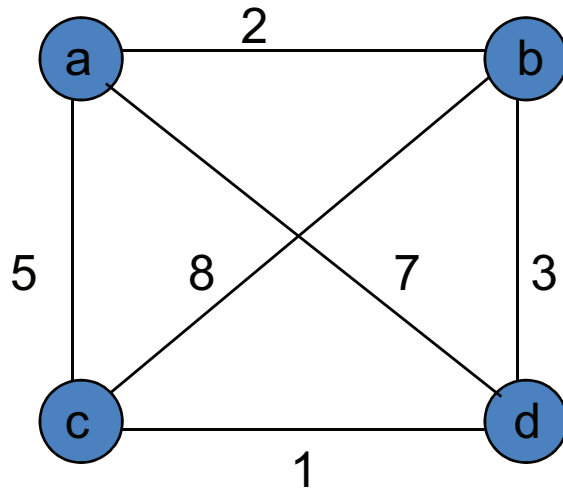


# Branch and Bound

## Traveling salesman problem

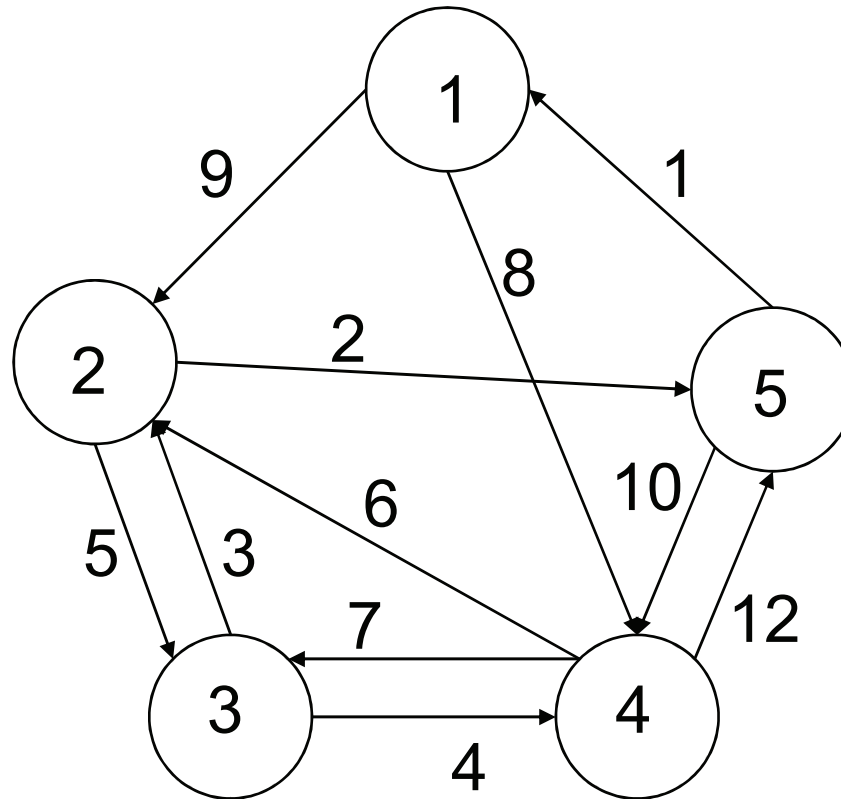


$\{ a, b, c, d \}$   
represents 4 cities

The weights  
represent distances  
between cities

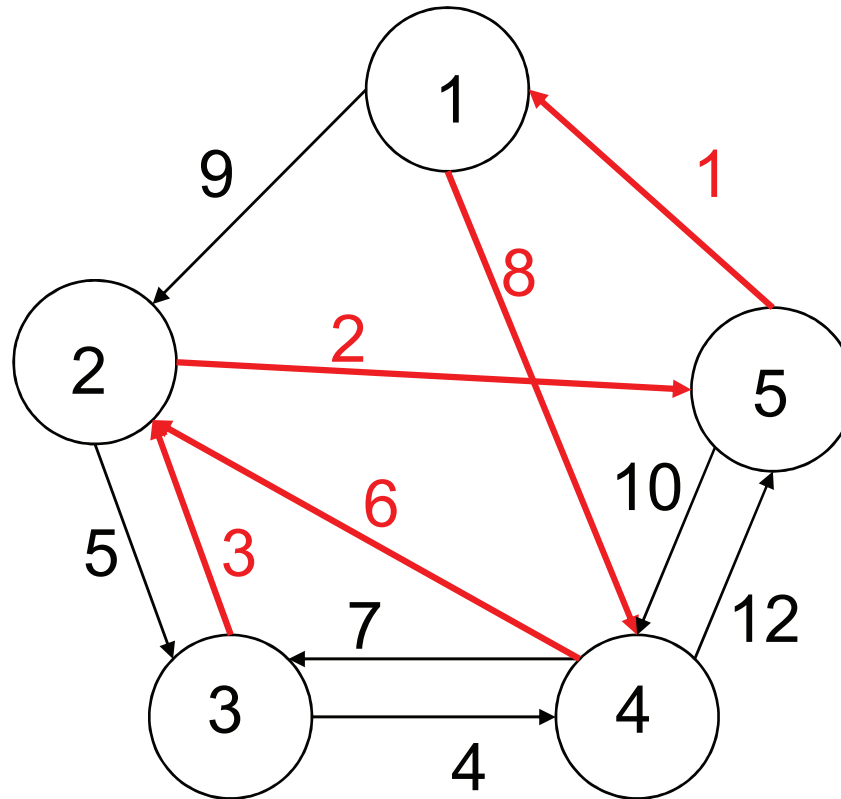
**Problem: find the shortest path from a city (say  $a$ ), visit all other cities exactly once, and return to the city where it started (city  $a$ ).**

