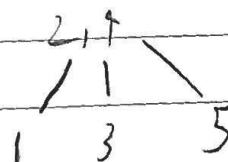
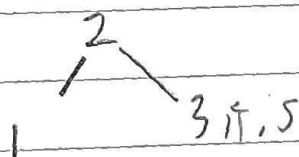
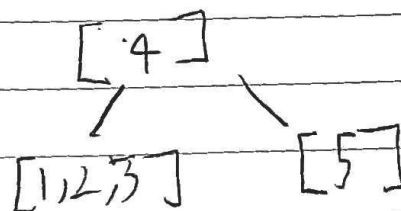
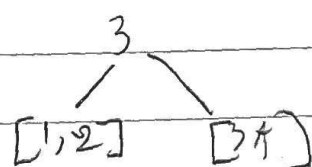


313 Assignment

1) All legal trees of degree 2 [1, 2, 3, 4, 5]

$$M-1=1, \quad M+1, 2M-1=3$$

4 possible trees



2) a) To find Minimum Node start from root then traverse to leftmost child/leaf, take leftmost element of the Node

Take a Node of the tree as a parameter, if value is NULL then return nil

If left and right child is NULL (so it's a leaf node) then return the minimum tree

Else call B-Tree recursively

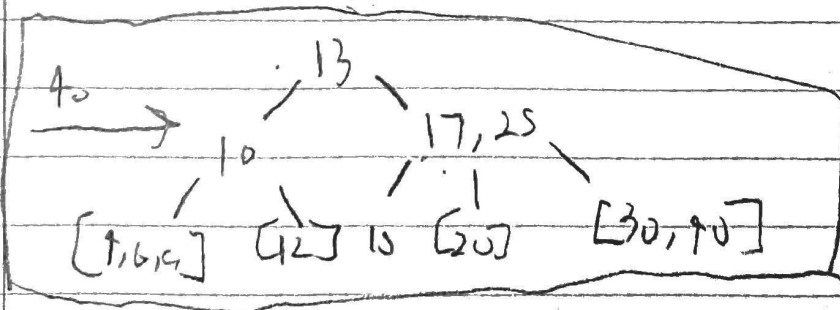
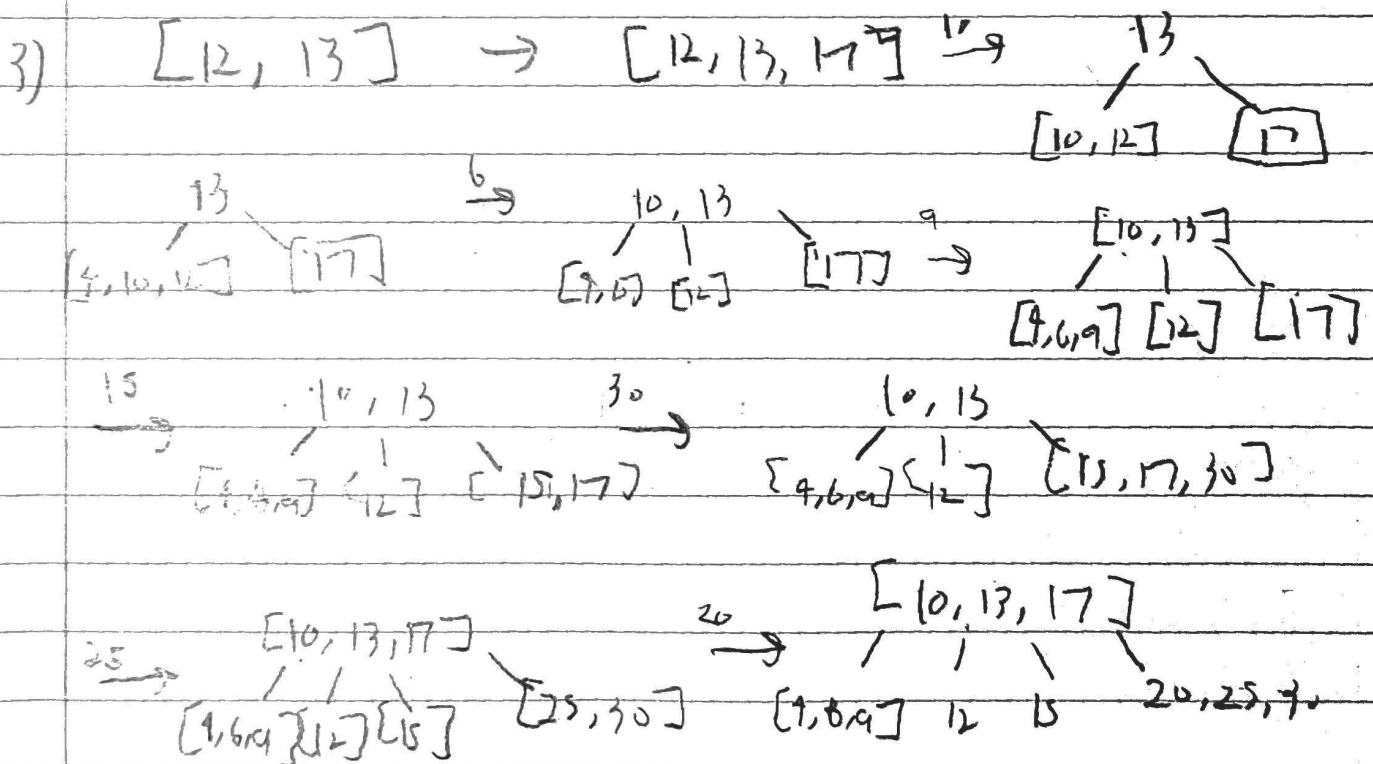
b) To find the predecessor find if the node is a leaf or not. If it's not then return largest element of child that precedes it, which has been found, if it's a leaf then return predecessor of given node

