

Class 5: Review*

February 20

readability

Q: is writing code like this really more *readable*?

question from homework

```
absAll :: [Int] -> [Int]  
absAll xs = map abs xs
```

```
absAll :: [Int] -> [Int]  
absAll = map abs
```

a lateral move

```
toRevDigits :: Int -> [Int]
sum :: [Int] -> Int
```

```
sumDigits :: [Int] -> Int
sumDigits [] = 0
sumDigits (x : xs) = sum (toRevDigits x)
                    + sumDigits xs
```

remember this?

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toRevDigits :: Int -> [Int]
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toRevDigits :: Int -> [Int]  
sum :: [Int] -> Int
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```
sumDigits :: [Int] -> Int  
sumDigits = sum . map sum . map toRevDigits
```

now look at this!

```
toRevDigits :: Int -> [Int]
sum :: [Int] -> Int
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```
sumDigits :: [Int] -> Int
sumDigits = sum . map sum . map toRevDigits
```

`[Int] -> [[Int]]`

now look at this!

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[[Int]] -> [Int]

now look at this!

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toRevDigits :: Int -> [Int]
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sumDigits :: [Int] -> Int
sumDigits = sum . map sum . map toRevDigits
```

┐
[Int] -> Int

now look at this!

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sumDigits :: [Int] -> Int
sumDigits = sum . map sum . map toRevDigits
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now look at this!

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toRevDigits :: Int -> [Int]
sum :: [Int] -> Int
```

```
sumDigits :: [Int] -> Int
sumDigits = sum . map (sum . toRevDigits)
```

but what about this?

A: concision is not the goal, though it is a side effect.
the goal is to reason about problems at a higher level,

so that our code can focus on communicating
what it does instead of *how* it does it.

question from homework

more (and less) pattern matching

Q: why does pattern matching work?

question from homework

```
exec :: [Instr] -> Stack -> Stack
exec (i : is) st = exec is (exec1 i st)
exec is          st = st
```

stylistic recommendation: avoid “catch all” cases that only match one case

```
exec :: [Instr] -> Stack -> Stack
exec (i : is) st = exec is (exec1 i st)
exec []          st = st
```

stylistic recommendation: avoid “catch all” cases that only match one case

```
eval :: Exp -> Nat
```

```
...
```

```
eval (Branch e1 e2 e3) = case eval e1 of  
    Z -> eval e2  
    S _ -> eval e3
```

```
...
```

case expressions

sum :: [Int] -> Int

sum [] = 0

sum (x : xs) = x + sum xs

case expressions

```
sum :: [Int] -> Int
sum xs = case xs of
    [] -> 0
    (x : xs) -> x + sum xs
```

case expressions

```
remove7 :: [[Int]] -> [[Int]]  
remove7 [] = []  
remove7 (xs : xss) = map (...) (xs : xss)
```

stylistic recommendation: avoid unnecessary pattern matching

```
remove7 :: [[Int]] -> [[Int]]  
remove7 xss = map (...) xss
```

stylistic recommendation: avoid unnecessary pattern matching


```
same :: [Int] -> Bool  
same [] = True  
same (x : xs) = all (== x) xs
```

cf. useful pattern matching

follow the types!

```
data WeatherForecast  
  = Forecast Weather  
  | Unknown
```


```
data Maybe a  
  = Nothing  
  | Just a
```

refactoring example

(example: refactoring to use Maybe)

$f :: a \rightarrow a \rightarrow a$

revisiting polymorphism

$f :: a \rightarrow a \rightarrow a$
 $f\ a1\ a2 = a1 \ \&\&\ a2$


revisiting polymorphism

```
f :: a -> a -> a    ✗  
f a1 a2 = case (typeof a1) of  
  Int -> a1 + a2  
  Bool -> a1 && a2  
  _ -> a1
```

revisiting polymorphism

$f :: a \rightarrow a \rightarrow a$

parametric polymorphism

$f :: a \rightarrow a \rightarrow a$
 $f\ a1\ a2 = a1$

$f :: a \rightarrow a \rightarrow a$
 $f\ a1\ a2 = a2$

parametric polymorphism

$f :: a \rightarrow a$

can we define this function?

$f :: a \rightarrow b$

can we define this function?

$f :: (a \rightarrow a) \rightarrow a$

can we define this function?

$f :: (a \rightarrow a) \rightarrow a \rightarrow a$

can we define this function?

$f :: [a] \rightarrow a$

can we define this function?

$f :: (b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow (a \rightarrow c)$

can we define this function?

(exercise: const and flip)

your questions

Q: can {functional programming,
an expressive type system}
improve other languages?

question from homework

Q: what is Haskell's test framework like?

A: so far, we've been using HUnit for writing unit tests.
eventually, we'll see QuickCheck for writing *property-based tests*.

Q: what's this?

`concat [] ~?= ([] :: [()])`

A: `()` is the *unit type*.

usually, Haskell can infer the type of the elements in the list.
here, it can't because the list is empty. so we provide one!

Q: what's this?

```
sum :: (Foldable t, Num a) => t a -> a
```

A: type class constraints!

we'll learn about these right after spring break :)