# Bound on TSP Tour



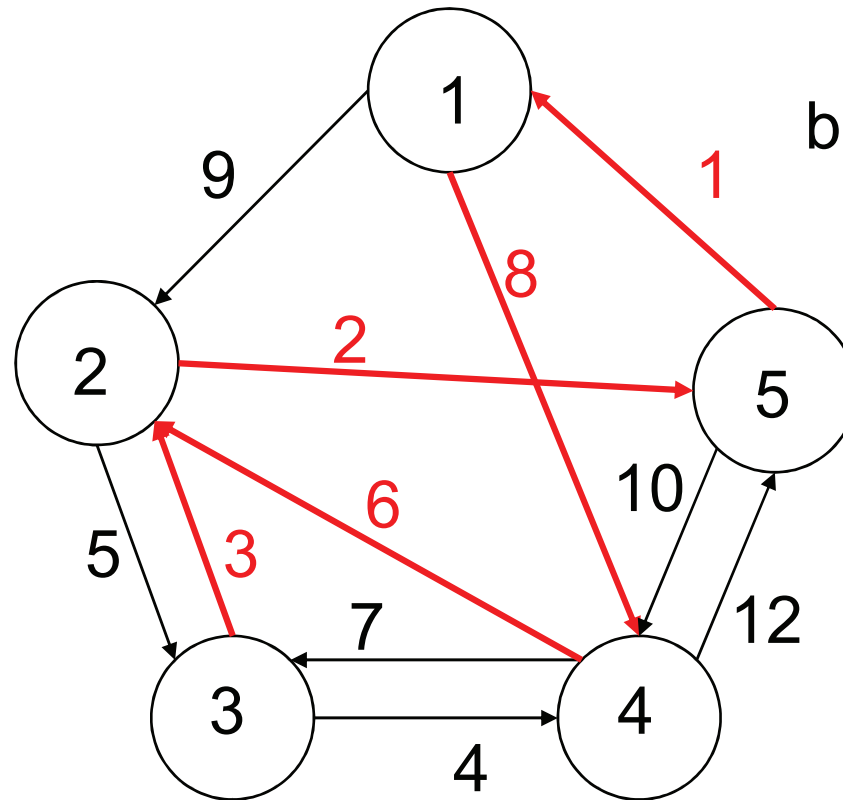
Every tour must leave every vertex and  
and arrive at every vertex.

# Bound on TSP Tour



What's the cheapest way to leave each vertex?

