Data Vault's XTS pattern on Snowflake



A problem that could occur when loading batch files to the enterprise data platform is the arrival of data *out of sequence*, yes data vault's satellite tables are change capture tables. That is, only new records are loaded to the target satellite table either if:

- We have never seen that record before for that *parent* key (hub or link).
- Or the record differs from the current record we have in the satellite for that *parent* key. The keyword here is *current*.

What do we do with a batch of records that arrive that are *older* than the current records in the target satellite table? Do we *rewind and replay* the records in the correct order? What of the downstream information marts, dashboards, reports... oh my!? Manual operations to fix the data order incur human a cost to identify the error and to correct the data in place. Alternatively, automation identifies that the batch is not the right sequence and human intervention is again needed to pause processing, and then process the batch in the correct order. In both scenarios there is a delay in analytic value.

Is there another way? Yes, there is, let's describe it below.

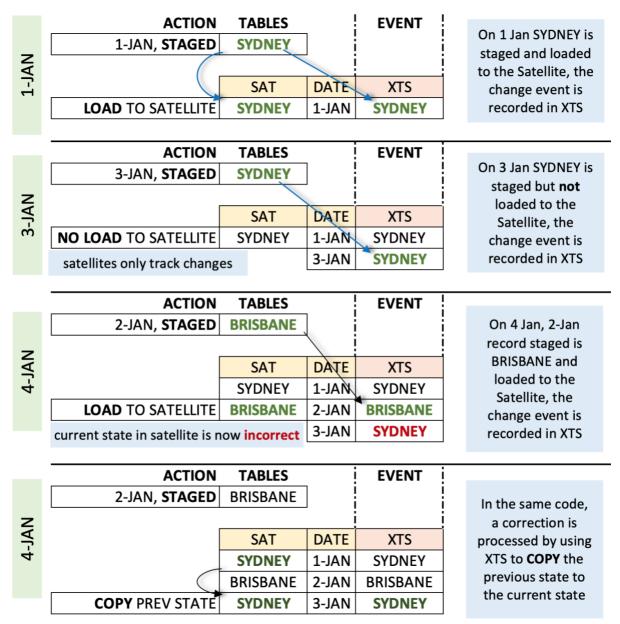


Figure 0-1 Load Date in Green on the left, Applied Date in Staged and Satellite

In Data Vault we have a pattern to solve this, with the help of the eXtended Record Tracking Satellites (or XTS). For a full discussion on this pattern please visit: bit.ly/3y4mUdV. XTS stores only the HashDiff for every record, the satellite has the complete change record.

What this article is designed to show is how the XTS pattern performs in Snowflake. With this Data Vault XTS *architecture* and *model* pattern the enterprise data platform is suddenly **self-healing**. It doesn't mean however that you should leave the platform to *excessively* heal itself! You should investigate as to why data is arriving out of sequence and look to resolve that issue!

In short this is what the XTS pattern looks like, XTS is used to monitor and manage adjacent satellite tables around a hub or a link.

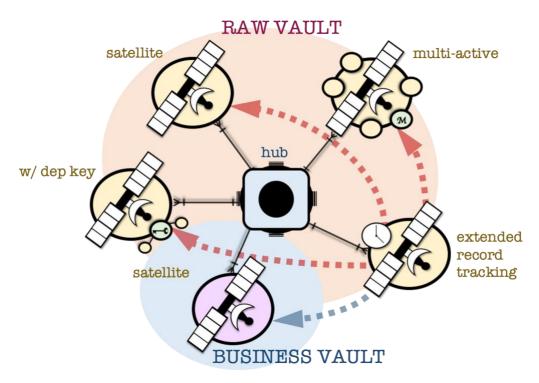
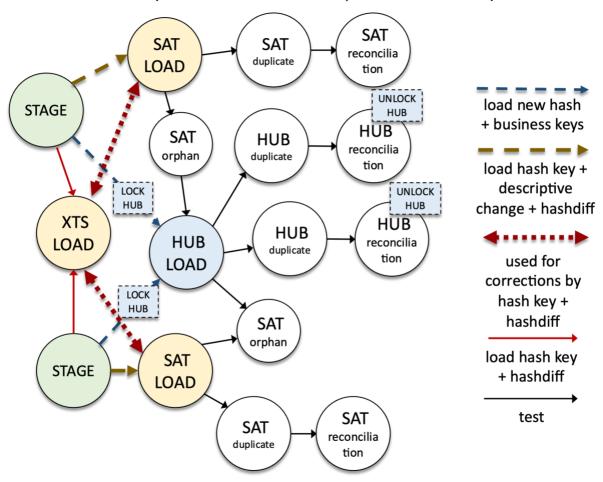


Figure 0-2 XTS can assist in correcting multiple satellite table types

As the modelled XTS is designed to facilitate timeline correction, its implementation within an automation architecture remains scalable and robust, observe:

one hub / link for shared satellites, one XTS for a hub / link



shared hub requires single-thread lock, XTS does not

Figure 0-3 Load DAGs

This extends on the article describing data vault test automation patterns described here: bit.ly/3dUHPIS

XTS satellite can be used with raw and business vault satellites; in addition, status tracking and record tracking satellites can also be used with this pattern and an XTS satellite can be hung off a hub or a link table tracking and assisting with satellite timeline correction events. All that is recommended is that you add a new data vault tag to your satellite tables: dv_xts_event to track if the record was either 'inserted' straight from staging or 'copied', the latter of which is the correction event.

