

The Power of Reformulation: From Validation to Planning in PDDL+

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PDDL+ planning problem

- A PDDL+ planning problem is a tuple $\Pi = \langle F, X, I, G, A, E, P \rangle$ where: F and X are the sets of **propositional** and **numeric variables**; I and G are the **initial** and **goal state**; A is the set of **actions** and E and P are the sets of **events** and **processes**.
- A PDDL+ plan is a tuple $\langle \pi, \langle t_s, t_e \rangle \rangle$ where: $\pi = \langle \langle a_1, t_1 \rangle, \dots, \langle a_n, t_n \rangle \rangle$, with $t_i \in \mathbb{Q}$, is a sequence of time-stamped actions and $\langle t_s, t_e \rangle$, with $t_s, t_e \in \mathbb{Q}$, is the **temporal envelope** within π is executed.

Processes and events

Processes and **Events** are used to model the **exogenous** and **environmental changes**.

- A **process** ρ is a tuple $\langle pre(\rho), eff(\rho) \rangle$ where $pre(\rho)$ is a logical formula involving conditions over F and X and $eff(\rho)$ is a set of **continuous numeric effects** having the form $\langle x, \xi \rangle$, where $x \in X$ and ξ is a mathematical expression defined over X and \mathbb{Q} ; if $pre(\rho)$ holds and time passes continuously then ρ contributes additively to the **first derivative** of x with ξ for each $\langle x, \xi \rangle \in eff(\rho)$.
- An **event** ε is a tuple $\langle pre(\varepsilon), eff(\varepsilon) \rangle$ where $pre(\varepsilon)$ is a logical formula and $eff(\varepsilon)$ is a set of conditional effects $c \triangleright e$ where c is a logical formula and e is a set of **numeric** and **Boolean assignments** having the form $\langle \{inc, dec, asgn\}, x, \xi \rangle$ and $\langle f := \{\perp, \top\} \rangle$, respectively; if $pre(\varepsilon)$ is triggered then the state changes instantaneously according to $eff(\varepsilon)$.

PDDL+ semantics: Intuitively, a PDDL+ problem consists of finding a number of time-stamped actions along with a potentially infinite timeline, whilst conforming to a number of processes and events that may change the state of the world in a **continuous** or an **instantaneous** manner as time goes by.

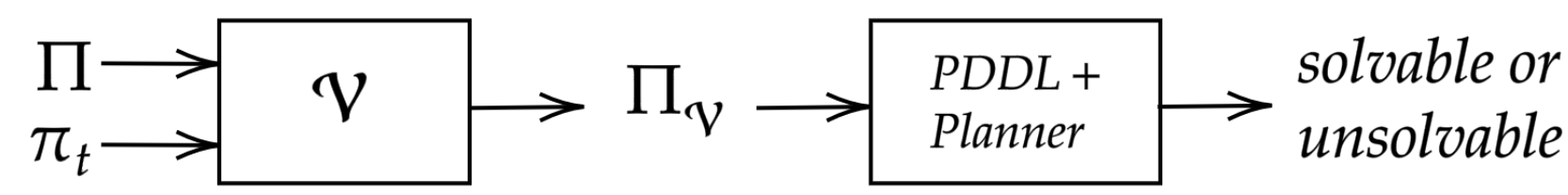
Research Question

Definition 1 ((Discrete) Validation Problem). Let Π be a PDDL+ problem, π_t be a plan for Π . The validation problem aims at establishing whether π_t is valid for Π . The discrete validation problem aims at establishing whether π_t is valid for Π under $\delta \in \mathbb{Q}$ discretisation.

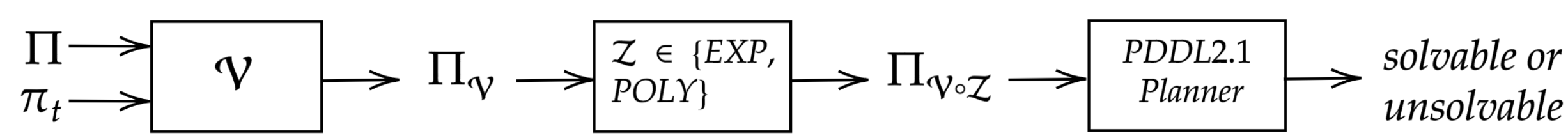
- Can reformulation be used to validate PDDL+ plans using planning techniques?

