

Longest Common Subsequence (LCS)

Subsequence:

Subset of letters in order from left to right.

Remember:

“Subsequence” is different from “**substring**”. Substring is continuous, no gap. But, subsequence can have gap.

Example : Subsequence

$X = \{A B C B D A B\}$

$Y = \{B D C A B A\}$

- CBD is a substring of X. BBA is a subsequence of X, but not a substring of X.
- DBBB is not a subsequence of X, because order violation.
- BDAB is both substring and subsequence of X, but it is only a subsequence of Y
- Substrings are also subsequence, but a subsequence may not be a substring.

Example : LCS

Common Subsequence:

Subsequence of both.

Longest Common Subsequence:

Same subsequence of maximum length.

$X = \{A B C B D A B\}$

$Y = \{B D C A B A\}$

➤ $X = A \text{BCBD} AB$, $Y = BD \text{CA} BA$ **LCS:** BCBA (4)

➤ Other examples: BCAB (4), BDAB (4).

➤ Maximum possible length = 4

Longest Common Subsequence Problem

Problem:

Given two sequences $X = \langle x_1, x_2, \dots, x_m \rangle$ and
 $Y = \langle y_1, y_2, \dots, y_n \rangle$ find a
maximum-length common subsequence of X and Y .

Solution (Idea):

Notice that the LCS problem has *optimal substructure*:
solutions of subproblems are parts of the final solution.

Subproblems: “Find LCS of pairs of *prefixes* of X and Y ”

Notations

- $X = \langle x_1, x_2, \dots, x_m \rangle$ and $Y = \langle y_1, y_2, \dots, y_n \rangle$: Sequences
- X_i, Y_j : Prefixes of X and Y of length i and j respectively
- $c[i,j]$: The length of LCS of X_i and Y_j
- $c[m,n]$: The length of LCS of X and Y

Optimal Substructure of an LCS

- Let $X = \langle x_1, x_2, \dots, x_m \rangle$ and $Y = \langle y_1, y_2, \dots, y_n \rangle$ be sequences and let $Z = \langle z_1, z_2, \dots, z_k \rangle$ be any LCS of X and Y .
- If $x_m = y_n$, then $z_k = x_m = y_n$ and Z_{k-1} is an LCS of X_{m-1} and Y_{n-1}
- If $x_m \neq y_n$, then $z_k \neq x_m$
implies that Z is an LCS of X_{m-1} and Y
- If $x_m \neq y_n$, then $z_k \neq y_n$
implies that Z is an LCS of X and Y_{n-1}

Optimal Substructure of an LCS (Book: Cormen; Page-392)

Let $X = \langle x_1, x_2, \dots, x_m \rangle$ and $Y = \langle y_1, y_2, \dots, y_n \rangle$ be sequences, and let $Z = \langle z_1, z_2, \dots, z_k \rangle$ be any LCS of X and Y .

1. If $x_m = y_n$, then $z_k = x_m = y_n$ and Z_{k-1} is an LCS of X_{m-1} and Y_{n-1} .
2. If $x_m \neq y_n$, then $z_k \neq x_m$ implies that Z is an LCS of X_{m-1} and Y .
3. If $x_m \neq y_n$, then $z_k \neq y_n$ implies that Z is an LCS of X and Y_{n-1} .

Proof (1) If $z_k \neq x_m$, then we could append $x_m = y_n$ to Z to obtain a common subsequence of X and Y of length $k + 1$, contradicting the supposition that Z is a *longest* common subsequence of X and Y . Thus, we must have $z_k = x_m = y_n$. Now, the prefix Z_{k-1} is a length- $(k - 1)$ common subsequence of X_{m-1} and Y_{n-1} . We wish to show that it is an LCS. Suppose for the purpose of contradiction that there exists a common subsequence W of X_{m-1} and Y_{n-1} with length greater than $k - 1$. Then, appending $x_m = y_n$ to W produces a common subsequence of X and Y whose length is greater than k , which is a contradiction.