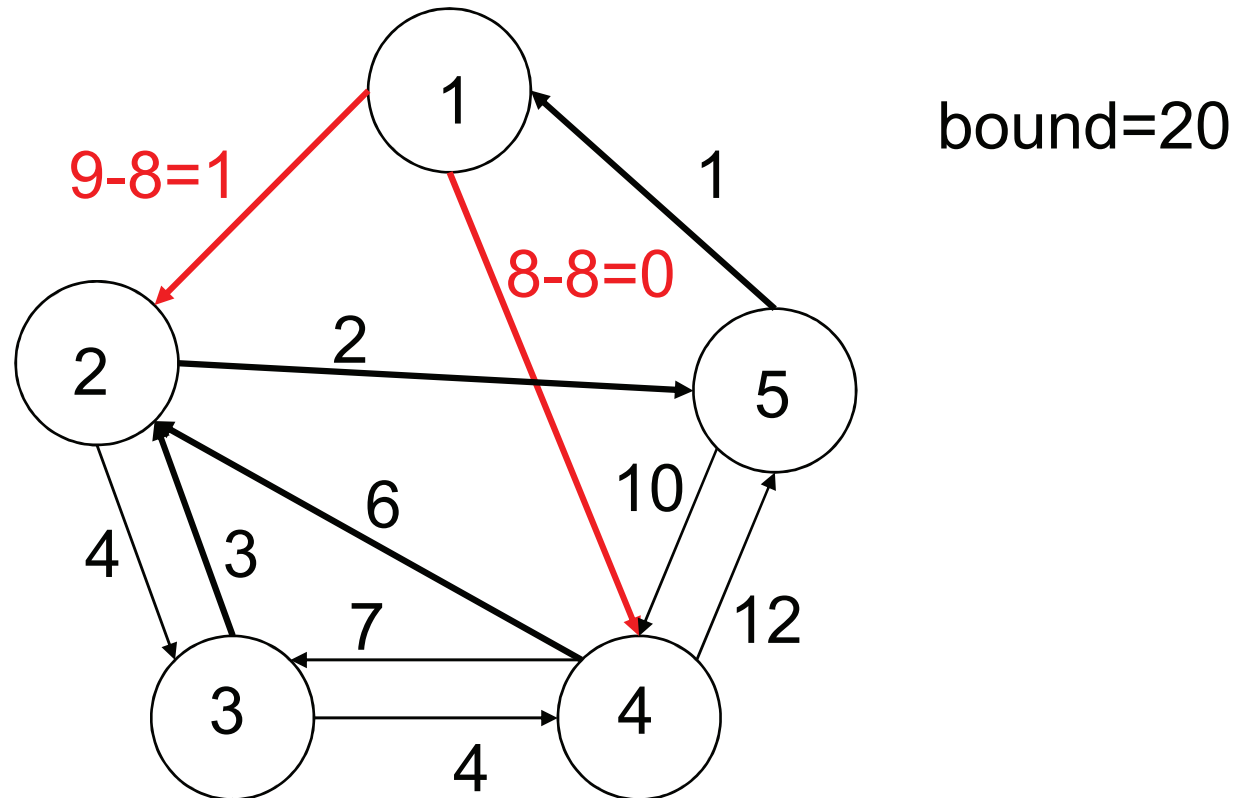
# Bound on TSP Tour



*rough draft*  
 $\text{bound} = 8 + 6 + 3 + 2 + 1$   
 $= 20$

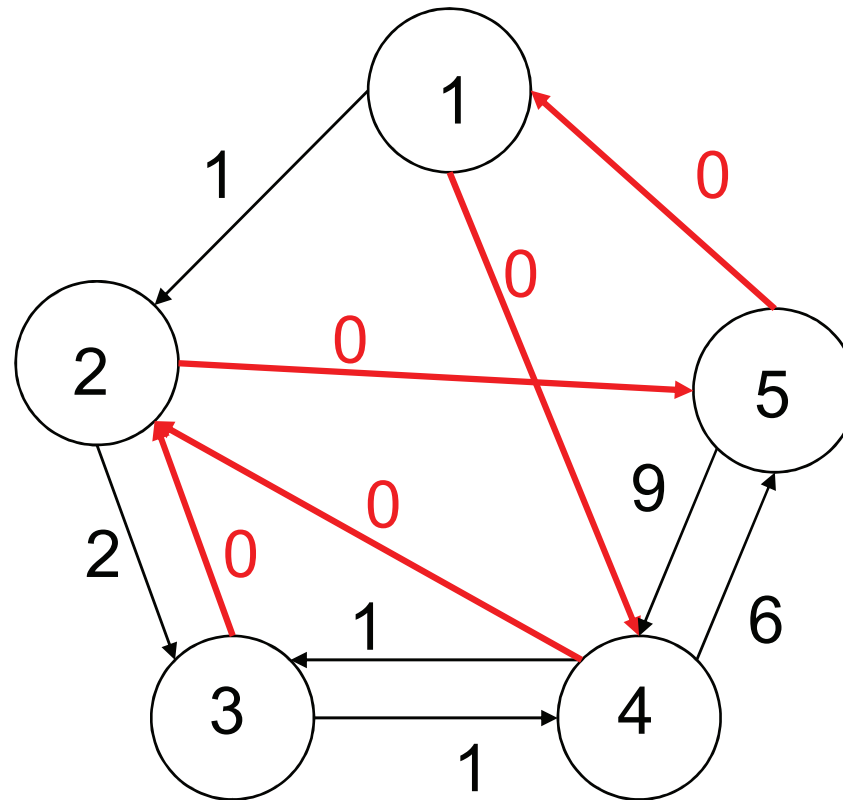
Save the sum of those costs in the bound (as a rough draft).  
Can we find a tighter lower bound?

# Bound on TSP Tour



For a given vertex, subtract the least cost departure from each edge leaving that vertex.

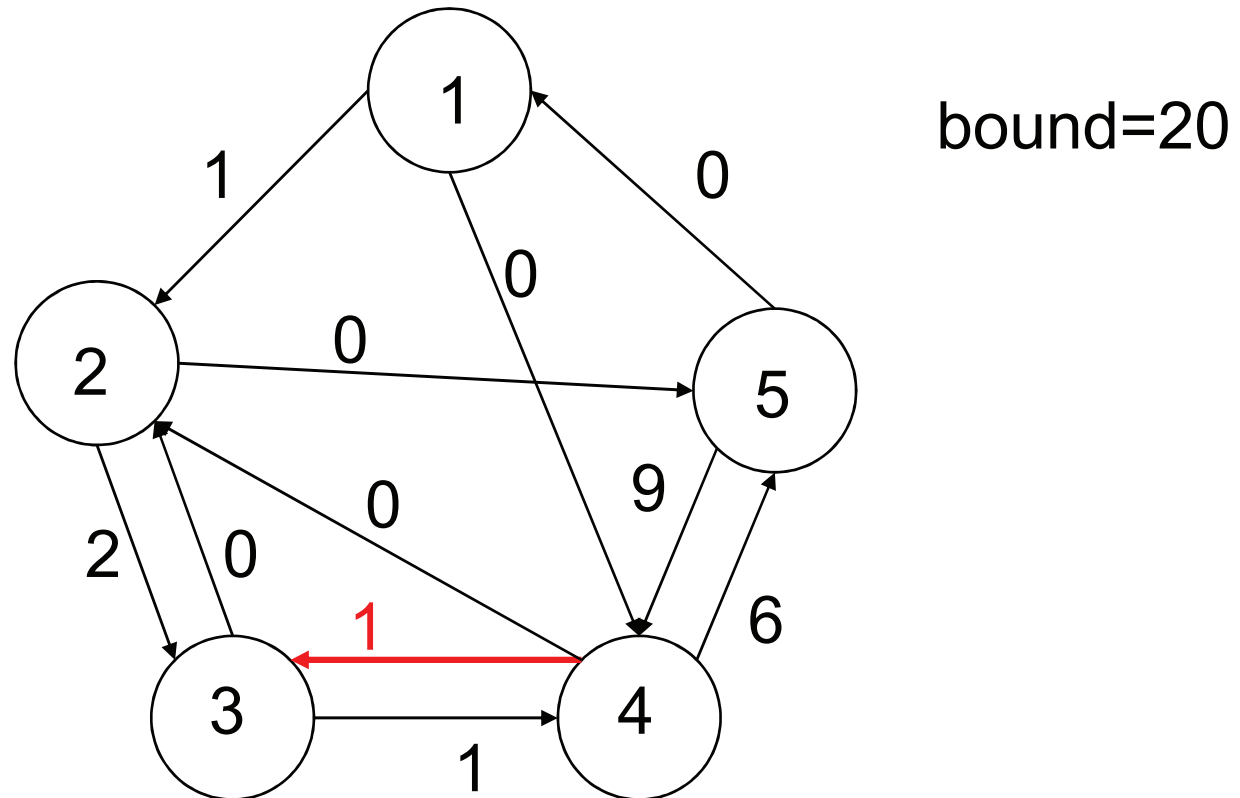
# Bound on TSP Tour



bound=20

Repeat for the other vertices.

