

# Balancing Privacy and Utility in Correlated Data

INRIA Montpellier, November 12th, 2025

Patricia Guerra-Balboa

PRIVACY  
AND SECURITY

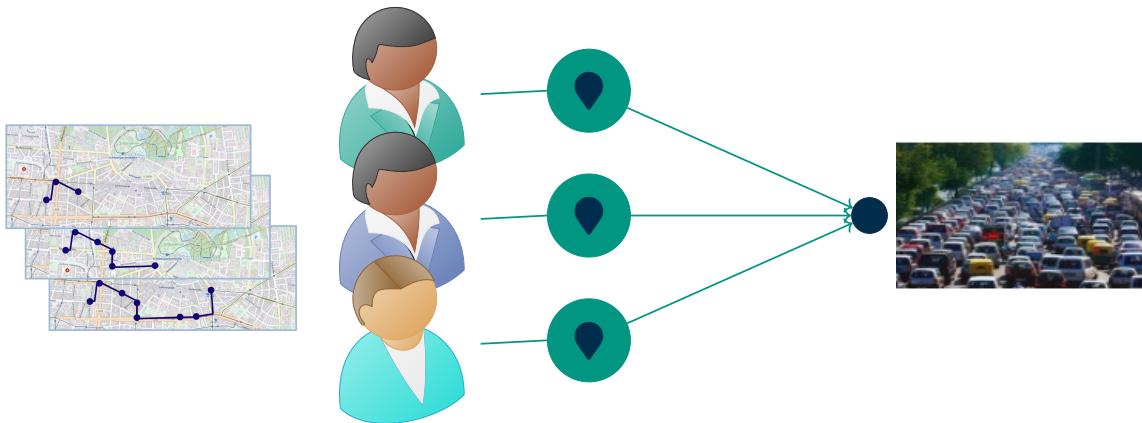
KASTEL

# Our General Goal

**Learn population-level information without harming  
individual's privacy**

# Our General Goal

Learn population-level information without harming individual's privacy

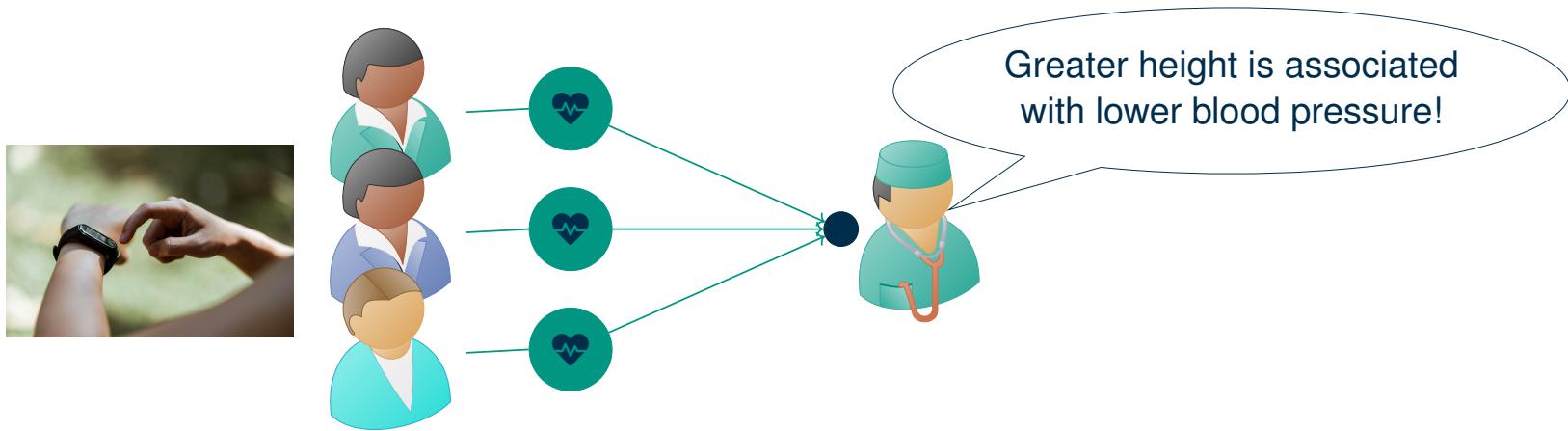


**Privacy Goal:** Protect Alice's location

**Utility Goal:** Number of cars per street

# Our General Goal

Learn population-level information without harming individual's privacy



Privacy Goal: Protect Alice's activity data

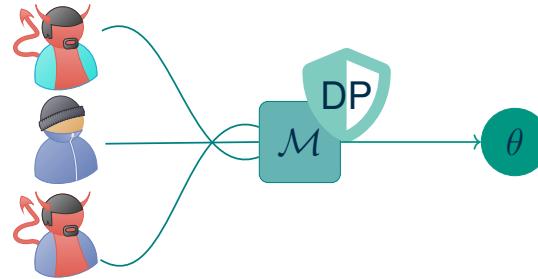
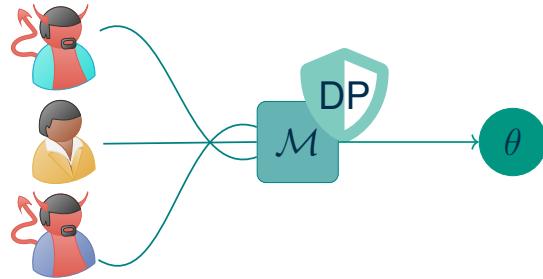
Utility Goal: Correlation between height and health

# The Best Tool Until Now: Differential Privacy

**Idea:** We want to bound participation risk.

# The Best Tool Until Now: Differential Privacy

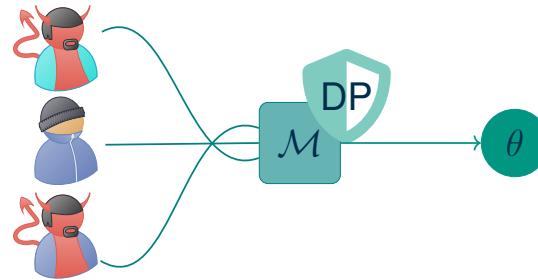
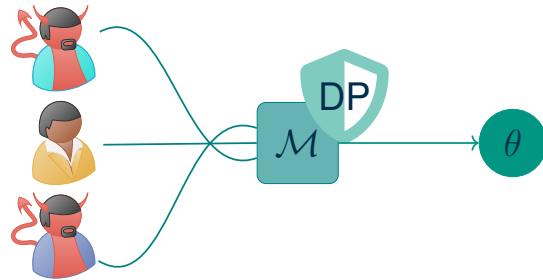
Idea: We want to bound participation risk.



- “Strongest” assumption: everybody’s record is known but the target.

# The Best Tool Until Now: Differential Privacy

Idea: We want to bound participation risk.



- “Strongest” assumption: everybody’s record is known but the target.

# The Best Tool Until Now: Differential Privacy

Idea: We want to bound participation risk.



$$\ln \frac{p_{\mathcal{M}}(\theta | x_1, \dots, x_{n-1}, \textcolor{brown}{x_n})}{p_{\mathcal{M}}(\theta | x_1, \dots, x_{n-1}, \textcolor{blue}{y_n})} \leq \epsilon$$

- “Strongest” assumption: everybody’s record is known but the target.
- The privacy leakage  $\epsilon$  controls the indistinguishability level between  $\textcolor{brown}{x_n}$ ,  $\textcolor{blue}{y_n}$ .

# The Best Tool Until Now: Differential Privacy

Idea: We want to bound participation risk.

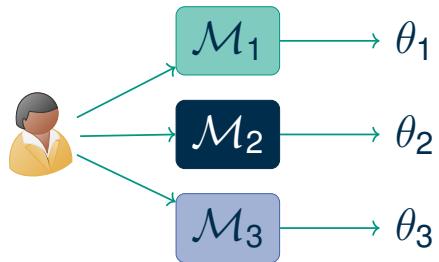


$$\ln \frac{p_{\mathcal{M}}(\theta | x_1, \dots, x_{n-1}, \textcolor{brown}{x_n})}{p_{\mathcal{M}}(\theta | x_1, \dots, x_{n-1}, \textcolor{blue}{y_n})} \leq \epsilon$$

- “Strongest” assumption: everybody’s record is known but the target.
- The privacy leakage  $\epsilon$  controls the indistinguishability level between  $\textcolor{brown}{x_n}$ ,  $\textcolor{blue}{y_n}$ .
- But at some cost! The smaller the  $\epsilon$  the less utility.

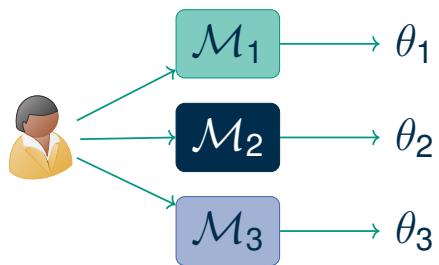
# Why DP Is The Best So Far?

