

UFL (2)

x_{ij} = demand of i sat. from j
 $y_j = 1$ if depot open in j

$$\Rightarrow \max \sum_i \sum_j x_{ij} c_{ij} - \sum_j y_j f_j$$

$$\text{s.t. } \sum_j x_{ij} = 1 \quad \forall i$$

$$1. \quad \boxed{\sum_i x_{ij} \leq m y_j + \nu_j}$$

$$x_{ij} \geq 0, \quad x_{ij} \leq 1, \quad y_j \in \{0, 1\}$$

Alternative for 1. $\boxed{x_{ij} \leq y_j}$ 2.

Lagrange:

$$\max \sum_i \sum_j x_{ij} c_{ij} - \sum_j y_j f_j + \sum_i u_i (1 - \sum_j x_{ij})$$

...

$$\Rightarrow \max \sum_j \left[\sum_i (x_{ij} c_{ij} - u_i x_{ij}) - y_j f_j \right] + \sum_i u_i$$

$$\sum_j \left[\sum_i (c_{ij} - u_i) x_{ij} - y_j f_j \right] + \sum_i u_i$$

$$w(\underline{u}) = \sum_j w_j(\underline{u}) + \sum_i u_i$$

$$w_j(\underline{u}) = \max \sum_i (c_{ij} - u_i) x_{ij} - y_j f_j$$

$$x_{ij} \leq y_j$$

$$x_{ij} \in [0, 1],$$

$$y_j \in \{0, 1\}$$

$$w_j(\underline{u}) = \max \{0, \sum_i \max \{0, c_{ij} - u_i\} - f_j\}$$

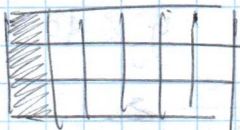


(c_{ij})



$$\boxed{c_{ij} - u_i}$$

(if > 0)

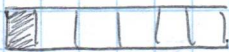


$$x_{ij} = 1$$

if $(c_{ij} - u_i) > 0$



Σ elements of column j



$$w_1(u_0) = \max \{0, \blacksquare_1 - f_1\}$$

:

$$w_j(u_0) = \max \{0, \blacksquare_j - f_j\}$$



$$w(u_0) = \Sigma_j w_j(u_0) + \Sigma_i (u_0)_i$$

$$(\underline{x}^0)_i = 1 - \Sigma_j x_{ij} \quad (\underline{u}_1 = \underline{u}_0 + \alpha_0 \underline{x}^0)$$

STSP (2)

1-tree: spanning tree on $V \setminus \{1\}$ plus two edges incident in 1

$$\min \sum_{e \in E} x_e c_e$$

$$\sum_{e \in \delta(i)} x_e = 2 \quad \forall i$$

CUT

$$\left[\begin{array}{l} \sum_{e \in \delta(S)} x_e \geq 2 \\ x_e \in \{0, 1\} \end{array} \quad S \subset V, 1 \leq |S| \leq n \right]$$

SEC

$$\left[\sum_{e \in E(S)} x_e \leq |S| - 1 \quad S \subset V, |S| \geq 2 \right]$$

starting from SEC:

$$\min \sum_{e \in E} x_e c_e$$

$$\sum_{e \in \delta(i)} x_e = 2 \quad \forall i$$

$$x_e \in \{0, 1\}$$

$$\sum_{e \in E(S)} x_e \leq |S| - 1 \quad S \subset V, |S| \geq 2, 1 \notin S$$

$$\sum_{e \in E} x_e = n$$

lagrangianize $\sum_{e \in \delta(i)} x_e = 2 \quad \forall i \in V \setminus \{1\}$

$$\min \left[\sum_{e \in E} x_e c_e + \sum_i u_i (2 - \sum_{e \in \delta(i)} x_e) \right]$$

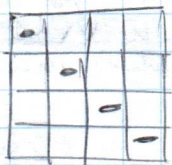
$$\sum_{e \in \delta(i)} x_e = 2$$

$$\sum_{e \in E(S)} x_e \leq |S| - 1 \quad S \subset V, |S| \geq 2, 1 \notin S$$

$$\sum_{e \in E} x_e = n$$

$$x_e \in \{0, 1\}$$

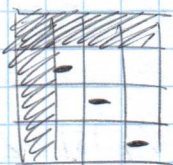
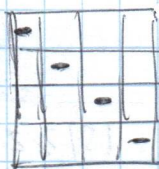
$$\sum_{e \in E} (c_e - u_i - u_j) x_e + 2 \sum_i u_i$$



(c_{ij})



$$c_{ij} - u_i - u_j$$



choose not a cycle
and add two
min cost from



$x_{ij} = 1$ if delectionated

$$W(\underline{u}^0) = \underset{\text{(new cost)}}{\text{(cost of the tree)}} + 2 \left(\sum_i u_i^0 \right)$$

$$(\underline{u}^0)_i = 2 - \left(\begin{array}{l} \# \text{ edges} \\ \text{going out} \\ \text{from } i \end{array} \right)$$