

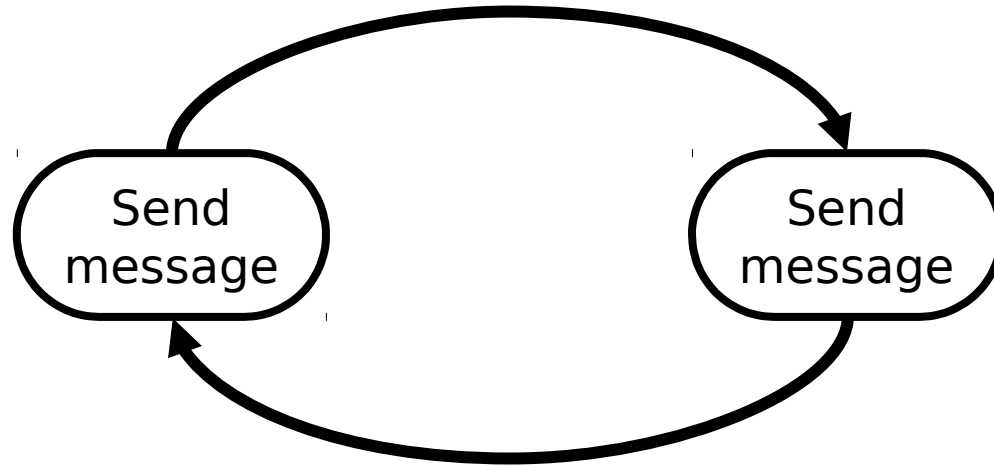


# Solving CSP

# Deadlock and Livelock

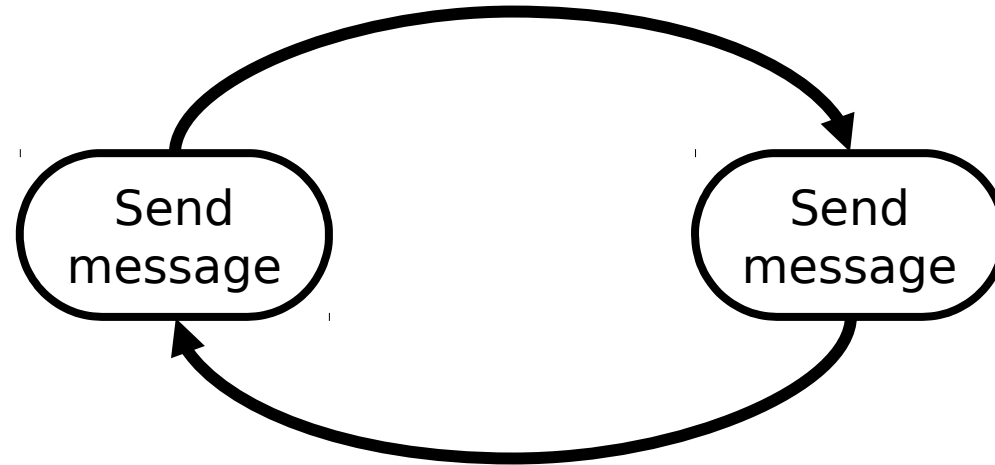
- Deadlock and livelock prevent the system continuing by starving its resources. It can be seen as a total system failure.
- Deadlock occurs when each process is locked out by another.
- Livelock occurs when each process tries to let another go first.
- There are ways to recover from deadlock and livelock, but the best solution is to avoid it altogether

# Deadlock



# Deadlock

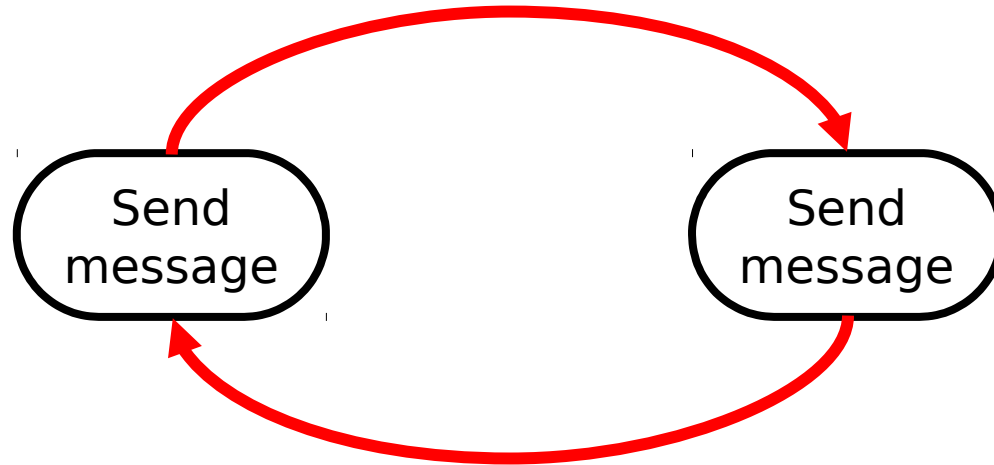
I want  
to go  
first!



I want  
to go  
first!

