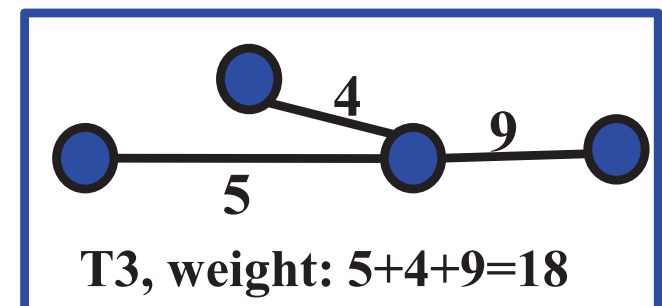
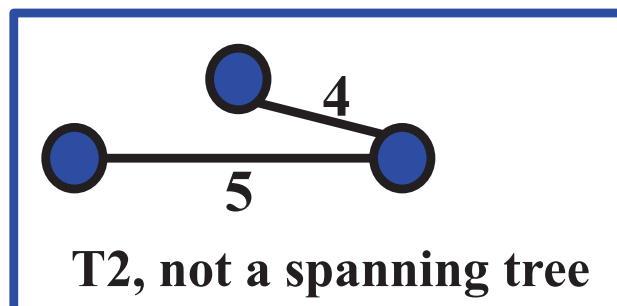
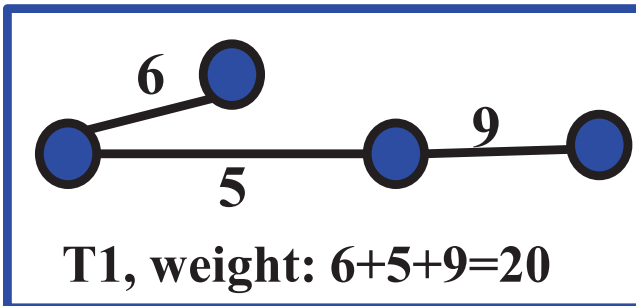
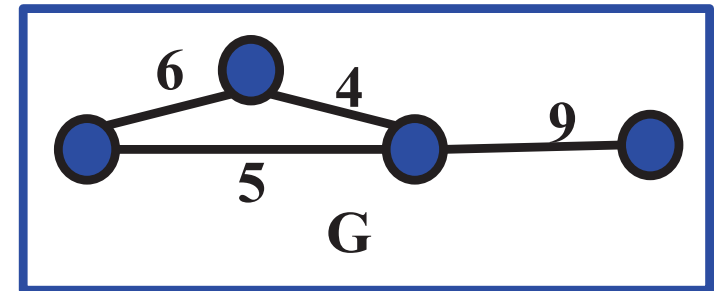


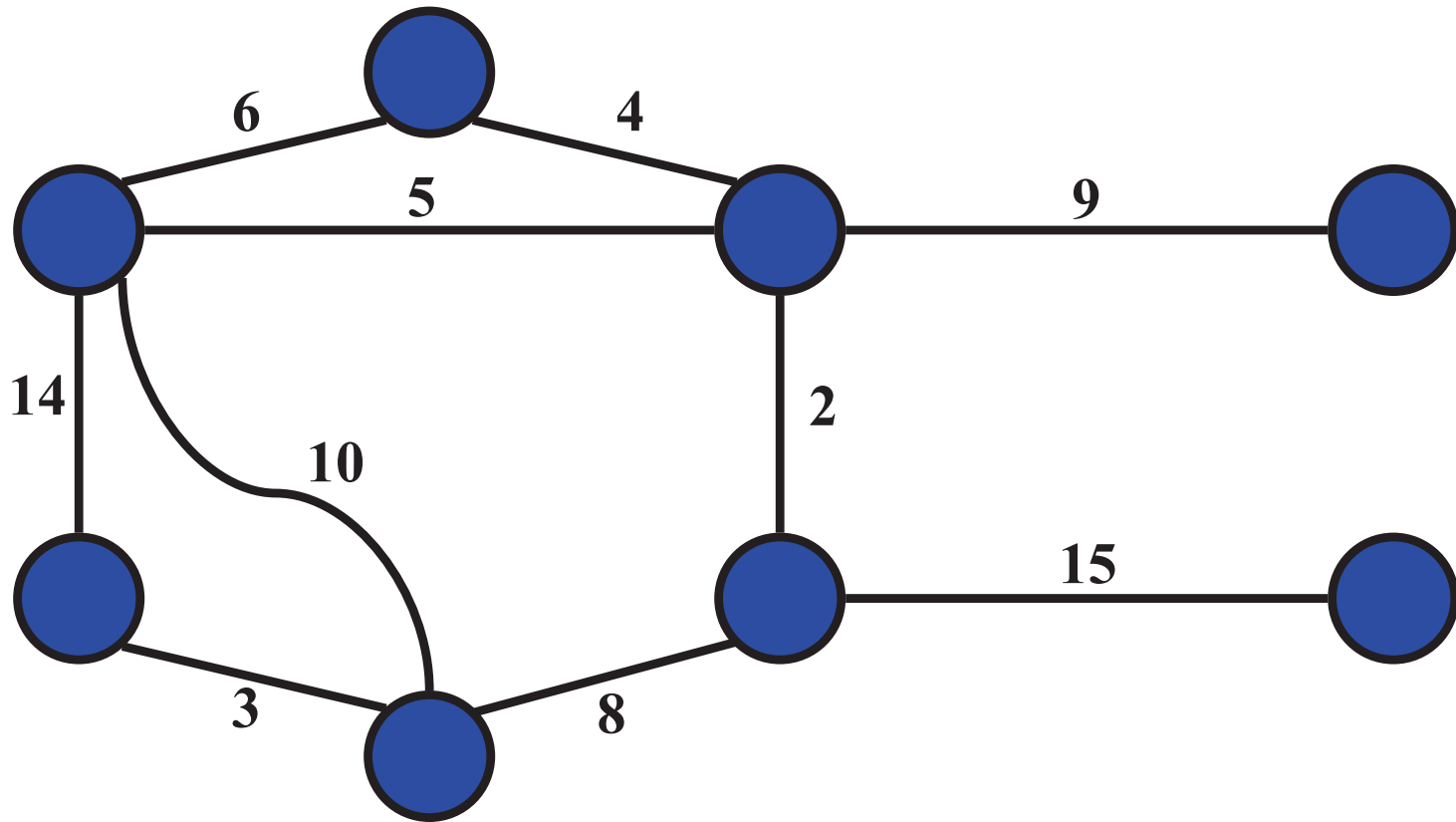
# Spanning Tree and Minimum Spanning Tree

- **Spanning tree** of a connected graph: A tree that connects *all* nodes of the graph.
- **Minimum Spanning tree** of a connected weighted graph: A tree that connects all nodes of the graph with total edge weight lowest possible.
- **Example:** For the graph G,
  - T1, T3 are spanning trees
  - T2 is not a spanning tree, because it does not connect all nodes of G
  - T1 is not minimum, because its total weight is 20
  - T2 is minimum with lowest possible total weight 19



# Minimum Spanning Tree

- Problem: given a connected, undirected, weighted graph, find a *minimum spanning tree*



# Two Algorithm

- Krushkal's algorithm
- Prim's algorithm

# Kruskal's Algorithm

- Take the minimum weight edge one after another if there is no cycle
- Use disjoint set data structure
- MakeSet() for each element
- Maintain two sets:
  - tree edges
  - other remaining edges
- Use FindSet() to check for cycle (If same set, then must be cycle, because if you connect any two vertices of one set, then it makes a cycle.)

Symbols used:

