

**CS570**  
**Analysis of Algorithms**  
**Spring 2015**  
**Exam I**

Name: \_\_\_\_\_

Student ID: \_\_\_\_\_

Email Address: \_\_\_\_\_

\_\_\_\_\_ **Check if DEN Student**

	Maximum	Received
Problem 1	20	
Problem 2	20	
Problem 3	14	
Problem 4	13	
Problem 5	20	
Problem 6	13	
Total	100	

**Instructions:**

1. This is a 2-hr exam. Closed book and notes
2. If a description to an algorithm or a proof is required please limit your description or proof to within 150 words, preferably not exceeding the space allotted for that question.
3. No space other than the pages in the exam booklet will be scanned for grading.
4. If you require an additional page for a question, you can use the extra page provided within this booklet. However please indicate clearly that you are continuing the solution on the additional page.

20 pts

Mark the following statements as **TRUE** or **FALSE**. No need to provide any justification.

[ **TRUE/FALSE** ]

For some graphs BFS and DFS trees can be the same.

[ **TRUE/FALSE** ]

The number of cycles in a bipartite graph may be odd.

[ **TRUE/FALSE** ]

Stable matching algorithm presented in class is based on the greedy technique.

[ **TRUE/FALSE** ]

To delete the  $i^{\text{th}}$  node in a binary min heap, you can exchange the last node with the  $i^{\text{th}}$  node, then check the nodes below the  $i^{\text{th}}$  node to see if the  $i^{\text{th}}$  node should move down the heap to “re-heapify” it.

[ **TRUE/FALSE** ]

Kruskal’s algorithm can fail in the presence of negative cost edges.

[ **TRUE/FALSE** ]

Consider an instance of the Stable Matching Problem in which there exists a man  $m$  and a woman  $w$  such that  $m$  is ranked first by  $w$ , and  $w$  is ranked first by  $m$ . Then  $(m, w)$  must appear in every stable matching.

[ **TRUE/FALSE** ]

If a connected undirected graph  $G$  has the same weights for every edge, then a minimum spanning tree can be found in linear time.

[ **TRUE/FALSE** ]

Given  $n$  numbers, one could construct a binary heap using the  $n$  numbers, then using the binary heap produce a sorted list of the numbers in  $O(n)$  time.

[ **TRUE/FALSE** ]

Binomial

In a **Fibonacci** heap, the insert operation has an amortized cost of  $O(1)$  time, but the worst case cost is higher.

[ **TRUE/FALSE** ]

Function  $10n^{10}2^n + 3^n \log(n)$  is  $O(n^{10}2^n)$ .

2) 20 pts

A pharmacist has  $W$  pills and  $n$  empty bottles. Let  $\{p_1, p_2, \dots, p_n\}$  denote the number of pills that each bottle can hold.

Describe a greedy algorithm, which, given  $W$  and  $\{p_1, p_2, \dots, p_n\}$ , determines the fewest number of bottles needed to store the pills. Prove that your algorithm is correct.

3) 14 pts

We are given a graph  $G = (V; E)$  with uniform cost edges and two vertices  $s$  and  $t$  within  $G$ . We are told that the length of the shortest path from  $s$  to  $t$  is more than  $n/2$  (where  $n$  is the number of vertices within  $G$ ). Give a linear time algorithm to find a vertex  $v$  (other than  $s$  and  $t$ ) such that every path from  $s$  to  $t$  contains  $v$ . Justify your solution.

4) 13 pts

Arrange the following functions in increasing order of asymptotic complexity. If  $f(n) = \Theta(g(n))$  then put  $f = g$ . Else, if  $f(n) = O(g(n))$ , put  $f < g$ .

$4n^2$ ,  $\log^2 n$ ,  $5n^2$ ,  $\log^3 n$ ,  $n^n$ ,  $3n$ ,  $n \log(n)$ ,  $2n$ ,  $2n+1$

5) 20 pts

Let  $G = (V, E)$  be an (undirected) graph with costs  $C_e \geq 0$  on the edges  $e \in E$ . Assume you are given a minimum-cost spanning tree  $T$  in  $G$ . Now assume that an edge connecting two nodes  $v, w \in V$  with cost  $c$  is deleted from graph  $G$  and let  $G'$  be the new graph. Assume that the graph  $G'$  is connected.

- If the removed edge doesn't appear in the MST  $T$ , will  $T$  still be the MST of  $G'$ ? Please justify your answer.
- If the removed edge appears in the MST  $T$ , give an  $O(V+E)$  algorithm to find a MST for the graph  $G'$ .
- Prove that the output produced by your algorithm in part b is an MST

6) 13 pts

Solve the following recurrences by giving tight  $\Theta$ -notation bounds. You do not need to justify your answers, but any justification that you provide will help when assigning partial credit.

i.  $T(n) = 5T(n/2) + n^2 \log n$

ii.  $T(n) = 4T(n/2) + n \log n$

iii.  $T(n) = (6006)^{1/2} * T(n/2) + n^{\sqrt{6006}}$

Additional Space