## Composition

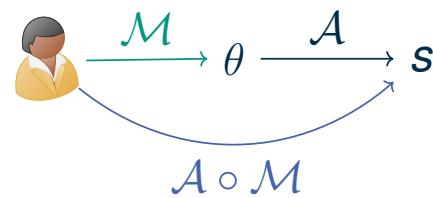


# Why DP Is The Best So Far?

## Composition

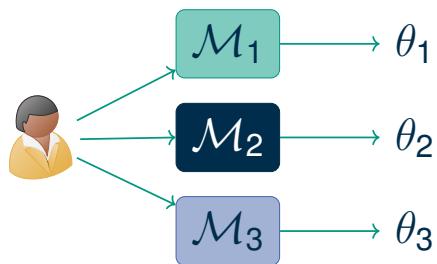


## Post-processing

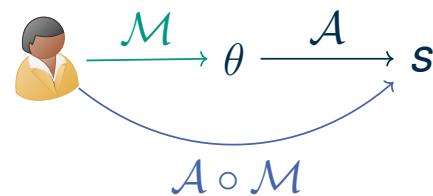


# Why DP Is The Best So Far?

## Composition



## Post-processing

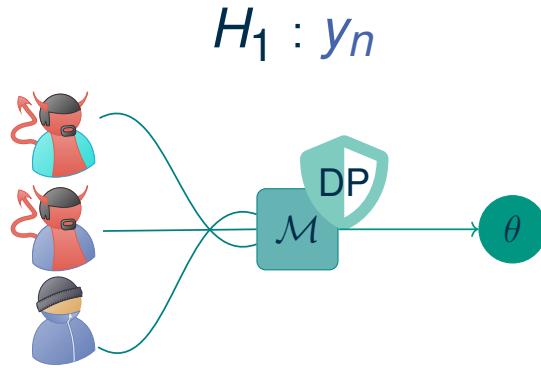
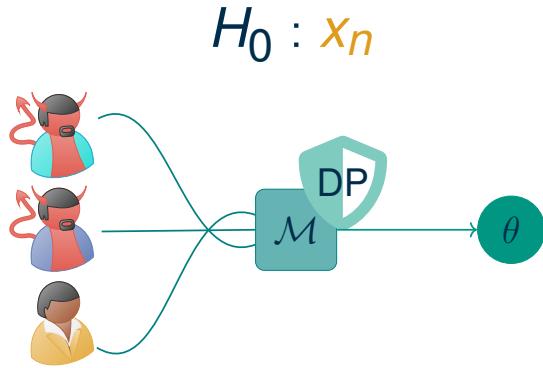


## Attack Mitigation

$$\mathcal{M} \text{- } \varepsilon\text{-DP} \Rightarrow \text{Adv} \leq f(\varepsilon)$$

# Membership Inference Attack Knowing $D_-$

The attacker receives  $\theta$  and aims to distinguish between:



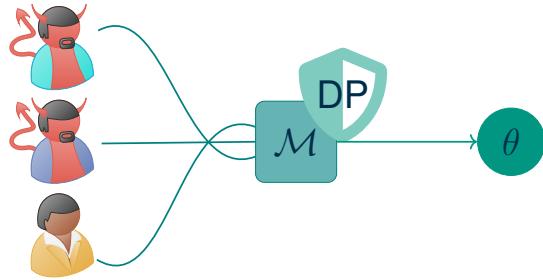
$D_-$  is known:

$$H_0 = D_{x_n} \text{ Vs. } H_1 = D_{y_n}$$

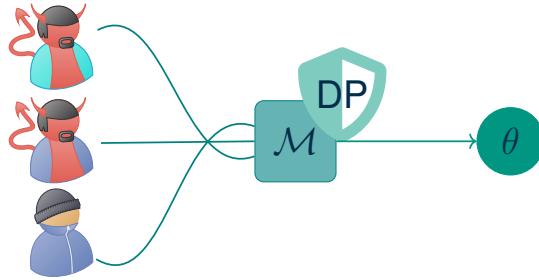
# Membership Inference Attack Knowing $D_-$

The attacker receives  $\theta$  and aims to distinguish between:

$$H_0 : x_n$$



$$H_1 : y_n$$



$D_-$  is known:

$$H_0 = D_{x_n} \text{ Vs. } H_1 = D_{y_n}$$

Type I error:

$$\alpha = \Pr_{A \circ \mathcal{M}}(y_n | D_{x_n})$$

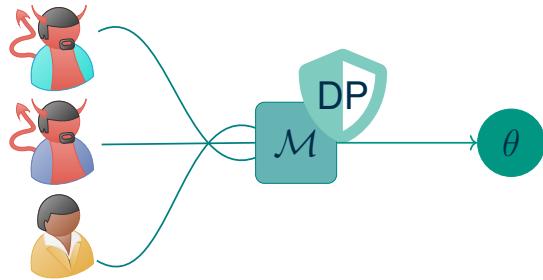
Type II error:

$$\beta = 1 - \Pr_{A \circ \mathcal{M}}(y_n | D_{y_n})$$

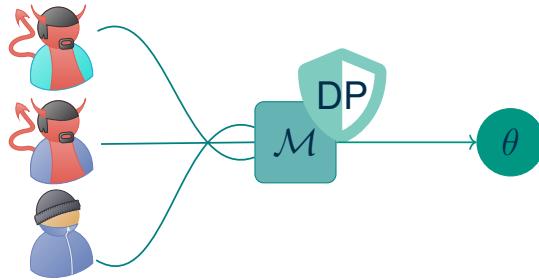
# Membership Inference Attack Knowing $D_-$

The attacker receives  $\theta$  and aims to distinguish between:

$$H_0 : x_n$$



$$H_1 : y_n$$



$D_-$  is known:

$$H_0 = D_{x_n} \text{ Vs. } H_1 = D_{y_n}$$

Type I error:

$$\alpha = \Pr_{A \circ \mathcal{M}}(y_n | D_{x_n})$$

Type II error:

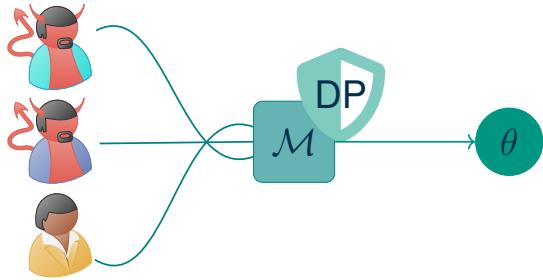
$$\beta = 1 - \Pr_{A \circ \mathcal{M}}(y_n | D_{y_n})$$

$A \circ \mathcal{M}$  is  $\varepsilon$ -DP  $\Rightarrow$

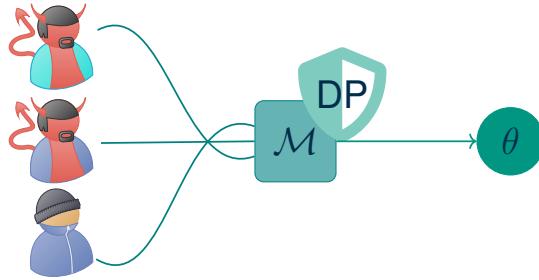
# Membership Inference Attack Knowing $D_-$

The attacker receives  $\theta$  and aims to distinguish between:

$H_0 : x_n$



$H_1 : y_n$



$D_-$  is known:

$$H_0 = D_{x_n} \text{ Vs. } H_1 = D_{y_n}$$

Type I error:

$$\alpha = \Pr_{A \circ M}(y_n | D_{x_n})$$

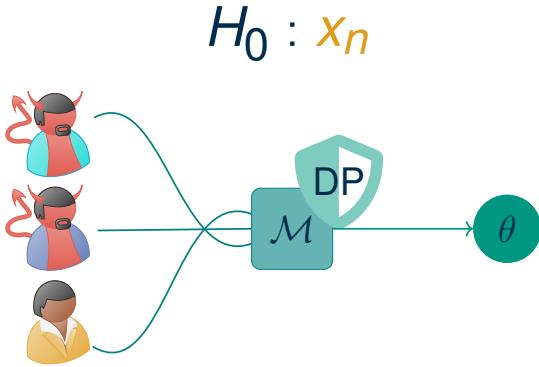
Type II error:

$$\beta = 1 - \Pr_{A \circ M}(y_n | D_{y_n})$$

$$A \circ M \text{ is } \varepsilon\text{-DP} \Rightarrow \begin{aligned} 1 - \beta &\leq e^\varepsilon \alpha \\ \alpha &\leq e^\varepsilon (1 - \beta) \end{aligned}$$

# Membership Inference Attack Knowing $D_-$

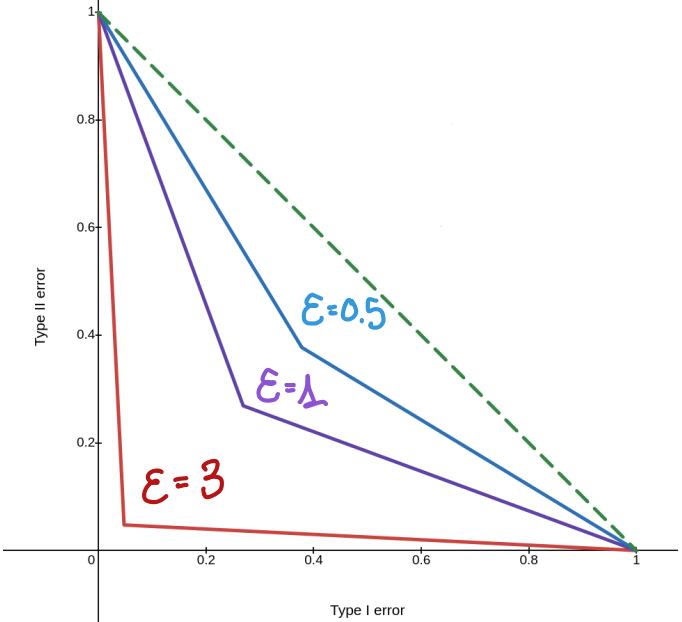
The attacker receives  $\theta$  and aims to distinguish between:



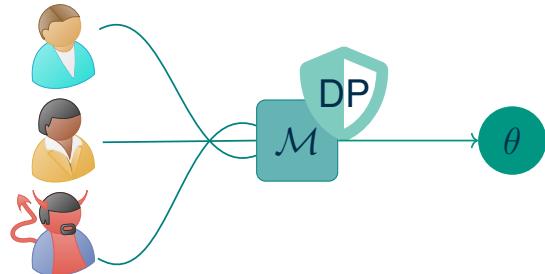
Type I error:  
 $\alpha = \Pr_{A \circ M}(y_n | D_{x_n})$