E: Edge set, V: Vertex set, T: Resulting tree, n: number of vertices

# Kruskal's Algorithm

```
Kruskal()
```

```
{
```

```
  T =  $\emptyset$ ;
```

```
  for each v  $\in$  V
```

```
    MakeSet(v);
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  sort E by increasing edge weight w
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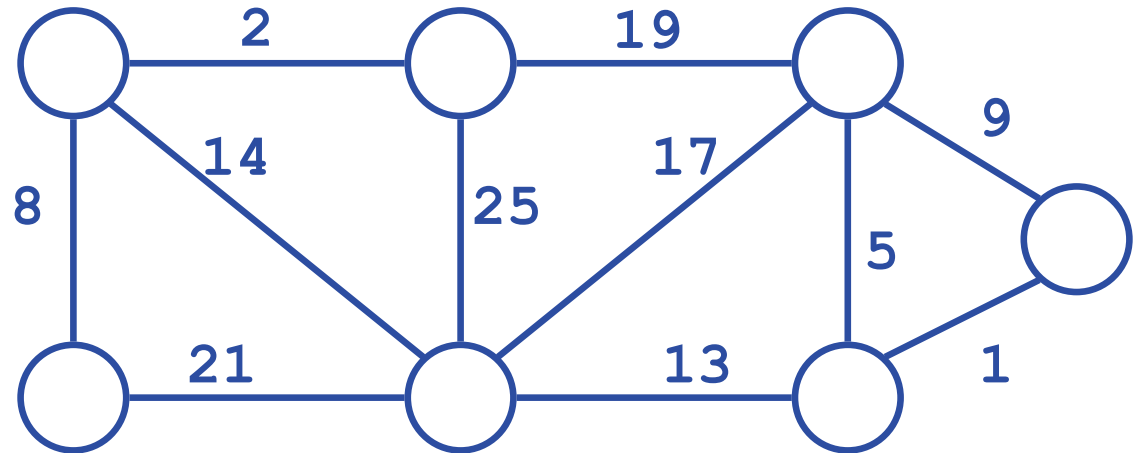
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*Run the algorithm:*



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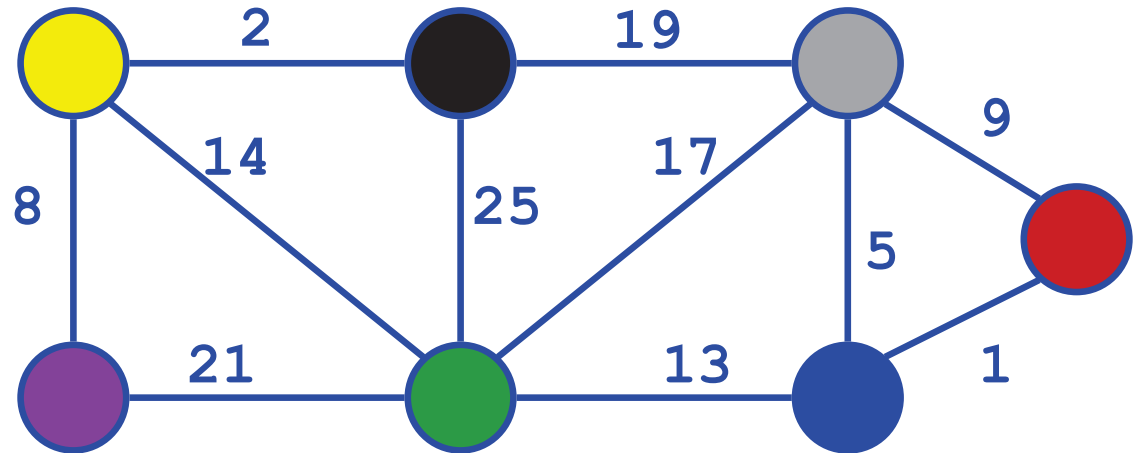
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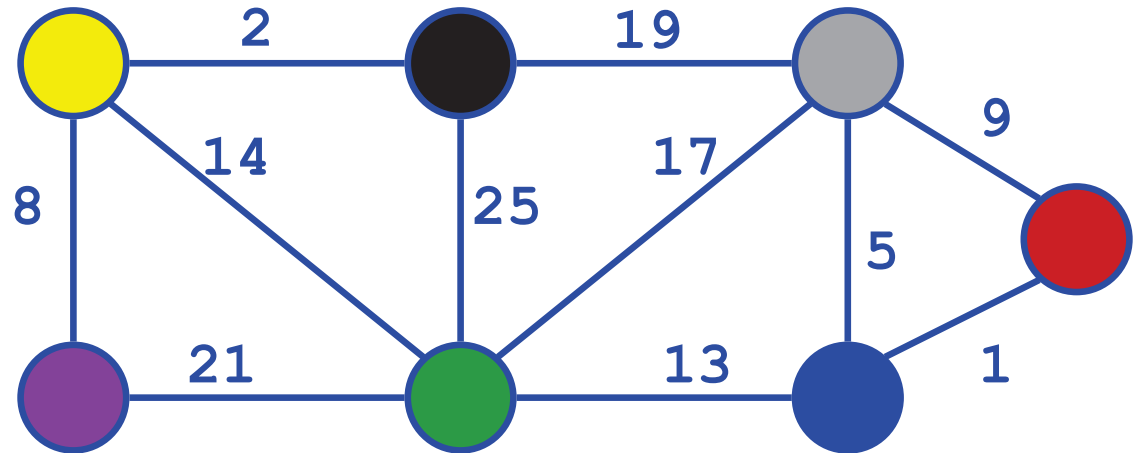
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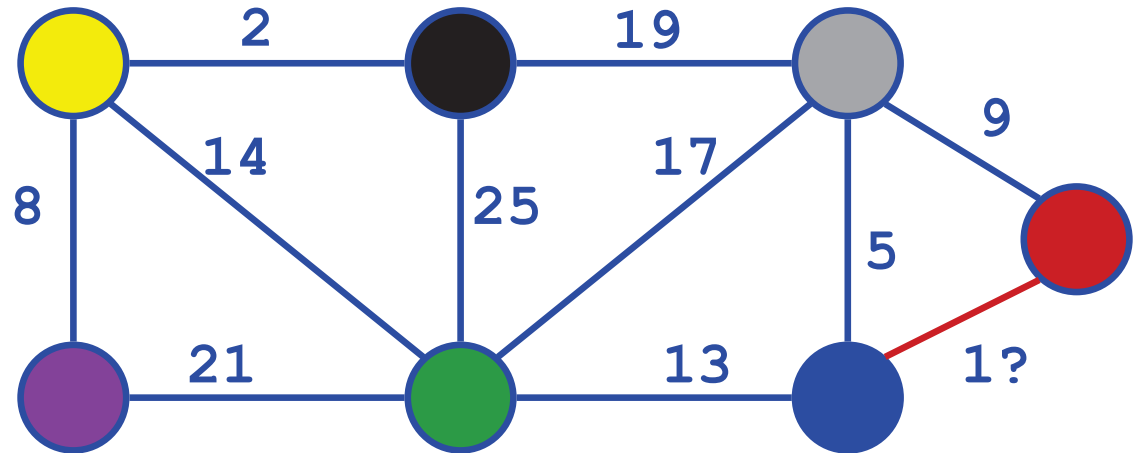
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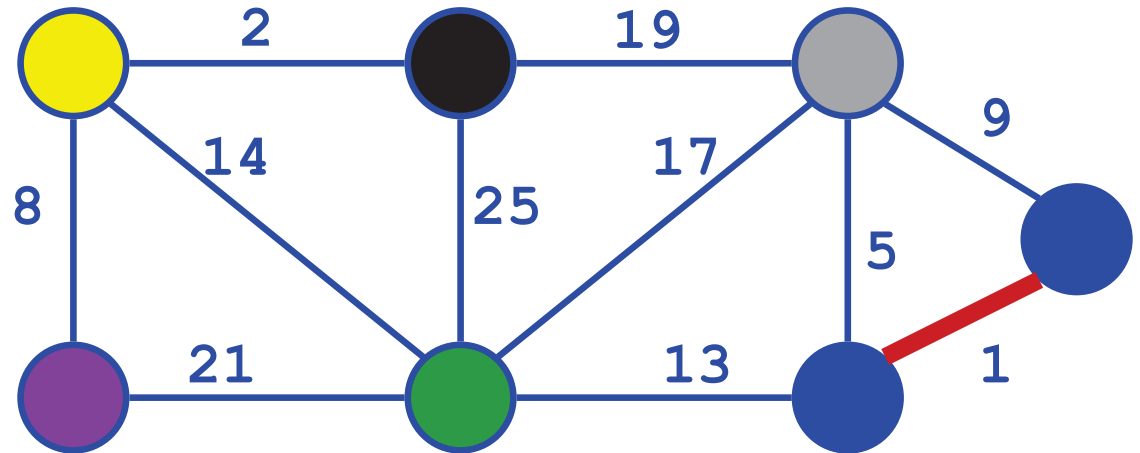
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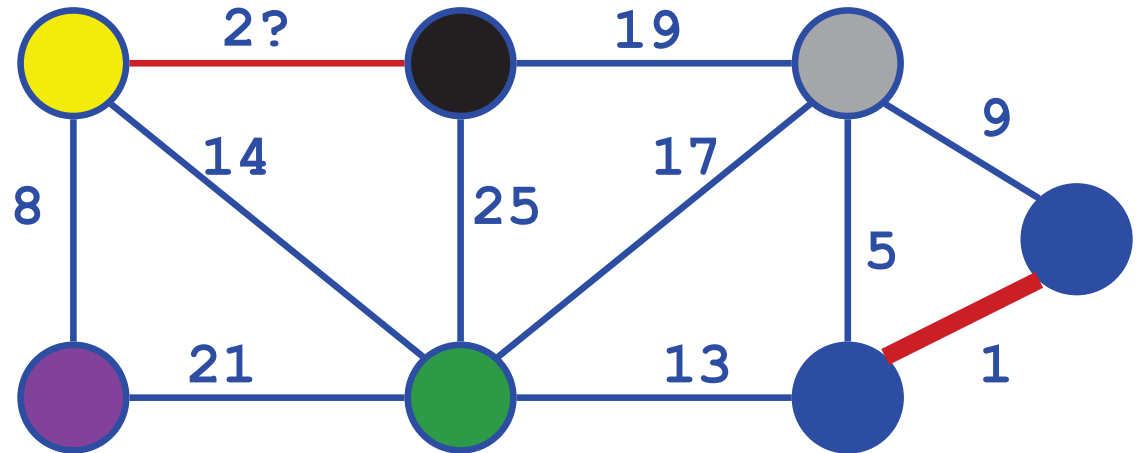
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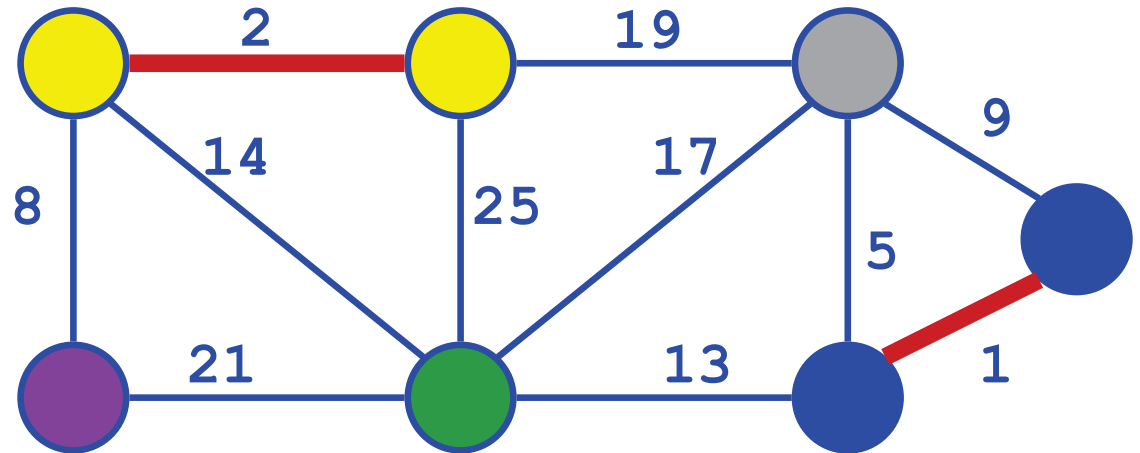
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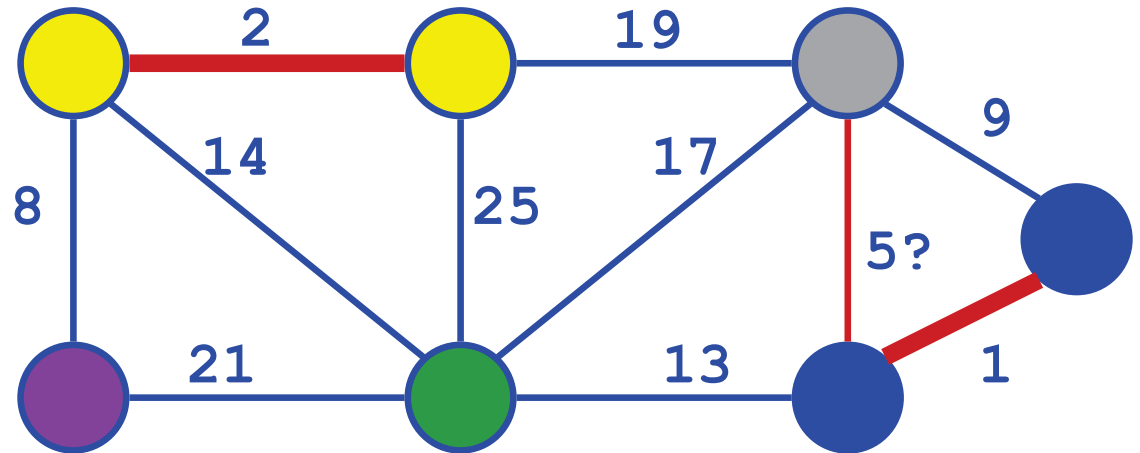
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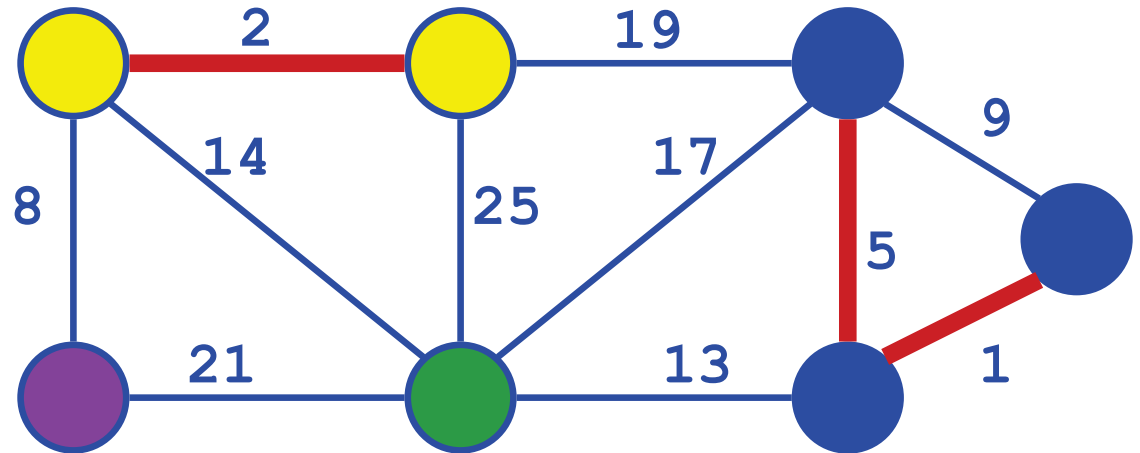
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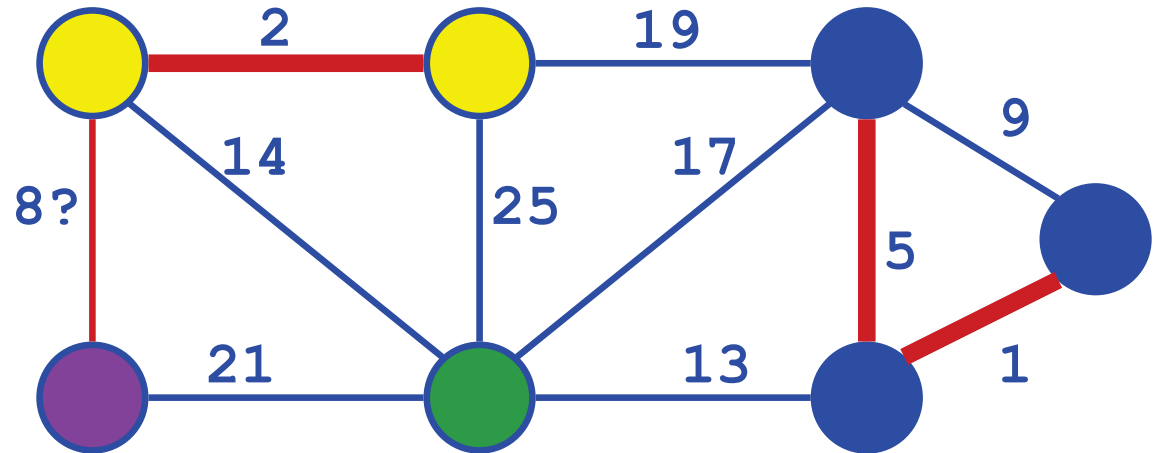
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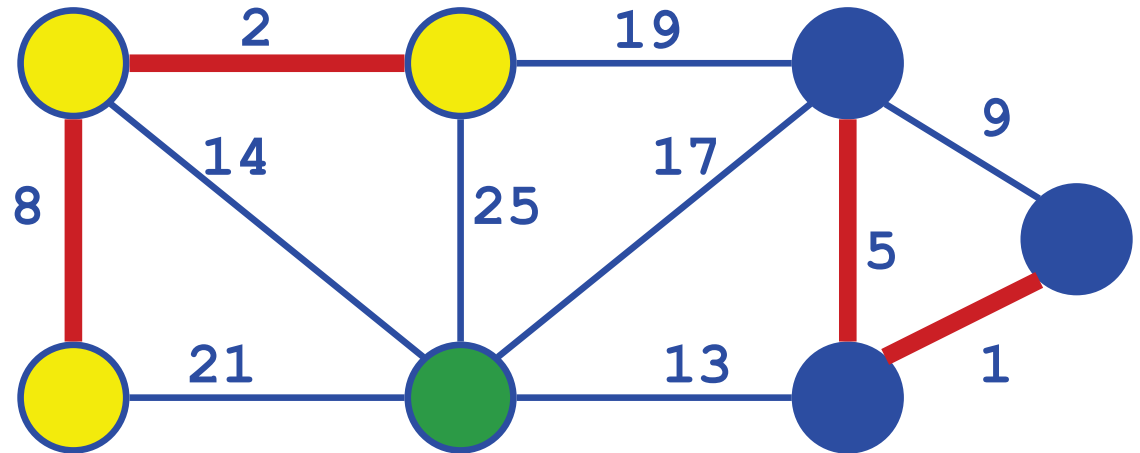
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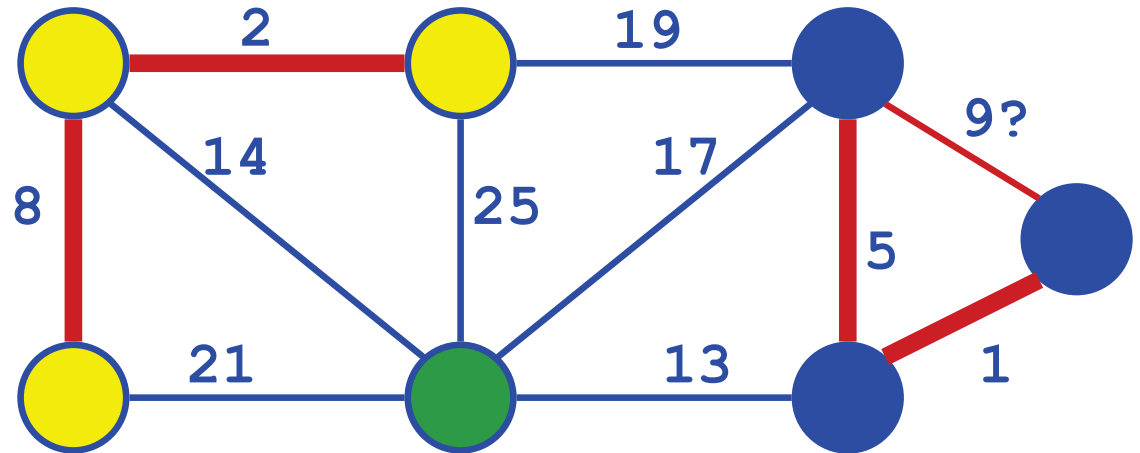
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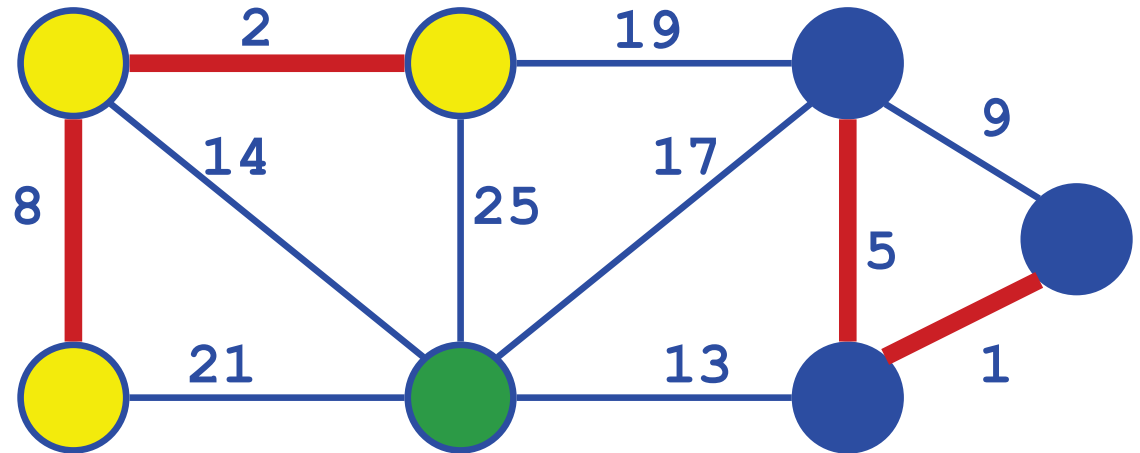
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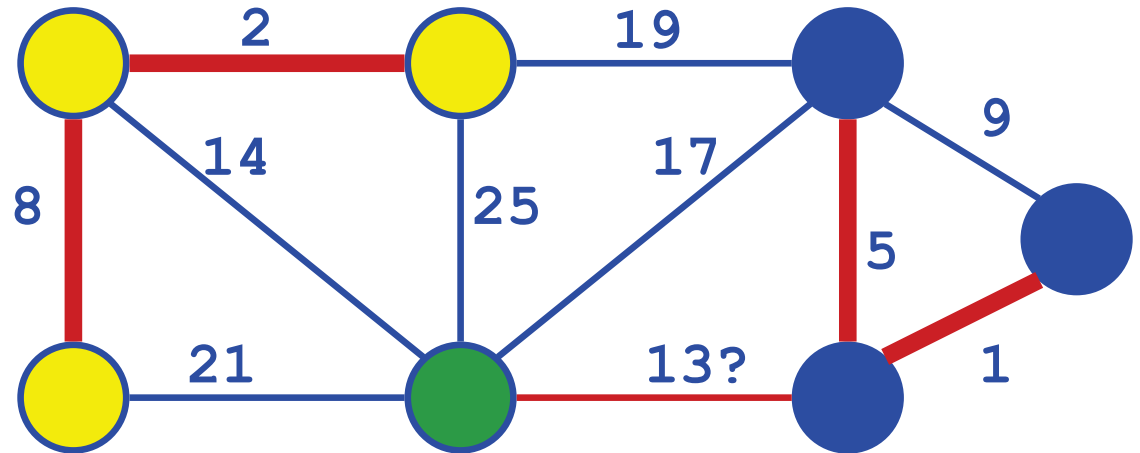
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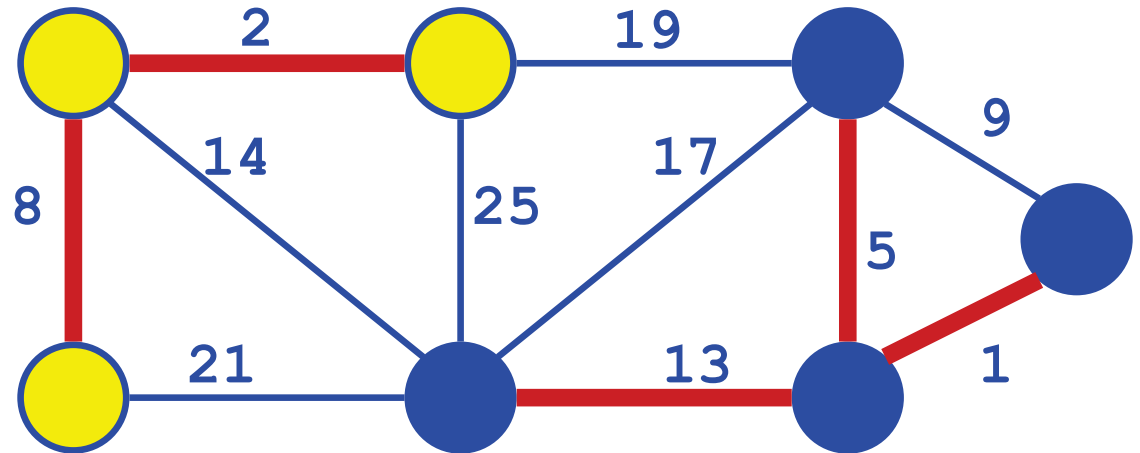
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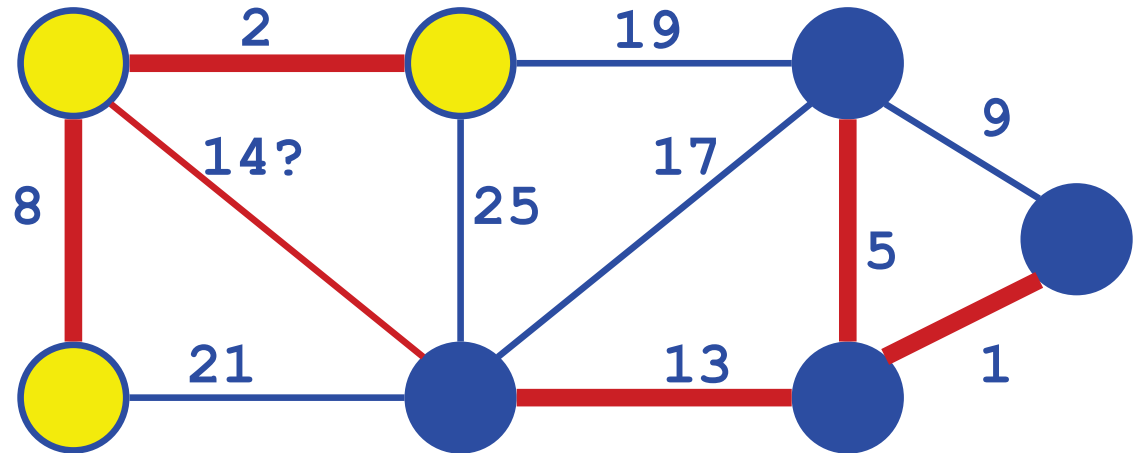
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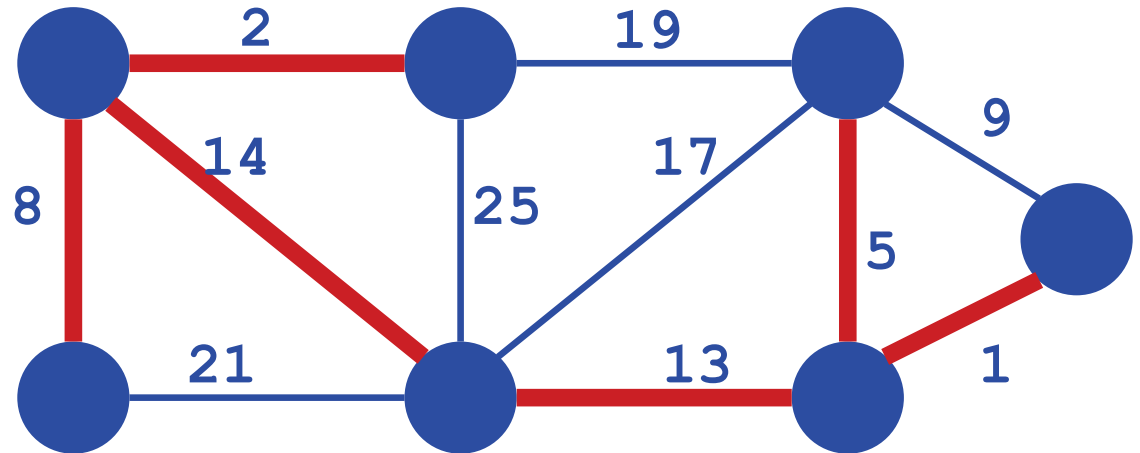