(2) If $z_k \neq x_m$, then Z is a common subsequence of X_{m-1} and Y . If there were a common subsequence W of X_{m-1} and Y with length greater than k , then W would also be a common subsequence of X_m and Y , contradicting the assumption that Z is an LCS of X and Y .

(3) The proof is symmetric to (2). ■

The Recurrence

$$c[i, j] = \begin{cases} 0 & \text{if } i = 0 \text{ or } j = 0 \\ c[i-1, j-1] + 1 & \text{if } i, j > 0 \text{ and } x_i = y_j, \\ \max(c[i, j-1], c[i-1, j]) & \text{if } i, j > 0 \text{ and } x_i \neq y_j \end{cases}$$

LCS Recursive Solution

- We start with $i = j = 0$ (empty substrings of x and y)
- Since X_0 and Y_0 are empty strings, their LCS is always empty (i.e. $c[0,0] = 0$)
- LCS of empty string and any other string is empty, so for every i and j : $c[0, j] = c[i, 0] = 0$

LCS Recursive Solution

- When we calculate $c[i,j]$, we consider two cases:
- **First case:** $x[i]=y[j]$: one more symbol matches, so the length of LCS X_i and Y_j equals to the length of LCS of smaller strings X_{i-1} and Y_{j-1} , plus 1
 - **Second case:** $x[i] \neq y[j]$: As symbols do not match, solution is not improved, and the length of $\text{LCS}(X_i, Y_j)$ is the same as before (maximum of $\text{LCS}(X_i, Y_{j-1})$ and $\text{LCS}(X_{i-1}, Y_j)$)

LCS Example

What is the Longest Common Subsequence of X and Y ?

$X = \text{ABCB}$

$Y = \text{BDCAB}$

LCS Example (0)

		j	0	1	2	3	4	5
			Y _j	B	D	C	A	B
i								
0	X _i							
1	A							
2	B							
3	C							
4	B							

$X = \text{ABCB}; \quad m = |X| = 4$

$Y = \text{BDCAB}; \quad n = |Y| = 5$

Allocate array $c[6,5]$

LCS Example (1)

		j	0	1	2	3	4	5
			Y _j	B	D	C	A	B
i	X _i							
0								
1	A		0					
2	B		0					
3	C		0					
4	B		0					

for $i = 1$ to m $c[i,0] = 0$

LCS Example (2)

		j	0	1	2	3	4	5
			Y _j	B	D	C	A	B
i	X _i	0	0	0	0	0	0	0
1	A	0						
2	B	0						
3	C	0						
4	B	0						

for $j = 0$ to n $c[0,j] = 0$

LCS Example (3)

		j	0	1	2	3	4	5
			Y _j	B	D	C	A	B
i	X _i							
0			0	0	0	0	0	0
1	A		0	0 ↑				
2	B		0					
3	C		0					
4	B		0					

case $i=1$ and $j=1$

$A \neq B$

but, $c[0,1] \geq c[1,0]$

so $c[1,1] = c[0,1]$, and $b[1,1] = \uparrow$

LCS Example (4)

		j	0	1	2	3	4	5
			Y _j	B	D	C	A	B
i								
0	X _i		0	0	0	0	0	0
1	A	0	0	0 ↑	0 ↑			
2	B							
3	C							
4	B							

case $i=1$ and $j=2$

$A \neq D$

but, $c[0,2] \geq c[1,1]$

so $c[1,2] = c[0,2]$, and $b[1,2] = \uparrow$

LCS Example (5)

		j	0	1	2	3	4	5
			Y _j	B	D	C	A	B
i								
0	X _i		0	0	0	0	0	0
1	A	0	0	0 ↑	0 ↑	0 ↑		
2	B	0						
3	C	0						
4	B	0						

case $i=1$ and $j=3$

$A \neq C$

but, $c[0,3] \geq c[1,2]$

so $c[1,3] = c[0,3]$, and $b[1,3] = \uparrow$

LCS Example (6)

		j	0	1	2	3	4	5
			Y _j	B	D	C	A	B
i	X _i	0	0	0	0	0	0	0
1	A	0	0	0 ↑	0 ↑	0 ↑	1 ↖	
2	B	0						
3	C	0						
4	B	0						