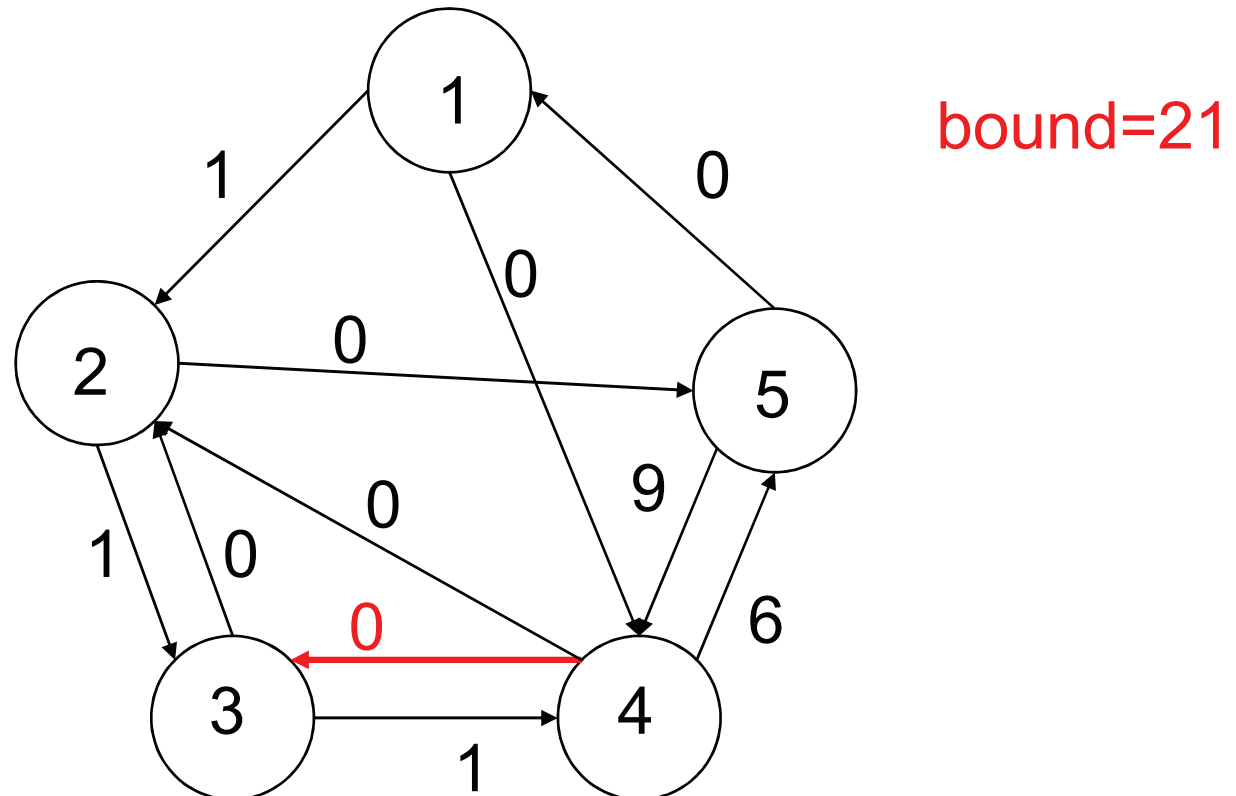
# Bound on TSP Tour



Does that set of edges now having 0 residual cost arrive at every vertex?

In this case, the edges never arrive at vertex 3.

# Bound on TSP Tour



We have to take an edge to vertex 3 from somewhere. Assume we take the cheapest. Subtract its cost from other edges entering vertex 3 and add the cost to the bound.

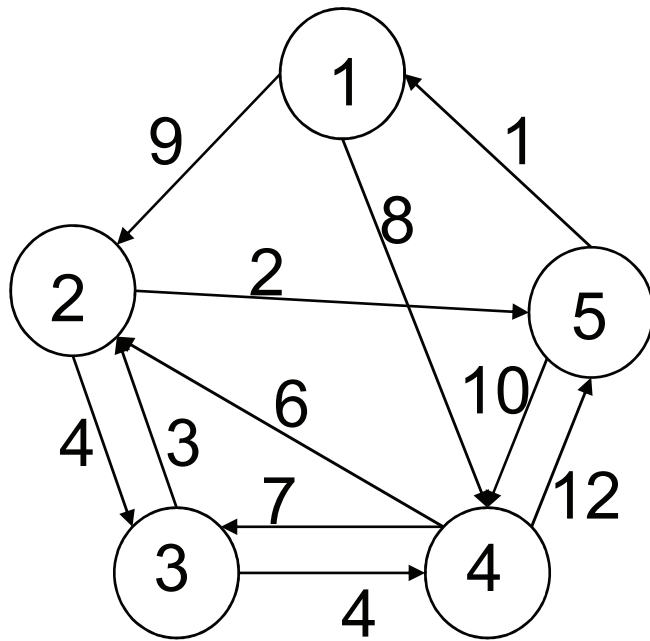
**We have just tightened the bound.**



# The Bound

- It will cost at least this much to visit all the vertices in the graph.
  - there's no cheaper way to get in and out of each vertex.
  - the edges are now labeled with the *extra* cost of choosing another edge.

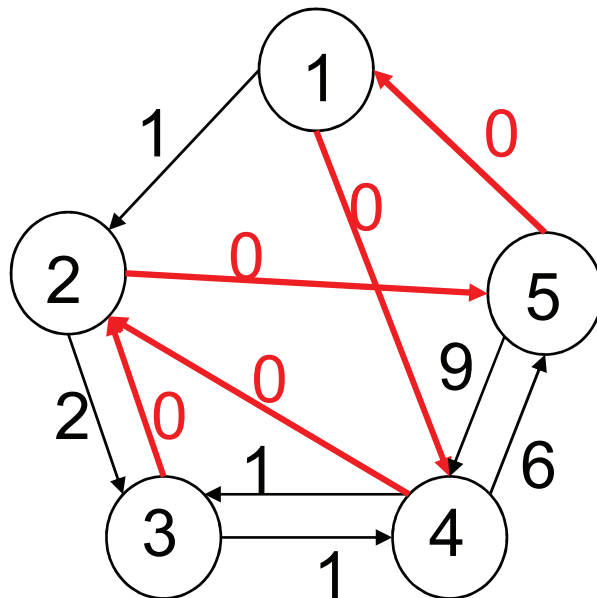
# Bound on TSP Tour



999	9	999	8	999
999	999	4	999	2
999	3	999	4	999
999	6	7	999	12
1	999	999	10	999

Algorithms do this using a cost matrix.

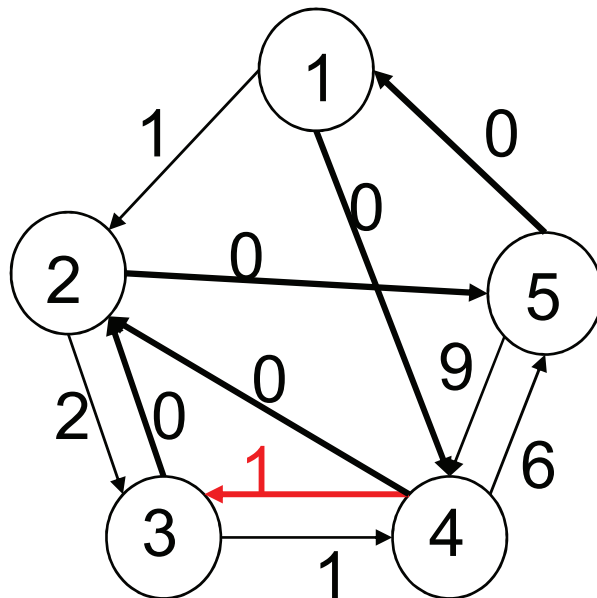
# Bound on TSP Tour



999	1	999	0	999	8
999	999	2	999	0	2
999	0	999	1	999	3
999	0	1	999	6	6
0	999	999	9	999	1
					<hr/>
					20

Reduce all rows.