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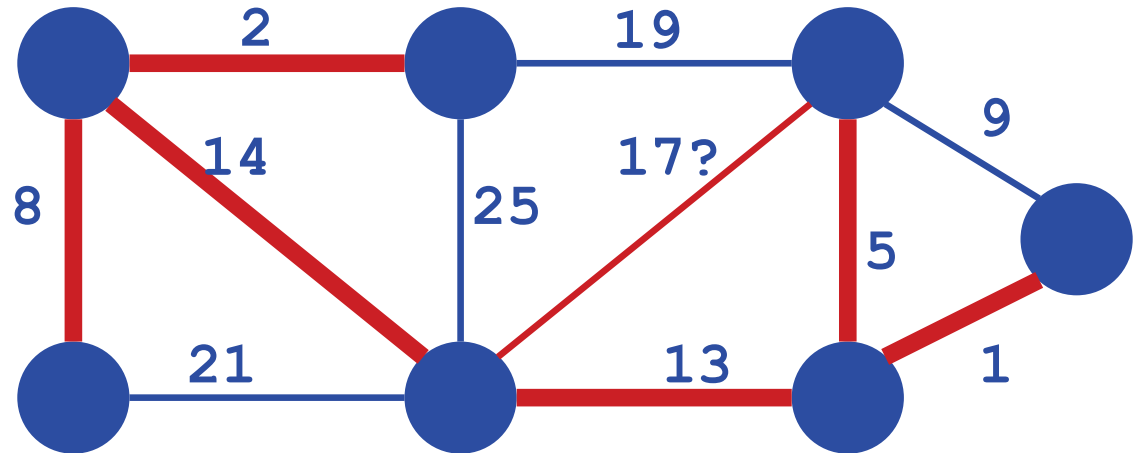
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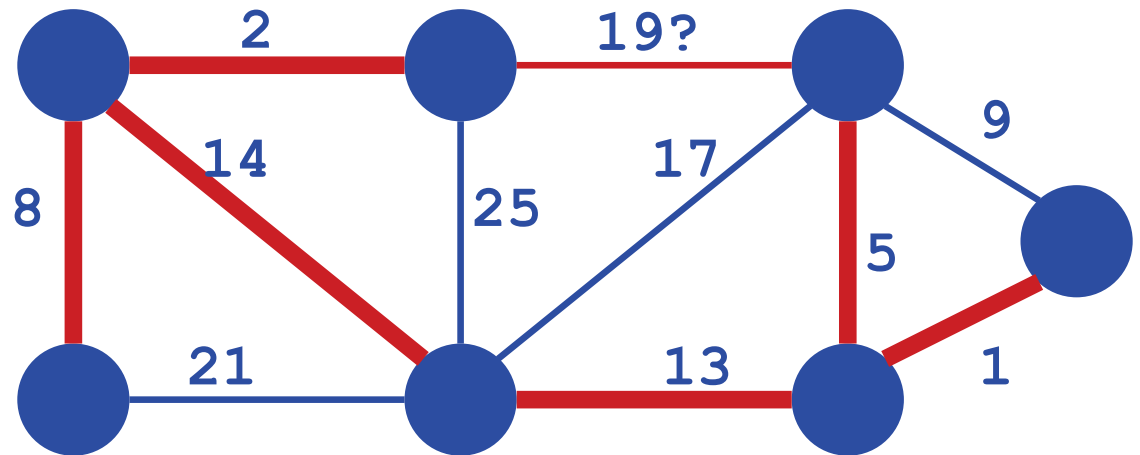
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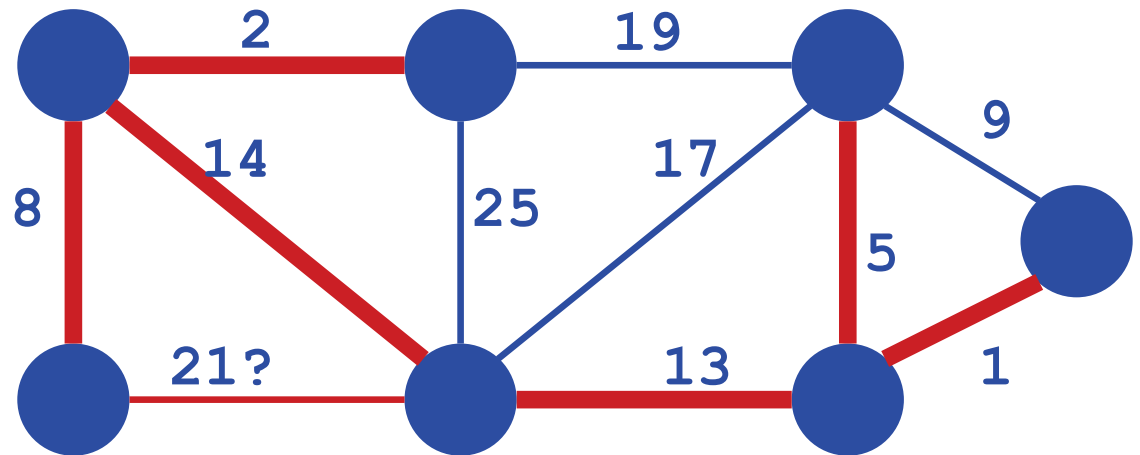
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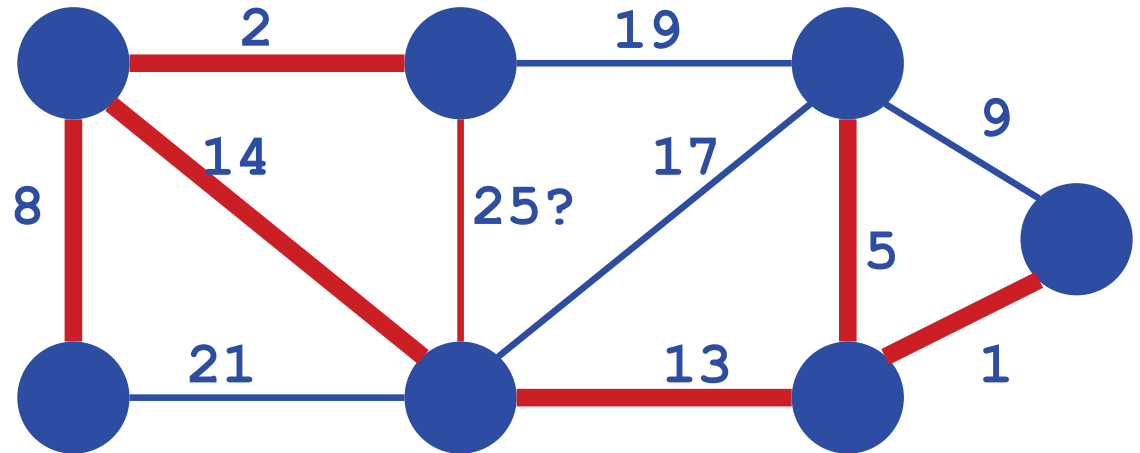
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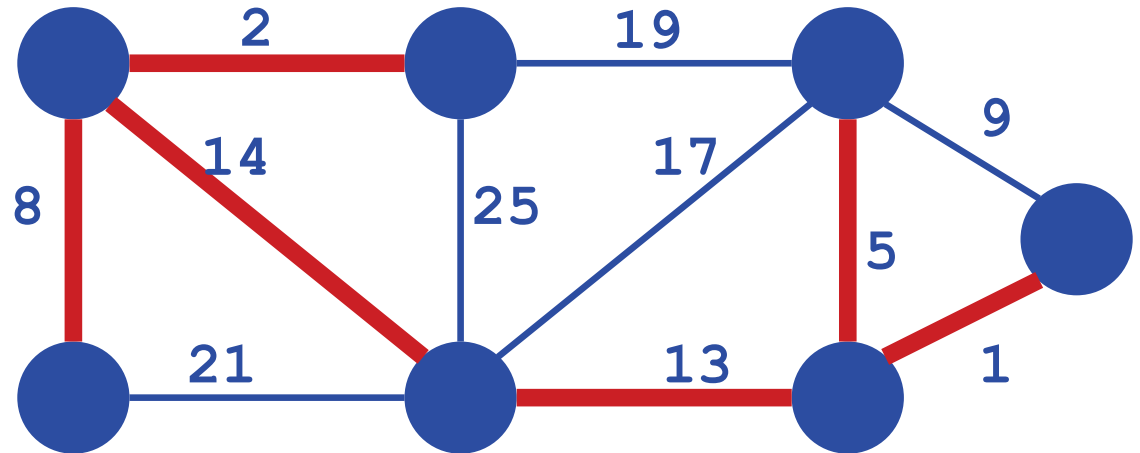
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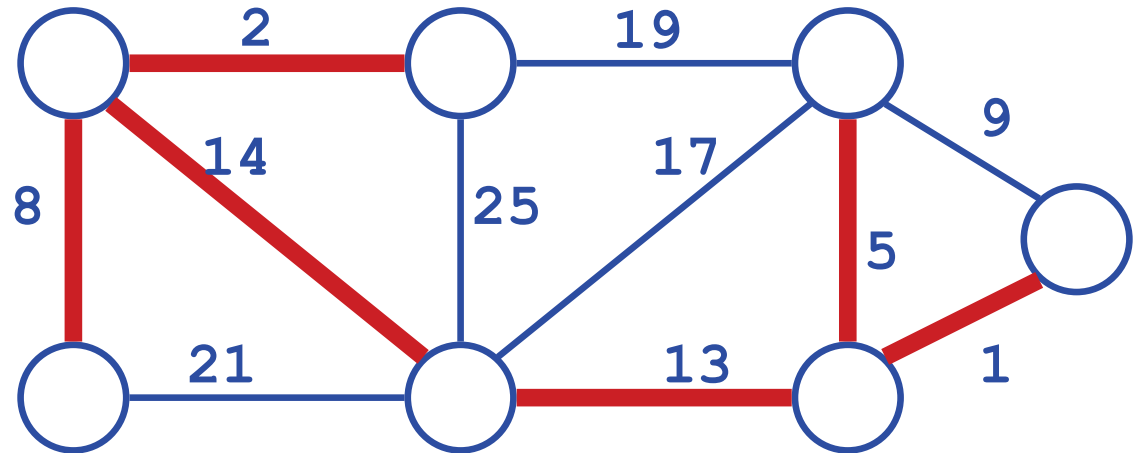
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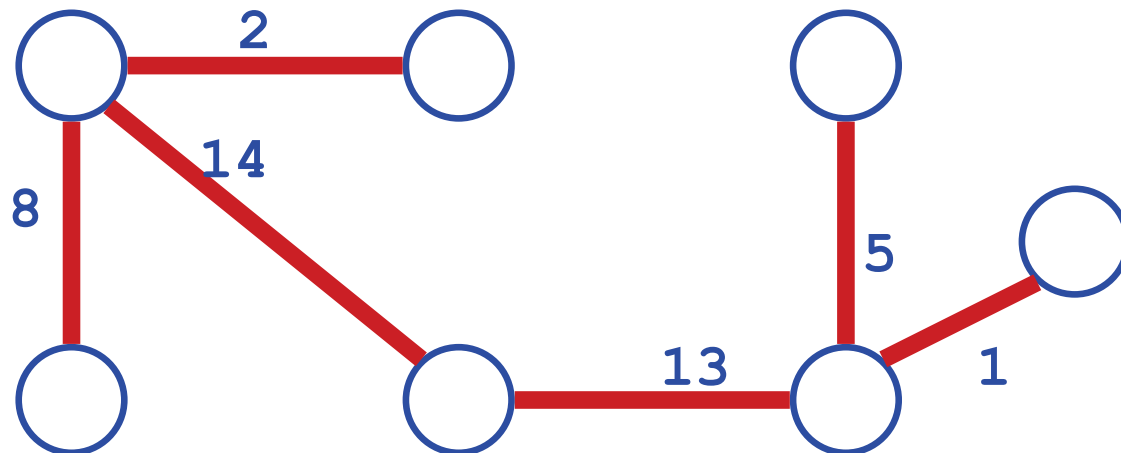
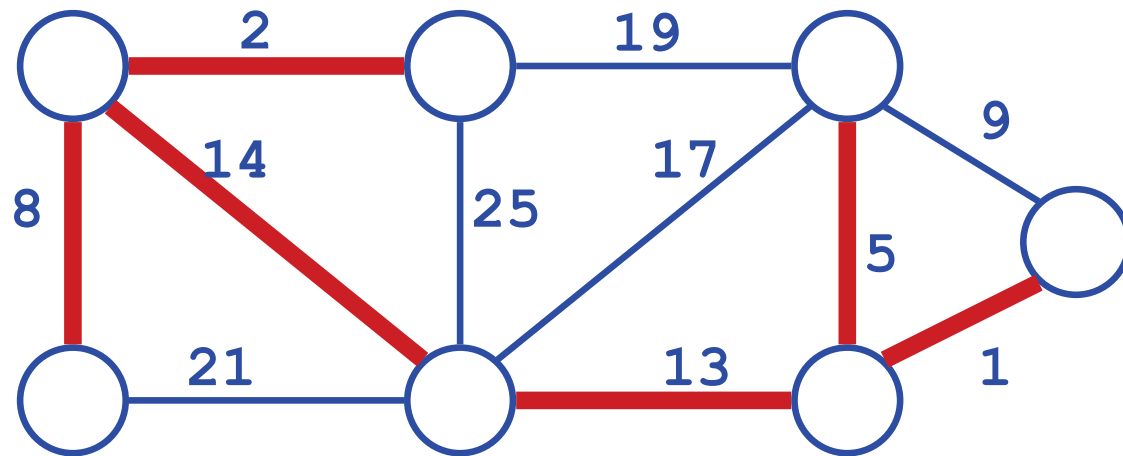
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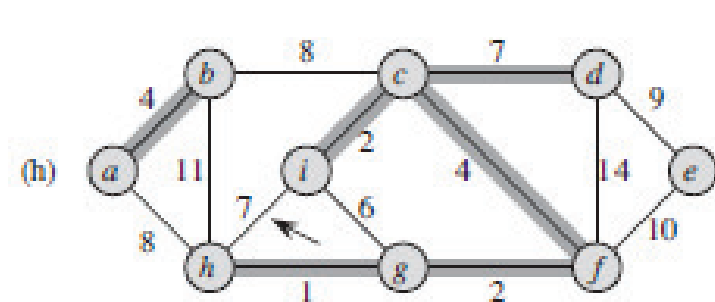
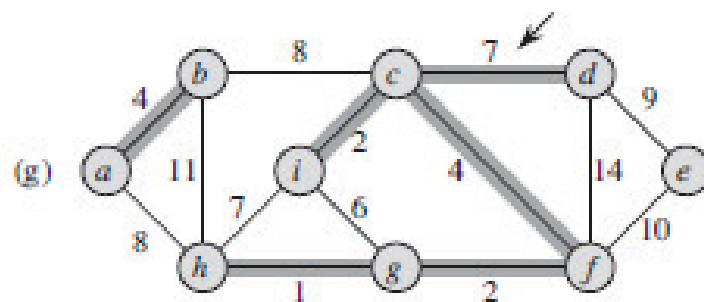
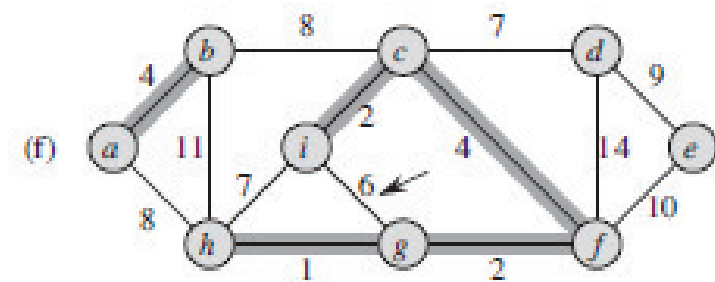
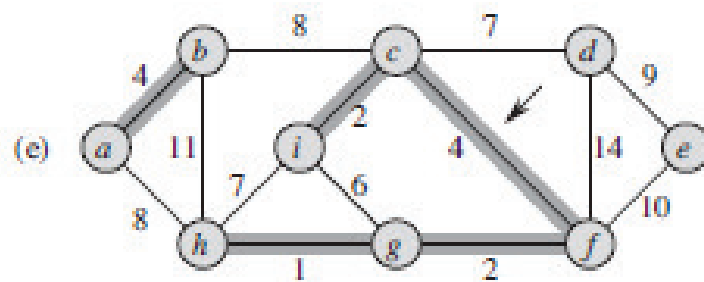
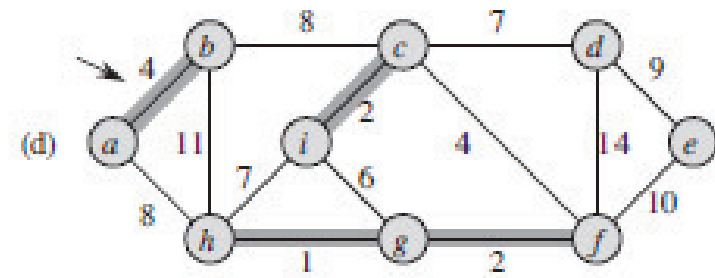
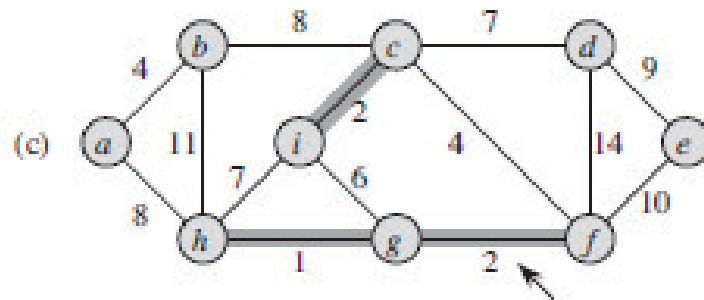
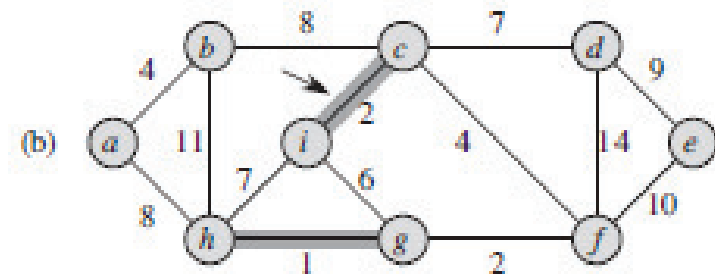
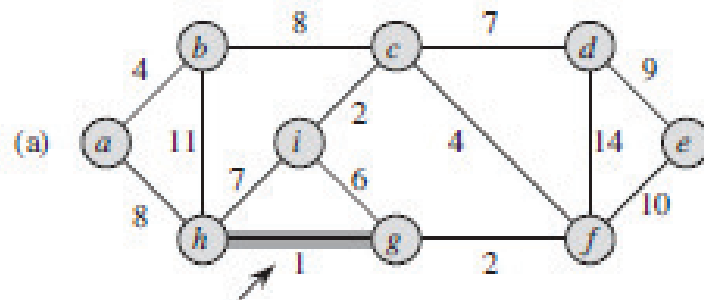
*Run the algorithm:*



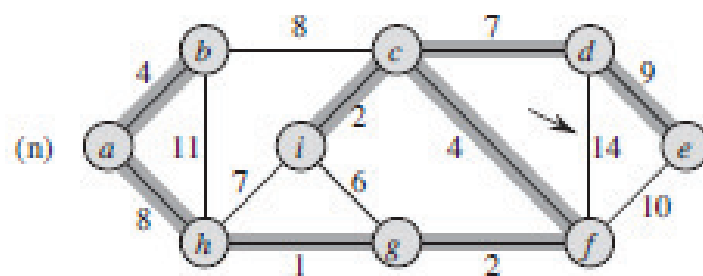
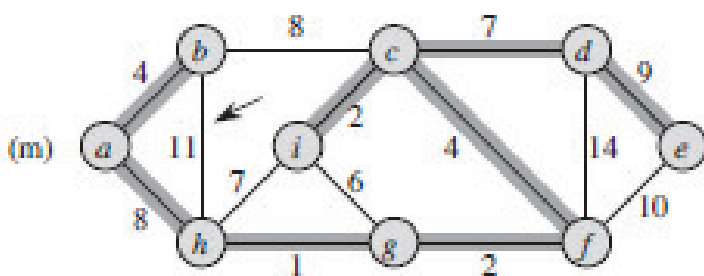
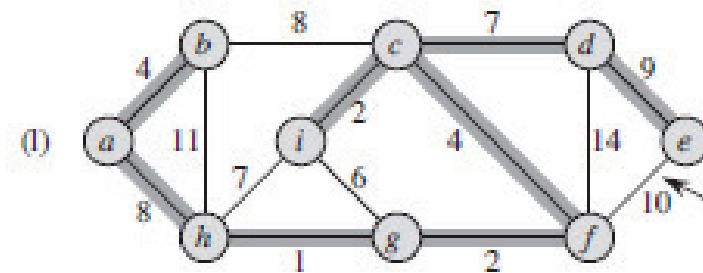
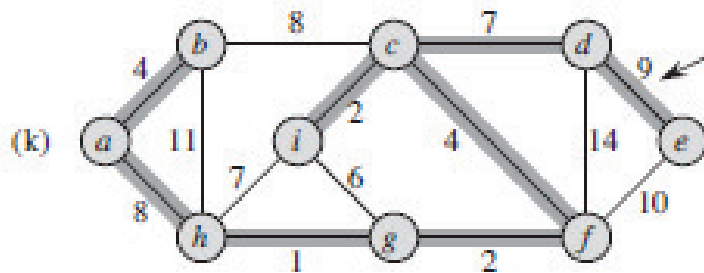
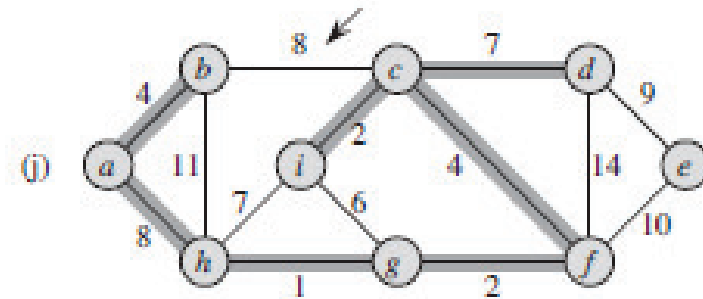
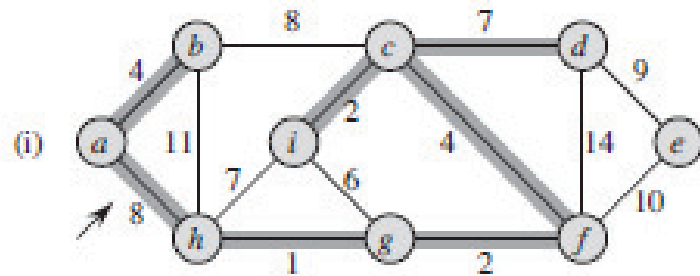
# Kruskal's Algorithm



# Another Example: Krushkal's Algorithm



# Another Example: Krushkal's Algorithm



# Prim's Algorithm

1. Initially, all node have weight  $\infty$
2. Keep all of them in a queue (heap)
3. Take minimum node from heap (ExtractMin())
4. For each neighbor of this node, if edge weight from this node is smaller than current weight , then update weight (DecreaseKey()) and add that edge in the tree.
5. Repeat step 3, 4 until finished.

n: number of vertex

# Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

## MST-Prim

