Provide planning information for migration
- Understand the new performance features

## Agenda

- Introduction
- Performance focus traditional workloads
- Performance focus enabling modern applications
- Migration
- Application enablement
- Reliability, availability, scalability, security

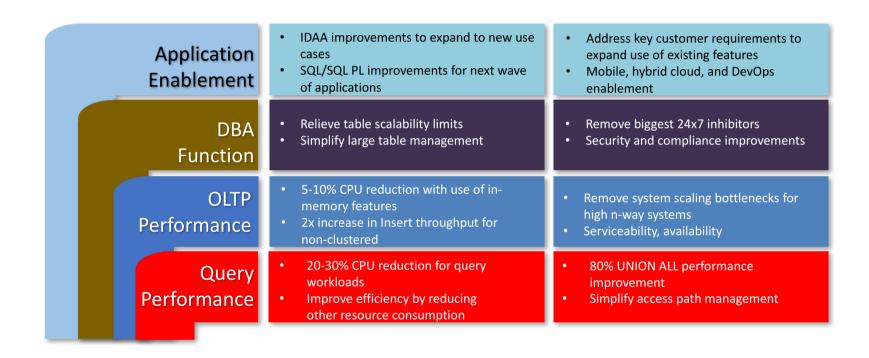
## **INTRODUCTION**

 $\ensuremath{\mathsf{DB2\ 12}}$  — The ultimate enterprise database for business-critical transactions and analytics





#### Goals



#### **Quick Hits**

- Scale and speed for the next era of mobile applications
  - Over 11 Million Inserts per second measured when hitting "sweet spot"
  - 6 trillion rows in a single table, with agile partition technology
- In Memory database
  - Up to 23% CPU reduction for index lookup with advanced in-memory techniques
- Next Generation application support
  - 360 million transactions per hour through RESTful web API
- Deliver analytical insights faster
  - Up to 25% CPU saving for traditional query workloads
  - Up to 2x speed up for modern (complex OLTP or real-time analytics) workloads

# PERFORMANCE FOCUS TRADITIONAL WORKLOADS



#### **Performance Enhancements**

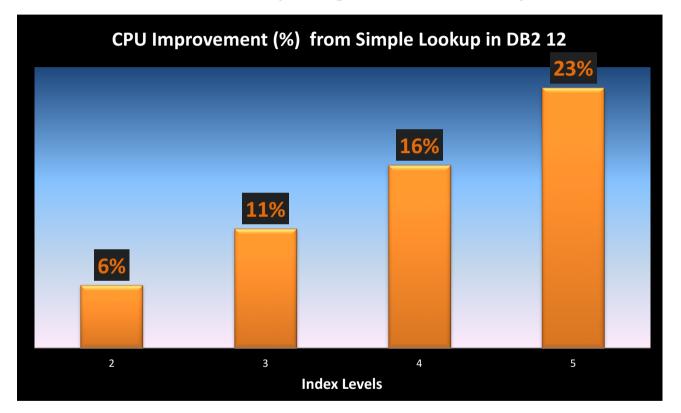
- In-memory contiguous buffer pools
  - Direct page access in-memory, greatly reduced GetPage overhead
  - Hash and LRU chains not maintained
  - Up to 8% CPU reduction measured for OLTP
  - PGSTEAL(NONE) improved in DB2 12 to avoid LRU and hash chain management overheads
  - Overflow area (10% of VPSIZE, max of 6400 buffers)
    - Used in case objects do not fit
    - Automatically managed by DB2 using FIFO steal algorithm
    - Allocated when buffer pool is allocated, but only backed when used
- In-memory index for fast traversal (see next slide)
- More granular Global Commit LSN and Read LSN
  - Potential huge improvement in lock avoidance (data sharing)
  - Help space reuse for LOB insert

#### **In-Memory Index Optimization**

- A new Index Fast Traverse Block (FTB) is introduced
  - Memory optimized structure for fast index lookups
  - Resides in memory areas outside of the buffer pool
    - New zparm INDEX MEMORY CONTROL
    - Default=AUTO (min of 500 MB or 20% of allocated buffer pool storage)
  - UNIQUE indexes only, key size 64 bytes or less
- DB2 automatically determines which indexes would benefit from FTB
- DISPLAY STATS command shows which indexes are using FTBs
- New SYSINDEXCONTROL catalog table
  - Specify time windows to control use of FTBs for an index
- New IFCIDs 389 and 477 to track FTB usage

## Simple Index Look-up: Faster & Cheaper

• Up to 23% CPU reduction for index look up using DB2 12 In-memory index tree



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#### Performance Enhancements ...

