Longest Common Subsequence (LCS)

Subsequence:

Subset of letters in order from left to right.

Remember:

"Subsequence" is different from "substring". Substring is continuous, no gap. But, subsequence can have gap.

Example: Subsequence

$$X=\{A B C B D A B\}$$
 $Y=\{B D C A B A\}$

- CBD is a substring of X. BBA is a subsequence of X, but not a substring of X.
- DBBB is not a subsequence of X, because order violation.
- ➤ BDAB is both substring and subsequence of X, but it is only a subsequence of Y
- Substrings are also subsequence, but a subsequence may not be a substring.

Example: LCS

Common Subsequence:

Subsequence of both.

Longest Common Subsequence:

Same subsequence of maximum length.

$$X = \{A B C B D A B\}$$

$$Y = \{B D C A B A\}$$

- \rightarrow X = A BCBD AB, Y = BD CA BA
- LCS: BCBA (4)
- > Other examples: BCAB (4), BDAB (4).
- Maximum possible length = 4

Longest Common Subsequence Problem

Problem:

Given two sequences $X = \langle x_1, x_2,, x_m \rangle$ and $Y = \langle y_1, y_2,, y_n \rangle$ find a maximum-length common subsequence of X and Y.

Solution (Idea):

Notice that the LCS problem has *optimal substructure*: solutions of subproblems are parts of the final solution.

Subproblems: "Find LCS of pairs of *prefixes* of *X* and *Y*"

Notations

• $X = \langle x_1, x_2, ..., x_m \rangle$ and $Y = \langle y_1, y_2, ..., y_n \rangle$: Sequences

• X_i , Y_j : Prefixes of X and Y of length i and j respectively

• c[i,j]: The length of LCS of X_i and Y_j

• c[m,n]: The length of LCS of X and Y

Optimal Substructure of an LCS

- Let $X = \langle x_1, x_2, \dots, x_m \rangle$ and $Y = \langle y_1, y_2, \dots, y_n \rangle$ be sequences and let $Z = \langle z_1, z_2, \dots, z_k \rangle$ be any LCS of X and Y.
- If $x_m = y_n$, then $z_k = x_m = y_n$ and Z_{k-1} is an LCS of X_{m-1} and Y_{n-1}
- If $x_m \neq y_n$, then $z_k \neq x_m$ implies that Z is an LCS of X_{m-1} and Y
- If $x_m \neq y_n$, then $z_k \neq y_n$ implies that Z is an LCS of X and Y_{n-1}

Optimal Substructure of an LCS (Book: Cormen; Page-392)

Let $X = (x_1, x_2, ..., x_m)$ and $Y = (y_1, y_2, ..., y_n)$ be sequences, and let $Z = (z_1, z_2, ..., z_k)$ be any LCS of X and Y.

- 1. If $x_m = y_n$, then $z_k = x_m = y_n$ and Z_{k-1} is an LCS of X_{m-1} and Y_{n-1} .
- 2. If $x_m \neq y_n$, then $z_k \neq x_m$ implies that Z is an LCS of X_{m-1} and Y.
- 3. If $x_m \neq y_n$, then $z_k \neq y_n$ implies that Z is an LCS of X and Y_{n-1} .

Proof (1) If $z_k \neq x_m$, then we could append $x_m = y_n$ to Z to obtain a common subsequence of X and Y of length k+1, contradicting the supposition that Z is a longest common subsequence of X and Y. Thus, we must have $z_k = x_m = y_n$. Now, the prefix Z_{k-1} is a length-(k-1) common subsequence of X_{m-1} and Y_{n-1} . We wish to show that it is an LCS. Suppose for the purpose of contradiction that there exists a common subsequence W of X_{m-1} and Y_{n-1} with length greater than k-1. Then, appending $x_m = y_n$ to W produces a common subsequence of X and Y whose length is greater than k, which is a contradiction.

(2) If $z_k \neq x_m$, then Z is a common subsequence of X_{m-1} and Y. If there were a common subsequence W of X_{m-1} and Y with length greater than k, then W would also be a common subsequence of X_m and Y, contradicting the assumption that Z is an LCS of X and Y.

(3) The proof is symmetric to (2).