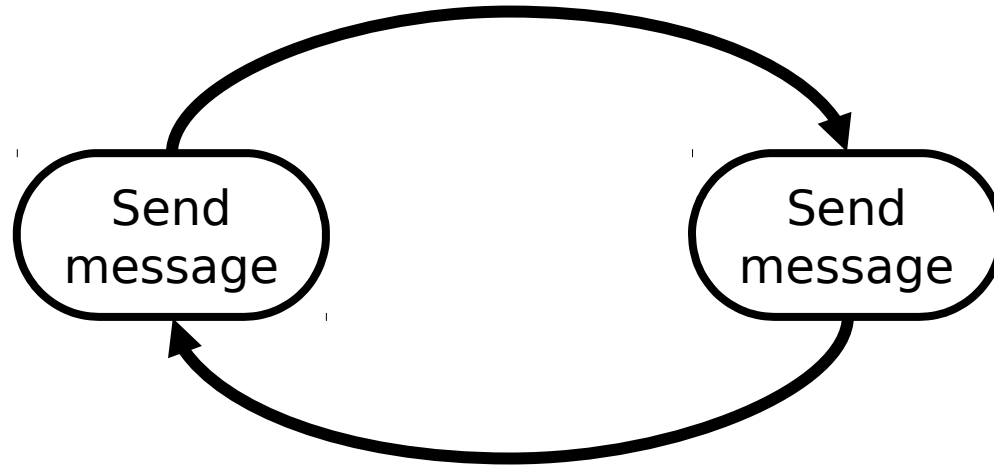
# Deadlock

I want  
to go  
first!



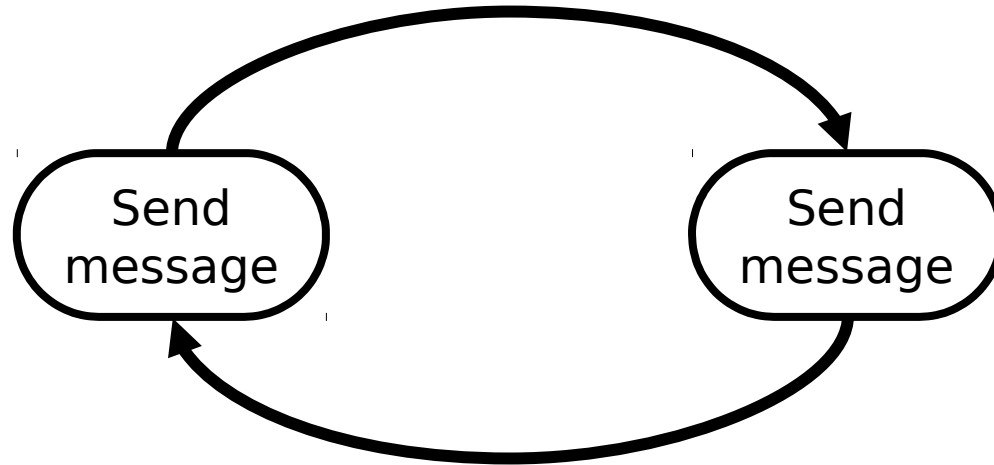
I want  
to go  
first!

# Livelock



# Livelock

You go  
first, I  
insist

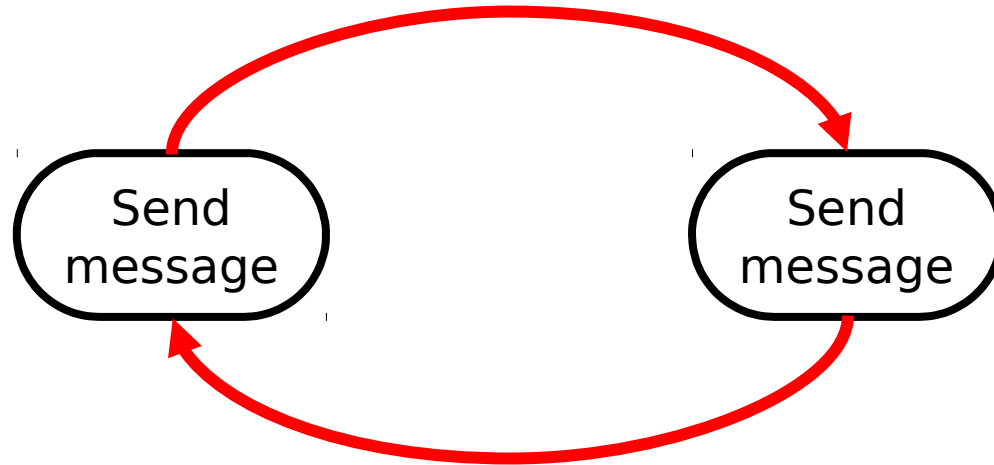


You go  
first, I  
insist



# Livelock

You go  
first, I  
insist



You go  
first, I  
insist

# Senders and Receivers

- We can define two types of processes; Senders and Receivers.
- Senders generate new messages and send them to receivers. They will wait for a response in a finite amount of time, if one is expected.
- Receivers will always wait to receive messages, and will always generate and send a response in a finite amount of time, if one is expected.