- Avoid scheduling unnecessary prefetch
  - Problem: when all pages are in memory, dynamic prefetch needlessly schedules prefetch
  - Wastes CPU, can cause "out of prefetch engine" condition
  - Attempts to solve this in the past failed solved in DB2 12
  - Up to 6.8% CPU reduction for OLTP, 4.5% for query

#### Performance Enhancements ...

- INSERT Algorithm 2 (see next slides)
- RLF control for static packages
- DB2 / DASD synergy enhancements retrofit to V10, V11
  - Exploit z/OS HyperWrite (PPRC log write accelerator)
    - Up to 30% log write latency reduction
  - Improved Integration with DS8870 Easy Tier multi-temperature management
    - Avoid I/O degradation after REORG

#### **INSERT Algorithm 2 Performance**

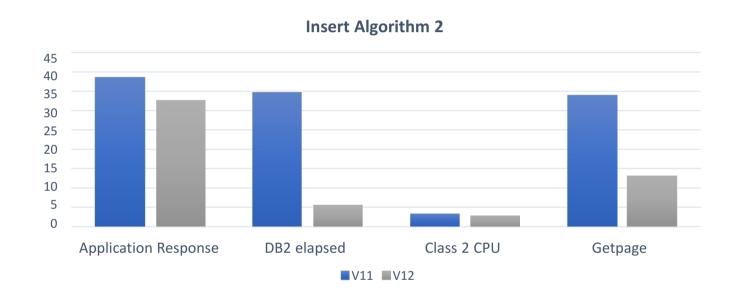
- Insert workloads are amongst the most prevalent and performance critical
- Performance bottleneck will vary across different insert workloads
  - Index maintenance?
  - Log write I/O?
  - Space search (page p-lock, page latch contention)
  - PPRC disk mirroring
  - Network latency
  - etc



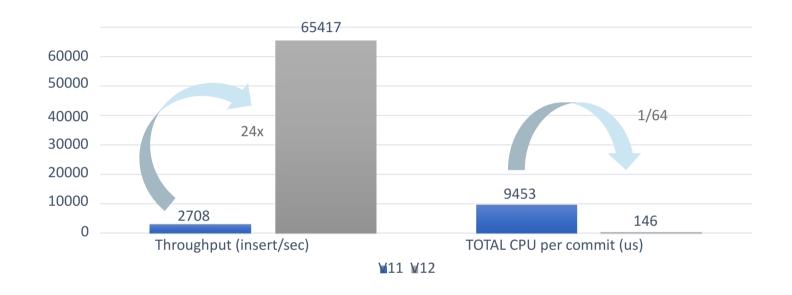
## **INSERT Algorithm 2 Performance ...**

- DB2 12 may potentially deliver significant improvement for non-clustered insert (e.g., journal table pattern where both <u>concurrent</u> insert activity and <u>space search is the constraint</u> on overall insert throughput
- Applies to UTS with MEMBER CLUSTER (both with/without APPEND)
- Implemented advanced new insert algorithm to streamline space search
  - Default is to use the new fast algorithm for qualifying table spaces
    - INSERT ALGORITHM zparm can change the default
    - INSERT ALGORITHM table space attribute can override zparm
- Your mileage will vary
  - Some insert workloads will see no improvement
  - Some specific insert workloads may see significant improvement
- Will shift the bottleneck to the next constraining factor

## **INSERT Algorithm 2 Performance – Shifting The Bottleneck ...**



#### **INSERT Algorithm 2 Performance - DB2 11 PMR Recreate ...**



UTS PBG with Member Cluster, RLL, with 400 bytes per row, one index, 800 concurrent threads, 10 insert per commit

#### Performance Enhancements ...

- Buffer Pool advisory mode to simulate larger buffer pools rollback to V11
  - Low CPU and real memory overhead
  - Statistics provided to indicate I/O savings
  - Retrofit to V11
- Streamlined Claim/Declaim processing
  - Avoid re-claim overhead across multiple commit scopes serially reusing persistent thread running RELEASE(DEALLOCATE)
  - Online REORG and other drainers can still break in