The Recurrence

$$c[i, j] = \begin{cases} 0 & \text{if } i = 0 \text{ or } j = 0 \\ c[i-1, j-1] + 1 & \text{if } i, j > 0 \text{ and } x_i = y_j, \\ \max(c[i, j-1], c[i-1, j]) & \text{if } i, j > 0 \text{ and } x_i \neq y_j \end{cases}$$

LCS Recursive Solution

- We start with i = j = 0 (empty substrings of x and y)
- Since X_0 and Y_0 are empty strings, their LCS is always empty (i.e. c[0,0]=0)
- LCS of empty string and any other string is empty, so for every i and j: c[0, j] = c[i, 0] = 0

LCS Recursive Solution

- \triangleright When we calculate c[i,j], we consider two cases:
- First case: x[i]=y[j]: one more symbol matches, so the length of LCS X_i and Y_j equals to the length of LCS of smaller strings X_{i-1} and Y_{i-1} , plus 1
- **Second case:** x[i] != y[j]: As symbols do not match, solution is not improved, and the length of $LCS(X_i, Y_j)$ is the same as before (maximum of $LCS(X_i, Y_{j-1})$ and $LCS(X_{i-1}, Y_j)$)

LCS Example

What is the Longest Common Subsequence of X and Y?

$$X = ABCB$$

 $Y = BDCAB$

LCS Example (0)

$$X = ABCB$$
; $m = |X| = 4$
 $Y = BDCAB$; $n = |Y| = 5$
Allocate array c[6,5]

LCS Example (1)

					_		
	j	0	1	2	3	4	5
i		Yj	В	D	C	A	В
0	Xi						
1	A	0					
2	В	0					
3	C	0					
4	В	0					

for
$$i = 1$$
 to m $c[i,0] = 0$

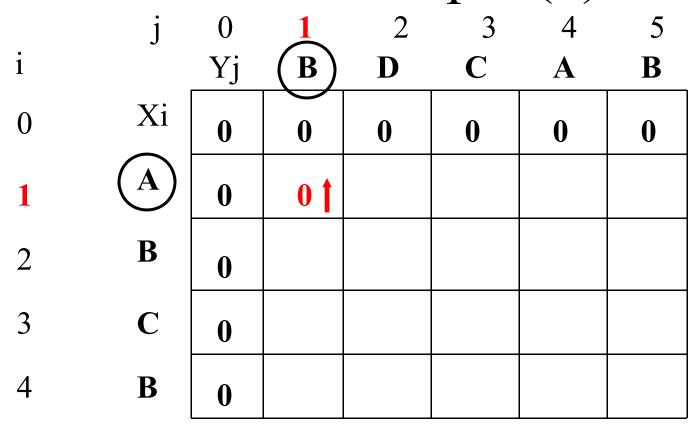
LCS Example (2)

	•	•	4			•	_
	J	0	1	2	3	4	5
i		Yj	В	D	C	A	В
0	Xi	0	0	0	0	0	0
1	A	0					
2	В	0					
3	C	0					
4	В	0					

for
$$j = 0$$
 to n $c[0,j] = 0$

$$c[0,j] = 0$$

LCS Example (3)



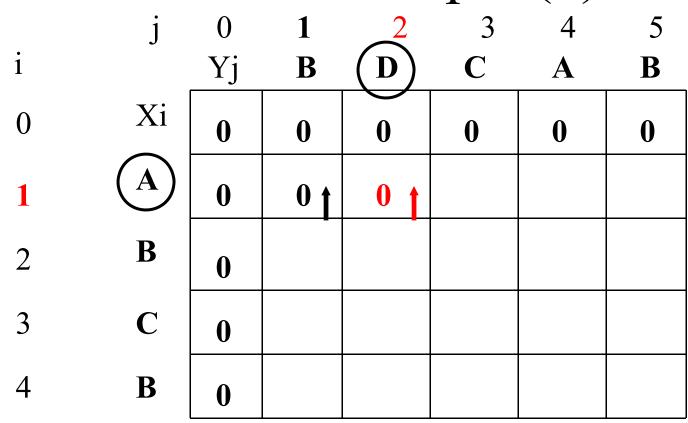
```
case i=1 and j=1

A!=B

but, c[0,1]>=c[1,0]

so c[1,1] = c[0,1], and b[1,1] = \uparrow
```

LCS Example (4)



