

Lambda Calculus

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Why is the λ -calculus important?

- Computer Science - Variable binding, scope
- Typed λ -calculi are the basis for modern type systems in programming
- Functional programming languages (Algol, Lisp, ML, Haskell)
- Recursion theory/computability
- Linguistics

How was the λ -calculus developed?

- Formal system of logic developed by Alonzo Church in 1932
- Used to address the *Entscheidungsproblem* – “Is every predicate first-order logic provable?”
- Similar developments - Godel numbers, Turing machines

Conversion Rules

- α -conversion: $\lambda x.[\dots x \dots] = \lambda y.[\dots y \dots]$.
“We can rename dummy variables.”
- Example: $\lambda a.a = \lambda b.b$
- β -conversion: $\lambda x.[\dots x \dots] T = [\dots T \dots]$.
“Evaluation by substitution.”
- Example: $(\lambda x.x)y = y$.
- η -conversion: $\lambda x.F(x) = F$.
“Extensionality - a function is defined by what it does.”
Controversial; often left out of compilers.
- Example: $\lambda y.\lambda x.yx = \lambda y.y$

Church numerals

- $0 := \lambda f. \lambda x. x$
- $1 := \lambda f. \lambda x. fx$
- $2 := \lambda f. \lambda x. f(fx)$
- $3 := \lambda f. \lambda x. f(f(fx))$
- ...