Distributed Systems

ECE428

Lecture 9

Adopted from Spring 2021

Today’s agenda

- Wrap up Multicast
  - Chapter 15.4
  - Tree-based multicast and Gossip
- Mutual Exclusion
  - Chapter 15.2

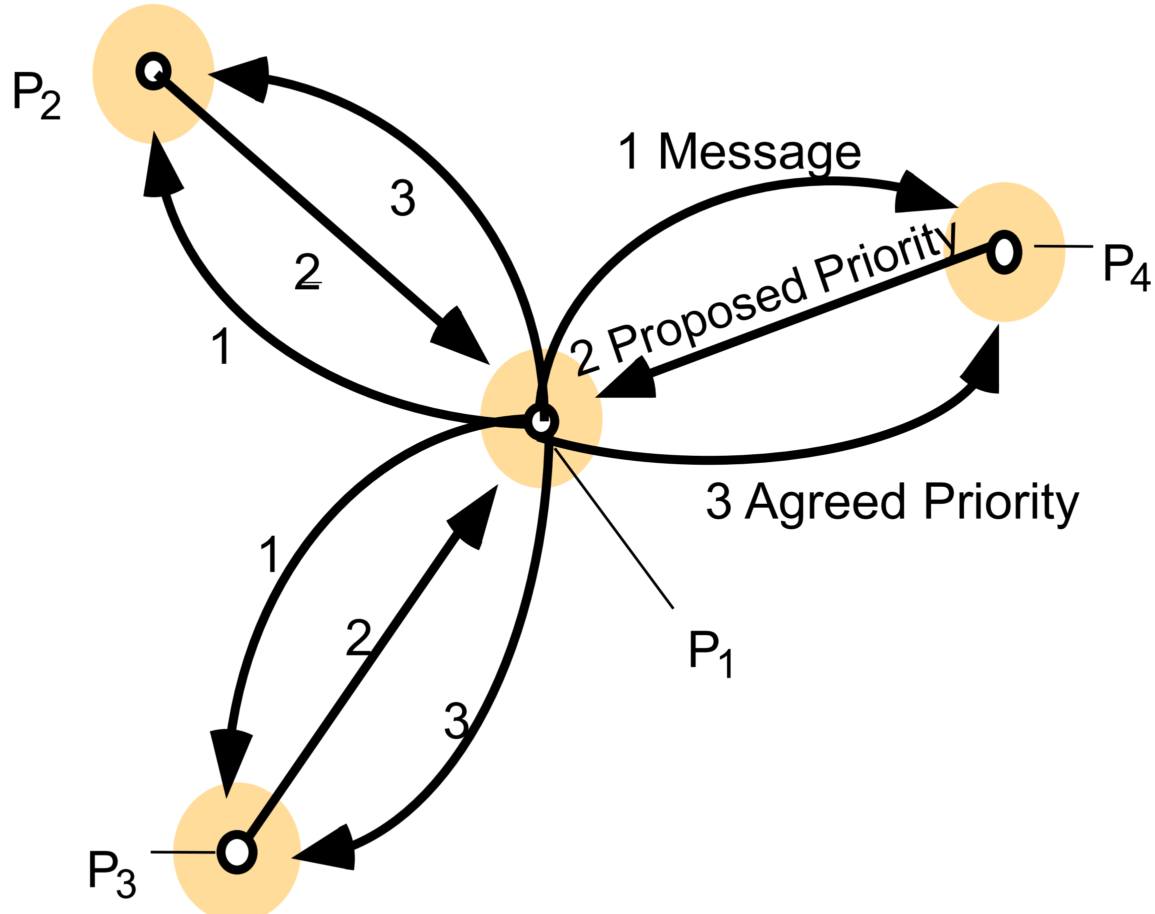
Recap: Ordered Multicast

- FIFO ordering: If correct process issues multicast(*g*,*m*) and then multicast(*g*,*m’*), then every correct process that delivers *m’* will have already delivered m.

- Causal ordering: If multicast(*g*,*m*) $\rightarrow$ multicast(*g*,*m’*) then any correct process that delivers *m’* will have already delivered *m*.
  - Note that $\rightarrow$ counts multicast messages delivered to the application, rather than all network messages.

- Total ordering: If a correct process delivers message *m* before *m’*, then any other correct process that delivers *m’* will have already delivered *m*.