#### Other Performance Enhancements ...

- Improvements in pool management to simplify and remove the scalability inhibitors
  - EDM pools
  - LOB/XML storage
- Other Improvements
  - Reducing DGTT DECLARE overhead
  - Remove log force write from identity column and sequence caching in data sharing



#### **System Scaling Enhancements**

- Large n-way scaling
  - Improved efficiency on LPARs with high number of CPs
  - Log latch contention reduction: testing shows up to 41% CPU reduction and 6% throughput improvement for high contention cases
  - Buffer Pool scaling improvements:
    - LC23 reduction, PLO avoidance
    - 5-30% CPU improvement when accessing hot pages
  - IRLM latch contention reduction
  - EDM DBD and skeleton pool scalability improvements
- Optimizations for new hardware
  - Exploitation of z13 decompression enhancement
  - Internal structure changes for cache efficiency, more processor prefetch
- Raise total buffer pool size limit to 16 TB
- >4G size active log datasets

## High level performance expectation

#### System and OLTP performance

- 2-3% CPU reduction without Index In-Memory feature
- 5-10% CPU reduction by exploiting Index In-Memory feature
- Further reduction is possible with contiguous buffer pools, and / or persistent RELEASE(DEALLOCATE)



- Wide range of improvement
  - Typically 0-20% without new access path
  - Typically 10-40% with new access path
  - Up to 90% reduction is observed in our evaluations

#### Concurrent insert against UTS and MEMBER CLUSTER

- 5-10% CPU reduction
- Throughput improvement if current bottleneck is space search or page contentions



#### **Instrumentation Enhancements**

- More granular wait times for IFCIDs 316 (dynamic) and 401 (static)
  - Accumulated wait time due to global contention for locks (broken out by type)
- Enhance IFCIDs 53/58 statement level section for PREPARE
  - Similar to INSERT/UPDATE/DELETE
- Enhance SQL performance tracing adding RDI Section Number in IFCIDs 53/58
- Add batch job STEP name in correlation header
- Add REFRESH TABLE to counts in DSNDQXST (SQL data section)
- Add workfile, tempfile usage information to Accounting trace
- Enhanced precision for IFCID 199 dataset I/O stats microseconds

#### Access path (plan) stability

- Dynamic SQL plan stability
- Static plan stability usability
- Preserve local dynamic statement cache at rollback
- Integrated RUNSTATS with optimizer
  - Automated update of statistics profiles by optimizer
- Statistics profile support
  - Automated update for Index DDL
  - Inline stats usage of profiles
- Simplify creation of all tables required by Explain
  - New ADMIN\_EXPLAIN\_MAINT stored procedure

#### **Dynamic SQL Plan Stability**

#### • Problem:

- Unstable performance of repeating dynamic SQL statements
- Environmental changes can result in change in access path or performance regression, and this can be tough to manage
  - RUNSTATS
  - Applying software maintenance
  - DB2 release migration
  - zparm changes
  - Schema changes
- Static SQL has several advantages
  - Access path established at BIND time
  - Static plan management gives advanced management functions
- Objective: extend static SQL advantages to dynamic SQL

#### **Dynamic SQL Plan Stability ...**

- **Base infrastructure** 
  - Opaque parameter CACHEDYN STABILIZATION
  - Capture via -START DYNQUERYCAPTURE STBLGRP(APP1) ...
    - Command with / without monitoring
    - Global variable
  - FRFF
  - EXPLAIN (current, invalid)
  - Invalidation
  - LASTUSED (identify stale statements)
  - Instrumentation (query hash, explain, cache + catalog hit ratio)
  - APPLCOMPAT is part of matching criteria
- Save guery and cache structures to catalog for subsequent reuse on cache mis



## **Dynamic SQL Plan Stability ...**

- Key limitations
  - Literal concentration and temporal stabilization not currently included
  - REBIND support not included
    - No PLANMGMT/SWITCH/APREUSE

