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# 1 Sem Outline

Week (dates)	Lecture
1	Computer Networks and the Internet
2	Principles of Nw Apps: HTTP, SMTP, DNS
3	Application Layer: P2P, CDN, Sockets
4	Networking at UQ
5	Transport Layer: UDP
6	Transport Layer: TCP
7	Network Layer: Data Plane
8	Network Layer: Control Place
9	Link Layer
11	Wireless and Mobile
12	Security
13	Multimedia

Table 1: Week Outline

### 2 Lecture 1

• billions of connected computing devices

• transmission rate: bandwidth

• Packet Switches: Forward packets

- routers and switches

Internet: "network of networks" (Interconnected ISPs)

 Protocols control sending, receiving (e.g. TCP, IP, HTTP, Skype, 802.11)

Internet standards

**RFC:** Request for comments

IETF: Internet Engineering Task Force

#### 2.1 Network Structure

### Network Edge

- hosts: clients and servers
- servers often in data centers
- Access networks, physical media: wired, wireless communication links
- network core:
  - interconnected routers
  - network of networks

#### 2.2 Access Network

#### 2.2.1 Digital Subscriber Line (DSL)

- use existing telephone line to central office DSLAM
  - data over DSL phone line goes to Internet
  - voice over DSL phone line goes to telephone net
- < 2.5 Mbps upstream transmission rate (typically < 1 Mbps)</li>
- < 24 Mbps downstream transmission rate (typically < 10 Mbps)</li>

### 2.2.2 Cable Network

frequency division multiplexing: different channels transmitted in different frequency bands

#### HFC: hybrid fiber coax

- asymmetric: up to 30Mbps downstream transmission rate, 2 Mbps upstream transmission rate
- network of cable, fiber attaches homes to ISP router
  - homes share access network to cable head-end

 unlike DSL, which has dedicated access to central office

#### wireless LANS:

- within building (30 meters)
- 802.11b/g/n (WiFi): 11,54,450 Mbps transmission rate

#### wide-area wireless access:

- provided by telco (cellular) operator, 10's km
- between 1 and 10 Mbps
- 3G, 4G, LTE

# 2.3 Sending

- takes application message
- breaks into smaller chunks, known as packets, of length L bits
- ullet transmites packet into access network at transmission rate R
  - link transmission rate, aka link capacity, aka link bandwidth

### **Note 1: Packet Transmission Delay**

$$\begin{array}{ccc} & & \text{time} \\ \text{packet} & & \text{needed} \\ \text{trans-} & & = \frac{\text{to} & \text{trans-}}{\text{mit} & L\text{-bit}} = \frac{L \text{ (bits)}}{R \text{ (bits/sec)}} \\ \text{delay} & & \text{packet into} \\ & & & \text{link} \end{array}$$

# 2.4 Physical Media

- bit: propagates between transmitter/receiver pairs
- physical link: what lies between transmitter and receiver
- guided media: signals propagate in solid media (copper, fiber, coax)
- unguided media: signals propagate freely, e.g. radio
- twisted pair (TP): two insulated copper wires
  - Category 5: 100 Mbps, 1 Gbps Ethernet
  - Category 6: 10 Gbps

#### 2.4.1 Coax

- two concentric copper conductors
- bidirectional
- broadband: multiple channels on cable, HFC

### 2.4.2 Fiber Optic Cable

- glass fiber carrying light pulses, each pulse a bit
- high-speed operation: high-speed point-topoint transmission (e.g. 10's - 100's Gbps transmission rate)
- low error rate
  - repeaters spaced far apart
  - immune to electromagnetic noise

#### 2.4.3 Radio

- signal carried in electromagnetic spectrum
- no physical "wire"
- bidirectional
- propagation environment effects:
  - reflection
  - obstruction by objects
  - interference

#### **Radio Link Types:**

- terrestrial microwave: up to 45 Mbps channels
- LAN (e.g. WiFi) 54 Mbps
- wide-area (e.g. cellular) 4G cellular: 10 Mbps
- satellite
  - Kbps to 45 Mbps channel (or multiple smaller channels)
  - 270 msec end-end delay
  - geosynchronous versus low altitude

# 2.5 Packet-switching

#### 2.5.1 Store-and-forward

L bits per packet

Source to destination: R bps

- takes  $\frac{L}{R}$  seconds to transmit (push out) L-bit packet into link at R bps
- store and forward: entire packet must arrive at router before it can be transmitted on next link

# Note 2: End-End delay

$$\mathrm{delay} = 2\frac{L}{R}$$

(assuming zero propagation delay)

# 2.5.2 Packet switching versus circuit switching

Is packet switching a "slam dunk winner?"