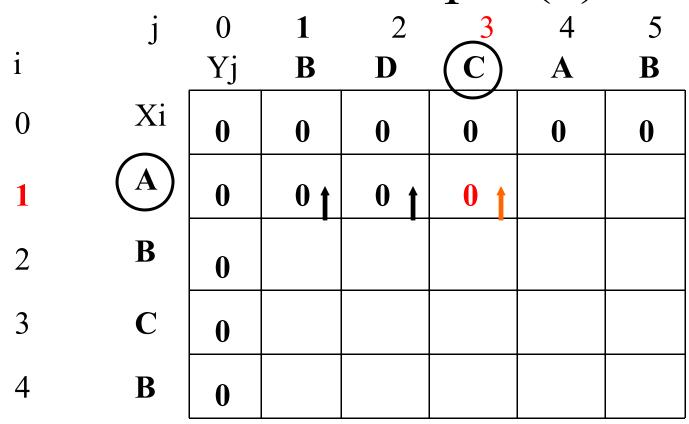
```
case i=1 and j=2

A!= D

but, c[0,2]>=c[1,1]

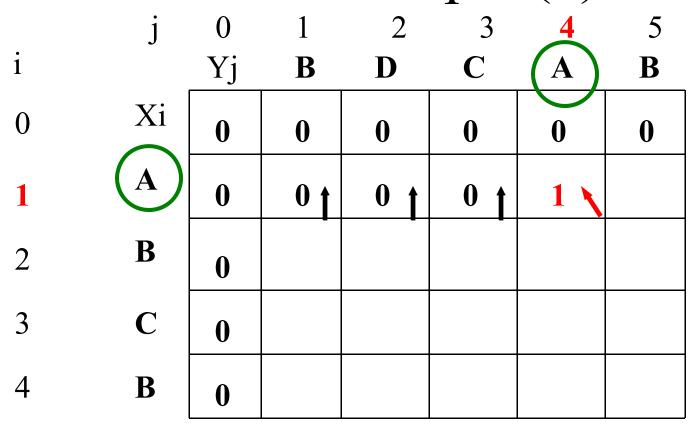
so c[1,2] = c[0,2], and b[1,2] = \uparrow
```

LCS Example (5)



case i=1 and j=3
A!= C
but, c[0,3]>=c[1,2]
so c[1,3] = c[0,3], and b[1,3] =
$$\uparrow$$

LCS Example (6)

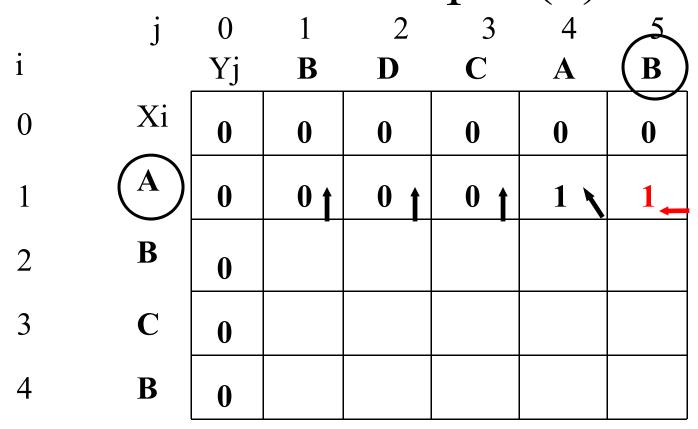


case i=1 and j=4

$$A = A$$

so c[1,4] = c[0,2]+1, and b[1,4] =

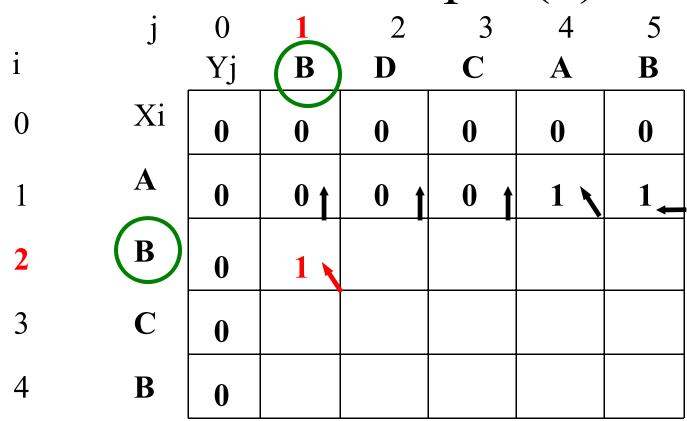
LCS Example (7)



case i=1 and j=5
A!= B
this time
$$c[0,5] < c[1,4]$$

so $c[1,5] = c[1,4]$, and $b[1,5] = \leftarrow$

LCS Example (8)