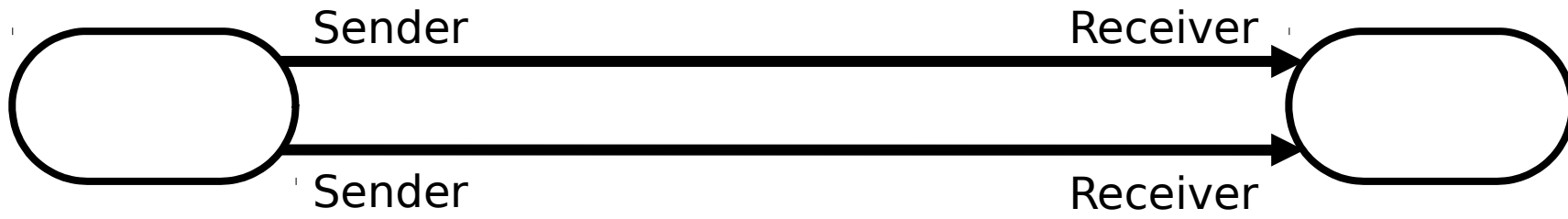
# Senders and Receivers



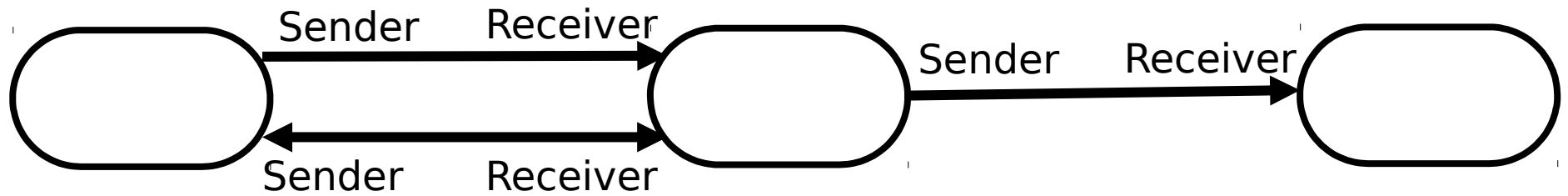
# Senders and Receivers



# Senders and Receivers



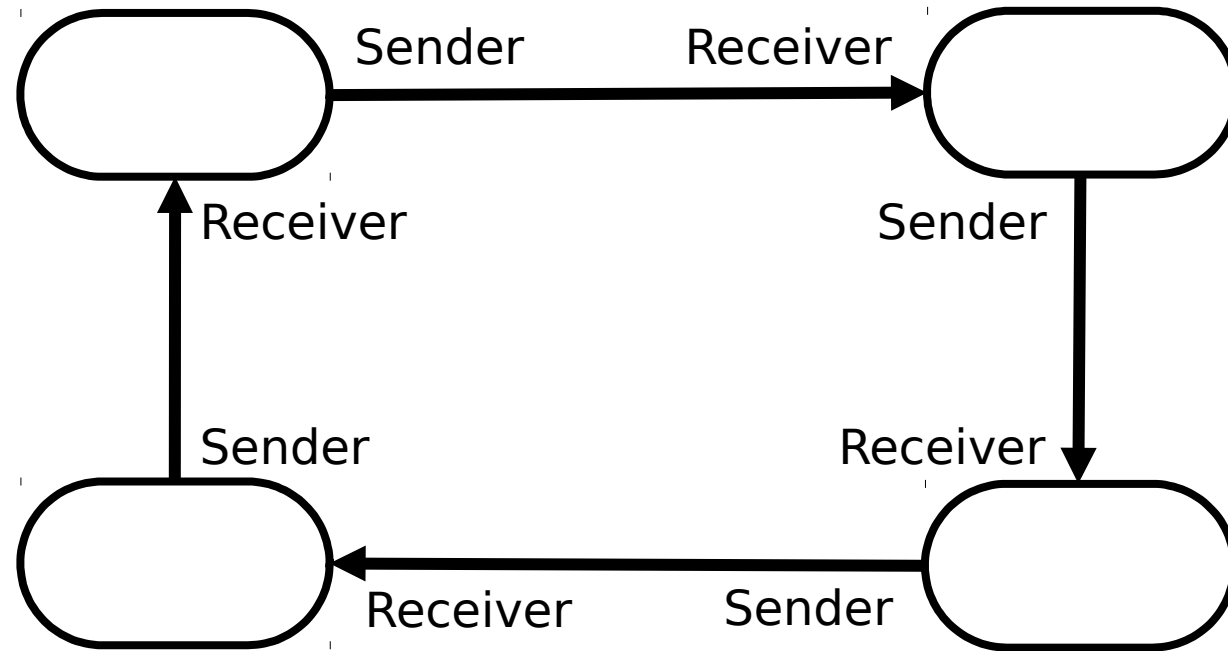
# Senders and Receivers



# Senders and Receivers

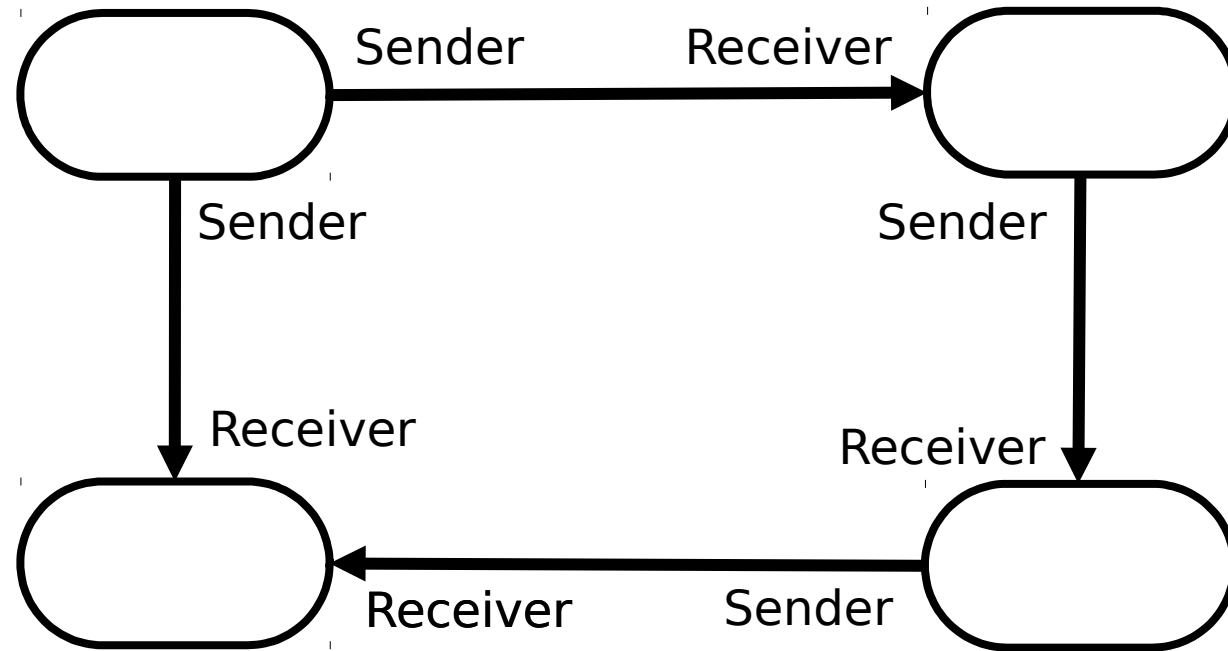
- Senders and Receivers allow us to avoid deadlock.
- As long as no Senders and Receivers interact in a loop, deadlock cannot occur.
- Livelock *might* still be a problem but its actually quite hard to get that to occur without trying to (famous last words...).