- great for bursty data (resource sharing, simpler, no call setup)
- excessive congestion possible: packet delay and loss (protocols needed for reliable data transfer, congestion control)

### 2.6 Packet Loss

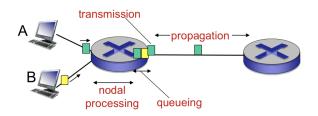


Figure 1: Packet Delay Algorithm Explanation

### Note 3: Packet Delay Algorithm

$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

#### 2.6.1 Nodal Processing

 $d_{\mathsf{proc}}$ 

- check bit errors
- · determine output link
- typically < msec</li>

#### 2.6.2 Queuing Delay

 $d_{\mathsf{queue}}$ 

- time waiting at output link for transmission
- depends on congestion level of router

#### 2.6.3 Transmission Delay

 $d_{\mathsf{trans}}$ 

- L: packet length (bits)
- R: link bandwidth(bps)
- $d_{\text{trans}} = \frac{L}{R}$

#### 2.6.4 Propagation Delay

 $d_{\mathsf{prop}}$ 

- d: length of physical link
- s: propagation speed ( $\approx 2 \times 10^8$  m/sec)
- $d_{prop} = \frac{d}{s}$

# 2.7 Throughput

Rate (bits/time unit) at which bits transferred between sender/receiver

**Instantaneous:** rate at given point in time **Average:** rate over longer period of time

#### Note 4: Bottleneck Link

Link on end-end path that constrains end-end throughput

# 2.8 Layering

#### 2.8.1 Why Layering?

Dealing with complex systems:

- Explicit structure allows identification, relationship of complex system's pieces (layered reference model for discussion)
- Modularization eases maintenance, updating system
  - change of implementation of layer's service transparent to rest of system
  - e.g. change in gate procedure doesn't affect rest of system
- layering considered harmful?

#### 2.8.2 Internet Protocol Stack

**Application:** supporting network applications (FTP, SMTP, HTTP)

**Transport:** process-process data transfer (TCP, UDP)

**Network:** routing of datagrams from source to destination (IP, routing protocols)

**Link:** data transfer between neighboring network elements (Ethernet, 802.111 (WiFi), PPP)

Physical: bits "on the wire"

#### 2.8.3 ISO/OSI Reference Model

Internet stack "missing" these layers. These services, if needed, must be implemented in application.

#### Application:

**Presentation:** allow applications to interpret meaning of data, e.g. encryption, compression, machine-specific conventions

**Session:** synchronization, check-pointing, recovery of data exchange

**Transport:** 

**Network:** 

Link:

Physical:

# 2.9 Security

• Malware can get in host from:

Virus: self-replicating infection by receiving/executing object (e.g. e-mail attachment)

**Worm:** self-replicating infection by passively receiving object that gets itself executed

- Spyware malware can record keystrokes, web sites visited, upload info to collection site
- Infected host can be enrolled in **botnet**, used for spam. DDoS attacks

#### 2.9.1 DoS: Denial of Service

**Denial of Service (DoS):** attackers make resources (server, bandwidth) unavailable to legitimate traffic by overwhelming resource with bogus traffic

- 1. select target
- 2. break into hosts around the network (botnet)
- send packets to target from compromised hosts

### 2.9.2 Sniffing

- broadcast media (shared Ethernet, wireless)
- promiscuous network interface reads/records all packets (e.g. including passwords) passing by

#### 2.9.3 IP Spoofing

Send packet with false source address

# 3 Lecture 2

# 3.1 Application Architectures

#### 3.1.1 Client-Server

**Server:** Always-on host, Permanent IP address **Clients:** Do not communicate directly with each other, May have dynamic IP addresses

#### 3.1.2 Peer-to-Peer (P2P)

- No always-on server
- Peers request service from other peers, provide service in return to other peers
- Self Scalability new peers bring new service capacity, as well as new service demands