ISIS algorithm for total ordering



Proposed Priority: *higher than all priorities proposed by the process and agreed priorities received by the process so far.*

Agreed Priority: *Maximum of all proposed priority for the message*

Proof of total order with ISIS

- Consider messages, m_1 and m_2, and two processes, p and p'.
- Suppose that p delivers m_1 before m_2.
- When *p* delivers m_1, it is at head of the queue. m_2 is either:
  - Already in *p*'s queue, and deliverable, so
    - Final_priority(m_1) < Final_priority(m_2)
  - Already in *p*'s queue, and not deliverable, so
    - Final_priority(m_1) < Proposed_priority(m_2) <= Final_priority(m_2)
  - Not yet in *p*'s queue:
    - same as above, since proposed priority > priority of any delivered message
- Suppose *p*' delivers m_2 before m_1, by the same argument:
  - Final_priority(m_2) < Final_priority(m_1)
  - Contradiction!

Ordered Multicast

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Implementing causal order multicast

- Similar to FIFO Multicast
  - What you send with a message differs.
  - Updating rules differ.

- Each receiver maintains a vector of per-sender sequence numbers (integers)
  - Processes P*1* through P*N.*
  - P*i* maintains a vector of sequence numbers P*i*[1…N] (initially all zeroes).
  - P*i*[*j*] is the latest sequence number P*i* has received from P*j.*

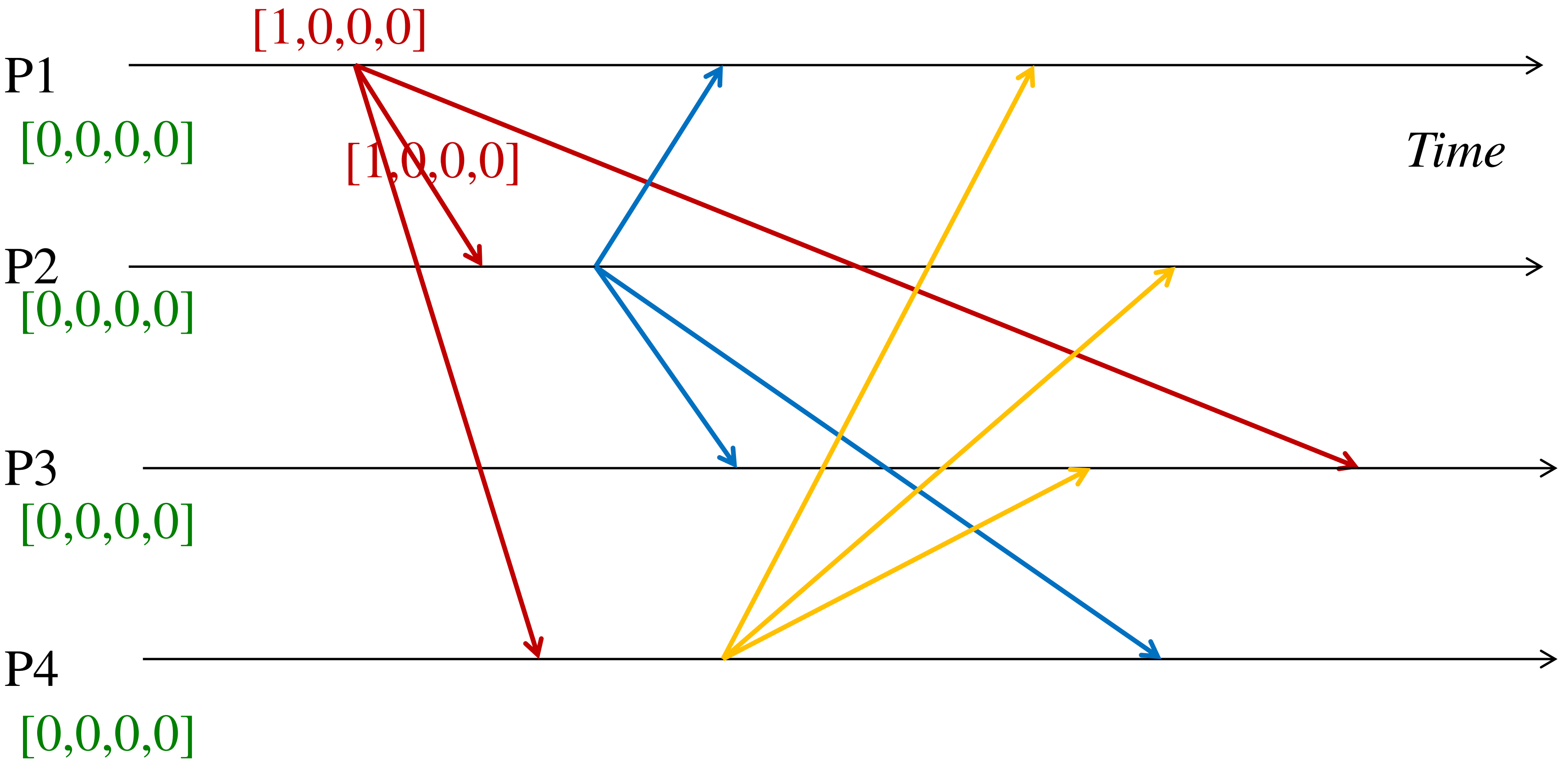
… these are NOT vector logical clocks!