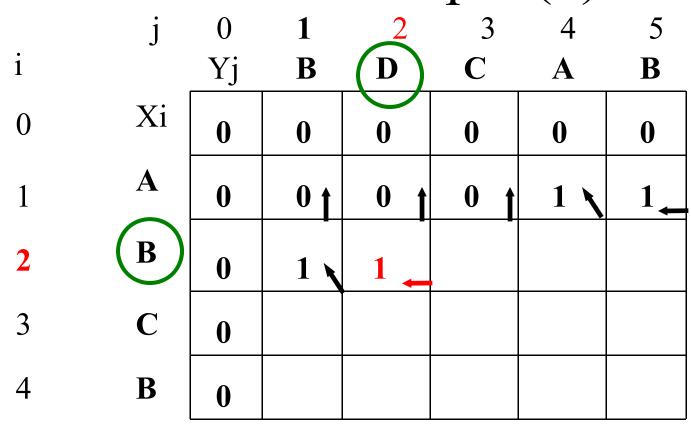


case i=2 and j=1

$$B = B$$

so c[2, 1] = c[1, 0]+1, and b[2, 1] =

LCS Example (9)



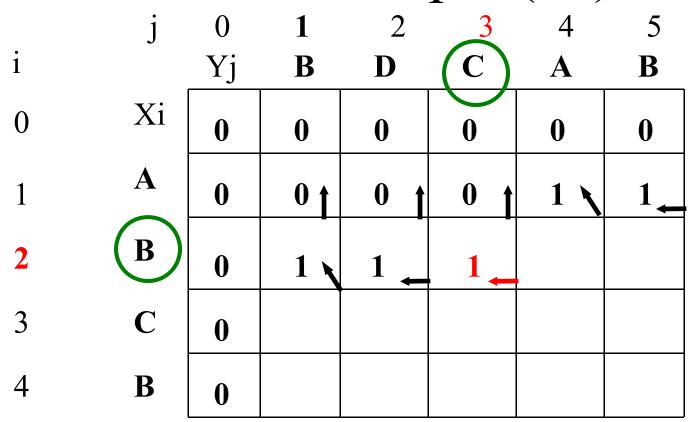
```
case i=2 and j=2

B!= D

and c[1, 2] < c[2, 1]

so c[2, 2] = c[2, 1] and b[2, 2] = \leftarrow
```

LCS Example (10)



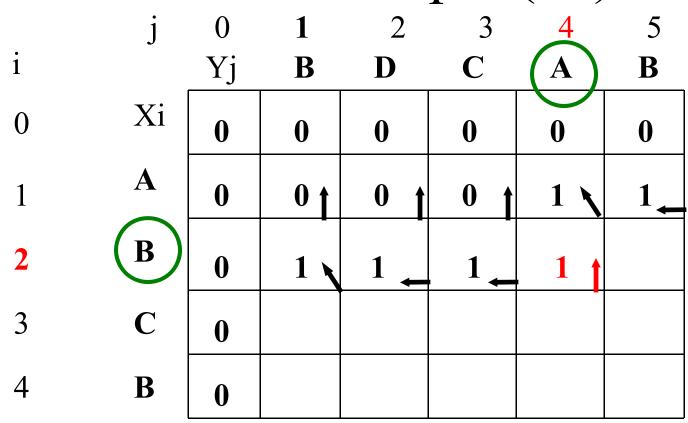
```
case i=2 and j=3

B!= D

and c[1, 3] < c[2, 2]

so c[2, 3] = c[2, 2] and b[2, 3] = \leftarrow
```

LCS Example (11)