```
for all node  $u$   
     $\text{weight}[u] = \infty$ ;  
 $Q = \text{all node}$ ;  
 $\text{weight}[\text{root}] = 0$ ;  
 $p[\text{root}] = \text{NULL}$ ;  
while ( $Q$  not empty)
```

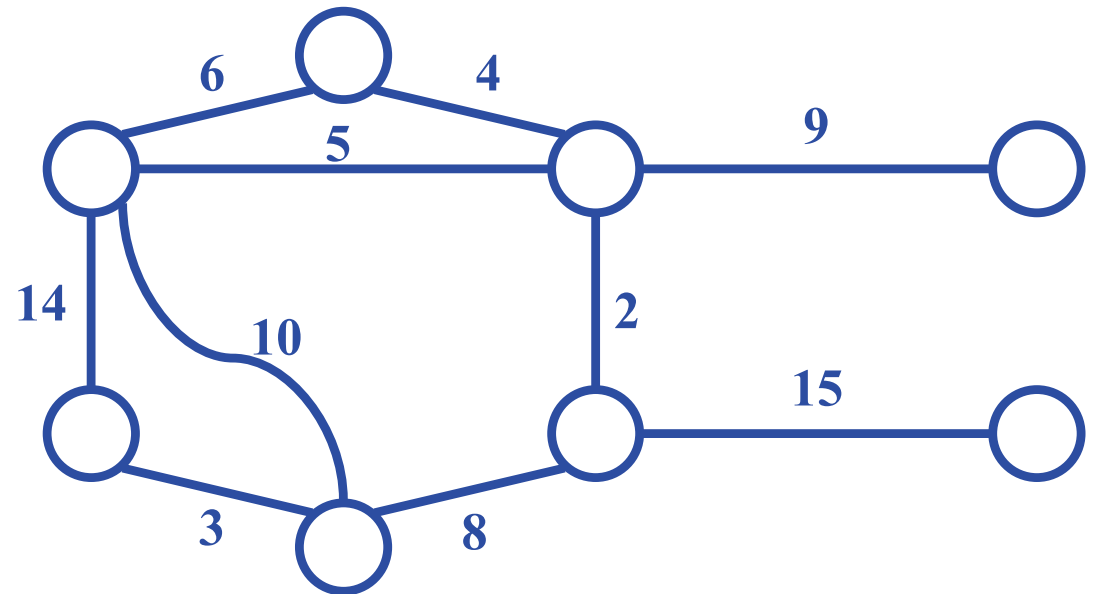
```
     $u = \text{ExtractMin}(Q)$  ;
```