Case 1: If node is not leaf then return MAX of left subtree

Case 2: If the node is the leaf and $i > 1$ then return $i-1^{th}$ leaf of node

Case 3: If the node has no left child, then find predecessor of node by traversing bottom to top until $i > 0$

Use a Temp variable to store root. then find PREDECESSOR



4) We will want to use either heap-sort or merge-sort which is $O(\log n)$ then we can't take most appeared search which will be $O(n)$ time.

Function (S, n)

Sort (S, S+n)

Most = S[0]

Current = 1

Current max = 1

for i = 1 to n-1

if S[i] == S[i-1]

{

Counter += 1

}

else {

Counter = 1

if Counter > Counter_max

Counter_max = Counter

Most = S[i]

}

}

return Most

- 5) 1. Sort an array A with size K in sorted order, the $A[i]$ with most votes is winner.
2. for each vote do a binary-search on Array A to find candidate with the particular vote ID and increment br 1
3. Once vote record is over, declare candidate with highest vote count in array A .

$O(K \log K)$ to sort the list,
 Binary search will take $O(\log K)$ time
 $O(K \log K) \rightarrow O(n \log K) = O(n \log K)$

Election (S, K)

```

Array  $A$  of size  $K$  = Candidates
Array  $S$  = Votes
for  $i$  in  $S$ 
  BinarySearch into  $A$ 
  Count  $++$  1
for  $i$  in  $A$ 
  if Count > MaxCount
    Candidate =  $A[i]$ 
return Candidate
  
```

6) $O(n)$ for $\left(\left\lfloor \frac{n+1}{2} \right\rfloor\right)$

Majority Element

Keep count of every student in a map and see if any candidate exceeded it.

Majority (A) -

Maj_index = 0

count = 1

for i in len(A)

if A[Maj_index] == A[i]:

count += 1

else:

count -= 1

if (count == 0):

Maj_index = i

count = 1

return A[Maj_index]

for i in len(A)

if A[i] == A[Maj_index]:

count += 1

if count > len(A)/2:

return True

else:

return False