

# The Fractional Knapsack Problem

- Given: A set  $S$  of  $n$  items, with each item  $i$  having
  - $p_i$  - a positive profit
  - $w_i$  - a positive weight
- Goal: Choose items, allowing fractional amounts( $x_i$ ), to maximize total profit but with weight at most  $m$ .

$$\text{maximize } \sum_{1 \leq i \leq n} p_i x_i$$

$$\text{subjected to } \sum_{1 \leq i \leq n} w_i x_i \leq m$$

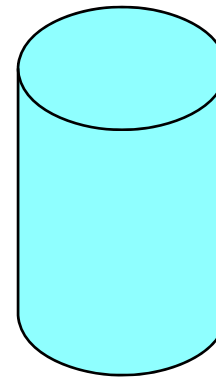
$$\text{and } 0 \leq x_i \leq 1, \quad 1 \leq i \leq n$$

## Example-4.4: Sahni (Page-218)

Greedy decision property:-

Select items in **decreasing** order of **profit/weight**.

Items:	1	2	3
$w_i$ :	18	15	10
$p_i$ :	25	24	15
Value: ( $p_i / w_i$ )	1.39	1.6	1.5



Knapsack = 20

Solution:

- 15 of  $i_2$
- 5 of  $i_3$
- 0 of  $i_1$

- Solution vector

$$(x_1, x_2, x_3) = (0, 1, 1/2)$$

- Profit  $= 25 * 0 + 24 * 1 + 15 * 1/2$   
 $= 0 + 24 + 7.5$   
 $= 31.5$

## Algorithm-4.3: Sahni (Page-220)

Greedy algorithm for the fractional Knapsack problem

Algorithm GreedyKnapsack(m,n)

//P[1:n] and w[1:n] contain the profits and weights

// respectively of the n objects ordered such that

// $p[i]/w[i] \geq p[i+1]/w[i+1]$ .

//m is the knapsack size and x[1:n] is the solution

// Vector.

```
{
    for i=1 to n do x[i]=0; // Initialize x.
    U=m;
    for i=1 to n do
    {
        if ( w[i]>U ) then break;
        x[i]=1; U=U-w[i];
    }
    if ( i <=n) then x[i]= U/w[i];
}
```

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Thank You

Stay Safe