Type II error:  
 $\beta = 1 - \Pr_{A \circ M}(y_n | D_{y_n})$

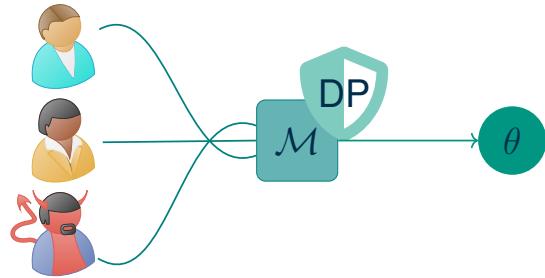
$$A \circ M \text{ is } \varepsilon\text{-DP} \Rightarrow 1 - \beta \leq e^\varepsilon \alpha \Rightarrow \alpha \leq e^\varepsilon (1 - \beta) \Rightarrow \beta \geq \max\{1 - e^\varepsilon \alpha, e^\varepsilon (1 - \alpha)\}$$



# What about other attacker models?



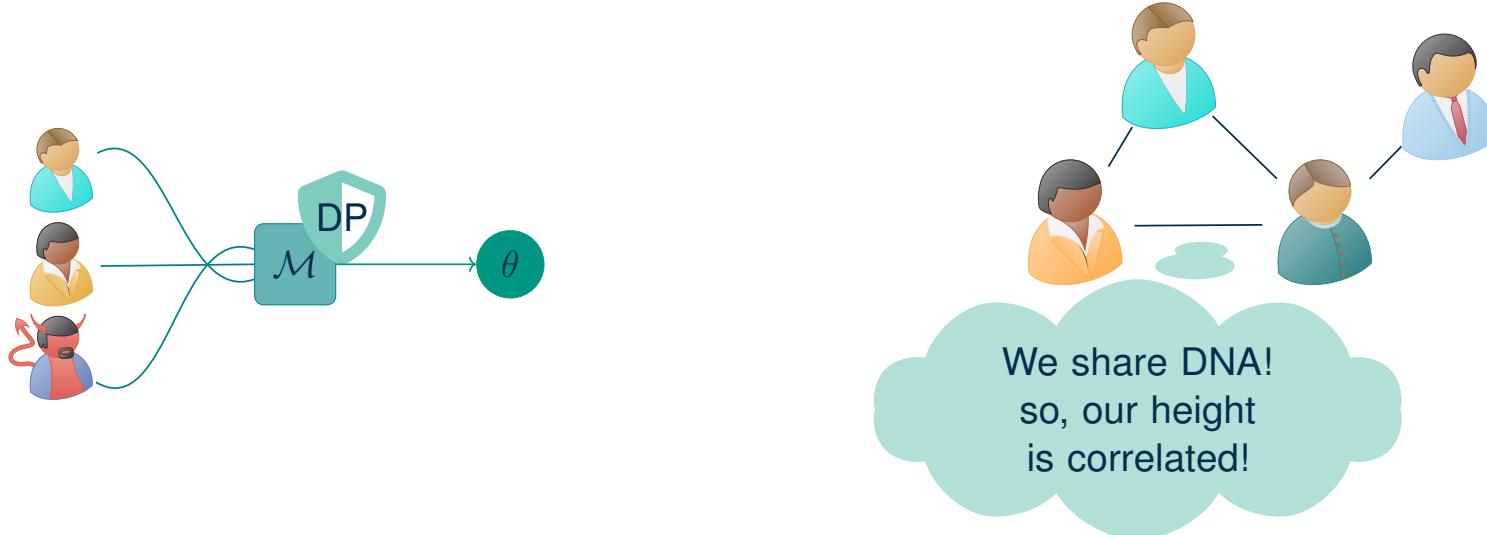
# What about other attacker models?



Statistical Independence

The strongest attacker is the worst-case one, and we have at least the same protection.

# What about other attacker models?



Statistical Independence

Dependencies between  
Records

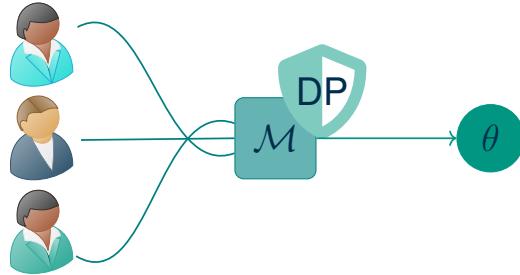
The strongest attacker is the worst-case one, and we have at least the same protection.

DP interpretation does not hold anymore.

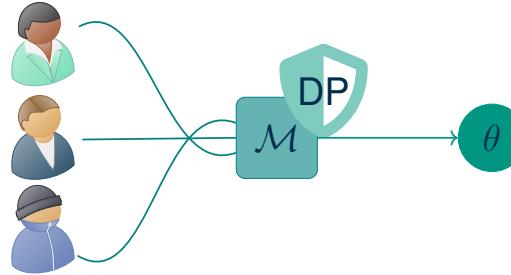
# Membership Inference Attack With Dependencies

The attacker receives  $\theta$  and aims to distinguish between:

$$H_0 : x_n$$



$$H_1 : y_n$$



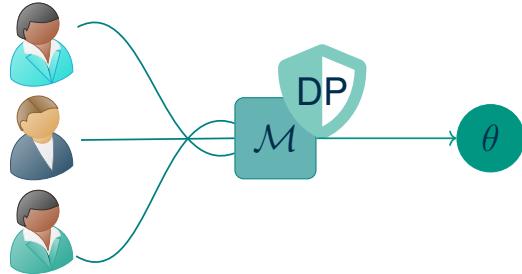
$D_-$  is **unknown**:

$$H_0 = \{D : x_n \in D\} \text{ Vs. } H_1 = \{D : y_n \in D\}$$

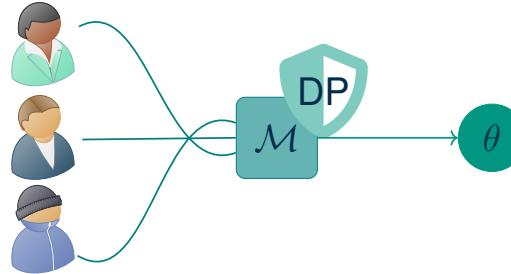
# Membership Inference Attack With Dependencies

The attacker receives  $\theta$  and aims to distinguish between:

$$H_0 : x_n$$



$$H_1 : y_n$$



$D_-$  is **unknown**:

$$H_0 = \{D : x_n \in D\} \text{ Vs. } H_1 = \{D : y_n \in D\}$$

Type I error:

$$\alpha = \Pr_{A \circ \mathcal{M}} (y_n | x_n)$$

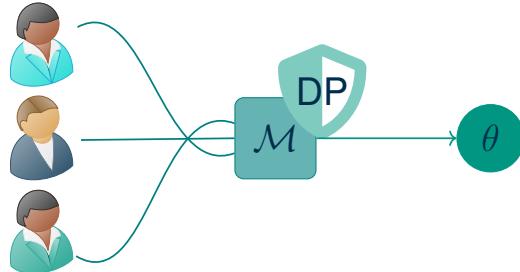
Type II error:

$$\beta = 1 - \Pr_{A \circ \mathcal{M}} (y_n | y_n)$$

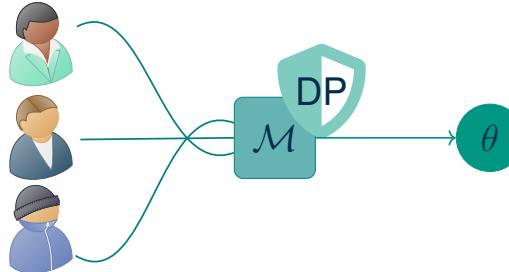
# Membership Inference Attack With Dependencies

The attacker receives  $\theta$  and aims to distinguish between:

$$H_0 : x_n$$



$$H_1 : y_n$$



$D_-$  is **unknown**:

$$H_0 = \{D : x_n \in D\} \text{ Vs. } H_1 = \{D : y_n \in D\}$$

Type I error:

$$\begin{aligned}\alpha &= \Pr_{A \circ \mathcal{M}}(y_n | x_n) \\ &= \sum_{D_-} \Pr_{A \circ \mathcal{M}}(y_n | D_{x_n}) \pi(D_- | x_n)\end{aligned}$$

Type II error:

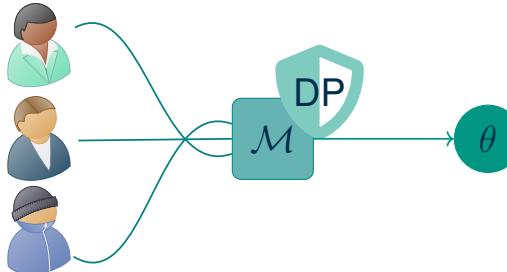
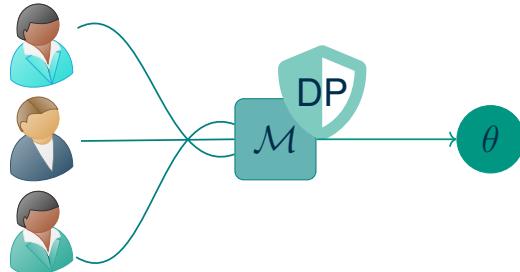
$$\begin{aligned}\beta &= 1 - \Pr_{A \circ \mathcal{M}}(y_n | y_n) \\ &= 1 - \sum_{D_-} \Pr_{A \circ \mathcal{M}}(y_n | D_{y_n}) \pi(D_- | y_n)\end{aligned}$$

# Membership Inference Attack With Dependencies

The attacker receives  $\theta$  and aims to distinguish between:

$$H_0 : x_n$$

$$H_1 : y_n$$



$D_-$  is **unknown**:

$$H_0 = \{D : x_n \in D\} \text{ Vs. } H_1 = \{D : y_n \in D\}$$

Type I error:

$$\begin{aligned}\alpha &= \Pr_{A \circ \mathcal{M}}(y_n | x_n) \\ &= \sum_{D_-} \Pr_{A \circ \mathcal{M}}(y_n | D_{x_n}) \pi(D_- | x_n)\end{aligned}$$

Type II error:

$$\begin{aligned}\beta &= 1 - \Pr_{A \circ \mathcal{M}}(y_n | y_n) \\ &= 1 - \sum_{D_-} \Pr_{A \circ \mathcal{M}}(y_n | D_{y_n}) \pi(D_- | y_n)\end{aligned}$$