Implementing causal order multicast

- *CO-multicast(g,m) at Pj:*
  set P*j*[*j*] = P*j*[*j*] + 1
  piggyback entire vector P*j*[1…N] with m.
  B-multicast(g,{m, P*j[1…N]*})
- *On B-deliver({m, V[1..N]}) at Pi from Pj:* If Pi receives a multicast from Pj with sequence vector V[1…N], buffer it until both conditions are true:
  1. This message is next one P*i* is expecting from P*j*, i.e.,
     V[*j*] = P*i*[*j*] + 1
  2. All multicasts, anywhere in the group, which happened-before m have been received at P*i*, i.e.,
     For all $k \neq j$: V[*k*] ≤ P*i*[*k*]

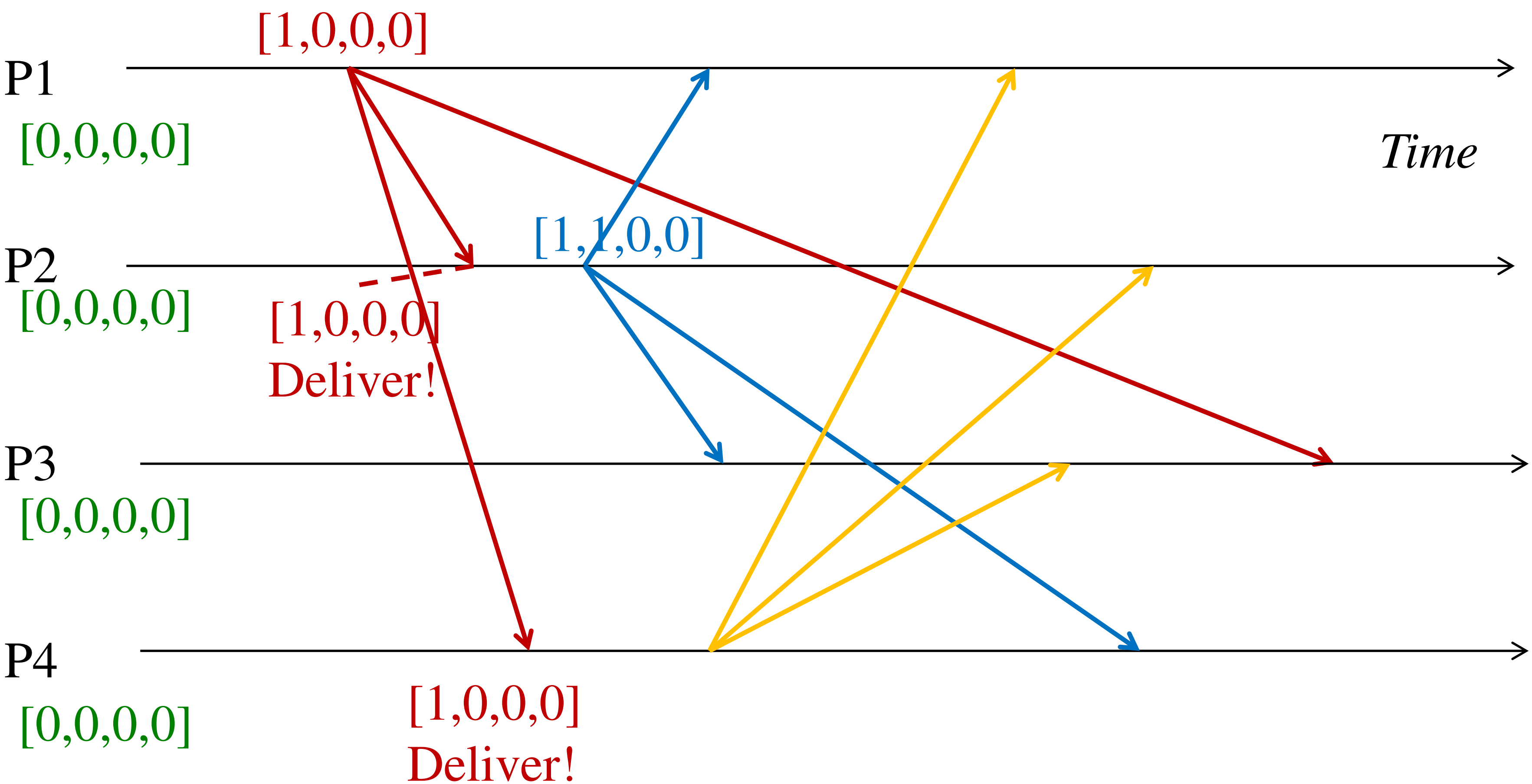
  When these conditions satisfied, CO-deliver(m), and set P*i*[*j*]=V[*j*]

