

1.

89/100

Packet Length = 1000 bytes (L)

Distance = 2500 km (D)

Speed = 2.5×10^8 m/s (S)

transmission rate = 2 Mbps (R)

$$\text{Time taken (t)} = \frac{L}{R} + \frac{D}{S}$$

(transmission delay + propagation delay)

$$\text{propagation delay} = \frac{D}{S} \Rightarrow \frac{2500 \text{ km}}{2.5 \times 10^8 \text{ m/s}}$$

$$\Rightarrow \frac{2500 \times 10^3}{2.5 \times 10^8} \text{ m/s} \quad (1 \text{ km} = 1000 \text{ m})$$

$$\Rightarrow \frac{2.5 \times 10^6}{2.5 \times 10^8} \Rightarrow \frac{1}{100} = 0.01 \text{ s} \Rightarrow \boxed{10 \text{ ms}}$$

$$\text{transmission delay} = \frac{1000 \text{ bytes}}{2 \text{ Mbps}} \Rightarrow \frac{8 \times 1000 \text{ bits}}{2 \times 1000 \times 1000 \text{ bits/sec}} \quad (1 \text{ Mb} = 1,000,000 \text{ bytes})$$

$$\Rightarrow 0.004 \Rightarrow 4 \text{ ms}$$

B) part just general form of p

$$\text{total delay} = 4 + 10 \text{ ms} \Rightarrow 14 \text{ ms}$$

$$9/10(2.5+1.5+2.5+)$$

this delay doesn't depend on both packet length and transmission rate as we don't consider there in propagation delay

2.

Data transfer rate,

$$R_1 \Rightarrow 500 \text{ Kbps}$$

$$R_2 \Rightarrow 2 \text{ Mbps}$$

$$R_3 \Rightarrow 1 \text{ Mbps}$$

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(a) if there is no traffic, we can say that minimum bottle neck value as an throughput for the file transfer, so, among all the data Rates (R_i) can be considered as minimum throughput.

(b) file size \Rightarrow 4 million bytes.

throughput \Rightarrow 500 Kbps

$$\text{time (t)} \Rightarrow \frac{\text{file size}}{\text{throughput}} \Rightarrow \frac{4,000,000 \times 8}{500 \text{ Kbps}}$$

(converting bytes to bits \Rightarrow 1 byte \Rightarrow 8 bits)

$$\Rightarrow \frac{4000 \text{ Kbp} \times 8}{500}$$

$$\Rightarrow 8 \times 8 \Rightarrow \boxed{64 \text{ sec}}$$

(c) If throughput \Rightarrow 100 Kbps

$$\Rightarrow \frac{4000 \text{ Kbps} \times 8}{100} \Rightarrow \boxed{320 \text{ sec}}$$

(3) Yes. It can be done through by putting reliable - data transfer. It can be Implemented by the Application programmer by keeping acknowledgement and re-transmission of data using RDUP, QUIC protocols with UDP reliable data transfer.

(4) - Circuit-switching (Ncs) \Rightarrow bandwidth - 20Mbps

share link cap - 100Mbps

- packet switching \Rightarrow share link = 100Mbps

min. transmission = 20Mbps

min. trans time = 20%

(a)

$$\text{maximum users} = \frac{\text{total cap}}{\text{min band for user}}$$

$$\Rightarrow \frac{100}{20} \Rightarrow \boxed{5 \text{ users}}$$

(b) probability for fraction of time (P) = 0.2

probability when specific user is not

$$\text{busy} = (1 - P)$$

therefore, prob of all users not transmitting

$$\Rightarrow (1 - P)^{N-1}$$

$$\therefore \text{probability} = N \cdot P + (1 - P)^{N-1}$$

N is not necessary

Consider for specific user

$$= \boxed{N \cdot (0.2) + (0.8)^{N-1}}$$

(c) If 9 users are in total:

$$\text{probability} \Rightarrow 9 \cdot 0.2 + (0.8)^{9-1}$$

$$\Rightarrow 9 \cdot 0.2 + 0.167 \Rightarrow \boxed{0.3006}$$

(d) In packet switching, probability of more than 5 users transmitting, and the rest users are idle then probability can be taken as $(p)^5 (1-p)^{9-5}$.

25/35(5+6+1)

lets assume six(6) members are using the channel at same time probability

$$\Rightarrow 6 \cdot (p)^6 (1-p)^{9-6}$$

$$\Rightarrow 6 \cdot p \cdot 6 \cdot (0.2)^6 (0.8)^3$$

$$\Rightarrow 0.0001$$

It is the summation from 6 to

$$P(7) \Rightarrow 7 \cdot p^7 (1-p)^{9-7}$$

$$\Rightarrow 7 \cdot (0.2)^7 (0.8)^2$$

$$= 0.00005$$

$$P(5) = 0.006$$

So, ~~It~~ varies based on the number of users using the packet switching at the same time.

(5)

$F = 4 \text{Gb}$, 7 peers, upload rate = 94 Mbps

Upload rates: $U_1 = 28$, $U_2 = 14$, $U_3 = 29$, $U_4 = 22$

$U_5 = 16$, $U_6 = 15$, $U_7 = 11$ Mbps.

Download rate: $d_1 = 15$ Mbps, $d_2 = 36$, $d_3 = 21$, $d_4 = 28$

$d_5 = 31$, $d_6 = 20$, $d_7 = 38$.

(a)

Minimum time Needed =

$$\text{Max} \left(\text{No. of peers} \times \frac{\text{file size}}{\text{upload rate}}, \frac{\text{file size}}{\text{Minimum download Speed}} \right)$$

$$\Rightarrow \text{Max} \left(7 \times \frac{4 \text{Gb}}{94 \text{Mbps}}, \frac{4 \text{Gb}}{15 \text{Mb}} \right)$$

$$\Rightarrow \text{Max} \left(\frac{7 \times 4 \times 1000}{94}, \frac{4 \times 1000}{15} \right)$$

$$\Rightarrow \text{Max} (297.87, 266.66)$$

$$\text{Minimum time needed} = \boxed{297.87 \text{ sec}}$$

(b) the root cause of this specific minimum time is due to server distributing file to every client.

(c)

minimum time needed to distribute \Rightarrow

$$\Rightarrow \text{Max} \left(\frac{\text{file size}}{\text{upload speed}}, \frac{\text{file size}}{\text{min download Speed}}, N * \frac{\text{file size}}{\sum_i^N \text{upload Speed}} \right)$$

$$\Rightarrow \text{Max} \left(\frac{4 * 1000}{94}, \frac{4 * 1000}{15}, 7 * \frac{4 * 1000}{28 + 14 + 29 + 22 + 16 + 18 + 11} \right)$$

$$\Rightarrow \text{Max} \left(\overset{42.55}{229}, 266.66, 207.40 \right)$$

$$\Rightarrow 266.66.$$

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minimum time needed to distribute in

$$\text{peer-to-peer model} \Rightarrow \boxed{266.66}$$

(d) the root cause of this specific minimum time is due slower download Speed of the client d1 (15 mbps).