 Pers are intermittently connected and change IP addresses

### Note 5: App-layer protocol defines

- type of messages exchanged e.g. request, response
- message syntax what fields in messages and how fields are delineated
- message semantics meaning of information in fields
- rules for when and how processes send and respond to messages
- open protocols defined in RFCs, allows for interoperability (e.g. HTTP, SMTP)
- proprietary protocols e.g. Skype

# 3.2 Transport Service is needed

**Data Integrity:** Some programs need 100% reliable data transfer (e.g. file transfer, web transactions), others can tolerate loss (e.g. audio)

**Timing:** Some programs require low delay to be "effective" (e.g. online games)

**Throughput:** Some programs require minimum amount of throughput to be "effective" (e.g. multimedia), some use whatever they have available ("elastic apps")

Security: Encryption, Data Integrity

# 3.3 Transport Protocol Services

#### 3.3.1 TCP

**Reliable Transport** between sending and receiving process

Flow Control: sender won't overwhelm receiver Congestion Control: throttle sender when network overloaded

**Connection-Oriented:** setup required between client and server processes

**Does Not Provide:** timing, minimum throughput guarantee, security

#### 3.3.2 UDP

**Unreliable Data Transfer** between sending and receiving process

**Does Not Provide:** reliability, flow control, congestion control, timing, throughput guarantee, security, or connection setup

### 3.3.3 Securing TCP

#### **TCP and UDP**

- no encryption
- cleartext passwords sent into socket traverse Internet in cleartext

#### SSL

- provides encrypted TCP connection
- data integrity
- end-point authentication

#### SSL is at app layer

app use SSL libraries, that "talk" to TCP

#### **SSL** socket API

 cleartext passwords sent into socket traverse Internet encrypted

# 3.4 HTTP: Hypertext Transfer Protocol

- Web's application layer protocol
- client/server model. Client request website and server serves HTTP object in response
- Uses TCP
- HTTP is stateless. Server maintains no information about past client requests
- non-persistent HTTP: one object sent over one TCP connection, downloading multiple object required multiple connections
- persistent HTTP: multiple object sent over single TCP connection

Non-persistent HTTP issues:

- requires 2 RTTs per object
- OS overhead for each TCP connection
- browsers often open parallel TCP connections to fetch referenced objects

#### Persistent HTTP:

- server leaves connection open after sending response
- subsequent HTTP messages between same client/server sent over open connection
- client sends requests as soon as it encounters a referenced object
- as little as one RTT for all the referenced objects

#### 3.4.1 Method Types

**HTTP/1.0:** GET, POST, HEAD (asks server to leave requested object out of response)

**HTTP/1.1:** GET, POST, HEAD, PUT (uploads file in entity body to path specified in URL field), DELETE (deletes file specified in the URL field)

#### 3.4.2 Response Codes

**200 OK:** request succeeded, requested object later in this msg

301 Moved Permanently: requested object moved, new location specified later in this msg

**400 Bad Request:** request msg not understood by

**404 Not Found:** requested document not found on this server

**505 HTTP Version Not Supported** 

### 3.5 Cookies

Uses: authorization, shopping carts, recommendations, user session state (Web, email)

# 3.6 Web Caches (proxy server)

**Goal:** satify client request without involving origin server

- Browsers requests object from cache, if in cache the object is sent back otherwise cache requests object from origin
- Cache acts as both client and server
- Reduce response time for client request
- Reduce traffic

#### 3.6.1 Conditional GET

**Goal:** don't send object if cache has up-to-date cached version (lower link usage)

- Cache: specify date of cached copy in HTTP request If-modified-since: <date>
- Server: response contains no object if cached copy is up-to-date: HTTP/1.0 Not Modified

#### 3.7 Electronic Mail: SMTP

RFC 2821

- uses TCP to reliably transfer email message from client to server, port 25
- direct transfer: sending server to receiving 3.8.1 server
- three phases of transfer: handshaking, transfer of messages, closure
- command/response interaction
- messages must be in 7-bit ASCII
- uses persistent connections

- requires message to be in 7-bit ASCII
- uses CRLF.CRLF to determine end of message