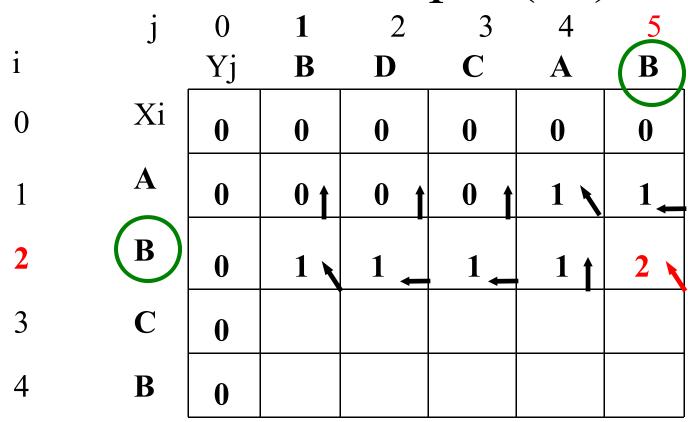
```
case i=2 and j=4

B!= A

and c[1, 4] = c[2, 3]

so c[2, 4] = c[1, 4] and b[2, 2] = \mathbf{1}
```

LCS Example (12)

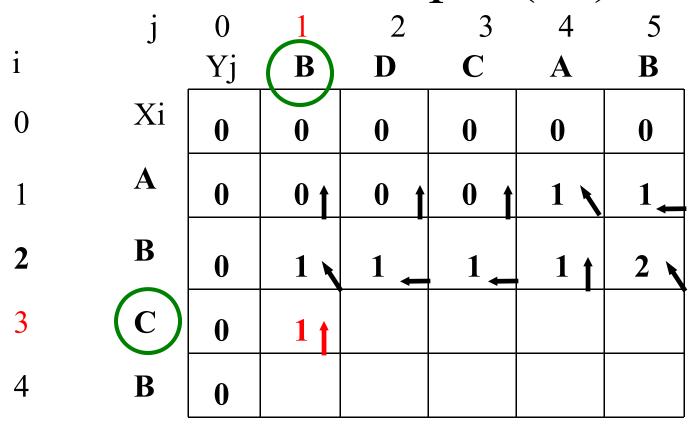


case i=2 and j=5

$$B = B$$

so c[2, 5] = c[1, 4]+1 and b[2, 5] =

LCS Example (13)



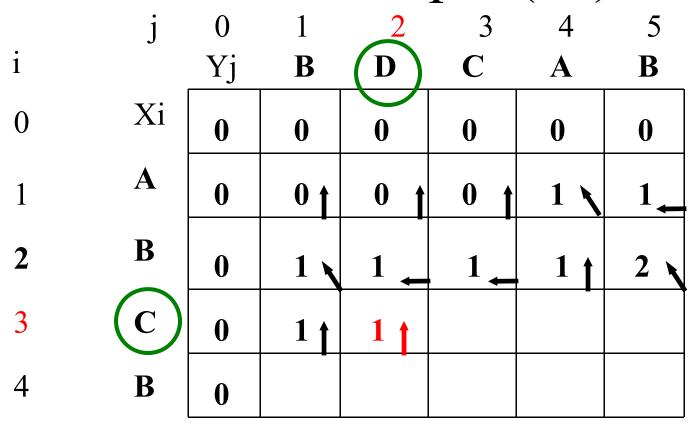
```
case i=3 and j=1

C != B

and c[2, 1] > c[3,0]

so c[3, 1] = c[2, 1] and b[3, 1] = 1
```

LCS Example (14)



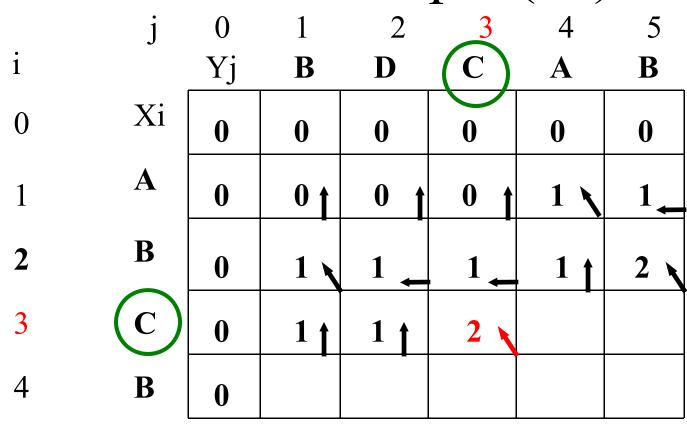
```
case i=3 and j=2

C != D

and c[2, 2] = c[3, 1]

so c[3, 2] = c[2, 2] and b[3, 2] = 1
```