case $i=1$ and $j=4$

$A = A$

so $c[1,4] = c[0,2] + 1$, and $b[1,4] = \nwarrow$

LCS Example (7)

		j	0	1	2	3	4	5
			Y _j	B	D	C	A	B
i	X _i							
0			0	0	0	0	0	0
1	A		0	0 ↑	0 ↑	0 ↑	1 ↖	1 ←
2	B		0					
3	C		0					
4	B		0					

case $i=1$ and $j=5$

$A \neq B$

this time $c[0,5] < c[1,4]$

so $c[1,5] = c[1,4]$, and $b[1,5] = \leftarrow$

LCS Example (8)

		j	0	1	2	3	4	5
			Y _j	B	D	C	A	B
i	X _i							
0			0	0	0	0	0	0
1	A		0	0 ↑	0 ↑	0 ↑	1 ↘	1 ←
2	B		0	1 ↘				
3	C		0					
4	B		0					

case $i=2$ and $j=1$

$B = B$

so $c[2, 1] = c[1, 0] + 1$, and $b[2, 1] = \nwarrow$

LCS Example (9)

		j	0	1	2	3	4	5
			Y _j	B	D	C	A	B
i	X _i							
0			0	0	0	0	0	0
1	A		0	0 ↑	0 ↑	0 ↑	1 ↘	1 ←
2	B		0	1 ↘	1 ←			
3	C		0					
4	B		0					

case $i=2$ and $j=2$

$B \neq D$

and $c[1, 2] < c[2, 1]$

so $c[2, 2] = c[2, 1]$ and $b[2, 2] = \leftarrow$

LCS Example (10)

		j	0	1	2	3	4	5
			Y _j	B	D	C	A	B
i								
0	X _i		0	0	0	0	0	0
1	A		0	0 ↑	0 ↑	0 ↑	1 ↘	1 ←
2	B		0	1 ↘	1 ←	1 ←		
3	C		0					
4	B		0					

case $i=2$ and $j=3$

$B \neq D$

and $c[1, 3] < c[2, 2]$

so $c[2, 3] = c[2, 2]$ and $b[2, 3] = \leftarrow$

LCS Example (11)

		j	0	1	2	3	4	5
			Y _j	B	D	C	A	B
i								
0	X _i		0	0	0	0	0	0
1	A		0	0 ↑	0 ↑	0 ↑	1 ↘	1 ←
2	B		0	1 ↘	1 ←	1 ←	1 ↑	
3	C		0					
4	B		0					

case $i=2$ and $j=4$

$B \neq A$

and $c[1, 4] = c[2, 3]$

so $c[2, 4] = c[1, 4]$ and $b[2, 2] = \uparrow$

LCS Example (12)

		j	0	1	2	3	4	5
		Yj		B	D	C	A	B
i	Xi							
0			0	0	0	0	0	0
1	A		0	0 ↑	0 ↑	0 ↑	1 ↘	1 ←
2	B		0	1 ↘	1 ←	1 ←	1 ↑	2 ↘
3	C		0					
4	B		0					

case $i=2$ and $j=5$

$B = B$

so $c[2, 5] = c[1, 4] + 1$ and $b[2, 5] = \nwarrow$

LCS Example (13)

		j	0	1	2	3	4	5
			Yj	B	D	C	A	B
i	Xi	0	0	0	0	0	0	0
1	A	0	0	0 ↑	0 ↑	0 ↑	1 ↘	1 ←
2	B	0	1	1 ↘	1 ←	1 ←	1 ↑	2 ↘
3	C	0	1	1 ↑				
4	B	0						

case $i=3$ and $j=1$

$C \neq B$

and $c[2, 1] > c[3, 0]$

so $c[3, 1] = c[2, 1]$ and $b[3, 1] = \uparrow$

LCS Example (14)