Causal order multicast execution



Self-deliveries omitted for simplicity.

Causal order multicast execution



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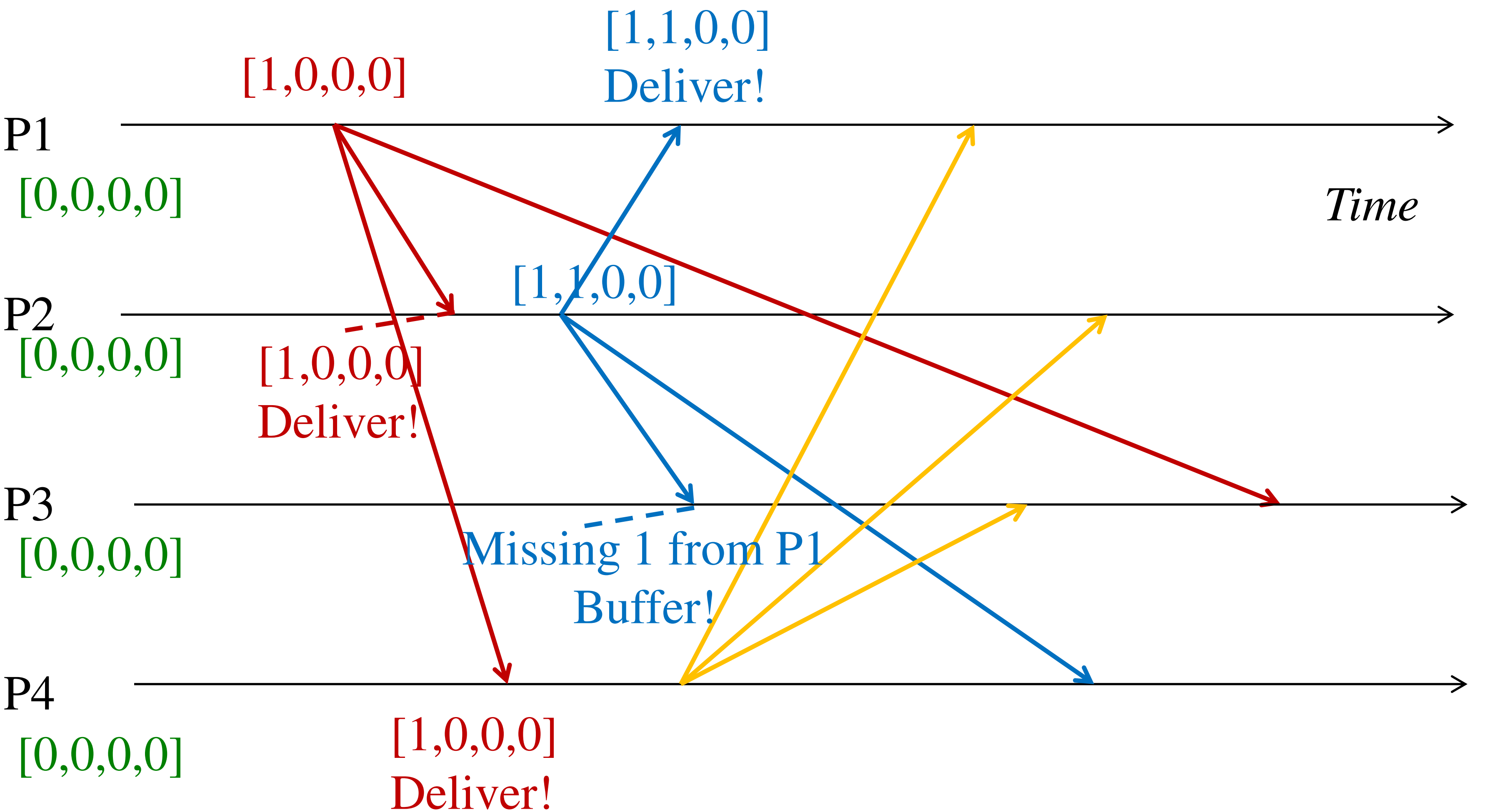