LCS Example (15)

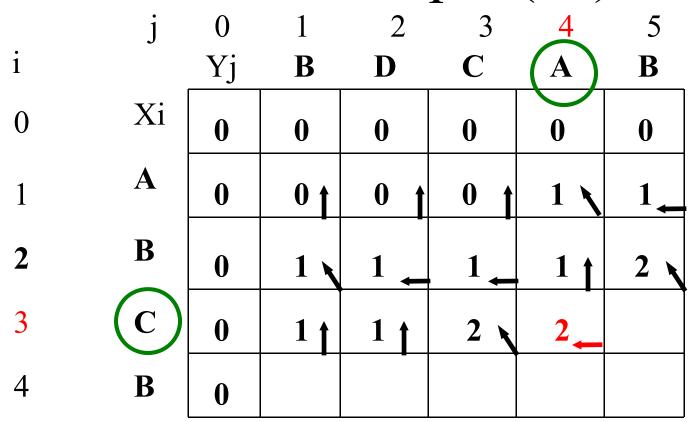


case i=3 and j=3

$$C = C$$

so c[3, 3] = c[2, 2]+1 and b[3, 3] =

LCS Example (16)

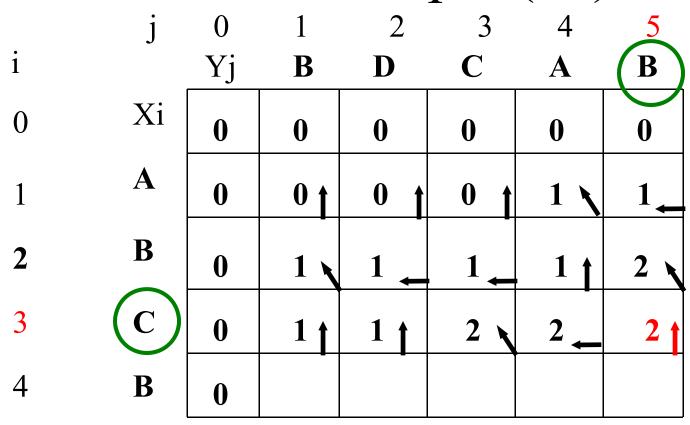


case i=3 and j=4

$$C != A$$

 $c[2, 4] < c[3, 3]$
so $c[3, 4] = c[3, 3]$ and $b[3, 3] =$

LCS Example (17)

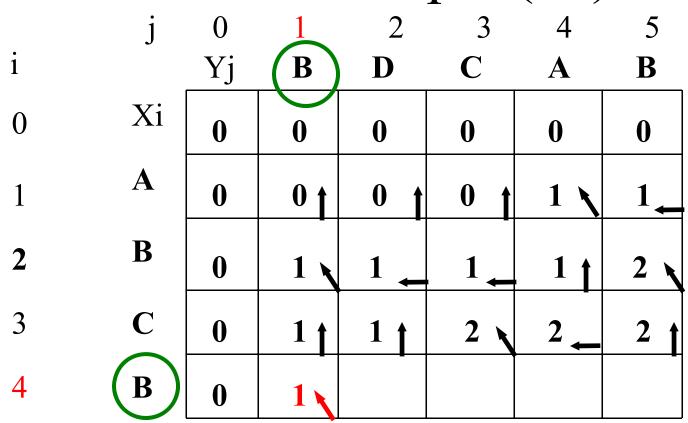


case i=3 and j= 5

$$C != B$$

 $c[2, 5] = c[3, 4]$
so $c[3, 5] = c[2, 5]$ and $b[3, 5] = 1$

LCS Example (18)

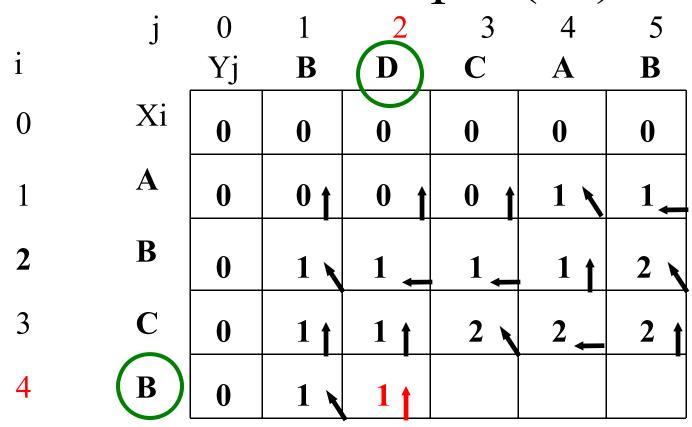


case i=4 and j=1

$$B = B$$

so c[4, 1] = c[3, 0]+1 and b[4, 1] =

LCS Example (19)