Difference to HTTP being, HTTP is server sending data, SMTP is client connection sending data

**SMTP:** protocol for exchanging email messages **RFC 822:** standard for text message format (To, From, Subject, Body)

#### 3.7.1 Mail Access Protocols

**SMTP:** delivery/storage to receiver's server

**POP:** Post Office Protocol (*RFC 1939*): authorization, download

- POP3 is stateless across sessions
- Two main modes; download and delete, download and keep (allows multiple clients to read the same email)

**IMAP:** Internet Mail Access Protocol (*RFC 1730*): more features, including manipulation of stored message on server

- · All messages stored on server
- Supports folders
- Keeps user state across sessions: names of folders and mappings between message IDs and folder name

HTTP: gmail, Hotmail, Yahoo, etc

# 3.8 DNS: Domain Name System

- Lookup between names (e.g. google.com) and IP addresses
- Distributed Database implemented in hierarchy of many name servers
- Application-layer protocol: hosts, name servers communicate to resolve names (address/name translation)

Why not centralize DNS? Single point of failure, traffic volume, doesn't scale

#### 3.8.1 DNS Services

- hostname to IP address translation
- host aliasing (canonical, alias names)
- mail server aliasing
- load distribution (many IP addresses correspond to one name)

#### 3.8.2 TLD, authoritative servers

#### top-level domain (TLD) servers:

- responsible for com, org, net, edu, aero, jos, io
- and top-level country domains au, uk, ca
- Network Solutions maintains servers for .com TLD
- Educause for .edu TLD

#### **Authoritative DNS servers:**

- organization's own DNS server(s), providing authorative hostname to IP mappings for organization's named hosts
- can be maintained by organization or service provider

#### 3.8.3 Local DNS name server

- does not strictly belong to hierarchy
- each ISP (residential ISP, company, university) has one (also called "default name server")
- when host makes DNS query, query is sent to its local DNS server
  - has local cache of recent name-toaddress translation pairs (but may be out of date!)
  - acts as proxy, forwards query into hierarchy

#### 3.8.4 DNS Name Resolution

**Iterated query:** contacted server replies with name of server to contact. So root dns sends the ip of the next dns server to contact

**Recursive query:** puts burden of name resolution on contacted name server. So root dns server contacts the next levels down which contacts next level down.

#### 3.8.5 Caching

Once (any) name server learns mapping, it **caches** mapping. Cache entries timeout (disappear) after some time (TTL). If name host changes IP address, the name servers might not update until TTLs expire.

update/notify mechanisms proposed IETF standard RFC 2136

#### Note 6: RR Format

(name, value, type, ttl)

type=A name is hostname, value is IP address

**type=NS** name is domain (e.g. google.com), value is hostname of authoritative name server for this domain

type=CNAME name is alias name for some "canonical" (the real) name (www.ibm.com is really servereast.backup2.ibm.com), value
is canonical name

**type=MX** value is name of mailserver associated with name

#### 3.8.7 Protocol

Query and reply messages both follow same format

Table 2: Protocol Layout

2 bytes

identification flags
# questions # answer RRs
# authority RRs # additional RRs
questions (variable # of questions)
answers (variable # of RRs)
authority (variable # of RRs)
additional info (variable # of RRs)

#### 3.8.8 Attacking DNS

#### **DDoS attacks**

- bombard root servers with traffic. Not successful to date, traffic filtering, local DNS servers cache protecting root DNS
- bombard TLD server. Potentially more dangerous

#### **Redirect Attacks**

- man-in-middle (Intercept queries)
- DNS Poisoning (Send bogus relies to DNS server, which caches)

#### **Exploit DNS for DDoS**

- send queries with spoofed source address: target IP
- requires amplification

#### 3.8.6 DNS Records

# 4 Acronyms

IP: Internet Protocol

TCP: UDP:

**HTTP:** Hypertext Transfer Protocol **SMTP:** Simple Mail Transfer Protocol **RDP:** Remote Desktop Protocol

VOIP: Voice over IP

RTT:

POP: Post Office Protocol

IMAP: Internet Mail Access Procotol

**DNS:** Domain Name System

SSN:

TLD: Top-level Domain
TTL: Time To Live
RR: Resource Records

DDoS: