

Lecture x

# More SQL: Complex Queries, Views, and Schema Modification

# AT71.01 Database Design

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## Outline (Chapter 5)

- More Complex SQL Retrieval Queries
- Views (Virtual Tables) in SQL
- Schema Change Statements in SQL

# MORE COMPLEX SQL RETRIEVAL QUERIES

## More Complex SQL Retrieval Queries

- Additional features allow users to specify more complex retrievals from database:
  - Nested queries, joined tables, outer joins, aggregate functions, and grouping

# Comparisons Involving NULL and Three-Valued Logic

- Meanings of NULL
  - Unknown value
  - Unavailable or withheld value
  - Not applicable attribute
- Each individual NULL value considered to be different from every other NULL value
- SQL uses a three-valued logic:
  - TRUE, FALSE, and UNKNOWN

# Comparisons Involving NULL and Three-Valued Logic (cont'd.)

**Table 5.1** Logical Connectives in Three-Valued Logic

(a)	<b>AND</b>	TRUE	FALSE	UNKNOWN
	TRUE	TRUE	FALSE	UNKNOWN
	FALSE	FALSE	FALSE	FALSE
	UNKNOWN	UNKNOWN	FALSE	UNKNOWN
(b)	<b>OR</b>	TRUE	FALSE	UNKNOWN
	TRUE	TRUE	TRUE	TRUE
	FALSE	TRUE	FALSE	UNKNOWN
	UNKNOWN	TRUE	UNKNOWN	UNKNOWN
(c)	<b>NOT</b>			
	TRUE	FALSE		
	FALSE	TRUE		
	UNKNOWN	UNKNOWN		

## Comparisons Involving NULL and Three-Valued Logic (cont'd.)

- SQL allows queries that check whether an attribute value is NULL
  - IS or IS NOT NULL

**Query 18.** Retrieve the names of all employees who do not have supervisors.

```
Q18:  SELECT  Fname, Lname
      FROM    EMPLOYEE
      WHERE   Super_ssn IS NULL;
```

## Nested Queries, Tuples, and Set/Multiset Comparisons

- **Nested queries**
  - Complete select-from-where blocks within WHERE clause of another query
  - **Outer query**
- Comparison operator IN
  - Compares value  $v$  with a set (or multiset) of values  $V$
  - Evaluates to TRUE if  $v$  is one of the elements in  $V$

## Nested Queries (cont'd.)

Q4A: **SELECT** **DISTINCT** Pnumber  
**FROM** PROJECT  
**WHERE** Pnumber **IN**  
( **SELECT** Pnumber  
**FROM** PROJECT, DEPARTMENT, EMPLOYEE  
**WHERE** Dnum=Dnumber **AND**  
Mgr\_ssn=Ssn **AND** Lname='Smith' )  
**OR**  
Pnumber **IN**  
( **SELECT** Pno  
**FROM** WORKS\_ON, EMPLOYEE  
**WHERE** Essn=Ssn **AND** Lname='Smith' );

## Nested Queries (cont'd.)

- Use tuples of values in comparisons
  - Place them within parentheses

**SELECT** **DISTINCT** Essn  
**FROM** WORKS\_ON  
**WHERE** (Pno, Hours) **IN** ( **SELECT** Pno, Hours  
**FROM** WORKS\_ON  
**WHERE** Essn='123456789' );

## Nested Queries (cont'd.)

- Use other comparison operators to compare a single value  $v$ 
  - $=$  ANY (or  $=$  SOME) operator
    - Returns TRUE if the value  $v$  is equal to some value in the set  $V$  and is hence equivalent to IN
  - Other operators that can be combined with ANY (or SOME):  $>$ ,  $>=$ ,  $<$ ,  $<=$ , and  $<>$

```
SELECT  Lname, Fname
FROM    EMPLOYEE
WHERE   Salary > ALL ( SELECT  Salary
                       FROM    EMPLOYEE
                       WHERE   Dno=5 );
```

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## Nested Queries (cont'd.)