		j	0	1	2	3	4	5
			Y _j	B	D	C	A	B
i								
0	X _i		0	0	0	0	0	0
1	A		0	0 ↑	0 ↑	0 ↑	1 ↘	1 ←
2	B		0	1 ↘	1 ←	1 ←	1 ↑	2 ↘
3	C		0	1 ↑	1 ↑			
4	B		0					

case $i=3$ and $j=2$

$C \neq D$

and $c[2, 2] = c[3, 1]$

so $c[3, 2] = c[2, 2]$ and $b[3, 2] = \uparrow$

LCS Example (15)

		j	0	1	2	3	4	5
			Y _j	B	D	C	A	B
i								
0	X _i		0	0	0	0	0	0
1	A		0	0 ↑	0 ↑	0 ↑	1 ↘	1 ←
2	B		0	1 ↘	1 ←	1 ←	1 ↑	2 ↘
3	C		0	1 ↑	1 ↑	2 ↘		
4	B		0					

case $i=3$ and $j=3$

$C = C$

so $c[3, 3] = c[2, 2] + 1$ and $b[3, 3] = \nwarrow$

LCS Example (16)

		j	0	1	2	3	4	5
			Y _j	B	D	C	A	B
i	X _i							
0			0	0	0	0	0	0
1	A		0	0 ↑	0 ↑	0 ↑	1 ↘	1 ←
2	B		0	1 ↘	1 ←	1 ←	1 ↑	2 ↘
3	C		0	1 ↑	1 ↑	2 ↘	2 ←	
4	B		0					

case $i=3$ and $j=4$

$C \neq A$

$c[2, 4] < c[3, 3]$

so $c[3, 4] = c[3, 3]$ and $b[3, 3] =$

LCS Example (17)

		j	0	1	2	3	4	5
			Y _j	B	D	C	A	B
i	X _i							
0		0	0	0	0	0	0	0
1	A	0	0 ↑	0 ↑	0 ↑	1 ↖	1 ←	
2	B	0	1 ↖	1 ←	1 ←	1 ↑	2 ↖	
3	C	0	1 ↑	1 ↑	2 ↖	2 ←	2 ↑	
4	B	0						

case $i=3$ and $j=5$

$C \neq B$

$c[2, 5] = c[3, 4]$

so $c[3, 5] = c[2, 5]$ and $b[3, 5] = \uparrow$

LCS Example (18)

		j					
		0	1	2	3	4	5
i	Yj		B	D	C	A	B
	Xi	0	0	0	0	0	0
	A	0	0 ↑	0 ↑	0 ↑	1 ↘	1 ←
	B	0	1 ↘	1 ←	1 ←	1 ↑	2 ↘
	C	0	1 ↑	1 ↑	2 ↘	2 ←	2 ↑
4	B	0	1 ↘				

case $i=4$ and $j=1$

$B = B$

so $c[4, 1] = c[3, 0] + 1$ and $b[4, 1] = \nwarrow$

LCS Example (19)

		j	0	1	2	3	4	5
			Y _j	B	D	C	A	B
i	X _i							
0			0	0	0	0	0	0
1	A		0	0 ↑	0 ↑	0 ↑	1 ↘	1 ←
2	B		0	1 ↘	1 ←	1 ←	1 ↑	2 ↘
3	C		0	1 ↑	1 ↑	2 ↘	2 ←	2 ↑
4	B		0	1 ↘	1 ↑			

case $i=4$ and $j=2$

$B \neq D$

$c[3, 2] = c[4, 1]$

so $c[4, 2] = c[3, 2]$ and $b[4, 2] = \uparrow$

LCS Example (20)

		j	0	1	2	3	4	5
			Yj	B	D	C	A	B
i	Xi							
0			0	0	0	0	0	0
1	A		0	0 ↑	0 ↑	0 ↑	1 ↘	1 ←
2	B		0	1 ↘	1 ←	1 ←	1 ↑	2 ↘
3	C		0	1 ↑	1 ↑	2 ↘	2 ←	2 ↑
4	B		0	1 ↘	1 ↑	2 ↑		