```
    for each  $v \in \text{Adjacent}[u]$ 
```

```
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
```

```
             $p[v] = u$ ;
```

```
             $\text{weight}[v] = w(u, v)$  ;
```



*Run on example graph*



# Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

## MST-Prim

```
for all node  $u$ 
```

```
     $\text{weight}[u] = \infty$ ;
```

```
 $Q = \text{all node}$ ;
```

```
 $\text{weight}[\text{root}] = 0$ ;
```

```
 $p[\text{root}] = \text{NULL}$ ;
```

```
while ( $Q$  not empty)
```

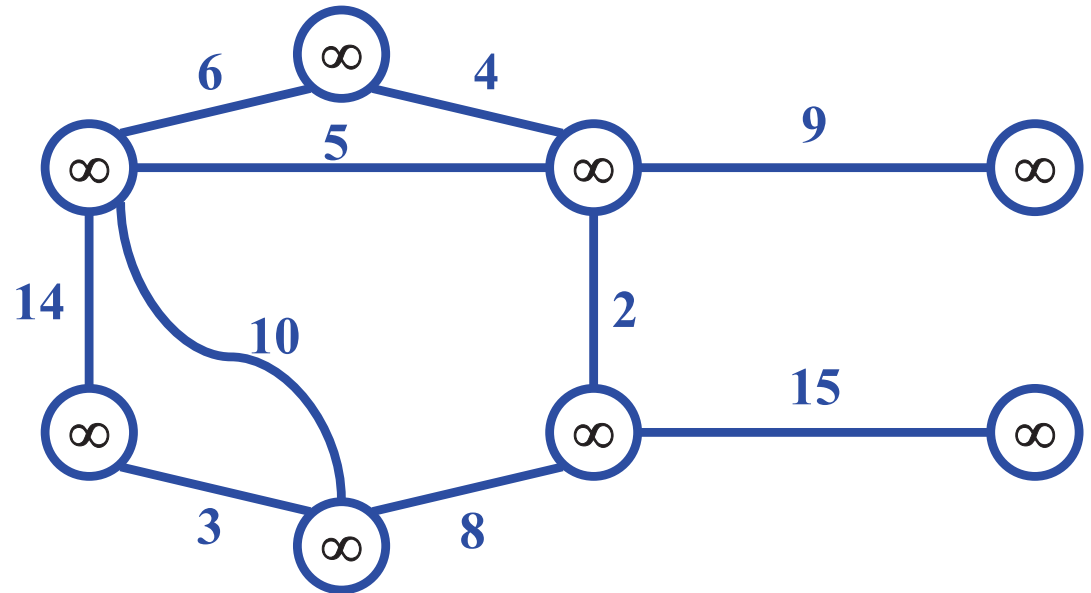
```
     $u = \text{ExtractMin}(Q)$  ;
```

```
    for each  $v \in \text{Adjacent}[u]$ 
```

```
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
```

```
             $p[v] = u$ ;
```

```
             $\text{weight}[v] = w(u, v)$  ;
```



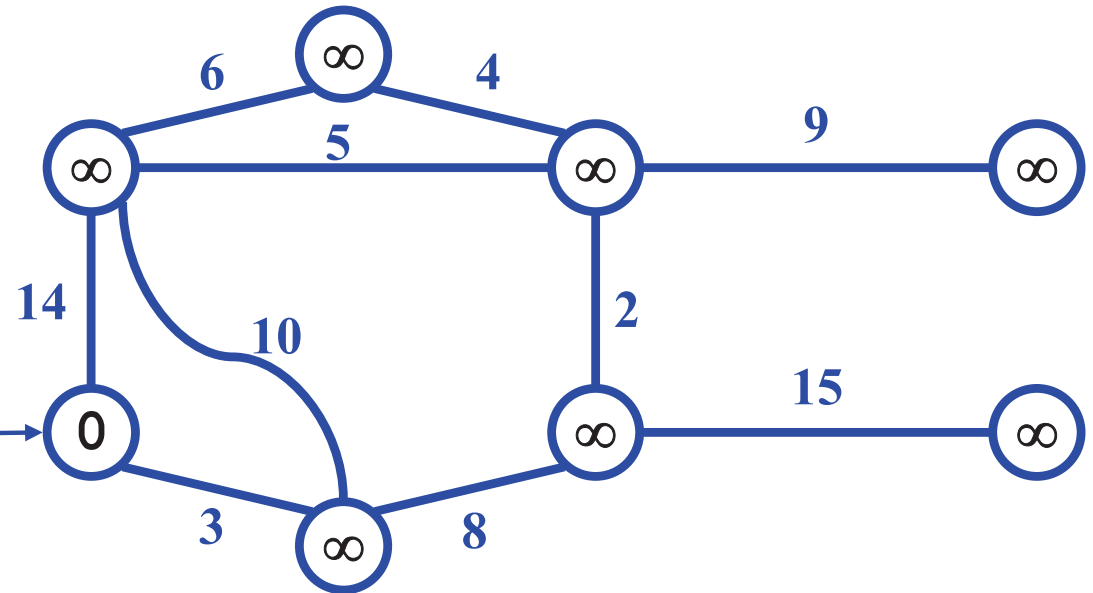
*Run on example graph*

# Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

## MST-Prim