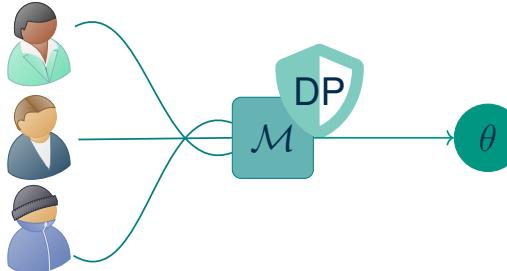
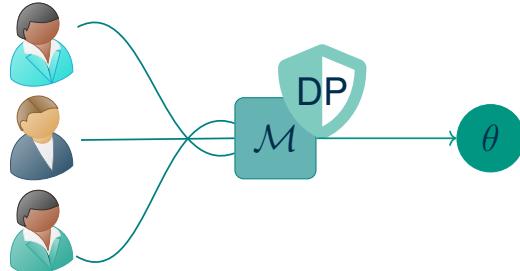
$$A \circ \mathcal{M} \text{ is } \varepsilon\text{-DP} \implies 1 - \beta \leq \sum_{D_-} e^\varepsilon \Pr_{A \circ \mathcal{M}}(y_n | D_{x_n}) \pi(D_- | x_n)$$

# Membership Inference Attack With Dependencies

The attacker receives  $\theta$  and aims to distinguish between:

$$H_0 : x_n$$

$$H_1 : y_n$$



$D_-$  is **unknown**:

$$H_0 = \{D : x_n \in D\} \text{ Vs. } H_1 = \{D : y_n \in D\}$$

Type I error:

$$\begin{aligned}\alpha &= \Pr_{A \circ \mathcal{M}}(y_n | x_n) \\ &= \sum_{D_-} \Pr_{A \circ \mathcal{M}}(y_n | D_{x_n}) \pi(D_- | x_n)\end{aligned}$$

Type II error:

$$\begin{aligned}\beta &= 1 - \Pr_{A \circ \mathcal{M}}(y_n | y_n) \\ &= 1 - \sum_{D_-} \Pr_{A \circ \mathcal{M}}(y_n | D_{y_n}) \pi(D_- | y_n)\end{aligned}$$

$$A \circ \mathcal{M} \text{ is } \varepsilon\text{-DP} \quad \Rightarrow \quad 1 - \beta \leq \sum_{D_-} e^\varepsilon \Pr_{A \circ \mathcal{M}}(y_n | D_{x_n}) \pi(D_- | x_n) = e^\varepsilon \sum_{D_-} \Pr_{A \circ \mathcal{M}}(y_n | D_{x_n}) \pi(D_- | y_n) \neq e^\varepsilon \alpha$$

# Standard DP Underestimates Participation Risk

Differential Privacy fails to measure privacy leakage under correlation

↔️ 😈 Empirically confirmed

⌚️ 🖊️ Theoretically exposed

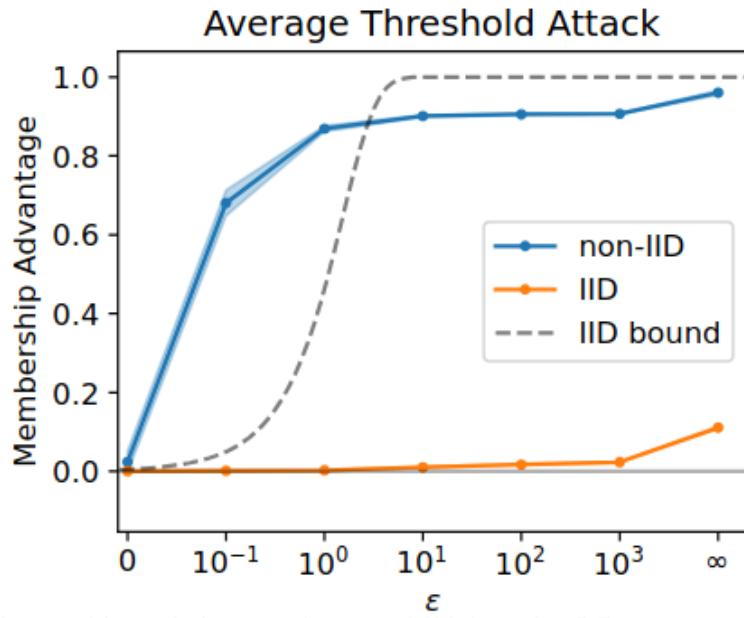


Figure: Humphries et al. 2023 MIA breaks DP guarantees.

# Standard DP Underestimates Participation Risk

Differential Privacy fails to measure privacy leakage under correlation

🔗 😈 Empirically confirmed

⌚ 📝 Theoretically exposed

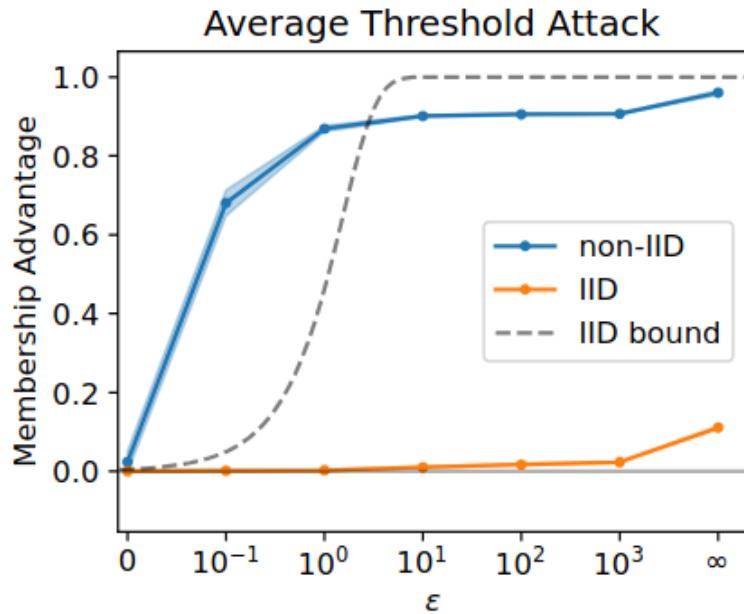
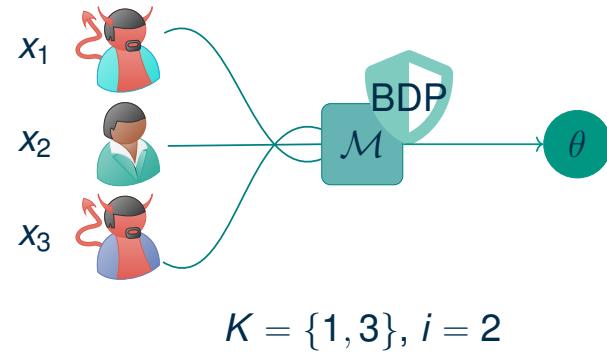


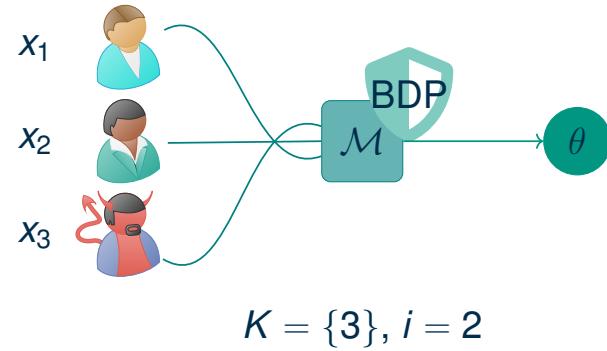
Figure: Humphries et al. 2023 MIA breaks DP guarantees.

New enhanced notion: Bayesian Differential Privacy

# Proposed Solution: Bayesian Differential Privacy



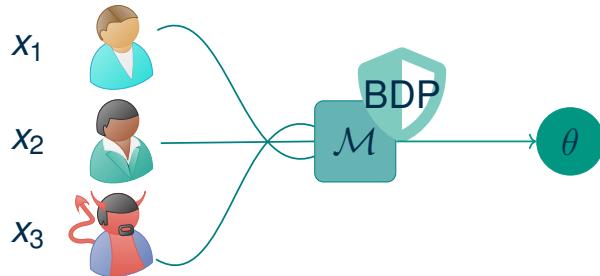
# Proposed Solution: Bayesian Differential Privacy



# Proposed Solution: Bayesian Differential Privacy

Bayesian DP leakage (Yang et al. 2017)

$$\text{BDPL}_{(K,i)} = \sup_{x_i, x'_i, \mathbf{x}_K, S} \ln \frac{\Pr_{\mathcal{M}}[Y \in S | \mathbf{X}_K = \mathbf{x}_K, X_i = x_i]}{\Pr_{\mathcal{M}}[Y \in S | \mathbf{X}_K = \mathbf{x}_K, X_i = x'_i]}, \text{ then } \varepsilon = \sup_{K,i} \text{BDPL}_{(K,i)}.$$



# Proposed Solution: Bayesian Differential Privacy

Bayesian DP leakage (Yang et al. 2017)

$$\text{BDPL}_{(K,i)} = \sup_{x_i, x'_i, \mathbf{x}_K, S} \ln \frac{\Pr_{\mathcal{M}}[Y \in S \mid \mathbf{X}_K = \mathbf{x}_K, X_i = x_i]}{\Pr_{\mathcal{M}}[Y \in S \mid \mathbf{X}_K = \mathbf{x}_K, X_i = x'_i]}, \text{ then } \varepsilon = \sup_{K,i} \text{BDPL}_{(K,i)}.$$

We check our MIA again, with  $H_0 : \textcolor{orange}{x_n} \in D, H_1 : \textcolor{blue}{y_n} \in D (K = \emptyset)$