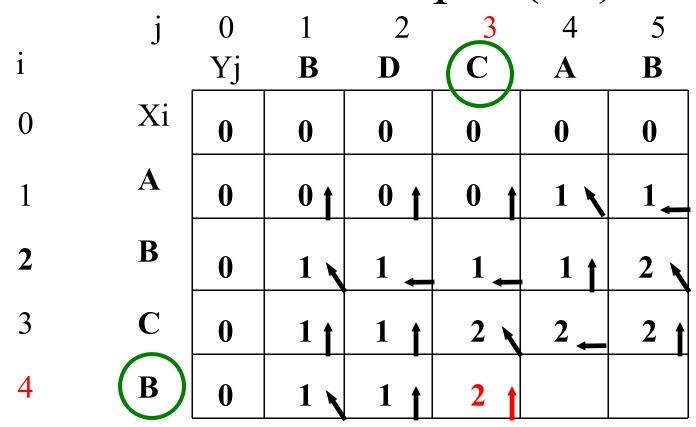


case i=4 and j=2
B!= D

$$c[3, 2] = c[4, 1]$$

so $c[4, 2] = c[3, 2]$ and $b[4, 2] = 1$

LCS Example (20)

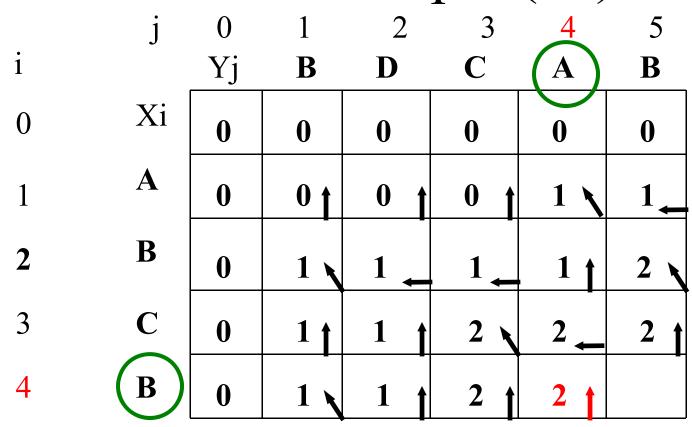


case i=4 and j=3
B!= C

$$c[3, 3] > c[4, 2]$$

so $c[4, 3] = c[3, 3]$ and $b[4, 3] = 1$

LCS Example (21)