# Senders and Receivers

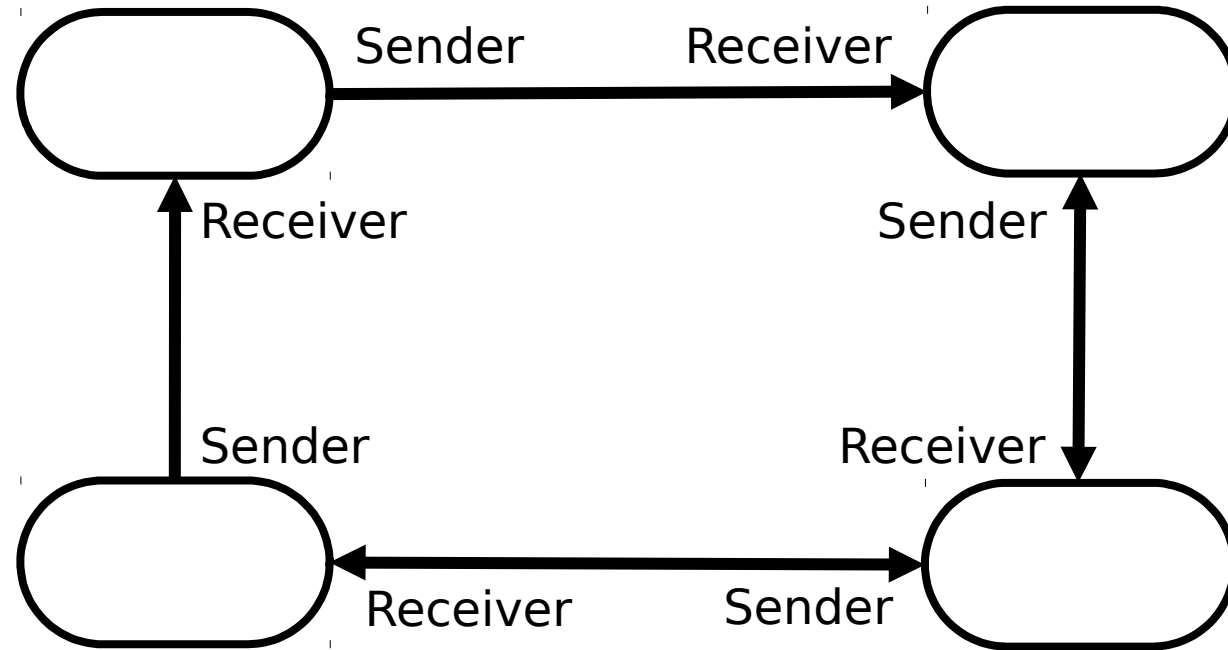




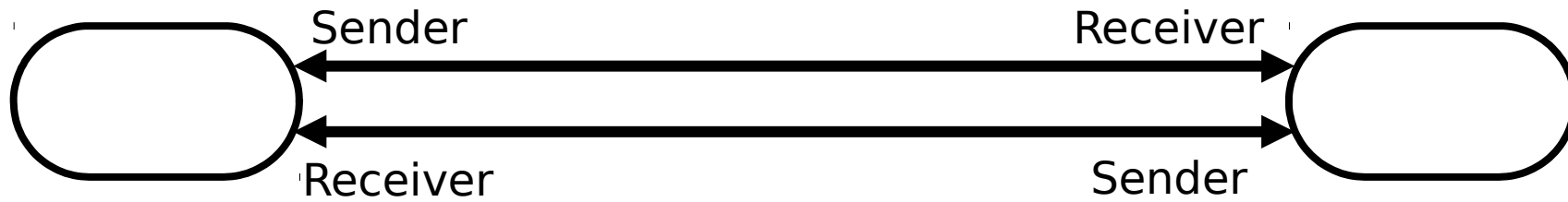
# Senders and Receivers



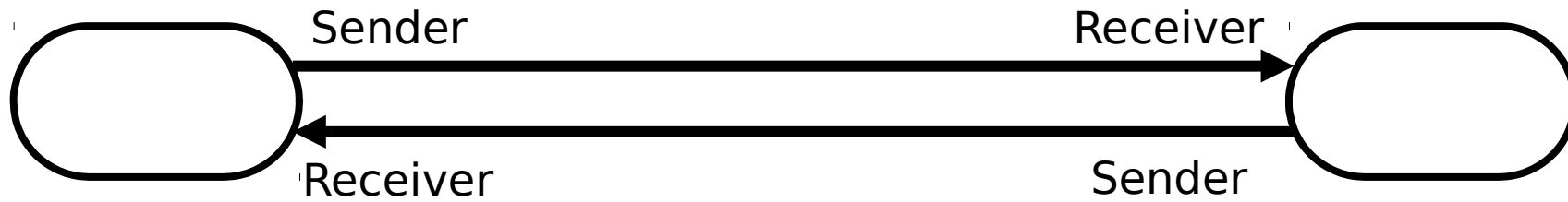
# Senders and Receivers



# Senders and Receivers



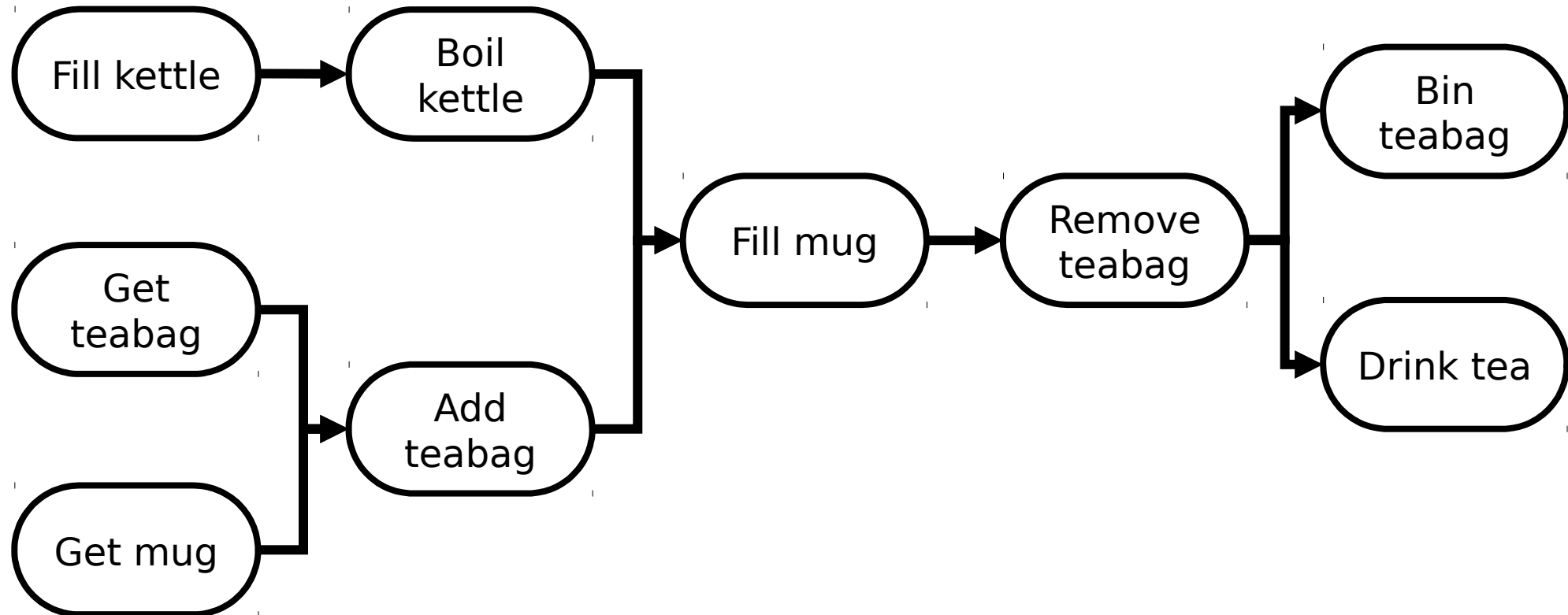
# Senders and Receivers



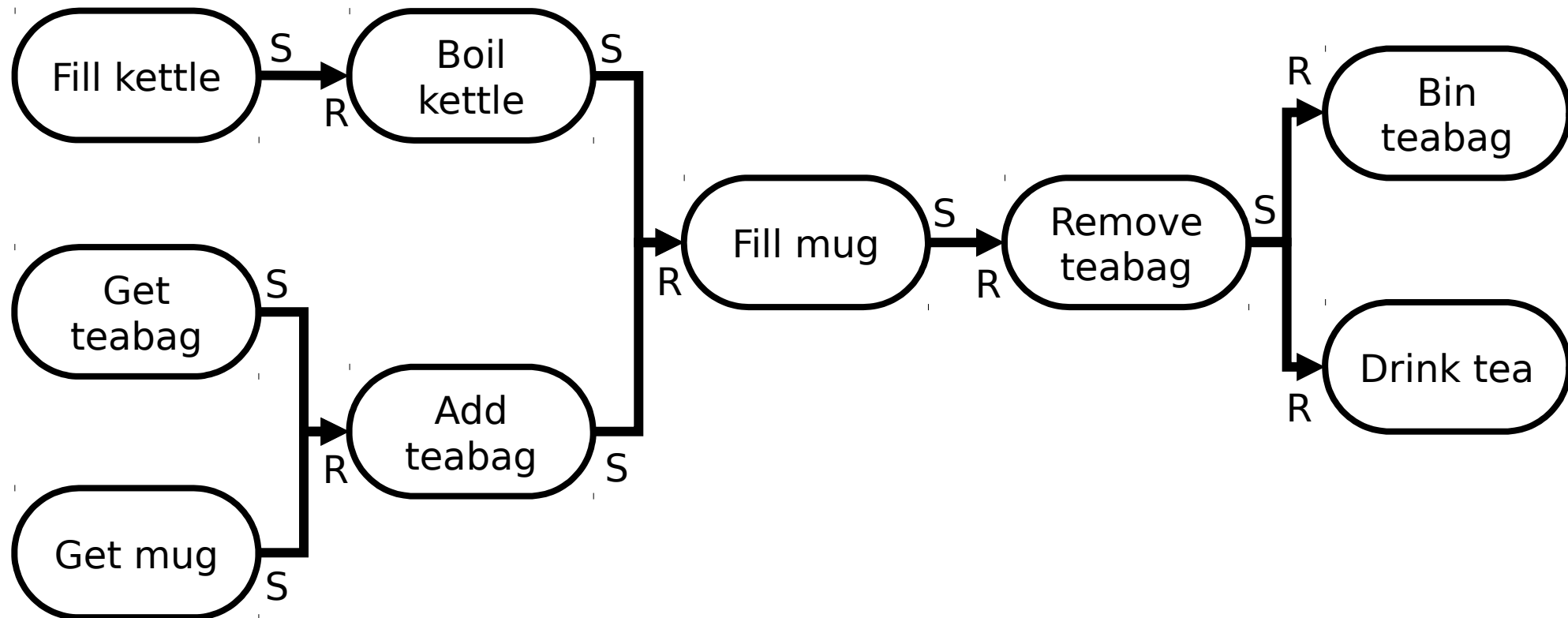
# Senders and Receivers

- The Sender Receiver model solves deadlock.
- But only if you've implemented your system as designed.
- Still possible to livelock.
- Still possible to make any other mistake.
- Still possible to avoid deadlock with careful management.
- Note that it may be referred to as the Client Server model in literature.