# Proposed Solution: Bayesian Differential Privacy

Bayesian DP leakage (Yang et al. 2017)

$$\text{BDPL}_{(K,i)} = \sup_{x_i, x'_i, \mathbf{x}_K, S} \ln \frac{\Pr_{\mathcal{M}}[Y \in S \mid \mathbf{X}_K = \mathbf{x}_K, X_i = x_i]}{\Pr_{\mathcal{M}}[Y \in S \mid \mathbf{X}_K = \mathbf{x}_K, X_i = x'_i]}, \text{ then } \varepsilon = \sup_{K,i} \text{BDPL}_{(K,i)}.$$

We check our MIA again, with  $H_0 : \textcolor{orange}{x_n} \in D, H_1 : y_n \in D (K = \emptyset)$

Type I error:

$$\alpha = \Pr_{A \circ \mathcal{M}}(y_n \mid \textcolor{orange}{x_n})$$

Type II error:

$$\beta = 1 - \Pr_{A \circ \mathcal{M}}(y_n \mid y_n)$$

# Proposed Solution: Bayesian Differential Privacy

Bayesian DP leakage (Yang et al. 2017)

$$\text{BDPL}_{(K,i)} = \sup_{x_i, x'_i, \mathbf{x}_K, S} \ln \frac{\Pr_{\mathcal{M}}[Y \in S \mid \mathbf{X}_K = \mathbf{x}_K, X_i = x_i]}{\Pr_{\mathcal{M}}[Y \in S \mid \mathbf{X}_K = \mathbf{x}_K, X_i = x'_i]}, \text{ then } \varepsilon = \sup_{K,i} \text{BDPL}_{(K,i)}.$$

We check our MIA again, with  $H_0 : \mathbf{x}_n \in D, H_1 : y_n \in D (K = \emptyset)$

Type I error:

$$\alpha = \Pr_{A \circ \mathcal{M}}(y_n \mid \mathbf{x}_n)$$

Type II error:

$$\beta = 1 - \Pr_{A \circ \mathcal{M}}(y_n \mid y_n)$$

$$A \circ \mathcal{M} \text{ is } \varepsilon\text{-BDP} \Rightarrow 1 - \beta \leq e^\varepsilon \alpha \\ \alpha \leq e^\varepsilon (1 - \beta)$$

# Proposed Solution: Bayesian Differential Privacy

Bayesian DP leakage (Yang et al. 2017)

$$\text{BDPL}_{(K,i)} = \sup_{x_i, x'_i, \mathbf{x}_K, S} \ln \frac{\Pr_M[Y \in S | \mathbf{X}_K = \mathbf{x}_K, X_i = x_i]}{\Pr_M[Y \in S | \mathbf{X}_K = \mathbf{x}_K, X_i = x'_i]}, \text{ then } \varepsilon = \sup_{K,i} \text{BDPL}_{(K,i)}.$$

We check our MIA again, with  $H_0 : \textcolor{orange}{x_n} \in D, H_1 : y_n \in D (K = \emptyset)$

Type I error:

$$\alpha = \Pr_{A \circ M}(y_n | \textcolor{orange}{x_n})$$

Type II error:

$$\beta = 1 - \Pr_{A \circ M}(y_n | y_n)$$

$A \circ M$  is  
 $\varepsilon$ -BDP

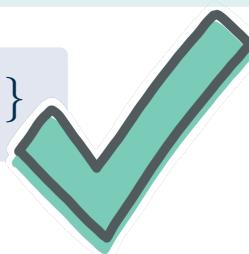


$$1 - \beta \leq e^\varepsilon \alpha$$



$$\alpha \leq e^\varepsilon (1 - \beta)$$

$$\beta \geq \max\{1 - e^\varepsilon \alpha, e^\varepsilon (1 - \alpha)\}$$



# Proposed Solution: Bayesian Differential Privacy

Bayesian DP leakage (Yang et al. 2017)

$$\text{BDPL}_{(K,i)} = \sup_{x_i, x'_i, \mathbf{x}_K, S} \ln \frac{\Pr_{\mathcal{M}}[Y \in S \mid \mathbf{X}_K = \mathbf{x}_K, X_i = x_i]}{\Pr_{\mathcal{M}}[Y \in S \mid \mathbf{X}_K = \mathbf{x}_K, X_i = x'_i]}, \text{ then } \varepsilon = \sup_{K,i} \text{BDPL}_{(K,i)}.$$

## Privacy

- ✓ Effective measure and resistance to correlation-based attacks.

# Proposed Solution: Bayesian Differential Privacy

## Bayesian DP leakage (Yang et al. 2017)

$$\text{BDPL}_{(K,i)} = \sup_{x_i, x'_i, \mathbf{x}_K, S} \ln \frac{\Pr_{\mathcal{M}}[Y \in S \mid \mathbf{X}_K = \mathbf{x}_K, X_i = x_i]}{\Pr_{\mathcal{M}}[Y \in S \mid \mathbf{X}_K = \mathbf{x}_K, X_i = x'_i]}, \text{ then } \varepsilon = \sup_{K,i} \text{BDPL}_{(K,i)}.$$

## Privacy

- ✓ Effective measure and resistance to correlation-based attacks.
- ✓ Good properties: post-processing & composition.
  - While other correlation-aware notions (General Pufferfish framework) don't!

# Proposed Solution: Bayesian Differential Privacy

## Bayesian DP leakage (Yang et al. 2017)

$$\text{BDPL}_{(K,i)} = \sup_{x_i, x'_i, \mathbf{x}_K, S} \ln \frac{\Pr_{\mathcal{M}}[Y \in S | \mathbf{X}_K = \mathbf{x}_K, X_i = x_i]}{\Pr_{\mathcal{M}}[Y \in S | \mathbf{X}_K = \mathbf{x}_K, X_i = x'_i]}, \text{ then } \varepsilon = \sup_{K,i} \text{BDPL}_{(K,i)}.$$

### Privacy

- ✓ Effective measure and resistance to correlation-based attacks.
- ✓ Good properties: post-processing & composition.
  - While other correlation-aware notions (General Pufferfish framework) don't!

### Utility

- ✗ Poor utility (methods based on group privacy).

# Proposed Solution: Bayesian Differential Privacy

## Bayesian DP leakage (Yang et al. 2017)

$$\text{BDPL}_{(K,i)} = \sup_{x_i, x'_i, \mathbf{x}_K, S} \ln \frac{\Pr_{\mathcal{M}}[Y \in S \mid \mathbf{X}_K = \mathbf{x}_K, X_i = x_i]}{\Pr_{\mathcal{M}}[Y \in S \mid \mathbf{X}_K = \mathbf{x}_K, X_i = x'_i]}, \text{ then } \varepsilon = \sup_{K,i} \text{BDPL}_{(K,i)}.$$

### Privacy

- ✓ Effective measure and resistance to correlation-based attacks.
- ✓ Good properties: post-processing & composition.
  - While other correlation-aware notions (General Pufferfish framework) don't!

### Utility

- ✗ Poor utility (methods based on group privacy).
- ✗ Computationally intractable methods (computing the Wasserstein distance).

# Proposed Solution: Bayesian Differential Privacy

## Bayesian DP leakage (Yang et al. 2017)

$$\text{BDPL}_{(K,i)} = \sup_{x_i, x'_i, \mathbf{x}_K, S} \ln \frac{\Pr_{\mathcal{M}}[Y \in S \mid \mathbf{X}_K = \mathbf{x}_K, X_i = x_i]}{\Pr_{\mathcal{M}}[Y \in S \mid \mathbf{X}_K = \mathbf{x}_K, X_i = x'_i]}, \text{ then } \varepsilon = \sup_{K,i} \text{BDPL}_{(K,i)}.$$

### Privacy

- ✓ Effective measure and resistance to correlation-based attacks.
- ✓ Good properties: post-processing & composition.
  - While other correlation-aware notions (General Pufferfish framework) don't!

### Utility

- ✗ Poor utility (methods based on group privacy).
- ✗ Computationally intractable methods (computing the Wasserstein distance).
- ✗ Limited applicability (lazy, binary, stationary Markov chains).

# Our Research Question

Can we reduce utility loss while still retaining the privacy guarantees of BDP?

**Our methodology:** Understanding how DP leakage relates to BDP leakage:

$$\varepsilon\text{-DP} \Rightarrow ??\text{-BDP}.$$

# Against Arbitrary Correlations It Is Impossible

Kifer and Machanavajjhala 2014:

Pufferfish (including BDP)

&

$\Rightarrow$

Free-lunch Privacy

$\Rightarrow$

No utility.

arbitrary correlation

# Against Arbitrary Correlations It Is Impossible

Kifer and Machanavajjhala 2014:

Pufferfish (including BDP)  
&  $\Rightarrow$  Free-lunch Privacy  $\Rightarrow$  No utility.  
arbitrary correlation

We express this in term of  $(\alpha, \beta)$ -accuracy for any numerical target query  $f$ :

$(\alpha, \beta)$ -accuracy

$$\Pr(|f(D) - \mathcal{M}(D)| \geq \alpha) \leq \beta$$

$1 - \beta = \text{confidence}$

$\alpha = \text{error interval}$



# Against Arbitrary Correlations It Is Impossible

Kifer and Machanavajjhala 2014:

Pufferfish (including BDP)  
&  $\Rightarrow$  Free-lunch Privacy  $\Rightarrow$  No utility.  
arbitrary correlation

We express this in term of  $(\alpha, \beta)$ -accuracy for any numerical target query  $f$ :

$(\alpha, \beta)$ -accuracy

$$\Pr(|f(D) - \mathcal{M}(D)| \geq \alpha) \leq \beta$$

$1 - \beta = \text{confidence}$

$\alpha = \text{error interval}$

**Our result (informal):**

$$\beta < \frac{1}{e^\varepsilon + 1} \Rightarrow \alpha > \frac{1}{2}\text{Range}(f).$$



# Against Arbitrary Correlations It Is Impossible

Kifer and Machanavajjhala 2014:

Pufferfish (including BDP)  
&  $\Rightarrow$  Free-lunch Privacy  $\Rightarrow$  No utility.  
arbitrary correlation