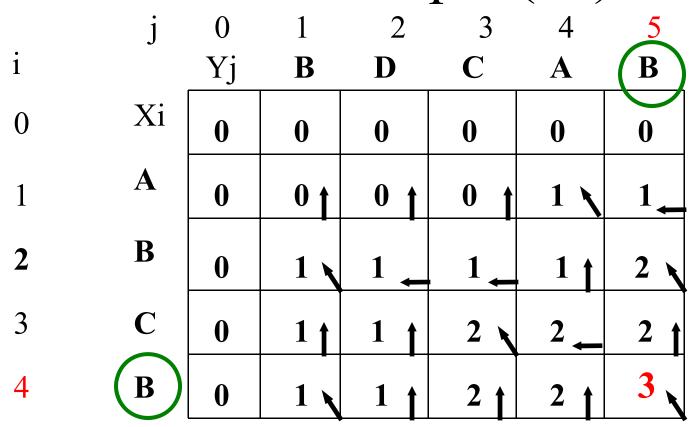


case i=4 and j=4
B!= A

$$c[3, 4] = c[4, 3]$$

so $c[4, 4] = c[3, 4]$ and $b[3, 5] = 1$

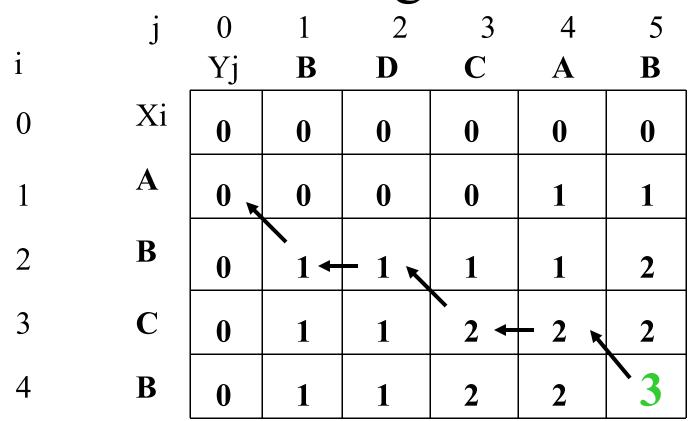
LCS Example (22)



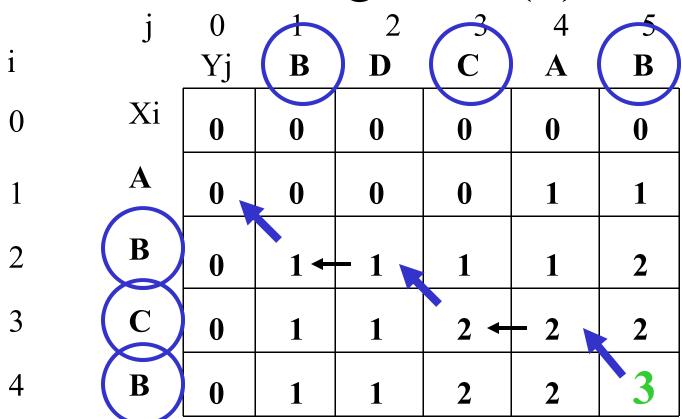
case i=4 and j=5
B= B
so c[4, 5] = c[3, 4]+1 and b[4, 5] =
$$\$$

```
LCS-Length(X, Y)
                                            Algorithm: LCS
m = length(X), n = length(Y)
for i = 1 to m
    c[i, 0] = 0
for j = 0 to n
    c[0, j] = 0
for i = 1 to m
       for j = 1 to n
            if (x_i = y_i)
                  c[i, j] = c[i - 1, j - 1] + 1
                  b[i, j] = "  "
            else if c[i-1, j] > = c[i, j-1]
                   c[i, j] = c[i - 1, j]
                   b[i, j] = "\uparrow"
            else c[i, j] = c[i, j - 1]
                   b[i, j] = "\leftarrow"
return c and b
```

Finding LCS



Finding LCS (2)



LCS (reversed order): B C B

LCS (straight order): B C B

Algorithm: Finding LCS

```
PRINT-LCS(b, X, i, j)

if i == 0 or j = 0

return

if b[i, j] == \text{``\}

PRINT-LCS(b, X, i - 1, j - 1)

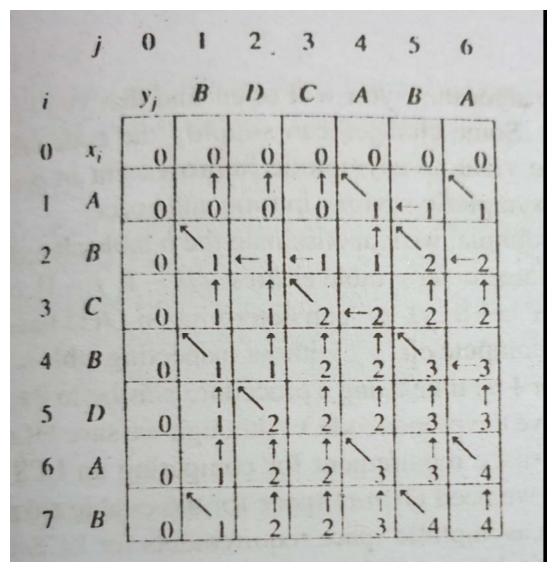
print x_i

elseif b[i, j] == \text{``\}

PRINT-LCS(b, X, i - 1, j)

else PRINT-LCS(b, X, i, j - 1, j)
```

Practice Example (Book: Cormen; Page-395)



Paris

LCS: BCBA

Length: 4

Thank You

Stay Safe