

A:

Step 1: Do not modify $f(n) = (\log(n))^2$; But modify this equation $g(n) = \log_2(n^{\log_2 n})^2$ by formula logarithm power rule $2^{\log_2 n^{\log_2 n}}$, use formula logarithm power rule again we can get $2^{\log_2 n * \log_2(n)}$

Step 2: The big O notation $f(n) = O(\log(n))$, and for all $n \geq n_0$ such that $C*f(n) \geq g(n)$; also $g(n) = O(\log(n))$ and for all $n \geq n_0$ such that $C*g(n) \geq f(n)$. These two equations have the properties of $f(n) = O(g(n))$ and $f(n) = \Omega(g(n))$; thus, $f(n)$ and $g(n)$ have the same asymptotic growth rate. Then they should be both, so $f(n) = \Theta(g(n))$.

B:

Given: $f(n) = n^{10}$; $g(n) = 2^{\sqrt[10]{n}}$

Step 1:

Do the logarithm for both functions. $\log f(n) = 2^{\log_2 n^{10}}$;
 $\log g(n) = 2^{\sqrt[10]{n}}$ we can compare the power of two functions as
 $\log f(n) = 10 \log n$; $\log g(n) = n^{0.1}$

Step 2:

Apply the L' Hopital' s rule, then it will be $\lim_{x \rightarrow \infty} \frac{g'(x)}{f'(x)} ==>$
 $\lim_{n \rightarrow \infty} \frac{0.1 * n^{-0.9}}{\frac{1}{n} \ln 2}$, if n is approach to infinite, then the value is
approaching to infinite as well.

Step 3:

So, $0 <= cg(n) <= f(n)$, for all $n \geq n_0$. The increasing rate infinite, $g(x)$ is the upper bound of $f(x)$; $g(x) = O(f(x))$ which is $g(x) = \Omega(f(x))$.

C:

Given $f(n) = n^{1+(-1)^n}$, $g(n) = n$.

Step 1:

There are two cases, When n is odd number then $f(n) = 1$, when n is even number $f(n) = n^2$.

Step 2:

If number n is odd number, $f(n) = 1$, $g(n) = n$. Then $g(n)$ is the upper bound of $f(n)$, then $g(n) = O(f(n))$.

If number is even number, $f(n) = n^2$, $g(n) = n$. Then $0 \leq c \cdot g(n) \leq f(n)$, so that, $f(n) = \Omega(g(n))$.

However in a single function $f(n)$, There is none a fixed n_0 could satisfied the condition of such that $0 \leq c \cdot g(n) \leq f(n)$ and $0 \leq f(n) \leq c \cdot g(n)$, for all $n \geq n_0$. So, it is neither.