

# More Fault Tolerant Consensus

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① TurpinCoan Sync Multi

② BenOr Async Stop

# Outline

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# TurpinCoan Sync Multi

- General agreement, over arbitrary finite set  $V$ ,  $|V| \geq 2$
- TurpinCoan: **two extra rounds** + binary Byz

# TurpinCoan Init aka Round 0 (Process $\#i$ )

- Initial choice:  $x \in V$ 
  - determined by other means, or
  - received from outside
- Proposal:  $y \in V \cup \perp = \perp$
- Candidate:  $z \in V \cup \perp = \perp$
- Vote:  $\hat{v} \in \{0, 1\} = 0$

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# TurpinCoan Round 1 (Process #i)

- **Send**  $x$  to all processes
- Let  $W \subseteq V \cup \perp =$  **multiset** of all **received** messages
- $|W| = N$  : **sync**,  $\perp$
- If  $\exists v \in V, |W|_v \geq N - F = 2F + 1$ , then  $y = v$
- Else, keep  $y = \perp$
- $y \in V \cup \perp$  is our **proposal**
- Note: **all non-faulty** processes select the **same**  $y \in V \cup \perp$ 
  - $aaab \Rightarrow y = a, aaac \Rightarrow y = a$
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## TurpinCoan Round 2 (Process $\#i$ )

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- $|W| = N$  : **sync**,  $\perp$
- If  $\exists v \in V, |W|_v \geq N - F = 2F + 1$ , then  $z = v, \hat{v} = 1$ 
  - We vote for candidate  $z \in V$
- Else if  $\exists v \in V, \arg \max_v |W|_v$ , arbitrary tie break, then  $z = v$  ( $\hat{v} = 0$ )
  - We do NOT vote for candidate  $z \in V$ , but this may be the final decision
- Else i.e.  $|W| \cap V = \emptyset$ . ( $z = \perp, \hat{v} = 0$ )
  - No candidate, no vote

# TurpinCoan Round 3, ... (Process $\#i$ )

- Binary Byz agreement on  $\hat{v} \in \{0, 1\}$ , for the candidate  $z \in V \cup \perp$
- If this Byz decision is 1 and  $z \in V$  (i.e.  $z \neq \perp$ ), then final decision  $z$
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The three loyal processes start with  $aab$ ,  $a, b \in V$ , the last process is faulty

- Variant agreement:  $z = a \in V$ 
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- Way around: stronger model, and weaker termination
- Stronger model: processes use **randomisation**
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# BenOr Init aka Round 0 (Process $\#i$ )

- Initial choice:  $x \in \{0, 1\}$ 
  - determined by other means, or
  - received from outside
- Proposal:  $y \in \{0, 1, \perp\} = \perp$
- Step:  $s \geq 0 = 0$ , unbounded
- Each step has two rounds

# BenOr Step $s$ , Round 1 (Process $\#i$ )

- **Send**  $(I, s, x)$  to all processes
- Let  $M =$  **multiset** of **first**  $N - F = 2F + 1$  **received** messages  $(I, s, *)$
- If all  $m \in M$  have same value  $v \in \{0, 1\}$ , then  $y = v$
- Else,  $y = \perp$

BenOr Step  $s$ , Round 1 (Process  $\#i$ )

- **Send**  $(I, s, x)$  to all processes
- Let  $M =$  **multiset** of **first**  $N - F = 2F + 1$  **received** messages  $(I, s, *)$
- If all  $m \in M$  have same value  $v \in \{0, 1\}$ , then  $y = v$
- Else,  $y = \perp$



BenOr Step  $s$ , Round 2 (Process  $\#i$ )

- **Send**  $(\Pi, s, y)$  to all processes
- Let  $M =$  **multiset** of **first**  $N - F = 2F + 1$  **received** messages  $(\Pi, s, *)$
- If **all**  $m \in M$  have same value  $v \in \{0, 1\}$ , then  $x = v$ , **decide**  $v$  (if not already), and **continue**
- If **at least**  $N - 2F = F + 1$   $m \in M$  have same value  $v \in \{0, 1\}$ , then  $x = v$ , but do not decide
- Else i.e.  $x = \text{random} \in \{0, 1\}$ .

BenOr Step  $s$ , Round 2 (Process  $\#i$ )

- **Send**  $(II, s, y)$  to all processes
- Let  $M =$  **multiset** of **first**  $N - F = 2F + 1$  **received** messages  $(II, s, *)$
- If **all**  $m \in M$  have same value  $v \in \{0, 1\}$ , then  $x = v$ , **decide**  $v$  (if not already), and **continue**
- If **at least**  $N - 2F = F + 1$   $m \in M$  have same value  $v \in \{0, 1\}$ , then  $x = v$ , but do not decide
- Else i.e.  $x = \text{random} \in \{0, 1\}$ .

BenOr Step  $s$ , Round 2 (Process  $\#i$ )

- **Send**  $(\Pi, s, y)$  to all processes
- Let  $M =$  **multiset** of **first**  $N - F = 2F + 1$  **received** messages  $(\Pi, s, *)$
- If **all**  $m \in M$  have same value  $v \in \{0, 1\}$ , then  $x = v$ , **decide**  $v$  (if not already), and **continue**
- If **at least**  $N - 2F = F + 1$   $m \in M$  have same value  $v \in \{0, 1\}$ , then  $x = v$ , but do not decide
- Else i.e.  $x = \text{random} \in \{0, 1\}$ .