case $i=4$ and $j=3$

$B \neq C$

$c[3, 3] > c[4, 2]$

so $c[4, 3] = c[3, 3]$ and $b[4, 3] = \uparrow$

LCS Example (21)

		j	0	1	2	3	4	5
			Yj	B	D	C	A	B
i	Xi							
0			0	0	0	0	0	0
1	A		0	0 ↑	0 ↑	0 ↑	1 ↘	1 ←
2	B		0	1 ↘	1 ←	1 ←	1 ↑	2 ↘
3	C		0	1 ↑	1 ↑	2 ↘	2 ←	2 ↑
4	B		0	1 ↘	1 ↑	2 ↑	2 ↑	

case $i=4$ and $j=4$

$B \neq A$

$c[3, 4] = c[4, 3]$

so $c[4, 4] = c[3, 4]$ and $b[3, 5] = \uparrow$

LCS Example (22)

		j	0	1	2	3	4	5
			Y _j	B	D	C	A	B
i								
0	X _i		0	0	0	0	0	0
1	A		0	0 ↑	0 ↑	0 ↑	1 ↘	1 ←
2	B		0	1 ↘	1 ←	1 ←	1 ↑	2 ↘
3	C		0	1 ↑	1 ↑	2 ↘	2 ←	2 ↑
4	B		0	1 ↘	1 ↑	2 ↑	2 ↑	3 ↘

case $i=4$ and $j=5$

$B = B$

so $c[4, 5] = c[3, 4] + 1$ and $b[4, 5] = \nwarrow$

Algorithm: LCS

```
LCS-Length(X, Y)
m = length(X), n = length(Y)
for i = 1 to m
    c[i, 0] = 0
for j = 0 to n
    c[0, j] = 0
for i = 1 to m
    for j = 1 to n
        if (  $x_i = y_j$  )
            c[i, j] = c[i - 1, j - 1] + 1
            b[i, j] = "↖"
        else if c[i - 1, j] >= c[i, j - 1]
            c[i, j] = c[i - 1, j]
            b[i, j] = "↑"
        else c[i, j] = c[i, j - 1]
            b[i, j] = "←"
return Paris c and b
```

Finding LCS

		j	0	1	2	3	4	5
i		Y _j	B	D	C	A	B	
0	X_i	0	0	0	0	0	0	
1	A	0	0	0	0	1	1	
2	B	0	1	1	1	1	2	
3	C	0	1	1	2	2	2	
4	B	0	1	1	2	2	3	

Finding LCS (2)

		j	0	1	2	3	4	5
		Yj		B	D	C	A	B
i	Xi							
0			0	0	0	0	0	0
1	A		0	0	0	0	1	1
2	B		0	1	1	1	1	2
3	C		0	1	1	2	2	2
4	B		0	1	1	2	2	3

LCS (reversed order): **B C B**

LCS (straight order): **B C B**

Algorithm: Finding LCS

```
PRINT-LCS( $b, X, i, j$ )  
  if  $i == 0$  or  $j == 0$   
    return  
  if  $b[i, j] == \nwarrow$   
    PRINT-LCS( $b, X, i - 1, j - 1$ )  
    print  $x_i$   
  elseif  $b[i, j] == \uparrow$   
    PRINT-LCS( $b, X, i - 1, j$ )  
  else PRINT-LCS( $b, X, i, j - 1$ )
```

Practice Example (Book: Cormen; Page-395)

		<i>j</i>	0	1	2	3	4	5	6
		<i>y_j</i>	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>	
<i>i</i>	<i>x_i</i>								
0									
1	<i>A</i>	0	0	0	0	0	0	0	
2	<i>B</i>	0	1	1	1	1	2	2	
3	<i>C</i>	0	1	1	2	2	2	2	
4	<i>B</i>	0	1	1	2	2	3	3	
5	<i>D</i>	0	1	2	2	2	3	3	
6	<i>A</i>	0	1	2	2	3	3	4	
7	<i>B</i>	0	1	2	2	3	4	4	

Paris

LCS : BCBA

Length : 4

39

Thank You

Stay Safe