#### **Static Plan Stability: Usability**

- **BIND REPLACE of same version to keep LASTUSED**
- **FREE PACKAGE improvements** 
  - Selectively FREE either original or previous
  - Choose to FREE if invalid only
  - Allow FREE of inactive package copies while application is running
- **REBIND PACKAGE improvements** 
  - SWITCH to include APREUSESOURCE option so as to choose ORIGINAL or PREVIOUS copy as source for APREUSE
  - Addresses issue of switch accidentally causing invalid copy to become current
  - APREUSESOURCE option avoids the two step process of
    - REBIND SWITCH followed by
    - REBIND APREUSE

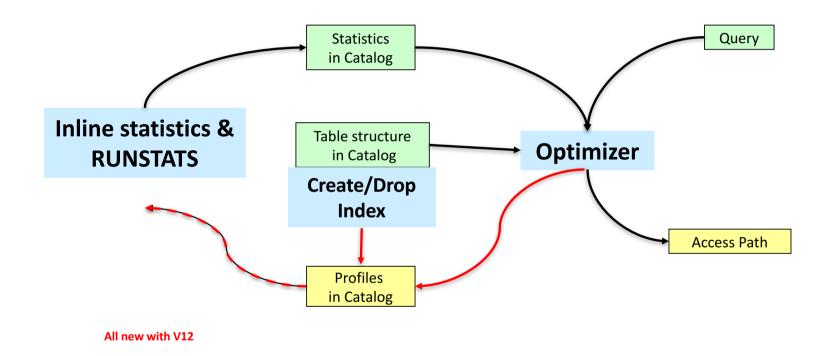
transactions and analytics



#### **RUNSTATS Enhancements for SQL Performance**

- RUNSTATS (for access path selection benefit)
  - CLUSTERRATIO formula improvements
  - Statement cache invalidation
    - Optional for RUNSTATS (new default)
    - For other utilities ONLY if objects in pending state before utility executed
  - Profile support for inline statistics
  - Automated COUNT for FREQVAL
    - Allow DB2 to collect the skewed values
      - Up to top 100 or until no skew for remaining values
  - Optimizer to automatically update statistics PROFILE with RUNSTATS recommendations
    - DSNACCOX to recommend RUNSTATS after profile update
      - Specify USE PROFILE on RUNSTATS to collect current statistics recommendations

## **Enhanced statistics profile management**



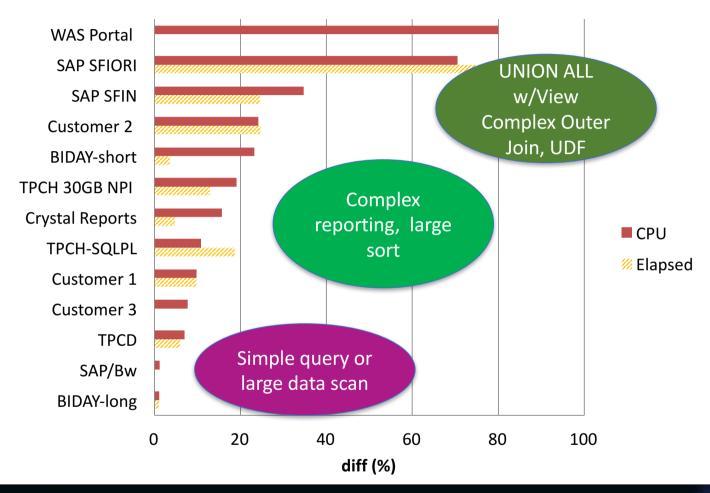
# PERFORMANCE FOCUS ENABLING NEW APPLICATIONS



## **Query Performance Enhancements**

- Up to 25% CPU improvement for traditional query workloads
- Up to 2x improvement for modern SQL applications
  - Performance improvements for next generation SAP applications
  - Real-time analytics, complex OLTP workloads
- 100% zIIP eligibility for parallel query child tasks
- Modern applications contain more complex SQL patterns (targeted in DB2 12), and more sorting, joins, stage 2 predicates etc.
  - These complex patterns are less common in traditional OLTP/batch

## Query Workload CPU and Elapsed time Improvement from V11 (%)



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## **High-level Performance Focus**

- Query (RDS) focus based upon new(er) workloads
  - Complex views or table UDFs
    - UNION ALL
    - Outer joins
    - Join predicates with (Stage 2) expressions
  - CASE expressions, CAST functions, scalar functions
- Query General Bottlenecks
  - Sort/workfile reductions
  - Reducing prepare cost and frequency
  - I/O performance
    - Reduce unnecessary prefetch scheduling