```
for all node  $u$ 
     $\text{weight}[u] = \infty$ ;
 $Q = \text{all node}$ ;
 $\text{weight}[\text{root}] = 0$ ;  $r \rightarrow$ 
 $p[\text{root}] = \text{NULL}$ ;
while ( $Q$  not empty)
     $u = \text{ExtractMin}(Q)$ ;
    for each  $v \in \text{Adjacent}[u]$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
             $\text{weight}[v] = w(u, v)$ ;
```



*Pick a start vertex  $r$*

# Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
for all node  $u$ 
```

```
     $\text{weight}[u] = \infty$ ;
```

```
Q = all node;
```

```
 $\text{weight}[\text{root}] = 0$ ;
```

```
p[root] = NULL;
```

```
while (Q not empty)
```

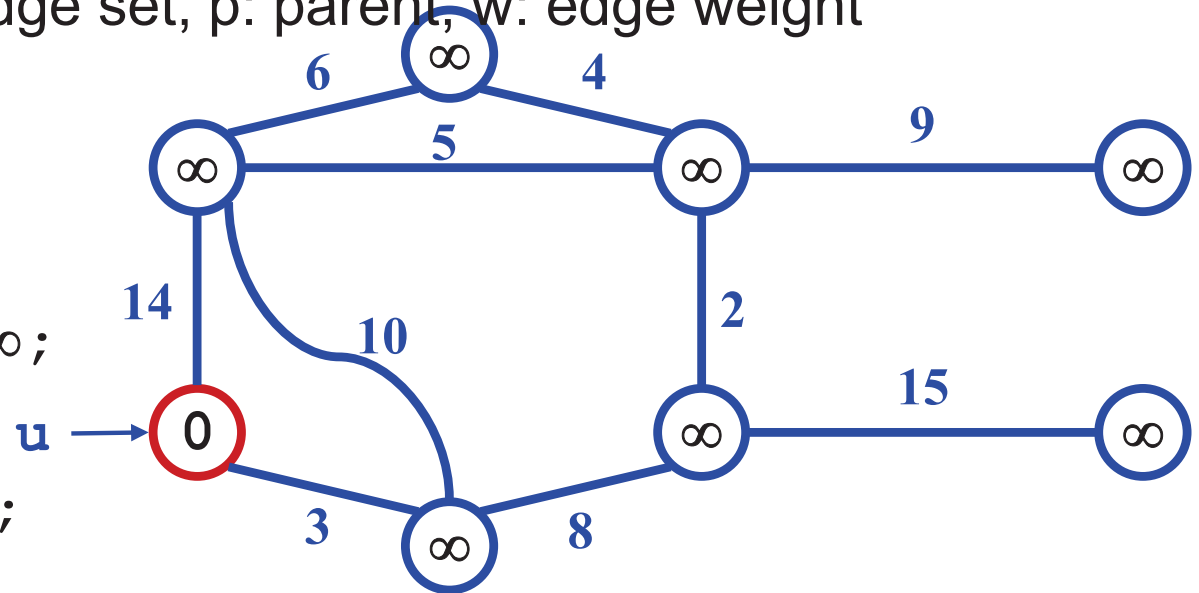
```
     $u = \text{ExtractMin}(Q)$  ;
```

```
    for each  $v \in \text{Adjacent}[u]$ 
```

```
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
```

```
            p[v] = u;
```

```
             $\text{weight}[v] = w(u, v)$  ;
```

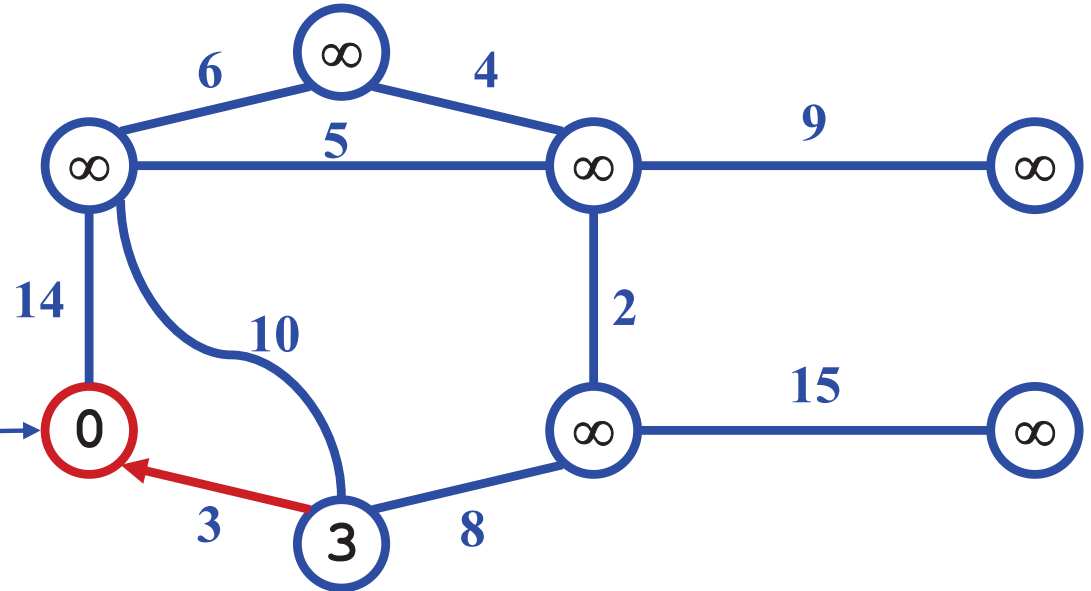


# Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

## MST-Prim

```
for all node  $u$ 
     $\text{weight}[u] = \infty$ ;
 $Q = \text{all node}$ ;
 $\text{weight}[\text{root}] = 0$ ;  $u \rightarrow$ 
 $p[\text{root}] = \text{NULL}$ ;
while ( $Q$  not empty)
     $u = \text{ExtractMin}(Q)$ ;
    for each  $v \in \text{Adjacent}[u]$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
             $\text{weight}[v] = w(u, v)$ ;
```



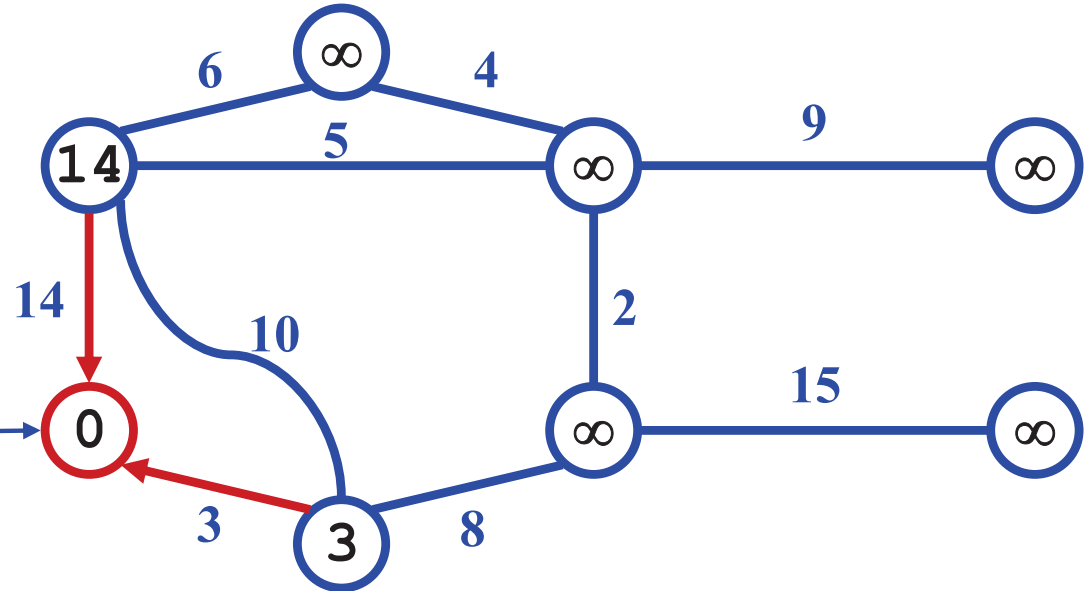
*Red arrows indicate parent pointers*

# Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
for all node  $u$ 
     $\text{weight}[u] = \infty$ ;
 $Q = \text{all node}$ ;
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 $p[\text{root}] = \text{NULL}$ ;
while ( $Q$  not empty)
     $u = \text{ExtractMin}(Q)$ ;
    for each  $v \in \text{Adjacent}[u]$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
             $\text{weight}[v] = w(u, v)$ ;
```

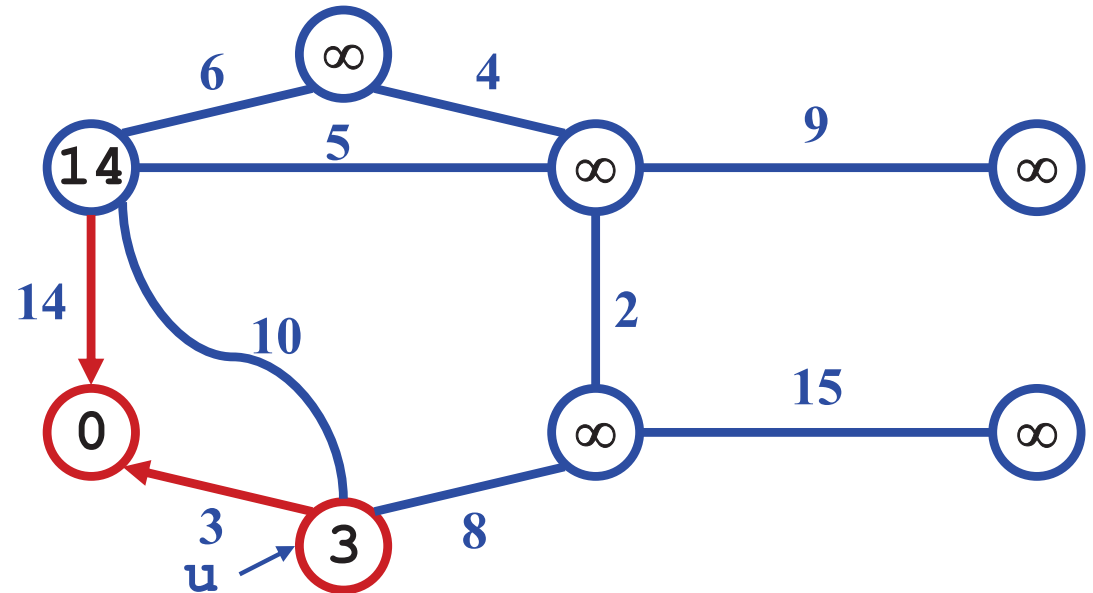


# Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

## MST-Prim

```
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     $\text{weight}[u] = \infty$ ;
 $Q = \text{all node}$ ;
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     $u = \text{ExtractMin}(Q)$ ;
    for each  $v \in \text{Adjacent}[u]$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
             $\text{weight}[v] = w(u, v)$ ;
```



# Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
for all node  $u$ 
```

```
     $\text{weight}[u] = \infty;$ 
```

```
Q = all node;
```

```
 $\text{weight}[\text{root}] = 0;$ 
```

```
p[root] = NULL;
```

```
while (Q not empty)
```

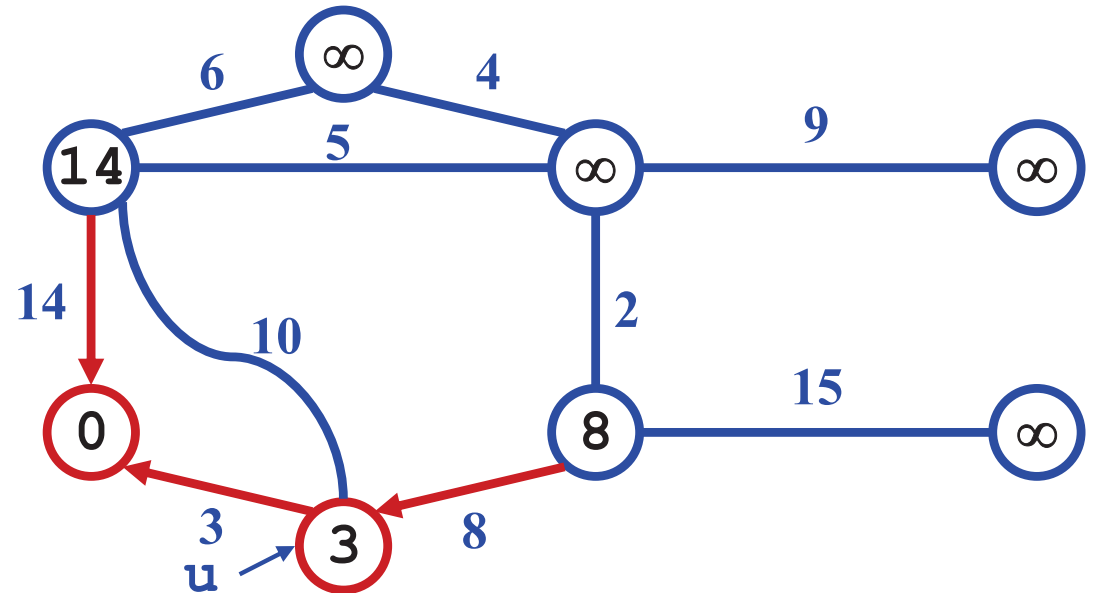
```
     $u = \text{ExtractMin}(Q);$ 
```

```
    for each  $v \in \text{Adjacent}[u]$ 
```

```
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
```

```
            p[v] = u;
```

```
             $\text{weight}[v] = w(u, v);$ 
```

