

Given: N cities with trip time T to each other and their own pods and populations

Aim: Find how many pods can escape at maximum condition.

Prepossessing: Regard City C_i as the starting point use Dijkstra algorithm to find all possible cities that can be reached from C_i within T days trip time. Append all the found cities into list L . Do this for all the N cities, so that we will get N list L .

Setup: Construct a bipartite graph with N vertexes on left and right, At the leftmost there is a vertex which is the Super source S , then on the left hand side there are N vertexes which represent populations of each cities. At the rightmost is the Super Sink S_i , on the right hand side there are N vertexes that represent number of pods of each cities. Connect with directed edge from S to all the vertexes on the left hand side, the capacity for each edge is equal to the connected city's(vertex) population. Then connect with directed edge from the vertexes on the right hand side to the rightmost, the capacity for each edge is equal to the connected city's(vertex) pods.

Solution: Connecting each city population with their own pods and connect the edge from left hand side to the cities (vertexes) on the right hand side that appeared in the corresponded List L . (For example: A city C_i , there are three cities C_j, C_k, C_m can be reached within X days then List L_i will be $[C_j, C_k, C_m]$, while connecting edges the C_i will only connect to C_j, C_k, C_m) The capacity for these edges are 0, because there are no restrictions about how many people pass the city, so that capacities are equal to 0. Now we finished graph. Then we find max flow in such a network using the Edmonds-karp algorithm. Then build the last residual network and find the minimal cut. The result will be the total capacity of the minimal cut.