#### **Query Performance Focus**

- Improve performance of
  - UNION ALL and outer join performance enhancements
    - Performance issues are similar with both types of guery pattern
    - Reduce materializations
      - Bypass workfile usage when materialization required
      - Trim unnecessary columns from materializations
    - Push predicates inside UNION ALL legs or OUTER JOIN query blocks
    - Push ORDER BY and FETCH FIRST into UNION ALL legs
    - Reorder OUTER JOIN tables to avoid materializations
  - Table UDFs
    - Improve merge similar to views
    - Indexability of join predicates

#### **Query Performance Focus ...**

- Improve performance of ...
  - Join predicates with Stage 2 expressions
  - Stage 2 to indexable
    - VARBINARY, COLLATION KEY IOE
  - Expression evaluation
    - CASE, SUBSTR, etc
    - Expression sharing (SELECT list only)
    - Caching deterministic UDF results
  - Enabling parallelism
    - 100% zIIP offload for parallel child tasks
    - Reduce cost and resource consumption
  - Sort
    - Reduce workfile usage for GROUP BY/DISTINCT
    - Reduce key length for GROUP BY/DISTINCT and sparse index
    - Continued progress towards in-memory for smaller sorts (begun in V9)

#### **Query Performance Focus ...**

#### **Adaptive index**

Simple example of targeted use case SELECT \* FROM TAB1 WHERE COL1 < ? AND COL2 < ? AND COL3 < ?;

```
INDEXES: IX1 (col1), IX2 (col2), IX3 (col3)
```

- Filtering of the above query is dependent on literals at execution
  - Common pattern for search screens (with BETWEENs or LIKEs)
  - Common that 1 index is filtering unless a highly skewed value is searched

## **Query Performance Focus ...**

- Adaptive index ...
  - Allow RID based plans (single index list PF or multi-index) to quickly determine filtering from index
    - Without requiring REOPT(ALWAYS)
    - For list prefetch or multi-index ORing
      - Earlier opportunity to fallback to tablespace scan if large % of table to be read
    - For multi-index ANDing
      - Reorder index legs from most to least filtering
      - Early-out for non-filtering legs, and fallback to rscan if no filtering
  - Optimizer to use uncertainty to determine risk of a single index plan
    - Quick evaluation done based upon literals used
    - Any further evaluation of filtering deferred until after 1 RID block retrieved
      - Ensuring that very short running queries do not incur overhead

## **MIGRATION**

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## **Migration Prerequisites – Hardware & Operating System**

#### • Processor requirements:

- z196, or higher, processors running z/OS V2.1, or later
- DB2 12 will probably require increased real storage for a workload compared to DB2 11 for z/OS

#### • Software Requirements:

- z/OS V2.1 Base Services, (5650-ZOS), or later
- DFSMS V2.1, or later
- Language Environment Base Services
- z/OS V2.1 Security Server (RACF), or later
- IRLM Version 2 Release 3 (Delivered with DB2 12)

#### Additional details:

http://www.ibm.com/common/ssi/rep\_ca/1/897/ENUS215-371/ENUS215-371.PDF

#### **Migration & Catalog**

- Single phase migration process
  - No ENFM phase
  - New function activated through new command
    - -ACTIVATE FUNCTION LEVEL with "V12R1Mnnn" where nnn >= 500
  - APPLCOMPAT rules, fallback rules continue to apply
- BSDS conversion to support 10 byte log RBA is pre-requisite
- No pre-V10 bound packages
  - Get rid of 31-bit runtime, some performance improvements
- BRF is deprecated
  - BRF page sets still supported, but zparm and REORG options are removed
- Temporal RTS tables
  - Defined in catalog, enablement is optional



#### **Online Migration Improvements**

- Pause statistics externalization during migration to reduce contention (part of fallback SPE)
  - Real time statistics
  - Stats feedback
  - SYSPACKAGE.LASTUSED
- Catalog/directory lock duration reduction APARs (impacts online migration, catalog REORGs)
  - PI43662 Accelerator resources released in a timely manner
  - PI43916 DB2 plan allocation locks released in timely manner
  - PI39053 Avoid locks from SET statement
  - PI40755 Dynamic SQL release prepare locks earlier

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