CS 6150: HW 5 - Randomized algorithms

Submission date: Wednesday, November 20, 2019, 11:59 PM

This assignment has 5 questions, for a total of 50 points. Unless otherwise specified, complete and reasoned arguments will be expected for all answers.

Question	Points	Score
Collecting coupons	14	
Brownian motion	10	
Trade-offs in sampling	6	
Satisfying ordering constraints	10	
Birthdays and applications	10	
Total:	50	

A cereal company has decided to give out superhero stickers with boxes of its cereal. There are n superheroes in total, and suppose that each cereal box you buy has a sticker of a uniformly random superhero. What is the expected number of boxes you need to buy so that you end up with at least one copy of all the n stickers?

There are many ways to do this analysis; let us see one of them. We would like to write down a recurrence for the expected value. Define f(n,k) to be the expected number of boxes you need to buy to end up with all the stickers, given that you have already seen k distinct stickers. Thus by definition, f(n,n) = 0, and the goal is to compute f(n,0).

(a) [6] Use the law of conditional expectations to prove that

$$f(n,k) = \frac{n-k}{n} (1 + f(n,k+1)) + \frac{k}{n} (1 + f(n,k)).$$

Simplify this to evaluate f(n,0). [Hint: you may use the identity $1 + \frac{1}{2} + \frac{1}{3} + \cdots + \frac{1}{n} = \log n + c$ for some $c \in (0,1)$.]

- (b) [2] Suppose n > 4. Prove that the probability that you need to buy $8n \log n$ boxes in order to see all the n stickers is $\leq 1/4$.
- (c) [6] Use a more direct computation to bound the probability above by $\frac{1}{n^4}$. [Hint: what is the probability that you buy $8n \log n 1$ boxes and you still have not seen a given sticker? Can you now use the union bound?]

[All logarithms above are natural logs. You might also find the inequality $1-x \le e^{-x}$ useful.]

- (a) [2] Let X_t be the random variable denoting the location of the particle at time t. For some $t \ge 1$, compute $\mathbf{E}[X_t]$.
- (b) [5] Compute $\mathbf{E}[X_t^2]$ for some integer $t \ge 1$, and use this to prove that with probability $\ge 3/4$, we have $|X_t| \le 2\sqrt{t}$.
- (c) [3] Part (b) shows that the magnitude of X_t after t steps of moving around is only $O(\sqrt{t})$. This raises the question: does it "move around" pretty uniformly in the interval say $(-\sqrt{t}, +\sqrt{t})$? Run experiments on the process with $t = 4 \cdot 10^4$. On average (over say 50 runs), how many times does the particle "cross the origin"? Repeat with $t = 9 \cdot 10^4$ and $t = 16 \cdot 10^4$ and report your answers.

Suppose we have a population of size 1 million, and suppose 52% of them vote +1 and 48% of them vote -1. Now, randomly pick samples of size (a) 20, (b) 100, (c) 400, and evaluate the probability that +1 is majority even in your sample (by running the experiment say 100 times and taking the average). Write down the values you observe for these probabilities in the cases (a-c).

Next, what is the size of the sample you need for this probability to become 0.9?

Given the constraints, the goal is to find an ordering that satisfies as many constraints as possible (for simplicity, assume in what follows that m is a multiple of 3). For large m, n, this problem becomes very difficult.

- (a) [6] As a baseline, let us consider a *uniformly random* ordering. What is the expected number of constraints that are satisfied by this ordering? [*Hint*: define appropriate random variables whose sum is the quantity of interest, and apply the linearity of expectation.]
- (b) [4] Let X be the random variable which is the number of constraints satisfied by a random ordering, and let E denote its expectation (which we computed in part (a)). Now, Markov's inequality tells us, for example, that $\Pr[X \geq 2E] \leq 1/2$. But it does not say anything that lets us argue that $\Pr[X \geq E]$ is "large enough" (which we need if we want to say that generating a few random orderings and picking the best one leads to many constraints being satisfied with high probability). Use the definition of X above to conclude that $\Pr[X \geq E] \geq 1/m$.
- - (a) [4] What is the expected number of pairs (i, j) with i < j such that person i and person j have the same birthday? For what value of n (as a function of m) does this number become 1?
 - (b) [6] This idea has some nice applications in CS, one of which is in estimating the "support" of a distribution. Suppose we have a radio station that claims to have a library of one million songs, and suppose that the radio station plays these songs by picking, at each step a uniformly random song from its library (with replacement), playing it, then picking the next song, and so on.
 Suppose we have a listener who started listening when the station began, and noticed that among

Suppose we have a listener who started listening when the station began, and noticed that among the first 200 songs, there was a repetition (i.e., a song played twice). Prove that with probability > 0.9, the station's claim of having a million song database is false.

[One recent application of this idea was in proving that "GANs", a recent ML technique to produce realistic data such as images, typically have a pretty small support size.]