

CAB301 ASSIGNMENT

Patrick Cummins n10627138



Table of Contents

Introduction	
Algorithm Analysis	4
Pseudocode	4
Theoretical Analysis	5
Empirical Analysis	7
Software Tests	c

Introduction

This project is a software management system for a tool library. The system manages the tools as well as the members of the library. It allows for staff to organise, add, remove both tools and members. It allows members to hire and return tools as well as other quality of life functionality. Chief amongst the features is the ability for a member to view the top three most borrowed tools, which is implemented using a *Heap Sort* algorithm.

The purpose of this report is to demonstrate the design and complexity of the implemented Heap Sort algorithm, explaining its implementation as well as its efficiency. Further, the report will outline the tests plans and results for the specified functionality, demonstrating that the implementation is correct.

The report will give the pseudo code of the HeapSort algorithm, present an analysis of the algorithm, showings its time complexity and efficiency.

The report will then demonstrate the testing plans and results.

Algorithm Analysis

Pseudocode

```
Heap Bottom-Up Pseudocode
HeapBottomUp(H[1...n])
// Constructs a heap from elements of a given array
// by the bottom-up algorithm
// Input: An array H[1...n] of orderable items
// Output: A heap H[1...n]
for i <- [n / 2] down to 1 do
       k <- i; v <- H[k]
       heap <- false
       while not heap and 2 * k \le n do
       i < -2*k
       if j < n // there are two children</pre>
               if H[j] < H[j + 1] j < -j + 1
       if v \ge H[j]
               heap <- true
       else H[k] <- H[j]; k <- j
    H[k] \leftarrow v
```

Heapsort Pseudocode

Max Key Delete Pseudocode

```
MaxKeyDelete(A[0...n-1], size)

// Exchange the root's key with the last key k of the heap temp <- A[0]

A[0] <- A[\text{size} - 1]

A[\text{size} - 1] <- \text{temp}

// Decrease the heap's size by 1
```

```
n <- size – 1
// Heapify the complete binary tree
heap <- false
k <- 0
v <- A[0]
while not heap and 2 * k + 1 <- n - 1
        j \leftarrow 2 * k + 1 // left child of k
        if j < n - 1 // \text{ key has two children}
                 if A[j] < A[j + 1]
                 j <- j + 1 // j is the larger child of k
        if v \ge A[j]
                 heap <- true
        else
                 A[k] \leftarrow A[j]
                 K <- j
A[k] \leftarrow v
return A[n]
```

Theoretical Algorithm Efficiency

```
Heap Bottom-Up
for i <- [n / 2] down to 1 do
       k < -i O(1)
       v <- H[k] O(1)
       heap <- false O(1)
       while not heap and 2 * k \le n do
              j < -2* k O(1)
              if j < n // there are two children</pre>
                      if H[j] < H[j + 1] O(1)
                             j < -j + 1 O(1)
              if v \ge H[j]
                                                                       O(log n)
                      heap <- true O(1)
              else
                      H[k] <- H[j] O(1)
                      k < -i O(1)
     H[k] <- v O(1)
Time Complexity = O(n log n)
Max Key Delete
temp <- A[0] O(1)
A[0] <- A[size - 1] O(1)
A[\text{size} - 1] \leftarrow \text{temp } O(1)
// Decrease the heap's size by 1
n < -size - 10(1)
// Heapify the complete binary tree
heap <- false O(1)
k < -00(1)
v < -A[0] O(1)
while not heap and 2 * k + 1 < -n - 1
       j < -2 * k + 1 // left child of k O(1)
       if j < n - 1 // \text{ key has two children}
              if A[j] < A[j + 1] O(1)
              i < -i + 1 // is the larger child of k O(1)
                                                                       O(log n)
       if v \ge A[j]
              heap <- true O(1)
       else
              A[k] <- A[j] O(1)
              K <- j O(1)
A[k] <- v O(1)
return A[n] O(1)
Time Complexity = O(log n)
```

Heap Sort

 $B[0...n-1] <- \emptyset O(1)$

HeapBottomUp() O(n log n)

for $v \leftarrow 0$ to 3 do

Use the MaximumKeyDeletion procedure to delete the root of the heap $O(\log n)$ B[0...n-1] <- v O(1)

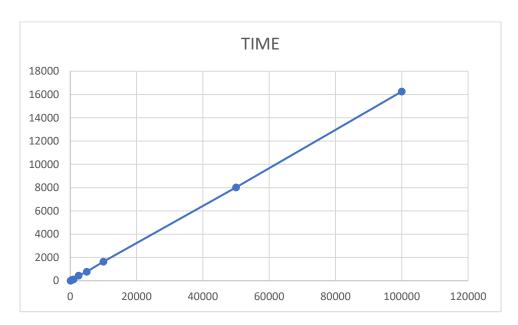


return B[0...n-1] O(1)

Time Complexity = O(n log n)

Empirical Algorithm Efficiency

To correctly test the algorithm empirically, multiple samples of arrays were created, loaded with random values. The program was setup to run the algorithm within a timer which would relay the time taken to complete the algorithm. The observed data was recorded and charted into a scatter plot seen below.



ARRAY SIZE	TIME m/s
100	10
500	71
1000	119
5000	778
10000	1641
50000	8021
100000	16260

An analysis of the data is as follows, taking two instances of t(2n) over t(n), their ratio was calculated to produce the efficiency of the algorithm.

$$\frac{t(10000)}{t(5000)} = \frac{1641}{778} = 2.10 \qquad \frac{t(100000)}{t(50000)} = \frac{16260}{8021} = 2.02$$
$$\frac{t(2n)}{t(n)} = \frac{(2n)\log(n)}{n\log n} = \frac{2\log(2n)}{\log n}$$

The mathematical results taken from the empirical tests show that the ratio of t(2n) and t(n) is greater than two, making it less efficient than O(n).

The ratio is less than four, making it more efficient than O(n²).

Therefore the efficiency class of the algorithm is O(n log n).

Software Tests

Staff Login

```
Please enter staff username...
staff
```

Add a new tool

Add pieces to a tool

Remove pieces from a tool

Add a new member

Remove a member

Find a member's number

Member Login

=== Member Login ===
Please enter first name
pat
Please enter last name
pat
Please enter PIN
1234_

Borrow tool

Display tools by category/type

Display of loois by Type	Line Trimmers
1: Gardening Tools 2: Flooring Tools 3: Fencing Tools 4: Measuring Tools 5: Cleaning Tools 6: Painting Tools 7: Electronic Tools 8: Electricity Tools 9: Automotive Tools Select option from menu (0 to exit):	Tool Name: tiny trimmer Available Quantity: 1 === Tool Name: small trimmer Available Quantity: 1 === Tool Name: medium trimmer Available Quantity: 1 === Tool Name: big trimmer Available Quantity: 1 ===
	Tool Name: huge trimmer Available Quantity: 1 === Pess any key to return to member menu

Return Tool

Display borrowed Tools

Top 3 Borrowed Tools

Please press any key to return to member menu

```
=== Top Three Most Borrowed Tools ===
#1: huge trimmer
Times Borrowed: 4
===
#2: big trimmer
Times Borrowed: 3
===
#3: medium trimmer
Times Borrowed: 2
===
```