# Bound on TSP Tour



$$\begin{pmatrix} 999 & 1 & 999 & 0 & 999 \\ 999 & 999 & 1 & 999 & 0 \\ 999 & 0 & 999 & 1 & 999 \\ 999 & 0 & 0 & 999 & 6 \\ 0 & 999 & 999 & 9 & 999 \end{pmatrix}$$

bound:  $20 + 1 = 21$

Then reduce column #3. Now we have a tight bound.

# Using this bound for TSP in B&B

start at ~~node 1~~ <sup>vertex</sup> in graph (arbitrary)

999	1	999	0	999
999	999	1	999	0
999	0	999	1	999
999	0	0	999	6
0	999	999	9	999

bound = 21

1to2

1to3

1to4

1to5

999	1	999	0	999
999	999	1	999	0
999	0	999	1	999
999	0	0	999	6
0	999	999	9	999

bound = 21+1

999	1	999	0	999
999	999	1	999	0
999	0	999	1	999
999	0	0	999	6
0	999	999	9	999

infeasible

999	1	999	0	999
999	999	1	999	0
999	0	999	1	999
999	0	0	999	6
0	999	999	9	999

bound = 21

999	1	999	0	999
999	999	1	999	0
999	0	999	1	999
999	0	0	999	6
0	999	999	9	999

infeasible

# Using this bound for TSP in B&B

start at ~~node 1~~ <sup>vertex</sup> in graph (arbitrary)

999	1	999	0	999
999	999	1	999	0
999	0	999	1	999
999	0	0	999	6
0	999	999	9	999

bound = 21

1to2

1to3

1to4

1to5

999	1	999	0	999
999	999	1	999	0
999	0	999	1	999
999	0	0	999	6
0	999	999	9	999

bound = 21+1

999	1	999	0	999
999	999	1	999	0
999	0	999	1	999
999	0	0	999	6
0	999	999	9	999

infeasible

999	1	999	0	999
999	999	1	999	0
999	0	999	1	999
999	0	0	999	6
0	999	999	9	999

bound = 21

999	1	999	0	999
999	999	1	999	0
999	0	999	1	999
999	0	0	999	6
0	999	999	9	999

infeasible

# Focus: going from 1 to 2

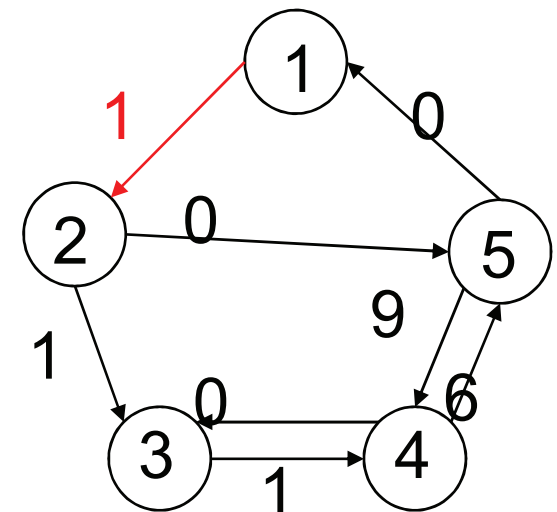
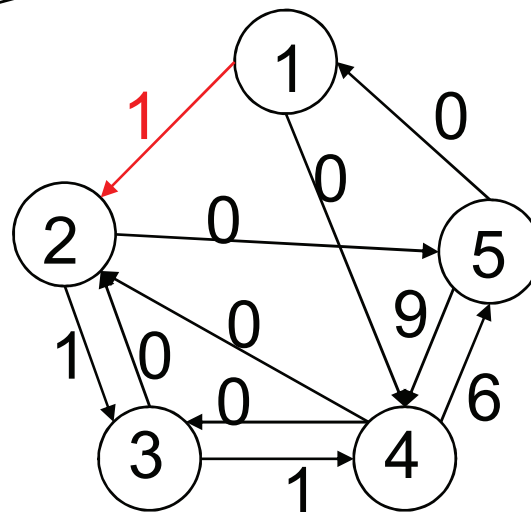
999	1	999	0	999
999	999	1	999	0
999	0	999	1	999
999	0	0	999	6
0	999	999	9	999

bound = 21

1to2

999	999	999	999	999
999	999	1	999	0
999	999	999	1	999
999	999	0	999	6
0	999	999	9	999

bound = 21+1



Add extra cost from 1 to 2, exclude edges from 1 or into 2.

# Focus: going from 1 to 2

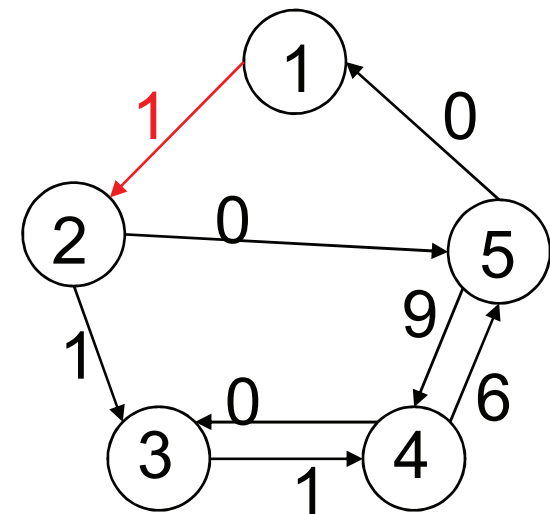
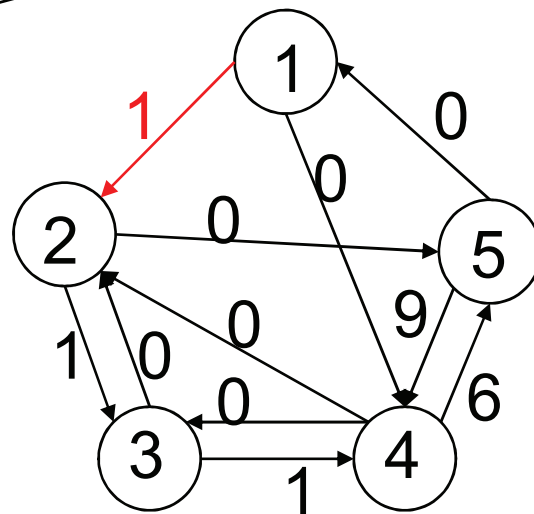
999	1	999	0	999
999	999	1	999	0
999	0	999	1	999
999	0	0	999	6
0	999	999	9	999