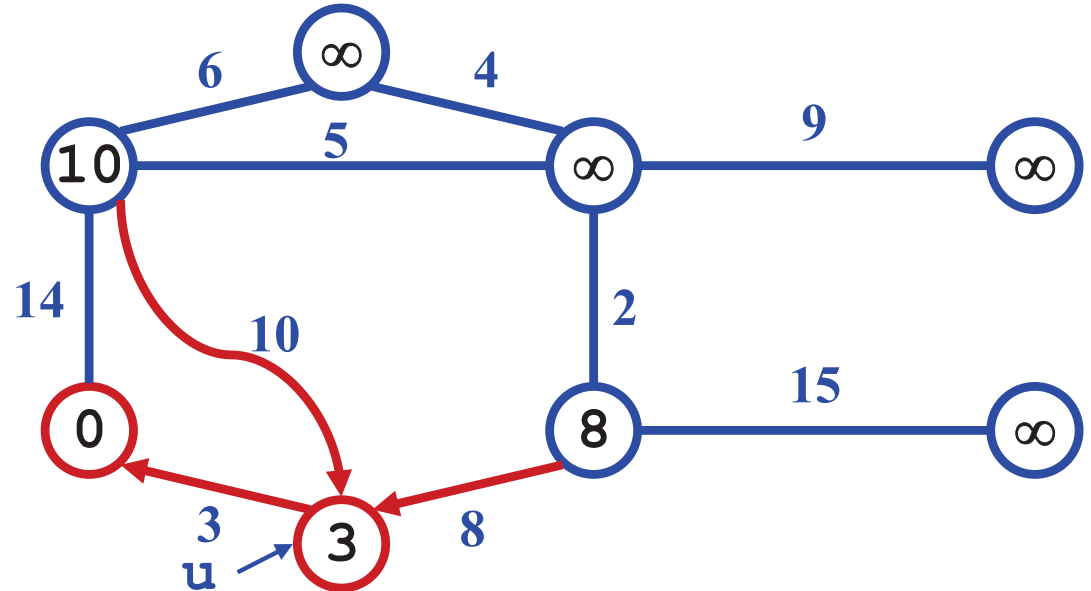


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             $p[v] = u$ ;
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```