We express this in term of  $(\alpha, \beta)$ -accuracy for any numerical target query  $f$ :

$(\alpha, \beta)$ -accuracy

$$\Pr(|f(D) - \mathcal{M}(D)| \geq \alpha) \leq \beta$$

$1 - \beta = \text{confidence}$

$\alpha = \text{error interval}$

**Our result (informal):**

$$\beta < \frac{1}{e^\varepsilon + 1} \Rightarrow \alpha > \frac{1}{2}\text{Range}(f).$$



# Against Arbitrary Correlations It Is Impossible

Kifer and Machanavajjhala 2014:

Pufferfish (including BDP)

&

$\Rightarrow$

Free-lunch Privacy

$\Rightarrow$

No utility.

arbitrary correlation

We express this in term of  $(\alpha, \beta)$ -accuracy for any numerical target query  $f$ :

$(\alpha, \beta)$ -accuracy

$$\Pr(|f(D) - M(D)| \geq \alpha) \leq \beta$$

$1 - \beta = \text{confidence}$

$\alpha = \text{error interval}$

**Our result (informal):**

$$\beta < \frac{1}{e^\varepsilon + 1} \Rightarrow \alpha > \frac{1}{2}\text{Range}(f).$$



# Few Correlated Records, Same Disaster

## Our result (informal)

Privacy decreases linearly proportional to number of correlated records:

$$\varepsilon\text{-DP} \Rightarrow m\varepsilon\text{-BDP}$$

## How does it impact utility?

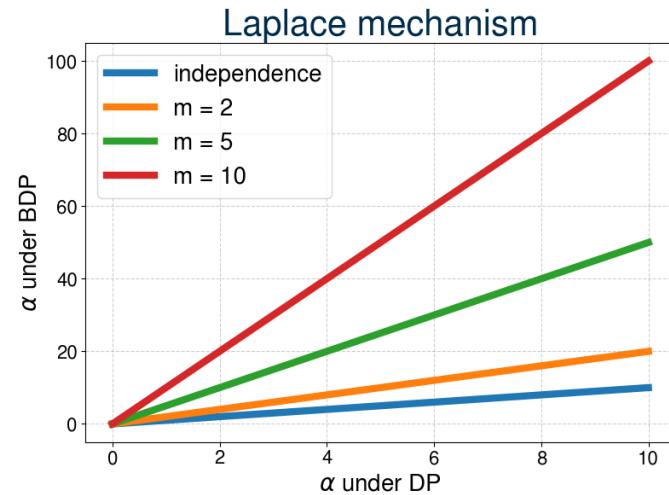


Figure: For the same confidence level, the upper bound on the query error  $\alpha$  increases sharply.

# Few Correlated Records, Same Disaster

## Our result (informal)

Privacy decreases linearly proportional to number of correlated records:

$$\varepsilon\text{-DP} \Rightarrow m\varepsilon\text{-BDP}$$

This result is tight! Even if  $\rho \rightarrow 0$ .

## How does it impact utility?

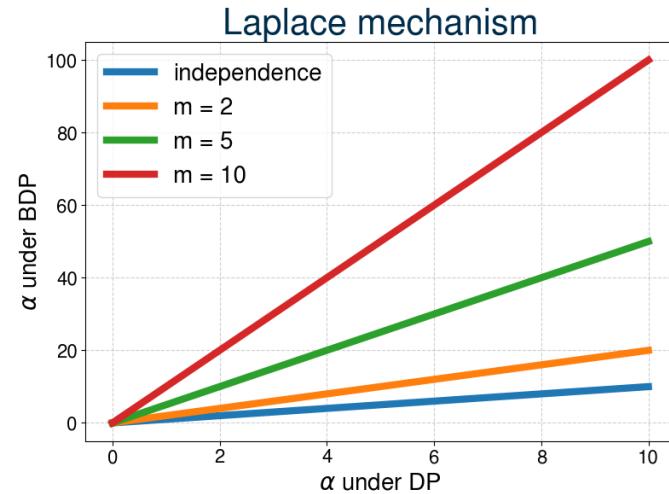


Figure: For the same confidence level, the upper bound on the query error  $\alpha$  increases sharply.

# Few Correlated Records, Same Disaster

## Our result (informal)

Privacy decreases linearly proportional to number of correlated records:

$$\varepsilon\text{-DP} \Rightarrow m\varepsilon\text{-BDP}$$

This result is tight! Even if  $\rho \rightarrow 0$ .

## Conclusion:

We need to target specific correlation models  $\pi$  to obtain utility

## How does it impact utility?

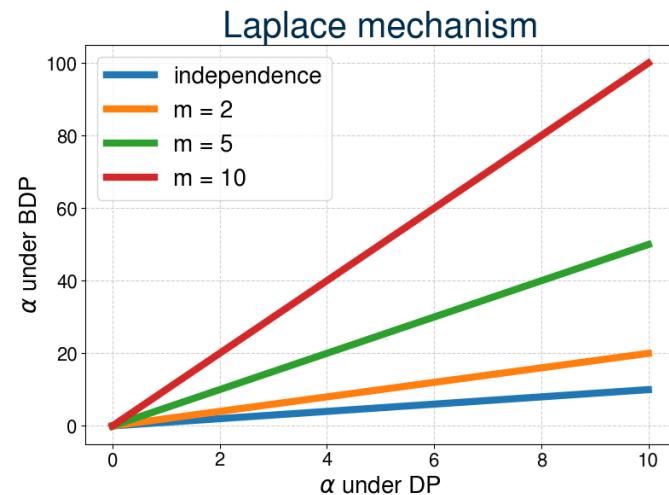


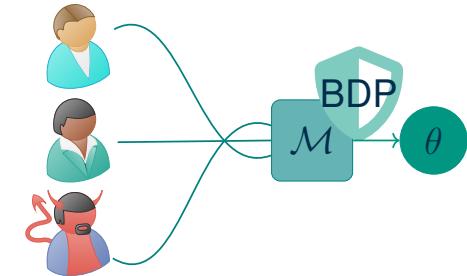
Figure: For the same confidence level, the upper bound on the query error  $\alpha$  increases sharply.

# New Strategy

Adjust the noise of DP mechanisms to obtain useful BDP mechanisms targeting specific priors  $\pi$ .

## Assumptions:

- Global setting: All data is collected by a trusted data curator that applies the mechanism.
- The attacker does not have more knowledge about  $\pi$  than the data curator.

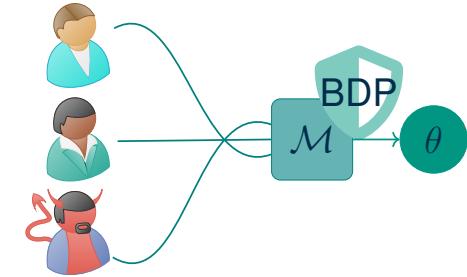


# New Strategy

Adjust the noise of DP mechanisms to obtain useful BDP mechanisms targeting specific priors  $\pi$ .

## Assumptions:

- Global setting: All data is collected by a trusted data curator that applies the mechanism.
- The attacker does not have more knowledge about  $\pi$  than the data curator.



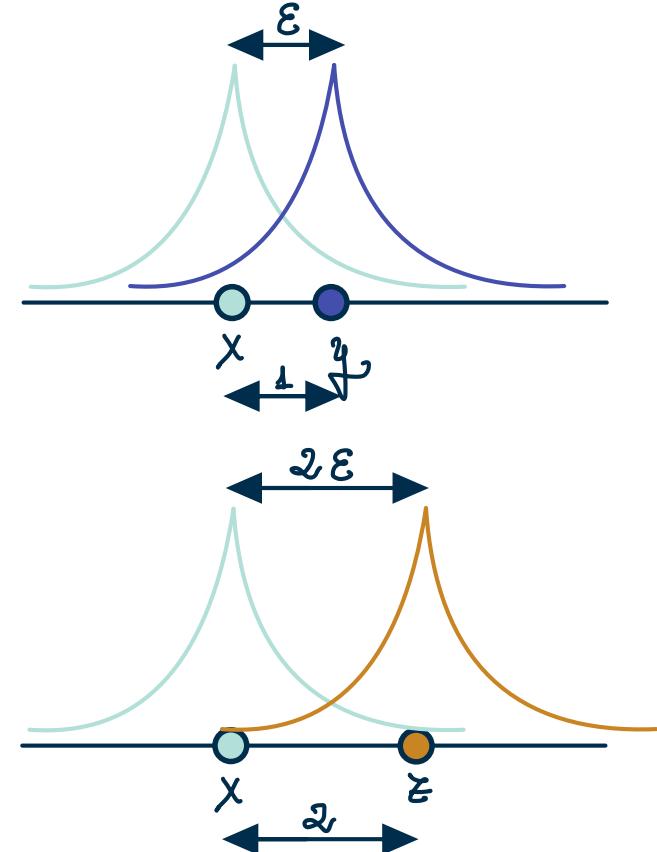
Multivariate Gaussian

Markov Chains

# Multivariate Gaussian Correlation

## Main Result (Informal)

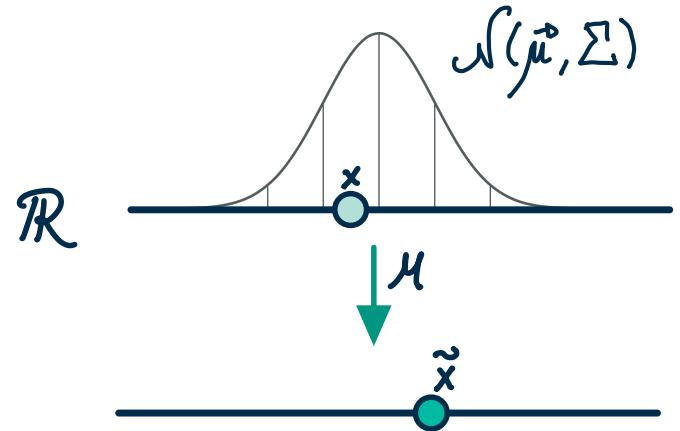
- Let  $\mathcal{M}$  be an  $\varepsilon\ell_1$ -private mechanism,