Methodology

Given Π and π_t , we **reformulate** the validation problem into a PDDL+ planning problem $\Pi_{\mathcal{V}}$ which is solvable (unsolvable) iff π_t is valid (invalid) w.r.t. Π .



We also reformulate $\Pi_{\mathcal{V}}$ through $\mathcal{Z} \in \{\text{POLY}, \text{EXP}\}$, which are translations for **discretising** a PDDL+ problem into a numeric one (Percassi et al., 2021), to produce a PDDL2.1 planning problem $\Pi_{\mathcal{V} \circ \mathcal{Z}}$ used for doing discrete PDDL+ validation.



Depending on the planning engine used to reason over $\Pi_{\mathcal{V}}$, it is possible to validate plans w.r.t. **discrete** or **continuous semantics**.

| | Continuous Semantics | Discrete Semantics | Polynomial Dynamic | Non-Polynomial Dynamic |
|-----|----------------------|--------------------|--------------------|------------------------|
| VAL | ✓ | ✗ | ✓ | ✗ |
| OUR | ✓ | ✓ | ✓ | ✓ |

Baseline translation \mathcal{V}_0

Given $\Pi = \langle F, X, I, G, A, E, P \rangle$ and $\pi_t = \langle \pi, \langle t_s, t_e \rangle \rangle$ plan for Π , \mathcal{V}_0 produces a new validating PDDL+ problem $\Pi_{\mathcal{V}_0}^{\pi_t} = \langle F \cup F_A \cup \{\top\}, X \cup \{\text{time}\}, I \cup \{a-d_0, \langle \text{time} := t_s \rangle, \top\}, G \wedge a-d_n \wedge \langle \text{time} = t_e \rangle, A_{\pi}, E, P \cup \{\rho_{\text{time}}\} \rangle$ such that $\rho_{\text{time}} = \langle \top, \{\langle \text{time}, 1 \rangle\} \rangle$ and:

$$F_A = \bigcup_{i=0}^{n=|\pi|} \{a-d_i\}$$

$$A_{\pi} = \bigcup_{\langle a_i, t_i \rangle \in \pi} \{ \langle pre(a_i) \wedge a-d_{i-1} \wedge \neg a-d_i \wedge \langle \text{time} = t_i \rangle, eff(a_i) \cup \{a-d_i\} \rangle \}$$

Key aspects: the only actions executable in $\Pi_{\mathcal{V}_0}^{\pi_t}$ are those from π_t . These actions can only be executed in the same ordering and at the same time-stamps prescribed by π_t .

Constrained translations \mathcal{V}_U and \mathcal{V}_D

\mathcal{V}_U : Given Π and π_t plan for Π , \mathcal{V}_U produces $\Pi_{\mathcal{V}_U}^{\pi_t} = \langle F \cup F_A \cup \{\top\}, X \cup \{\text{time}\}, I \cup \{a-d_0, \langle \text{time} := t_s \rangle, \top\}, G \wedge a-d_n \wedge \langle \text{time} = t_e \rangle, A_{\pi}, E, P_U \rangle$, where:

$$P_U = \bigcup_{\rho \in P \cup \{\rho_{\text{time}}\}} \{ \langle \langle \text{time} < t_e \rangle \wedge pre(\rho), eff(\rho) \rangle \}$$

The main idea of \mathcal{V}_U is to avoid some dead-ends by disallowing the occurrence of any process whenever time gets beyond what is prescribed by the plan. So, when time flows beyond t_e , then the processes P_U can not cause any change.