- Avoid potential errors and ambiguities
  - Create tuple variables (aliases) for all tables referenced in SQL query

**Query 16.** Retrieve the name of each employee who has a dependent with the same first name and is the same sex as the employee.

```
Q16:  SELECT  E.Fname, E.Lname
FROM    EMPLOYEE AS E
WHERE   E.Ssn IN ( SELECT  Essn
                  FROM    DEPENDENT AS D
                  WHERE   E.Fname=D.Dependent_name
                  AND E.Sex=D.Sex );
```

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# Correlated Nested Queries

- **Correlated** nested query
  - Evaluated once for each tuple in the outer query

# The EXISTS and UNIQUE Functions in SQL

- EXISTS function
  - Check whether the result of a correlated nested query is empty or not
- EXISTS and NOT EXISTS
  - Typically used in conjunction with a correlated nested query
- SQL function UNIQUE ( Q )
  - Returns TRUE if there are no duplicate tuples in the result of query Q

## Explicit Sets and Renaming of Attributes in SQL

- Can use explicit set of values in WHERE clause
- Use qualifier AS followed by desired new name
  - Rename any attribute that appears in the result of a query

```
Q8A:  SELECT  E.Lname AS Employee_name, S.Lname AS Supervisor_name
      FROM    EMPLOYEE AS E, EMPLOYEE AS S
      WHERE   E.Super_ssn=S.Ssn;
```

## Joined Tables in SQL and Outer Joins

- **Joined table**
  - Permits users to specify a table resulting from a join operation in the FROM clause of a query
- The FROM clause in Q1A
  - Contains a single joined table

```
Q1A:  SELECT  Fname, Lname, Address
      FROM    (EMPLOYEE JOIN DEPARTMENT ON Dno=Dnumber)
      WHERE   Dname='Research';
```



## Joined Tables in SQL and Outer Joins (cont'd.)

- Specify different types of join
  - NATURAL JOIN
  - Various types of OUTER JOIN
- NATURAL JOIN on two relations R and S
  - No join condition specified
  - Implicit EQUIJOIN condition for each pair of attributes with same name from R and S

## Joined Tables in SQL and Outer Joins (cont'd.)

- **Inner join**
  - Default type of join in a joined table
  - Tuple is included in the result only if a matching tuple exists in the other relation
- **LEFT OUTER JOIN**
  - Every tuple in left table must appear in result
  - If no matching tuple
    - Padded with NULL values for attributes of right table

## Joined Tables in SQL and Outer Joins (cont'd.)

- RIGHT OUTER JOIN
  - Every tuple in right table must appear in result
  - If no matching tuple
    - Padded with NULL values for the attributes of left table
- FULL OUTER JOIN
- Can nest join specifications

## Aggregate Functions in SQL

- Used to summarize information from multiple tuples into a single-tuple summary
- **Grouping**
  - Create subgroups of tuples before summarizing
- Built-in aggregate functions
  - **COUNT**, **SUM**, **MAX**, **MIN**, and **AVG**
- Functions can be used in the **SELECT** clause or in a **HAVING** clause

## Aggregate Functions in SQL (cont'd.)

- NULL values discarded when aggregate functions are applied to a particular column

**Query 20.** Find the sum of the salaries of all employees of the 'Research' department, as well as the maximum salary, the minimum salary, and the average salary in this department.

```
Q20:  SELECT  SUM (Salary), MAX (Salary), MIN (Salary), AVG (Salary)
      FROM    (EMPLOYEE JOIN DEPARTMENT ON Dno=Dnumber)
      WHERE   Dname='Research';
```

**Queries 21 and 22.** Retrieve the total number of employees in the company (Q21) and the number of employees in the 'Research' department (Q22).

```
Q21:  SELECT  COUNT (*)
      FROM    EMPLOYEE;
```

```
Q22:  SELECT  COUNT (*)
      FROM    EMPLOYEE, DEPARTMENT
      WHERE   DNO=DNUMBER AND DNAME='Research';
```

## Grouping: The GROUP BY and HAVING Clauses

- **Partition** relation into subsets of tuples
  - Based on **grouping attribute(s)**
  - Apply function to each such group independently
- **GROUP BY** clause
  - Specifies grouping attributes
- If NULLs exist in grouping attribute
  - Separate group created for all tuples with a NULL value in grouping attribute