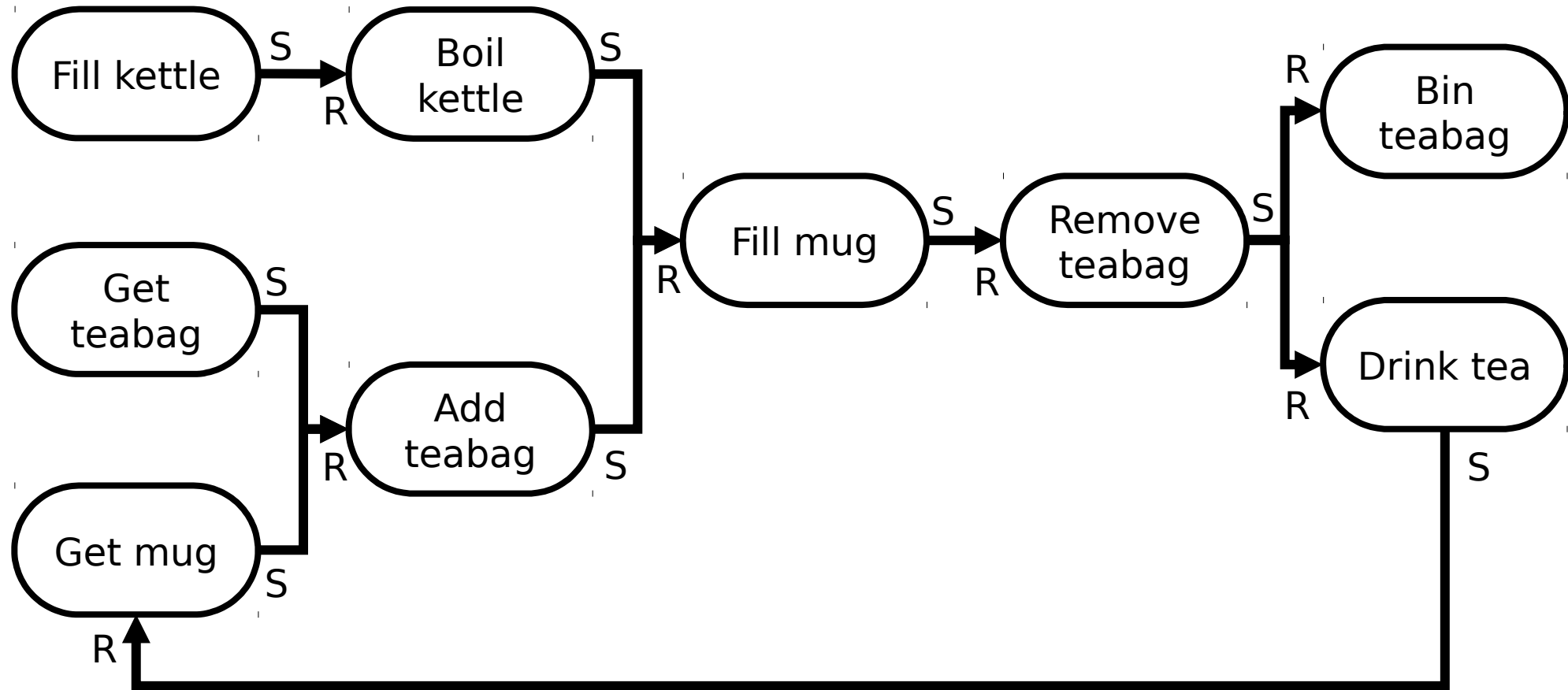
# Senders and Receivers



# Senders and Receivers

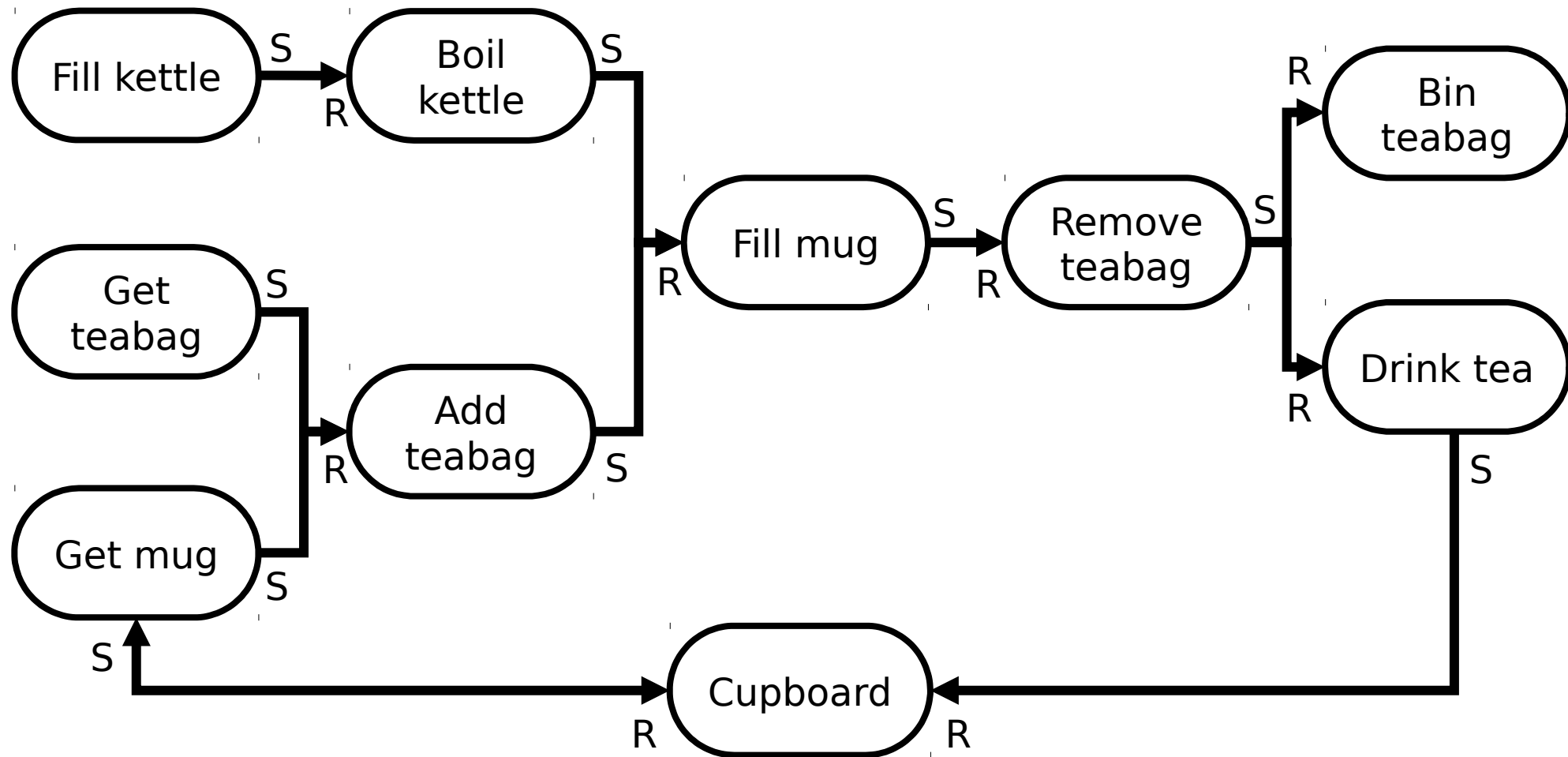


# Senders and Receivers





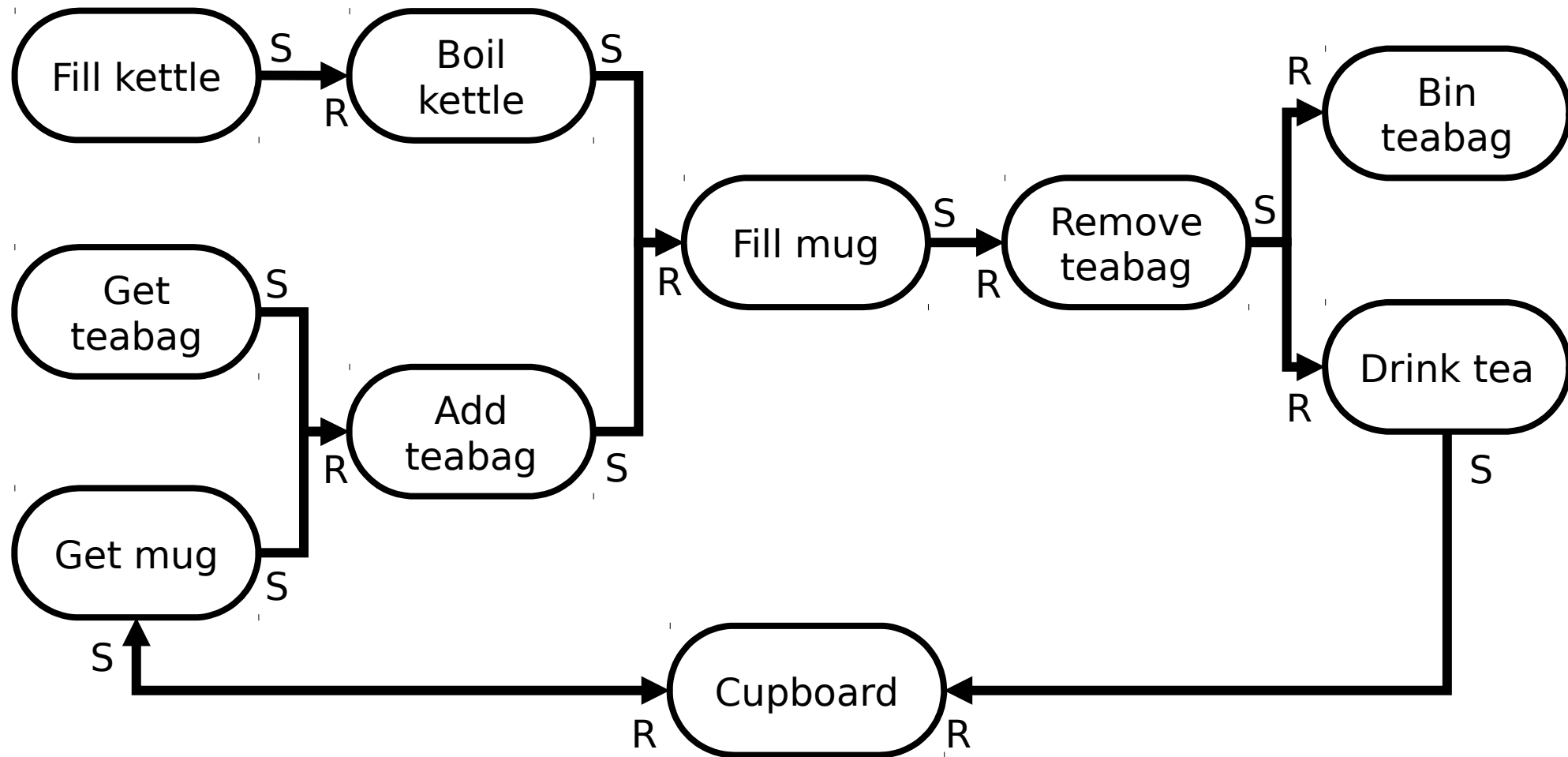
# Senders and Receivers



# The Resting Point

- It can help when checking the Sender/Receiver relationship to think of each process as having a Resting Point.
- This is the point at which the process will rest, waiting for an interaction and can be either listening for input, or waiting to send output.
- This may change as the process changes state and is intended only as a guide.

# The Resting Point



# Buffers

- The Cupboard process is acting as a *buffer*.
- A buffer is a place where messages can be stored to give the system some extra capacity.
- We can use this to swap the resting point between getting a mug and drinking tea

# A Note About Bi-directionality

- If we passed a channel to a process rather than one end, then we could read or write in either process.
- We have 1 channel, and can communicate in either direction.
- SPOT THE PROBLEM WITH THIS.