# Multivariate Gaussian Correlation

## Main Result (Informal)

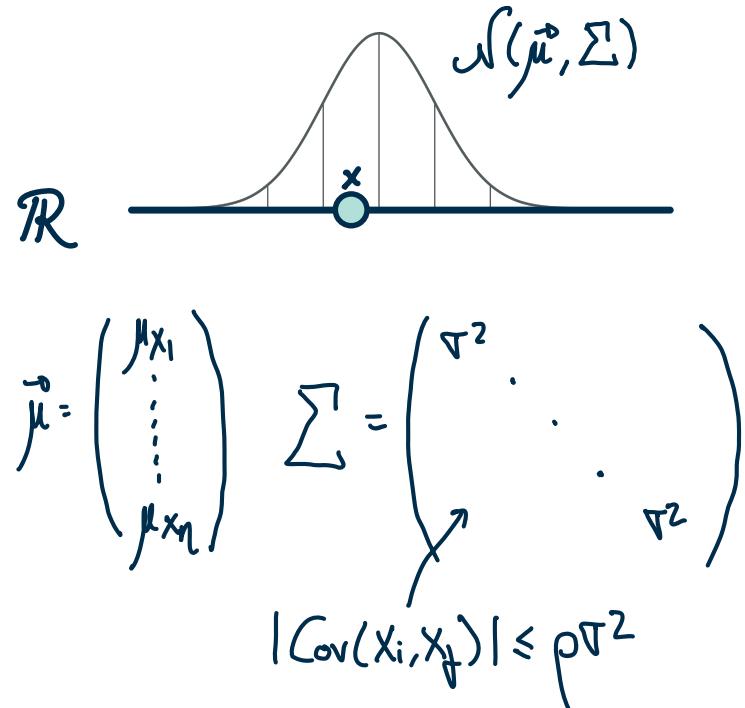
- Let  $\mathcal{M}$  be an  $\epsilon\ell_1$ -private mechanism,
- input data drawn from a multivariate Gaussian distribution



# Multivariate Gaussian Correlation

## Main Result (Informal)

- Let  $\mathcal{M}$  be an  $\varepsilon\ell_1$ -private mechanism,
- input data drawn from a multivariate Gaussian distribution
- $\rho(m - 2) < 1$  is the maximum correlation coefficient.

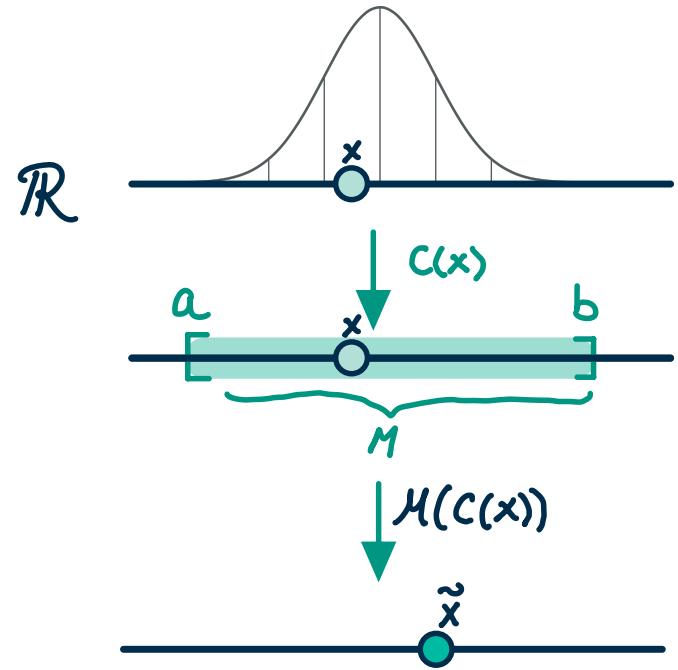


# Multivariate Gaussian Correlation

## Main Result (Informal)

- Let  $\mathcal{M}$  be an  $\varepsilon\ell_1$ -private mechanism,
- input data drawn from a multivariate Gaussian distribution
- $\rho(m - 2) < 1$  is the maximum correlation coefficient.

Then, using clipping as preprocessing step,  $c_i(D)_i = \max(a, \min(b, D_i))$ , we obtain  $\mathcal{M}'$ , satisfying



# Multivariate Gaussian Correlation

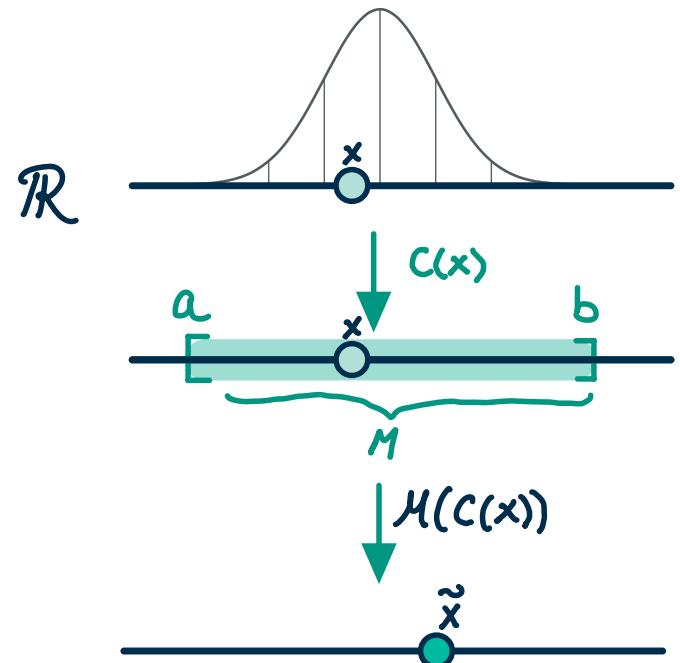
## Main Result (Informal)

- Let  $\mathcal{M}$  be an  $\varepsilon\ell_1$ -private mechanism,
- input data drawn from a multivariate Gaussian distribution
- $\rho(m - 2) < 1$  is the maximum correlation coefficient.

Then, using clipping as preprocessing step,  $c_l(D)_i = \max(a, \min(b, D_i))$ , we obtain  $\mathcal{M}_l$  satisfying

$$\text{BDPL}(\mathcal{M}_l) \leq \left( \frac{m^2}{4(\frac{1}{\rho} - m + 2)} + 1 \right) M\varepsilon.$$

where  $M$  is the diameter of the interval  $I = [a, b]$



# Multivariate Gaussian Correlation (Impact on Real Databases)

- **Use-case:** Sum queries with Laplace mechanism.  $\theta = f(D) + Z$  with  $Z \sim \text{Lap}(b)$ .

# Multivariate Gaussian Correlation (Impact on Real Databases)

- **Use-case:** Sum queries with Laplace mechanism.  $\theta = f(D) + Z$  with  $Z \sim \text{Lap}(b)$ .
- **Strategy:** We calibrate  $b$  to obtain BDP using our theorem.

# Multivariate Gaussian Correlation (Impact on Real Databases)

- **Use-case:** Sum queries with Laplace mechanism.  $\theta = f(D) + Z$  with  $Z \sim \text{Lap}(b)$ .
- **Strategy:** We calibrate  $b$  to obtain BDP using our theorem.
- **Utility metric:** We set  $\beta = 0.05$  (i.e., 95% confidence) and measure  $(\alpha, \beta)$ -accuracy, both theoretically (–) and empirically (×).

# Multivariate Gaussian Correlation (Impact on Real Databases)

- **Use-case:** Sum queries with Laplace mechanism.  $\theta = f(D) + Z$  with  $Z \sim \text{Lap}(b)$ .
- **Strategy:** We calibrate  $b$  to obtain BDP using our theorem.
- **Utility metric:** We set  $\beta = 0.05$  (i.e., 95% confidence) and measure  $(\alpha, \beta)$ -accuracy, both theoretically (—) and empirically (×).

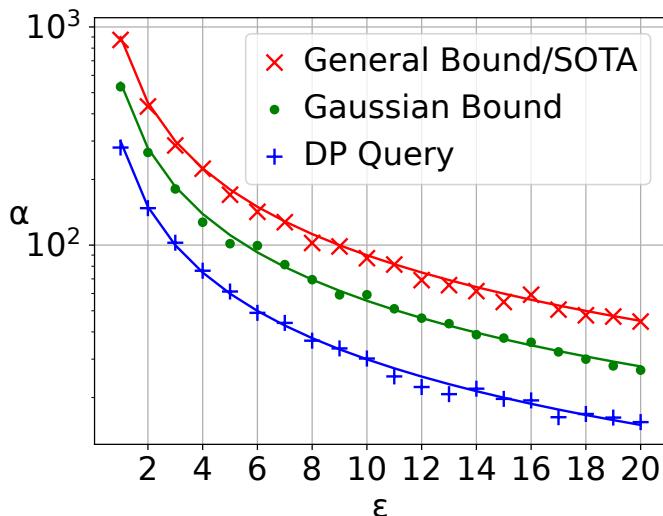


Figure: Galton,  $n = 897$   $m = 3$

## Key takeaway:

**Substantial utility gains** compared to the general bound!

- More experiments with different real and synthetic datasets in our paper show similar results.

# Markov Chain Correlation Model

## Main result (Informal)

- Let  $\mathcal{M}$  be an  $\varepsilon$ -DP mechanism,
- input data sampled from Markov chain with transition matrix  $P \in \mathbb{R}^{s \times s}$  and initial distribution  $w \in \mathbb{R}^s$  with the following properties:

$$(H1) \text{ For all } x, y \in \mathcal{S} \text{ we have } P_{x,y} > 0 \text{ and,} \quad (H2) \text{ } wP = w.$$

Then,  $\mathcal{M}$  is an  $(\varepsilon + 4 \ln \gamma)$ -BDP mechanism where  $\gamma = \frac{\max_{x,y \in \mathcal{S}} P_{xy}}{\min_{x,y \in \mathcal{S}} P_{xy}}$ .