## Grouping: The GROUP BY and HAVING Clauses (cont'd.)

- **HAVING** clause
  - Provides a condition on the summary information

**Query 28.** For each department that has more than five employees, retrieve the department number and the number of its employees who are making more than \$40,000.

```
Q28:  SELECT  Dnumber, COUNT (*)
        FROM    DEPARTMENT, EMPLOYEE
        WHERE   Dnumber=Dno AND Salary>40000 AND
              ( SELECT  Dno
                FROM    EMPLOYEE
                GROUP BY Dno
                HAVING   COUNT (*) > 5)
```

## Discussion and Summary of SQL Queries

```
SELECT <attribute and function list>
FROM <table list>
[ WHERE <condition> ]
[ GROUP BY <grouping attribute(s)> ]
[ HAVING <group condition> ]
[ ORDER BY <attribute list> ];
```

# VIEWS (VIRTUAL TABLES) IN SQL

## Views (Virtual Tables) in SQL

- Concept of a view in SQL
  - Single table derived from other tables
  - Considered to be a virtual table

# Specification of Views in SQL

- **CREATE VIEW** command
  - Give table name, list of attribute names, and a query to specify the contents of the view

```
V1:  CREATE VIEW  WORKS_ON1
      AS SELECT   Fname, Lname, Pname, Hours
      FROM        EMPLOYEE, PROJECT, WORKS_ON
      WHERE       Ssn=Essn AND Pno=Pnumber;

V2:  CREATE VIEW  DEPT_INFO(Dept_name, No_of_emps, Total_sal)
      AS SELECT   Dname, COUNT (*), SUM (Salary)
      FROM        DEPARTMENT, EMPLOYEE
      WHERE       Dnumber=Dno
      GROUP BY    Dname;
```

## Specification of Views in SQL (cont'd.)

- Specify SQL queries on a view
- View always up-to-date
  - Responsibility of the DBMS and not the user
- **DROP VIEW** command
  - Dispose of a view

# View Implementation, View Update, and Inline Views

- Complex problem of efficiently implementing a view for querying
- **Query modification** approach
  - Modify view query into a query on underlying base tables
  - Disadvantage: inefficient for views defined via complex queries that are time-consuming to execute

## View Implementation

- **View materialization approach**
  - Physically create a temporary view table when the view is first queried
  - Keep that table on the assumption that other queries on the view will follow
  - Requires efficient strategy for automatically updating the view table when the base tables are updated

## View Implementation (cont'd.)

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- **Incremental update strategies**
  - DBMS determines what new tuples must be inserted, deleted, or modified in a materialized view table

## View Update and Inline Views

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- Update on a view defined on a single table without any aggregate functions
  - Can be mapped to an update on underlying base table
- View involving joins
  - Often not possible for DBMS to determine which of the updates is intended



# View Update and Inline Views (cont'd.)

- Clause **WITH CHECK OPTION**
  - Must be added at the end of the view definition if a view is to be updated
- **In-line view**
  - Defined in the FROM clause of an SQL query

## SCHEMA CHANGE STATEMENTS IN SQL

# Schema Change Statements in SQL

- **Schema evolution commands**
  - Can be done while the database is operational
  - Does not require recompilation of the database schema

## The DROP Command

- **DROP command**
  - Used to drop named schema elements, such as tables, domains, or constraint
- **Drop behavior options:**
  - `CASCADE` and `RESTRICT`
- **Example:**
  - `DROP SCHEMA COMPANY CASCADE ;`

# The ALTER Command

- **Alter table actions** include:
  - Adding or dropping a column (attribute)
  - Changing a column definition
  - Adding or dropping table constraints
- **Example:**
  - `ALTER TABLE COMPANY.EMPLOYEE ADD COLUMN Job VARCHAR(12);`
- **To drop a column**
  - Choose either `CASCADE` or `RESTRICT`

## The ALTER Command (cont'd.)

- **Change constraints specified on a table**
  - Add or drop a named constraint

```
ALTER TABLE COMPANY.EMPLOYEE  
DROP CONSTRAINT EMPSUPERFK CASCADE;
```

# SUMMARY

## Summary

- Complex SQL:
  - Nested queries, joined tables, outer joins, aggregate functions, grouping
- CREATE ASSERTION and CREATE TRIGGER
- Views
  - Virtual or derived tables