\mathcal{V}_D : Given Π and π_t plan for Π , \mathcal{V}_D produces $\Pi_{\mathcal{V}_D}^{\pi_t} = \langle F \cup F_A \cup \{\top\}, X \cup \{\text{time}\}, I \cup \{a-d_0, \langle \text{time} := t_s \rangle, \top\}, G \wedge a-d_n \wedge \langle \text{time} = t_e \rangle \wedge \top, A_{\pi}, E \cup E_D, P_D \cup \{\rho_{\text{time}}\} \rangle$, where:

$$E_D = \bigcup_{\substack{\pi(t') \text{ in } \pi \\ \pi(t') \neq \langle \rangle \wedge t' \neq t_e}} \{ \langle \langle \text{time} > t' \rangle \wedge \neg a-d_{\text{last}(\pi(t'))} \wedge \top, \{\neg \top\} \rangle \}$$

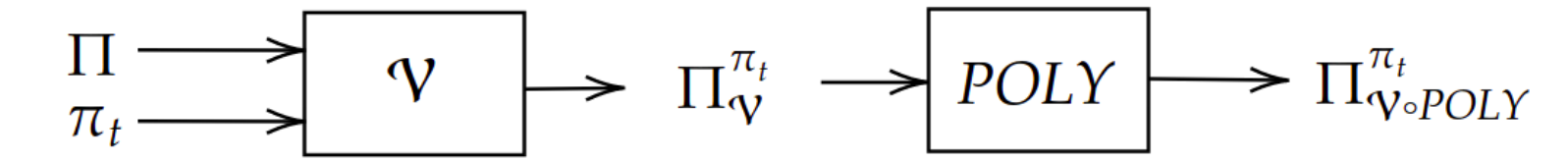
$$P_D = \bigcup_{\rho \in P} \{ \langle pre(\rho) \wedge \top, eff(\rho) \rangle \}$$

The main idea of \mathcal{V}_D is to prevent time from flowing unless all the actions prescribed by π_t up to the current instant have been executed.

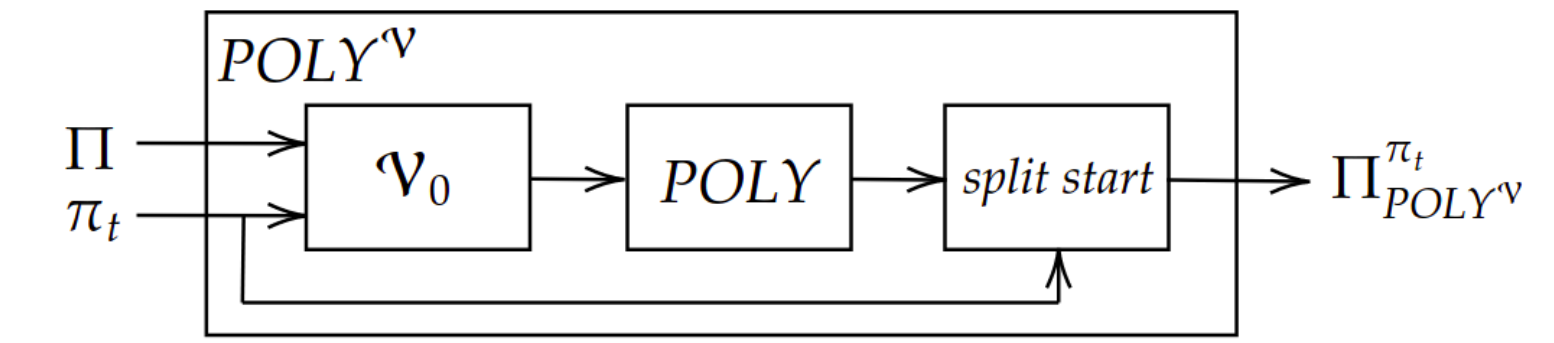
\mathcal{V}_{UD} : Jointly usage of \mathcal{V}_U and \mathcal{V}_D , i.e., $\mathcal{V}_{UD} = \mathcal{V}_U \circ \mathcal{V}_D$.

Can we use PDDL2.1 engines to validate PDDL+ plans?