bound = 21

1to2

999	999	999	999	999
999	999	1	999	0
999	999	999	1	999
999	999	0	999	6
0	999	999	9	999

bound = 21+1+1



No edges into vertex 4 w/ 0 reduced cost.



# Focus: going from 1 to 2

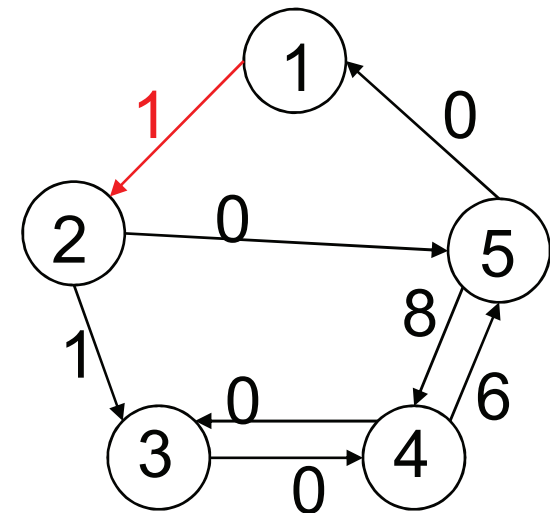
999	1	999	0	999
999	999	1	999	0
999	0	999	1	999
999	0	0	999	6
0	999	999	9	999

bound = 21

1to2

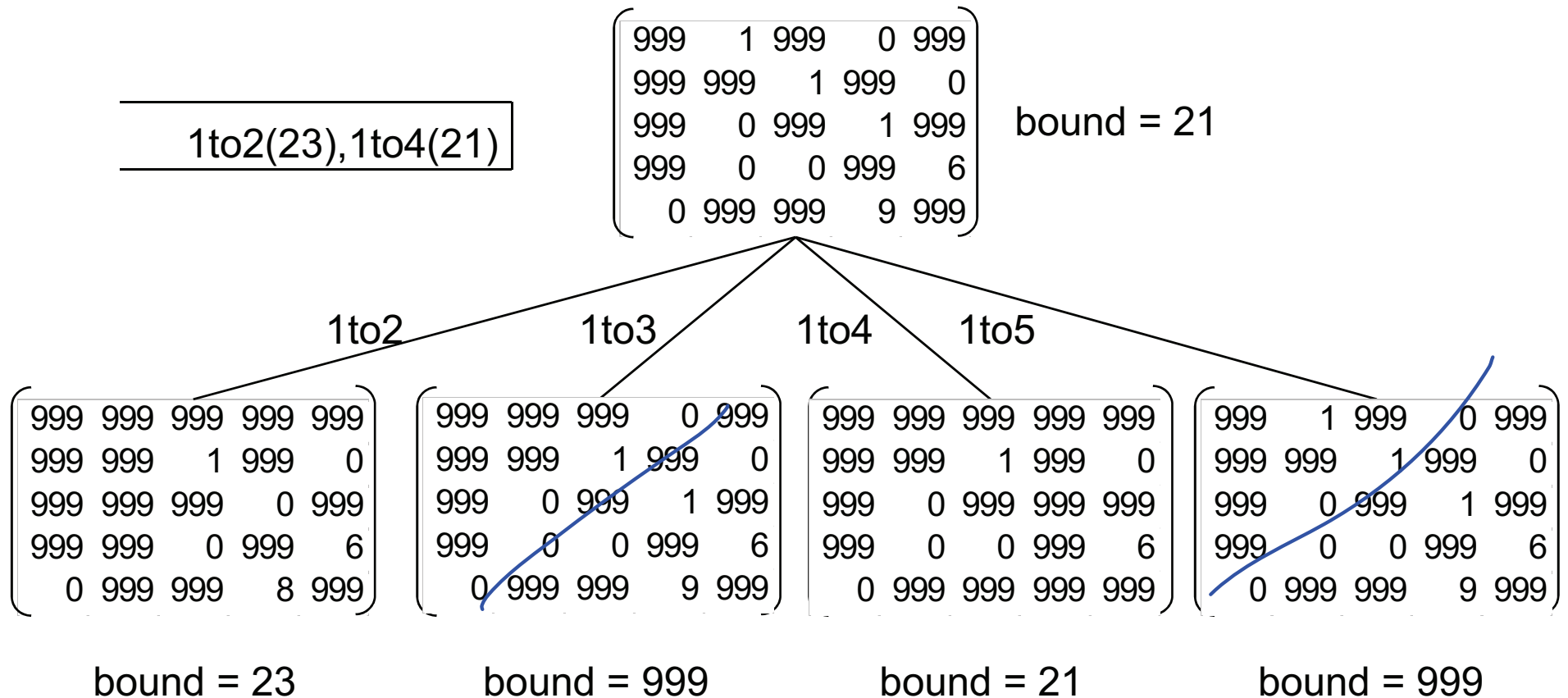
999	999	999	999	999
999	999	1	999	0
999	999	999	0	999
999	999	0	999	6
0	999	999	8	999

bound = 21+1+1 = 23



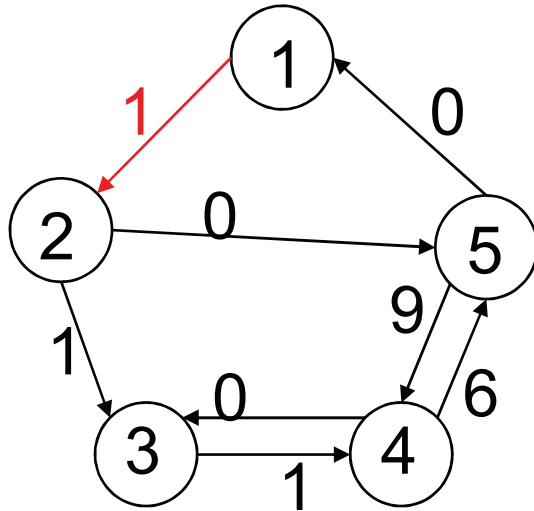
Add cost of reducing edge into vertex 4.