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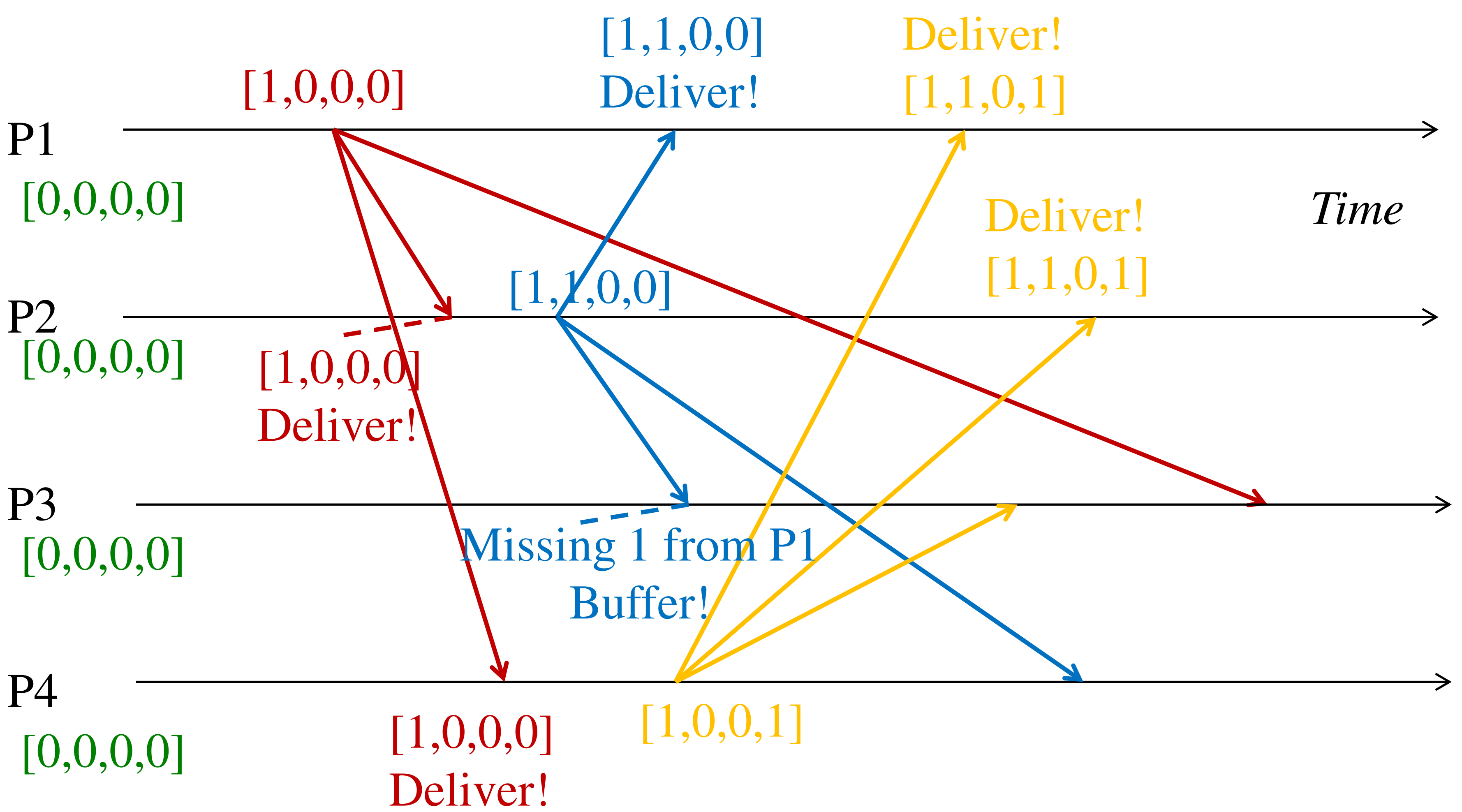