# Markov Chain Correlation Model

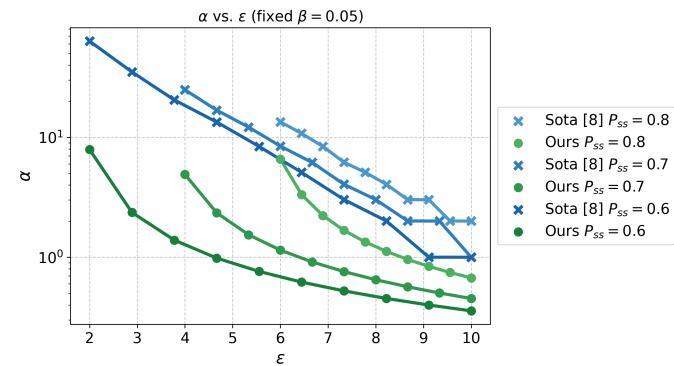
## Main result (Informal)

- Let  $\mathcal{M}$  be an  $\varepsilon$ -DP mechanism,
- input data sampled from Markov chain with transition matrix  $P \in \mathbb{R}^{s \times s}$  and initial distribution  $w \in \mathbb{R}^s$  with the following properties:

$$(H1) \text{ For all } x, y \in \mathcal{S} \text{ we have } P_{x,y} > 0 \text{ and,} \quad (H2) \text{ } wP = w.$$

Then,  $\mathcal{M}$  is an  $(\varepsilon + 4 \ln \gamma)$ -BDP mechanism where  $\gamma = \frac{\max_{x,y \in \mathcal{S}} P_{xy}}{\min_{x,y \in \mathcal{S}} P_{xy}}$ .

Previous mechanism	Ours
$P_{xy} > 0$	$P_{xy} > 0$
stationary	stationary
lazy	
binary	
symmetric	
$\varepsilon' > 0$	$\varepsilon' > 4 \ln(\gamma)$



# Markov Chain Correlation Model (Impact on Real Databases)

- **Use-case:** Counting queries with Laplace mechanism.  $\theta = f(D) + Z$  with  $Z \sim \text{Lap}(b)$ .

# Markov Chain Correlation Model (Impact on Real Databases)

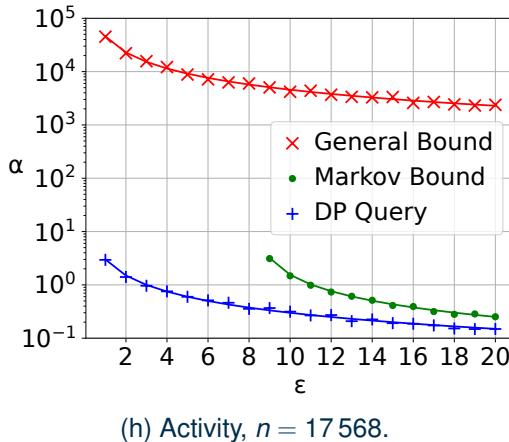
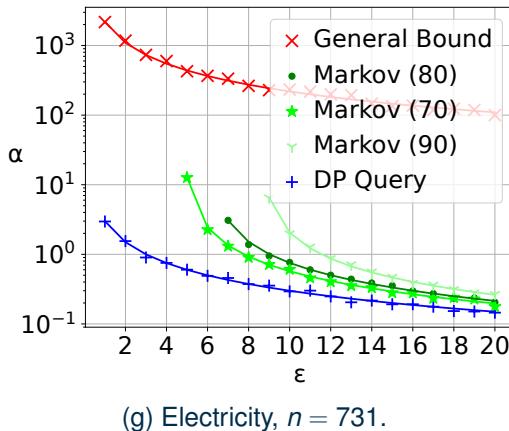
- **Use-case:** Counting queries with Laplace mechanism.  $\theta = f(D) + Z$  with  $Z \sim \text{Lap}(b)$ .
- **Strategy:** We calibrate  $b$  to obtain BDP using our theorem.

# Markov Chain Correlation Model (Impact on Real Databases)

- **Use-case:** Counting queries with Laplace mechanism.  $\theta = f(D) + Z$  with  $Z \sim \text{Lap}(b)$ .
- **Strategy:** We calibrate  $b$  to obtain BDP using our theorem.
- **Utility metric:** We set  $\beta = 0.05$  (i.e., 95% confidence) and measure  $-(\alpha, \beta)$ -accuracy,  $\times$  upper bound of a  $(1 - \beta)$  confidence interval for the absolute query error.

# Markov Chain Correlation Model (Impact on Real Databases)

- **Use-case:** Counting queries with Laplace mechanism.  $\theta = f(D) + Z$  with  $Z \sim \text{Lap}(b)$ .
- **Strategy:** We calibrate  $b$  to obtain BDP using our theorem.
- **Utility metric:** We set  $\beta = 0.05$  (i.e., 95% confidence) and measure  $-(\alpha, \beta)$ -accuracy,  $\times$  upper bound of a  $(1 - \beta)$  confidence interval for the absolute query error.



## Key takeaway:

- Substantial utility gains compared to the general bound!
- Markov bound independent of  $n$   $\Rightarrow$  huge improvement for large datasets.

# Conclusion

- ✓ We provide a **feasible method** to generate a **BDP mechanism** by **recalibrating** existing DP methods, tailored to **Gaussian** and **Markov** models.

# Conclusion

- ✓ We provide a **feasible method** to generate a **BDP mechanism** by **recalibrating** existing DP methods, tailored to **Gaussian** and **Markov** models.
- ✓ We offer **significantly better utility than prior results**.

# Conclusion

- ✓ We provide a **feasible method** to generate a **BDP mechanism** by **recalibrating** existing DP methods, tailored to **Gaussian** and **Markov** models.
- ✓ We offer **significantly better utility than prior results**.

## Key takeaway:

BDP becomes usable when correlations are structured.



# Conclusion

- ✓ We provide a **feasible method** to generate a **BDP mechanism** by **recalibrating** existing DP methods, tailored to **Gaussian** and **Markov** models.
- ✓ We offer **significantly better utility than prior results**.

## Key takeaway:

BDP becomes usable when correlations are structured.

## Future Work:

- Other distributions ?
- Can we build methods from scratch instead of recycling ?
- What if we calibrate directly to the attack advantage ?



# Conclusion

- ✓ We provide a **feasible method** to generate a **BDP mechanism** by **recalibrating** existing DP methods, tailored to **Gaussian** and **Markov** models.
- ✓ We offer **significantly better utility than prior results**.

## Key takeaway:

BDP becomes usable when correlations are structured.

## Future Work:

- Other distributions ?
- Can we build methods from scratch instead of recycling ?
- What if we calibrate directly to the attack advantage ?



Paper



Code

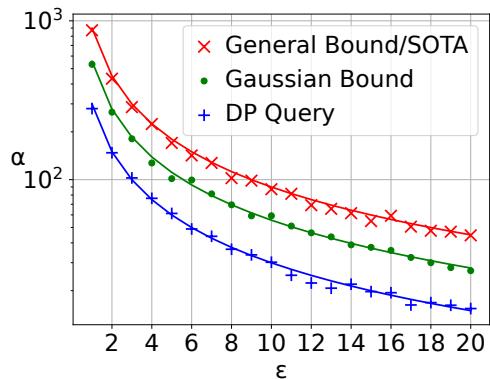
# Backup Slides

# Experiment Details

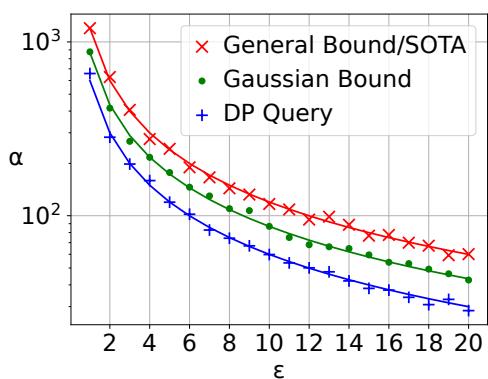
Database	n	m	Parameters	Sensitivity
Galton	897	3	$\rho = 0.275$	$\Delta q = 254cm$
FamilyIQ	868	2	$\rho = 0.4483$	$\Delta q = 120$
SyntheticIQ	20000	2	$\rho = 0.45$	$\Delta q = 120$
Activity	17568	<i>n</i>	$\gamma = 7.54$	$\Delta q = 1$
Activity Single Day	288	<i>n</i>	$\gamma = 7.54$	$\Delta q = 1$
Electricity	731	<i>n</i>	70 kWh, $\gamma = 3.29$	
			80 kWh, $\gamma = 4.49$	$\Delta q = 1$
			90 kWh, $\gamma = 8.43$	

Table: Data description.  $m$  is the max number of correlated records and  $n$  the total amount.

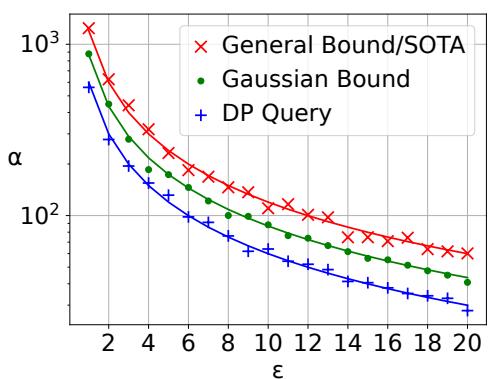
# Multivariate Gaussian More Results



(i) Galton,  $n = 897$   $m = 3$



(j) FamilyIQ,  $n = 868$ ,  $m = 2$ .



(k) SyntheticIQ,  $n = 20000$ ,  $m = 2$ .

Figure: Gaussian data results. Lines show theoretical error at  $\beta = 5\%$  and markers indicate empirical 95% upper bounds.