# Bounds for other choices.



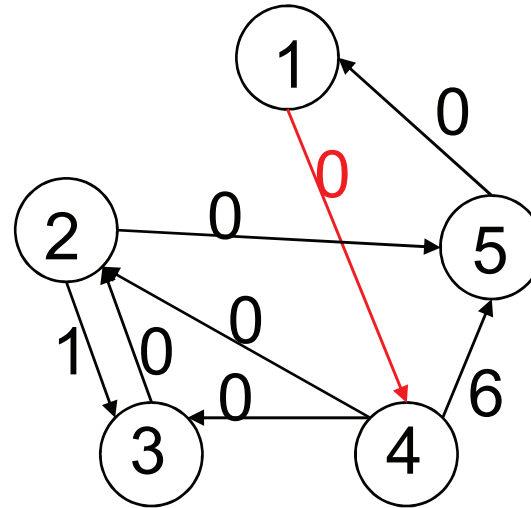
# Leaves us with Two Possibilities on Priority Queue

$1 \rightarrow 2$



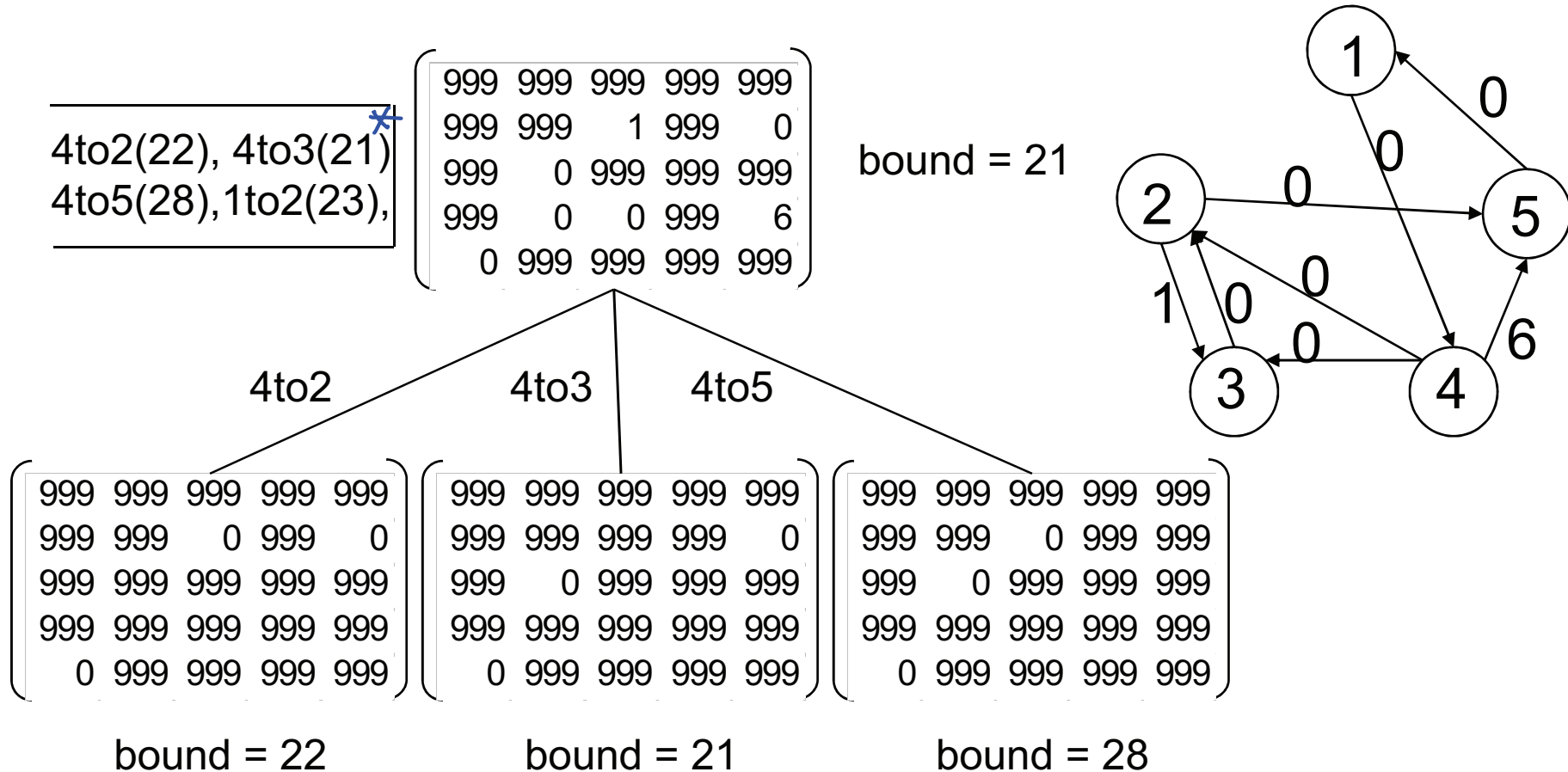
bound = 23

$1 \rightarrow 4$