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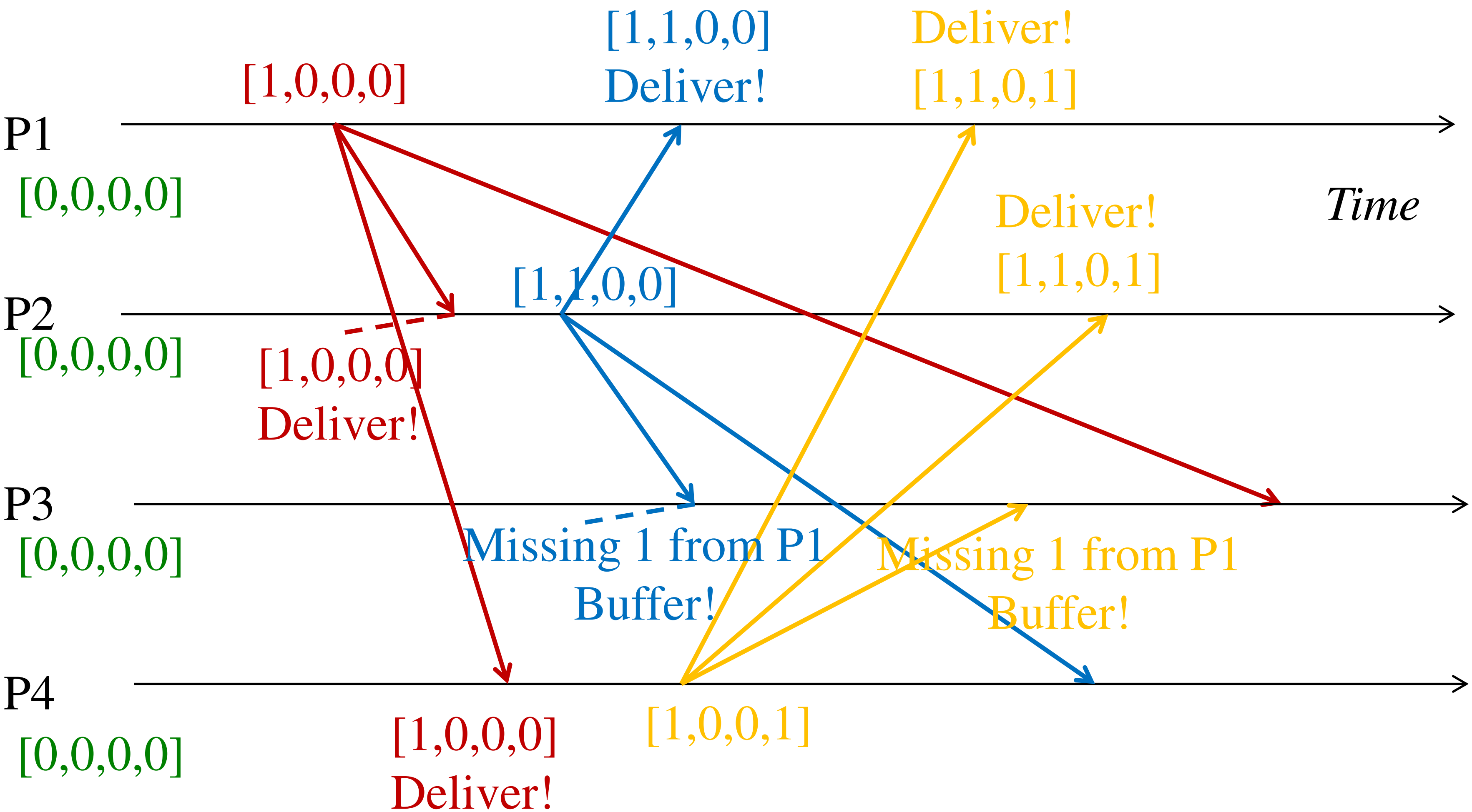

