



CSE3103: Database

Nazmus Sakib
Assistant Professor
Department of Computer Science and Engineering
Ahsanullah University of Science and Technology

Transaction Concept

- A **transaction** is a *unit* of program execution that accesses and possibly updates various data items.
- E.g., transaction to transfer \$50 from account A to account B:

```
read(A)
A := A - 50
write(A)
read(B)
B := B + 50
write(B)
```

- Two main issues to deal with:
 - Failures of various kinds, such as hardware failures and system crashes
 - Concurrent execution of multiple transactions

Required Properties of a Transaction

• Consider a transaction to transfer \$50 from account A to account B:

```
1.read(A)

2.A := A - 50

3.write(A)

4.read(B)

5.B := B + 50

6.write(B)
```

Atomicity requirement

- If the transaction fails after step 3 and before step 6, money will be "lost" leading to an inconsistent database state
 - Failure could be due to software or hardware
- The system should ensure that updates of a partially executed transaction are not reflected in the database

Durability requirement

• Once the user has been notified that the transaction has completed (i.e., the transfer of the \$50 has taken place), the updates to the database by the transaction must persist even if there are software or hardware failures.

Required Properties of a Transaction (Cont.)

• Consistency requirement :

- The sum of A and B is unchanged by the execution of the transaction.
- A transaction, when starting to execute, must see a consistent database.
- During transaction execution the database may be temporarily inconsistent.
- When the transaction completes successfully the database must be consistent.
- Erroneous transaction logic can lead to inconsistency

Required Properties of a Transaction (Cont.)

• **Isolation requirement** — if between steps 3 and 6 (of the fund transfer transaction), another transaction T2 is allowed to access the partially updated database, it will see an inconsistent database (the sum A + B will be less than it should be).

```
T1 T2

1.read(A)

2.A := A - 50

3.write(A)

read(B)

4.read(B)

5.B := B + 50

6.write(B
```

- Isolation can be ensured trivially by running transactions serially
 - That is, one after the other.
- However, executing multiple transactions concurrently has significant benefits.

ACID Properties

A **transaction** is a unit of program execution that accesses and possibly updates various data items. To preserve the integrity of data the database system must ensure:

- **Atomicity.** Either all operations of the transaction are properly reflected in the database or none are.
- **Consistency.** Execution of a transaction in isolation preserves the consistency of the database.
- **Isolation.** Although multiple transactions may execute concurrently, each transaction must be unaware of other concurrently executing transactions. Intermediate transaction results must be hidden from other concurrently executed transactions.
 - That is, for every pair of transactions T_i and T_j it appears to T_i that either T_j finished execution before T_i started, or T_i started execution after T_i finished.
- **Durability.** After a transaction completes successfully, the changes it has made to the database persist, even if there are system failures.

Transaction State

- **Active** the initial state; the transaction stays in this state while it is executing
- Partially committed after the final statement has been executed.
- Failed -- after the discovery that normal execution can no longer proceed.
- **Aborted** after the transaction has been rolled back and the database restored to its state prior to the start of the transaction. Two options after it has been aborted:
 - Restart the transaction
 - can be done only if no internal logical error
 - Kill the transaction
- **Committed** after successful completion.

Transaction State (Cont.)