HASHKEY	LOADATE	APPLIEDDATE	HASHDIFF	XTS_EVENT	CITYNAME
hash(1)	1-JAN	1-JAN	hash(SYDNEY)	insert	SYDNEY
hash(2)	1-JAN	1-JAN	hash(SYDNEY)	insert	SYDNEY
hash(2)	3-JAN	3-JAN	hash(BRISBANE)	insert	BRISBANE
hash(2)	4-JAN	2-JAN	hash(MELBOURNE)	insert	MELBOURNE
hash(3)	1-JAN	1-JAN	hash(SYDNEY)	insert	SYDNEY
hash(3)	3-JAN	3-JAN	hash(BRISBANE)	insert	BRISBANE
hash(4)	1-JAN	1-JAN	hash(SYDNEY)	insert	SYDNEY

HASHKEY	LOADATE	APPLIEDDATE	HASHDIFF	XTS_EVENT	CITYNAME
hash(4)	4-JAN	2-JAN	hash(BRISBANE)	insert	BRISBANE
hash(4)	4-JAN	3-JAN	hash(SYDNEY)	correction	SYDNEY
hash(5)	1-JAN	1-JAN	hash(SYDNEY)	insert	SYDNEY
hash(5)	3-JAN	3-JAN	hash(BRISBANE)	insert	BRISBANE
hash(5)	4-JAN	2-JAN	hash(BRISBANE)	insert	BRISBANE

Table 0-1 Refer to "Time Crime" article for scenario descriptions

All that aside, let's get into the loading statistics!

First Load (empty satellites), staged on the left, normal satellite in the middle and an XTS-influenced satellite on the right, all stats are based on XSMALL Warehouse size.

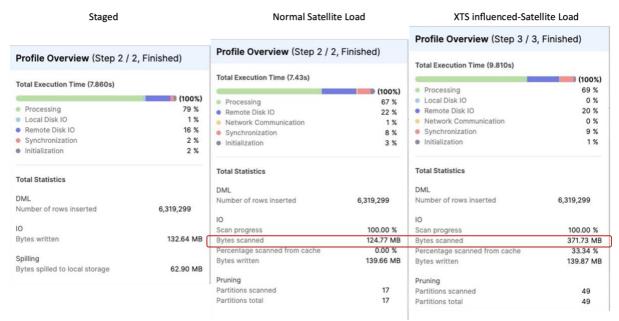


Figure 0-4 Note the bytes scanned, XTS includes both Staging and XTS to load a satellite

Now let's see what happens when we load the next day's data.



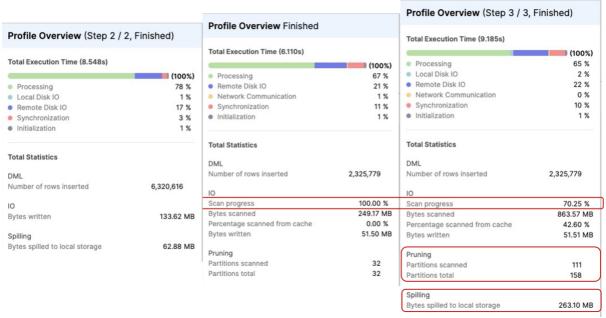


Figure 0-5 Now notice the amount of spillage under the XTS-influenced Satellite Load, including more micro-partitions scanned!

Yes, XTS involves more table scans, an option in applying this pattern would be to include a switch in your loading-pipeline, that is, if the staged applied date is older than the max applied date in the target satellite table then execute the XTS-influenced satellite loading pipeline instead. This switch-query will be free on Snowflake because the max value of every column is cached in Snowflake, essentially comparing the max applied date in staging versus the max applied date in the target satellite table.

So, you see, the Data Vault XTS pattern is a model and architecture pattern without the need to <u>refactor</u> anything! The optional loading pattern described above must be vetted by the data management team and above all the business. Every timeline correction event could affect reports, extracts, dashboards, and intelligence already acted upon and thus not only an alert should be created when such an event occurs, but the cost of the event calculated before publishing the correction. Yes, although this correction can happen in real time (and it is data driven), assess whether it is a pattern that suites your business & budget!

$$RV + BV + XTS = (self-healing) DV$$

Also consider, you're also loading content to an additional satellite table...

Day 1 Day 2

Profile Overview (Step 2 / 2	2, Finished)	Profile Overview (Step 2 / 2, Finished)		
Total Execution Time (7.860s)		Total Execution Time (8.548s)		
	(100%)		(100%)	
Processing	79 %	Processing	78 %	
Local Disk IO	1 %	 Local Disk IO 	1 %	
Remote Disk IO	16 %	Remote Disk IO	17 %	
Synchronization	2 %	Synchronization	3 %	
Initialization	2 %	Initialization	1 %	
Total Statistics		Total Statistics		
DML		DML		
Number of rows inserted	6,319,299	Number of rows inserted	6,320,616	
10		10		
Bytes written	132.64 MB	Bytes written	133.62 ME	
Spilling		Spilling		
Bytes spilled to local storage	62.90 MB	Bytes spilled to local storage	62.88 ME	

The views expressed in this article are that of my own, you should test implementation performance before committing to this implementation. The author provides no guarantees in this regard.