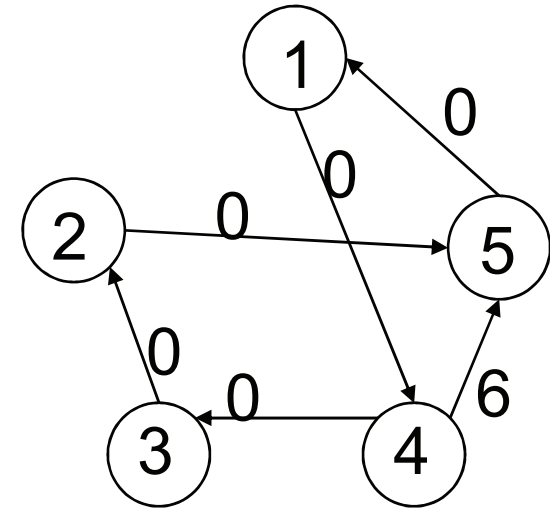
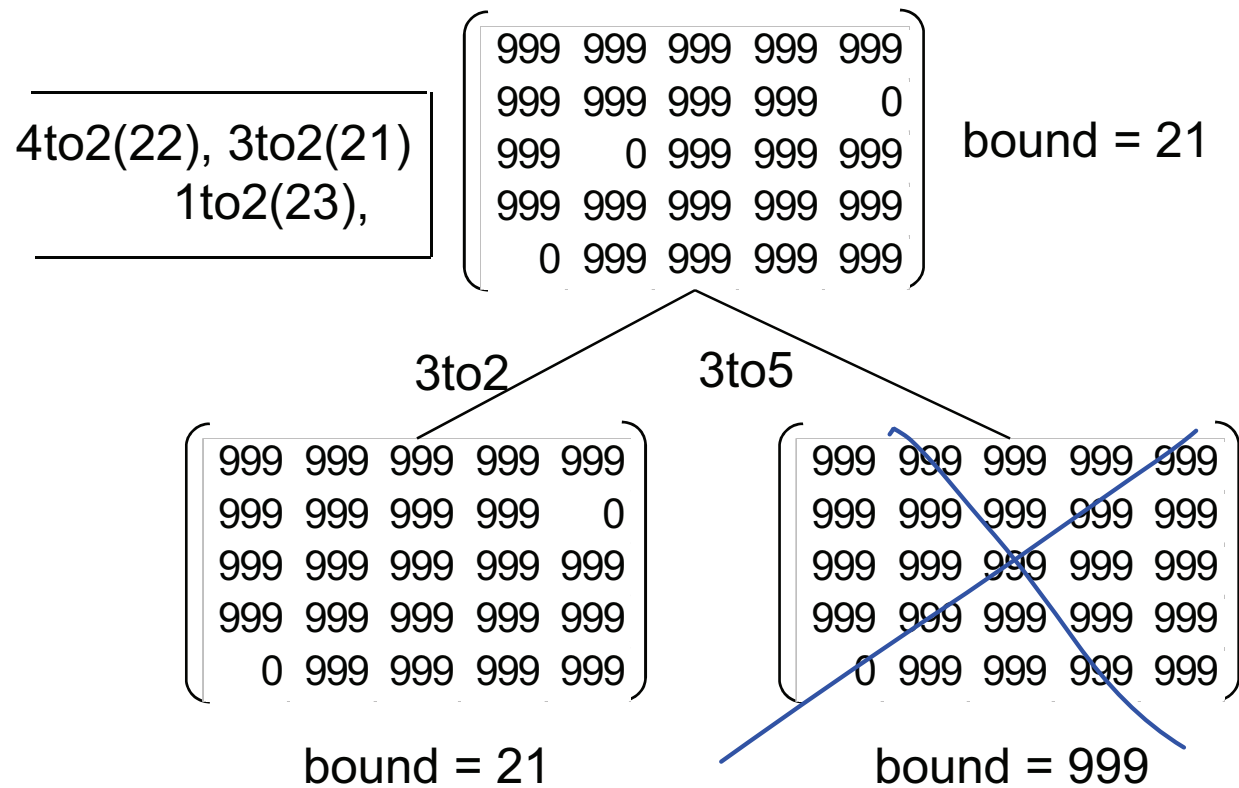


bound = 21

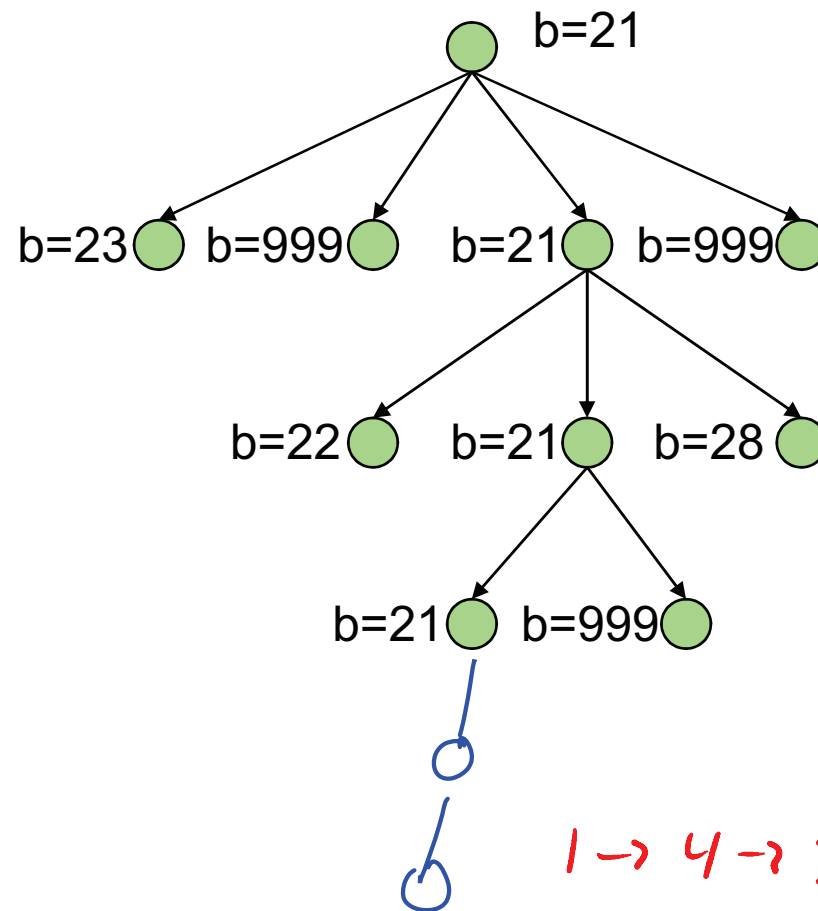
# Leaving Vertex 4



# Leaving Vertex 3



# Search Tree for This Problem



1 → 4 → 3 → 2 → 5 → 1