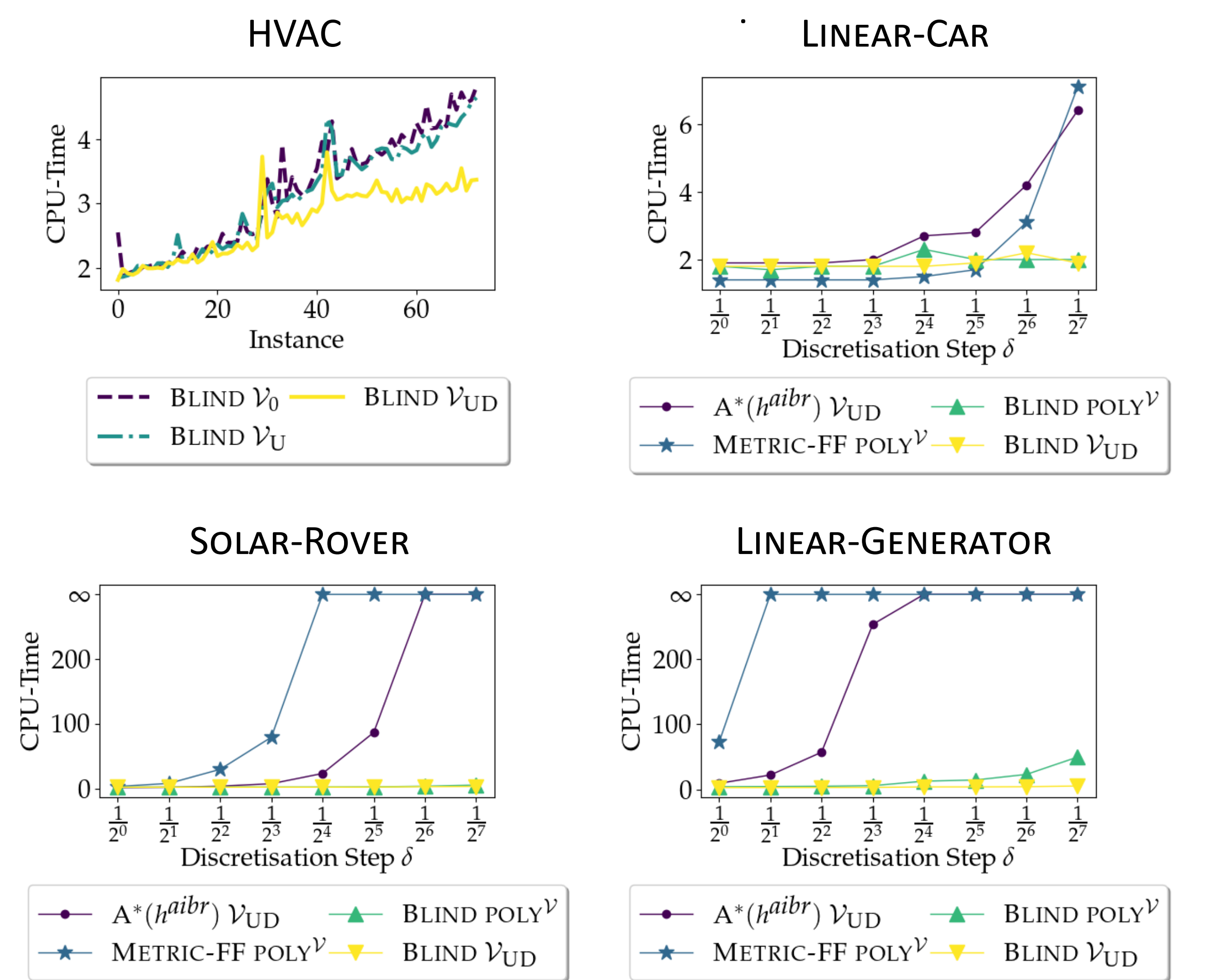
To **enable** the usage of numeric planners for doing PDDL+ validation we concatenate $\mathcal{V} \in \{\mathcal{V}_0, \mathcal{V}_U, \mathcal{V}_D, \mathcal{V}_{UD}\}$ with POLY. The resulting translation is denoted as $\mathcal{V} \circ \text{POLY}$ and the resulting PDDL2.1 validating task is denoted as $\Pi_{\mathcal{V} \circ \text{POLY}}^{\pi_t}$.



This chaining **obscures** the search properties provided by the constrained translations when $\mathcal{V} \in \{\mathcal{V}_U, \mathcal{V}_D\}$. We designed a variant of POLY, namely $\text{POLY}^{\mathcal{V}}$, which is **aware of the validation problem**. $\text{POLY}^{\mathcal{V}}$ is **equivalent** in terms of search properties to \mathcal{V}_{UD} . The resulting PDDL2.1 validating task is denoted as $\Pi_{\text{POLY}^{\mathcal{V}}}^{\pi_t}$.



Experimental Analysis



We can efficiently and effectively validate plans, even complex ones.

Selected Bibliography

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