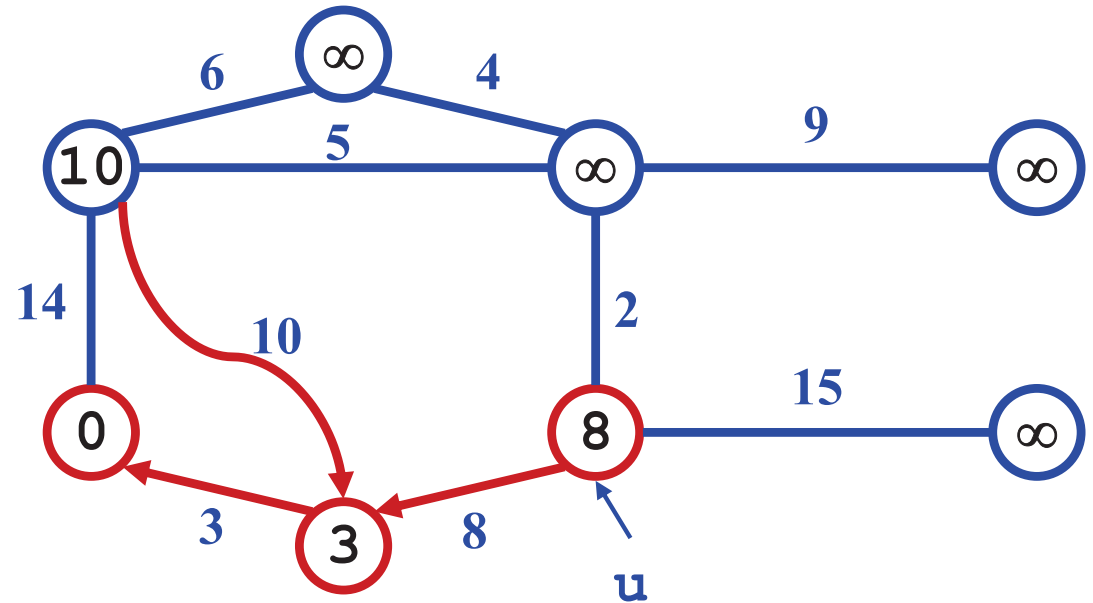


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```

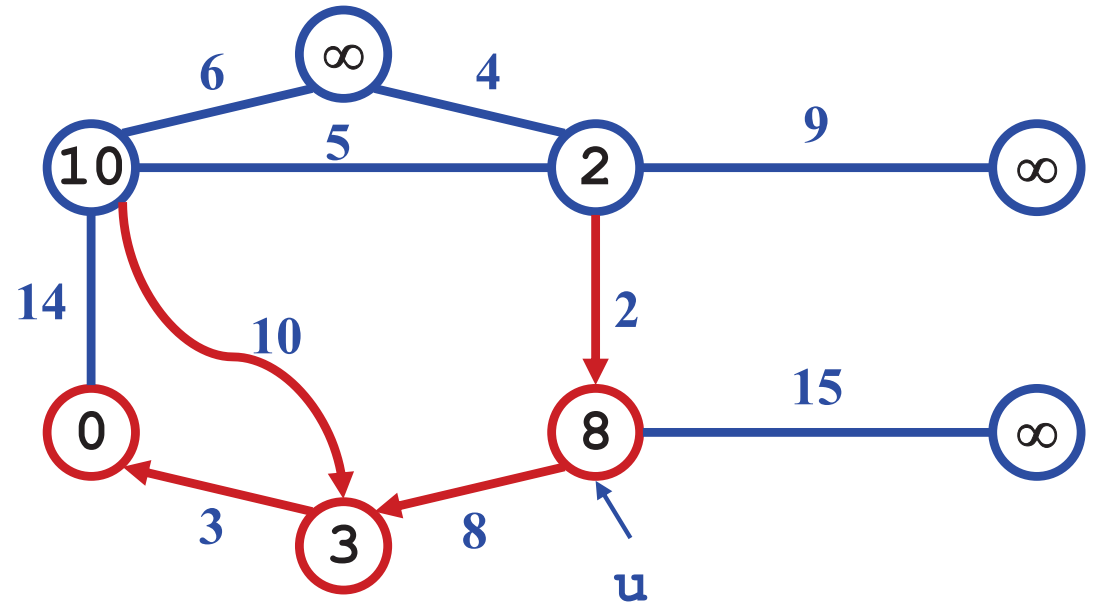


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             $p[v] = u$ ;
             $\text{weight}[v] = w(u, v)$ ;
```

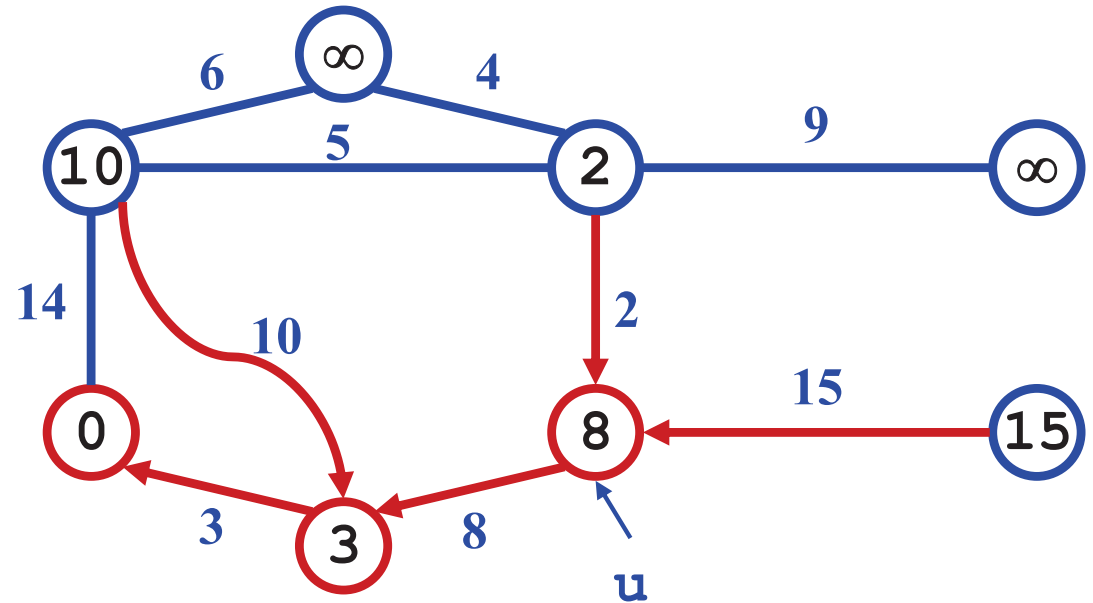


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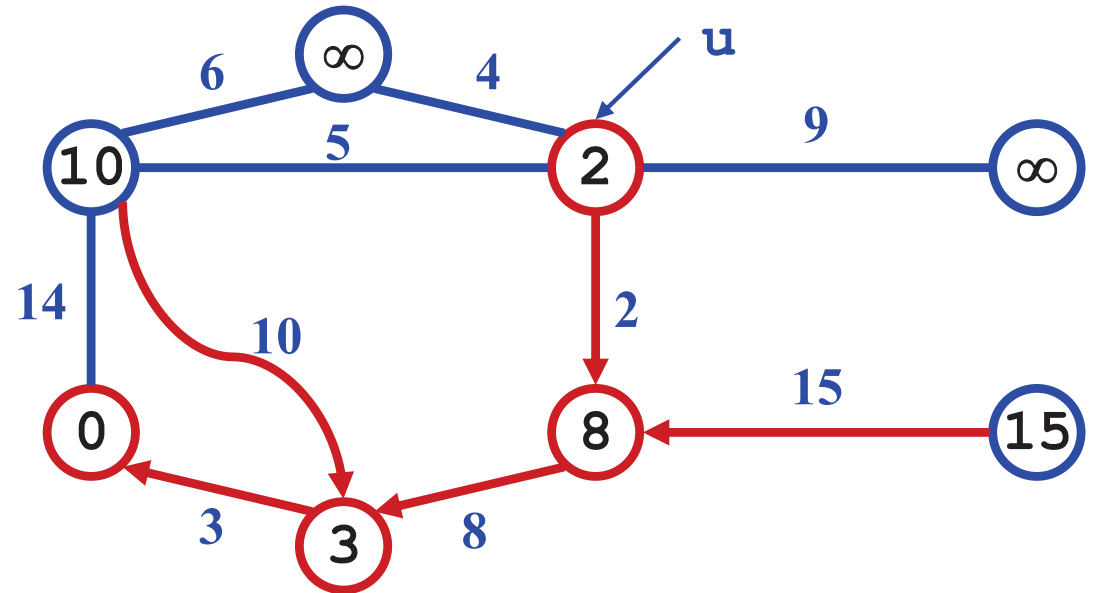


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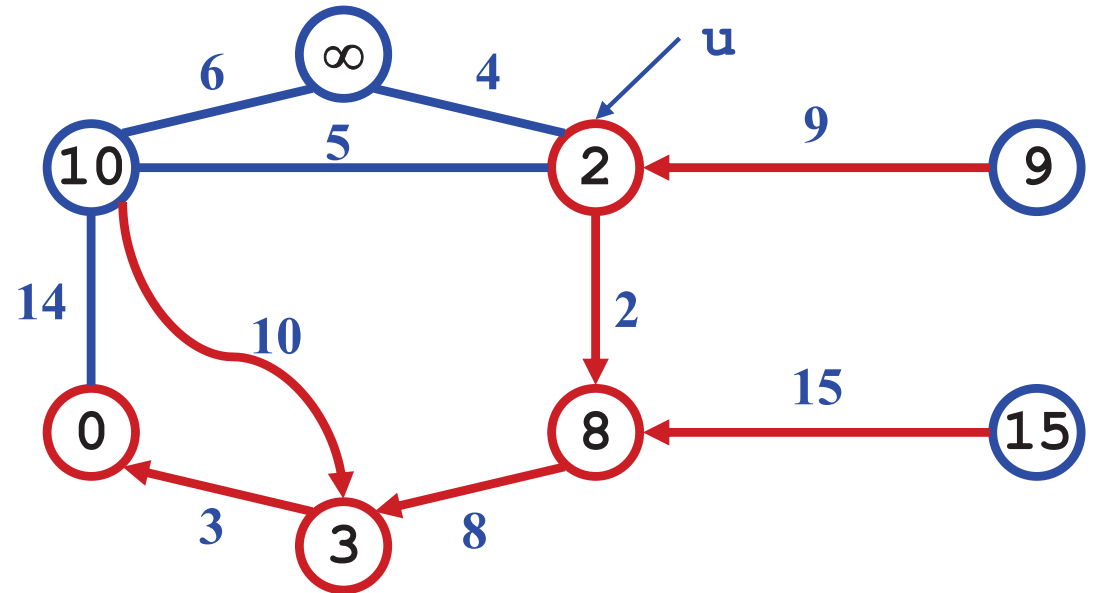


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```

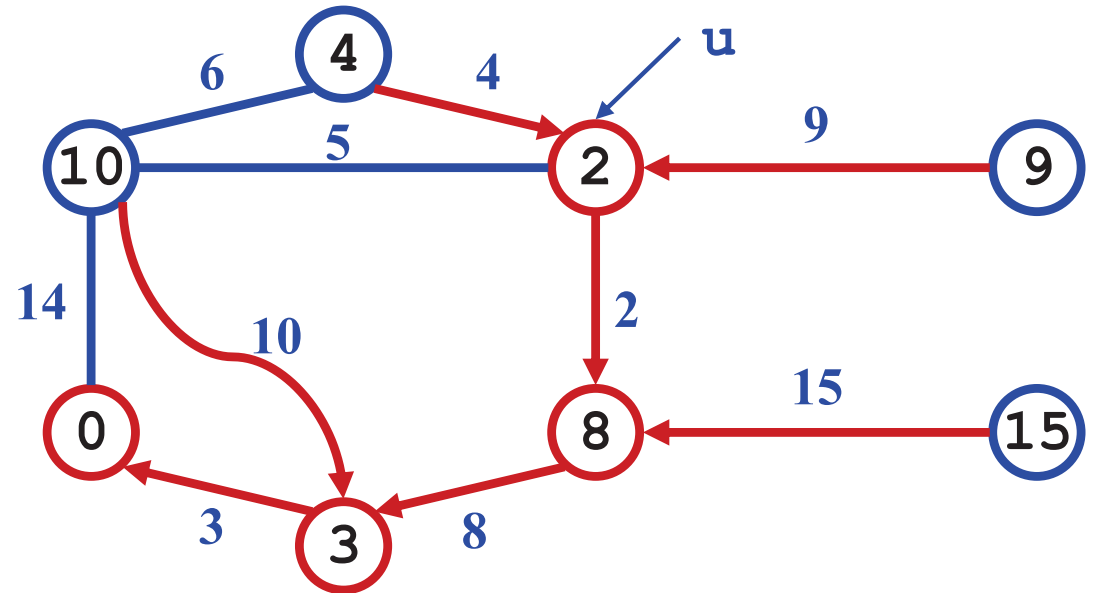


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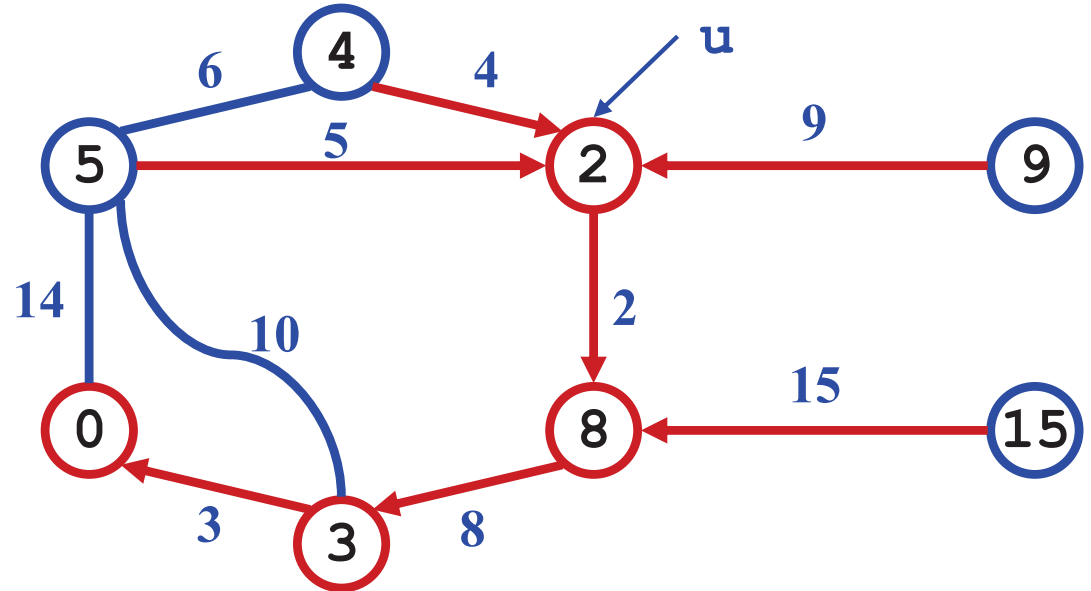


# Prim's Algorithm

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MST-Prim

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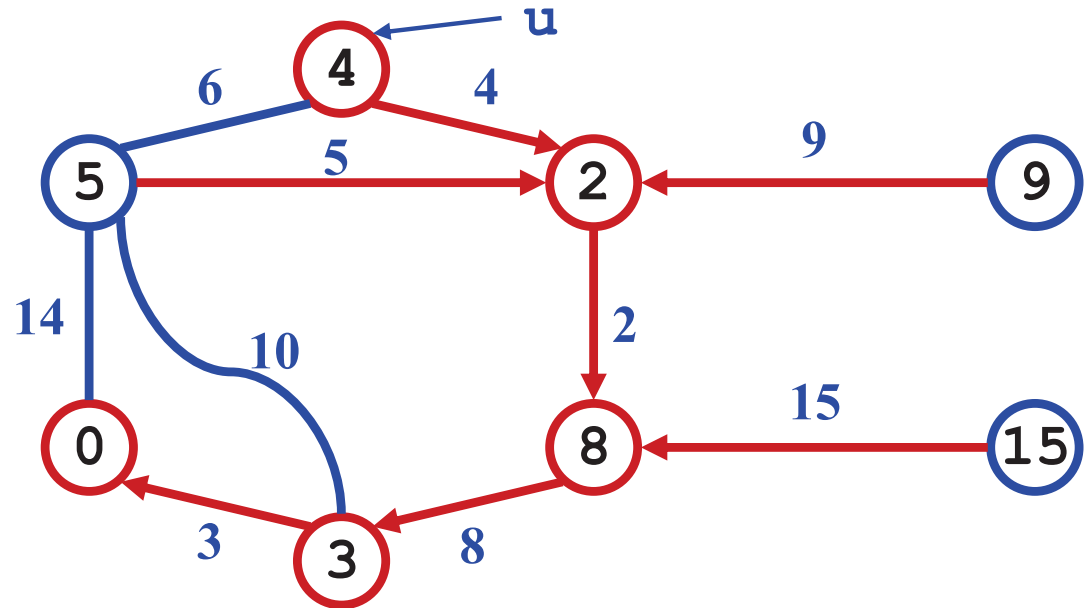


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MST-Prim

```
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```



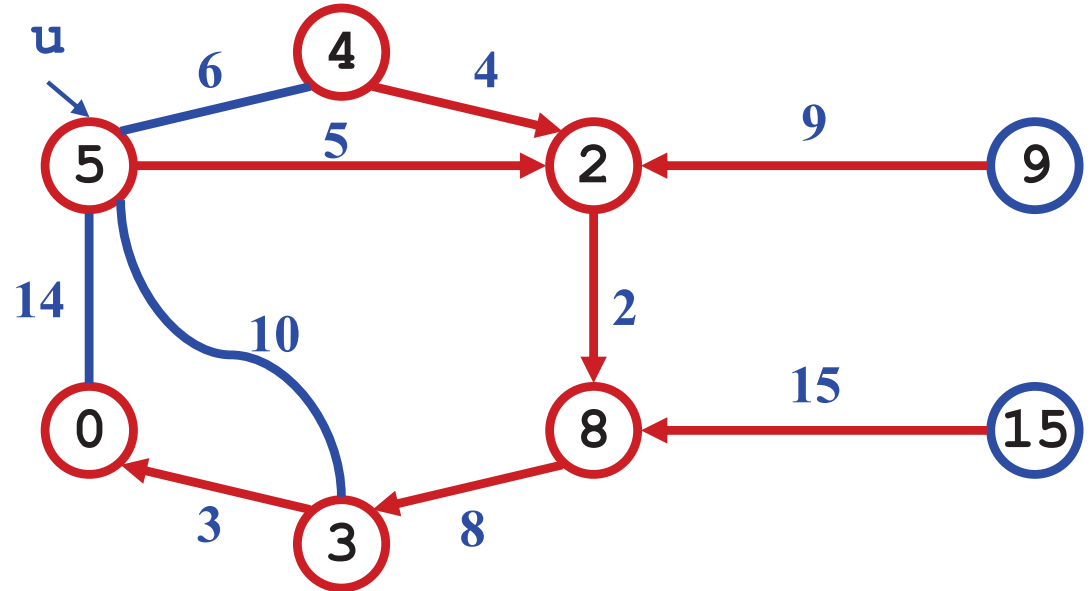


# Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

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```

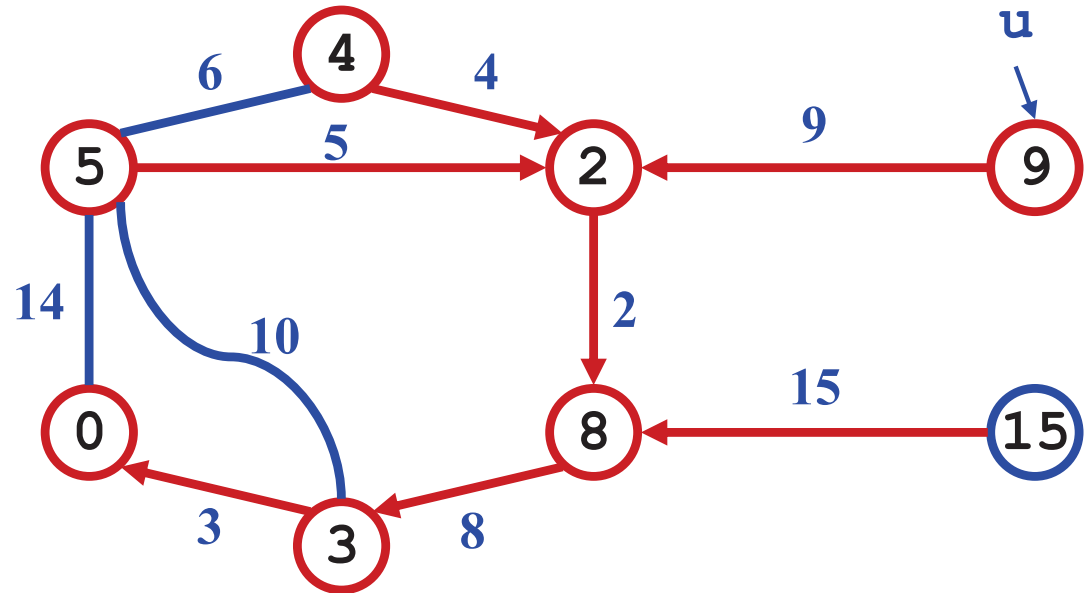


# Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
for all node  $u$ 
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        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
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```

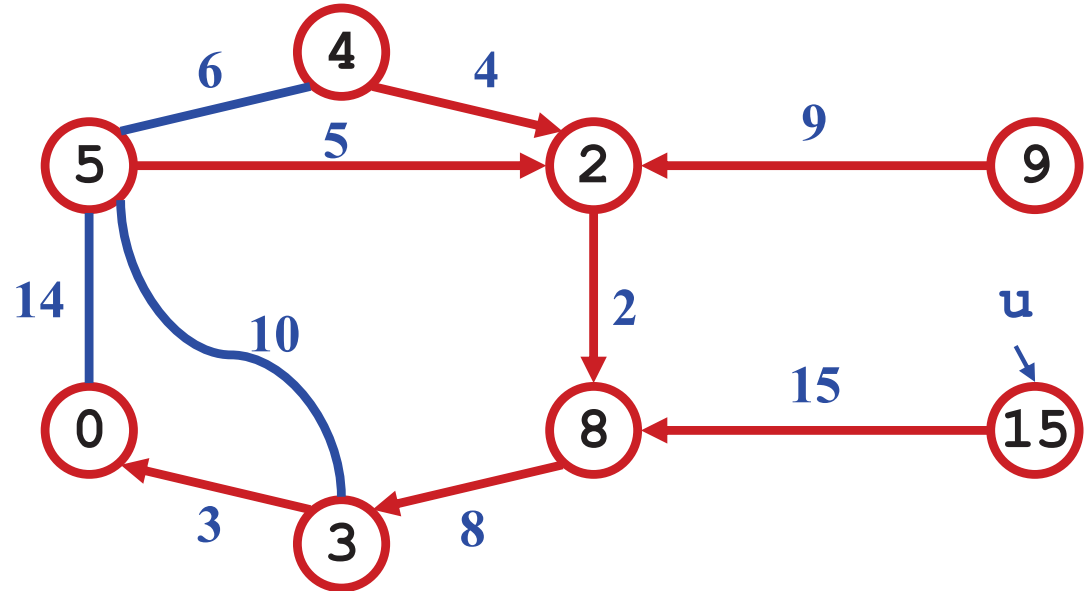


# Prim's Algorithm

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MST-Prim

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    for each  $v \in \text{Adjacent}[u]$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
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```



# Prim's Algorithm