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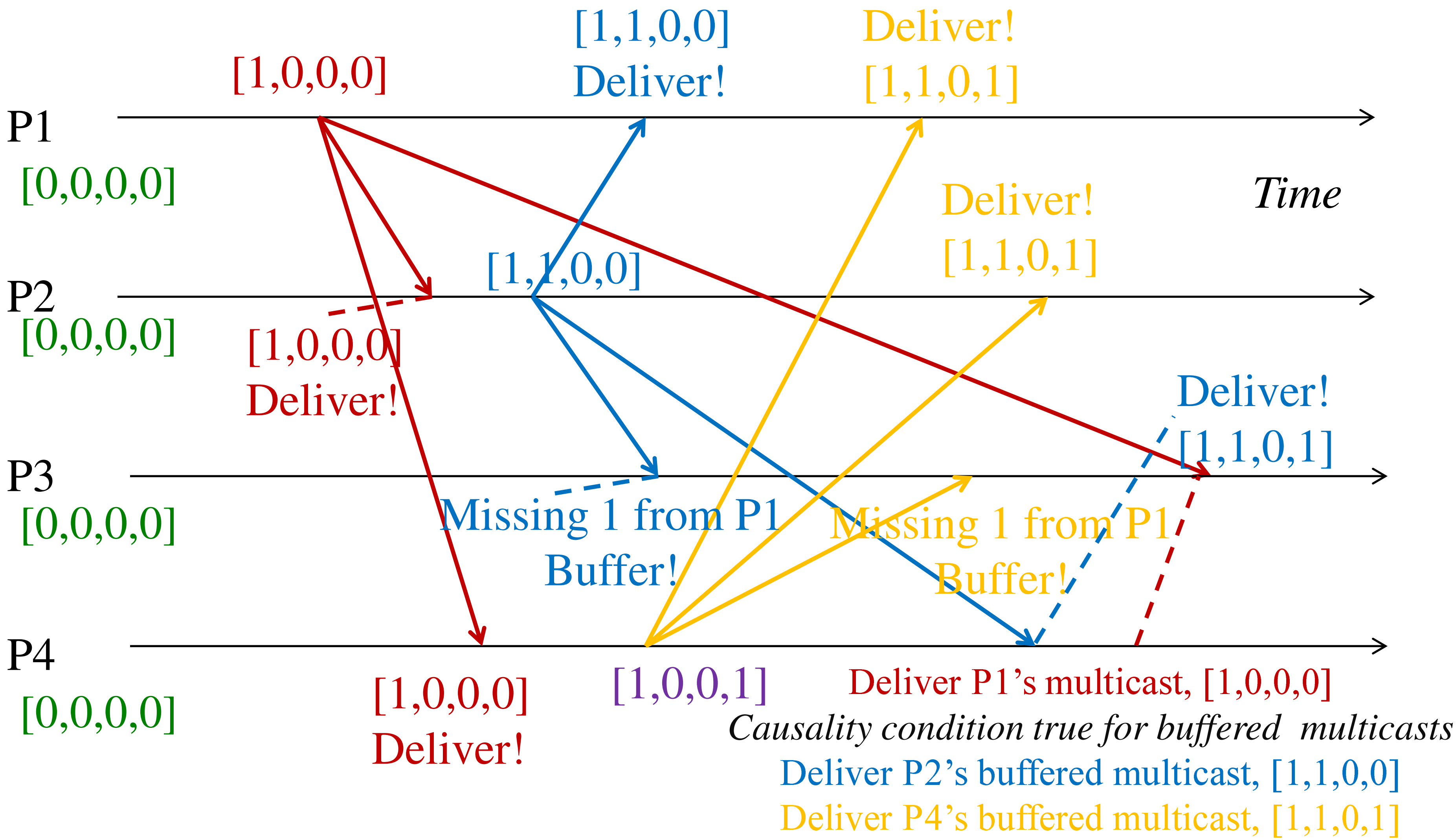
Causal order multicast execution



Causal order multicast implementation

- Only looks at multicast messages delivered to the application.

- Ignores causality created due to other network messages.

Ordered Multicast

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- Total ordering
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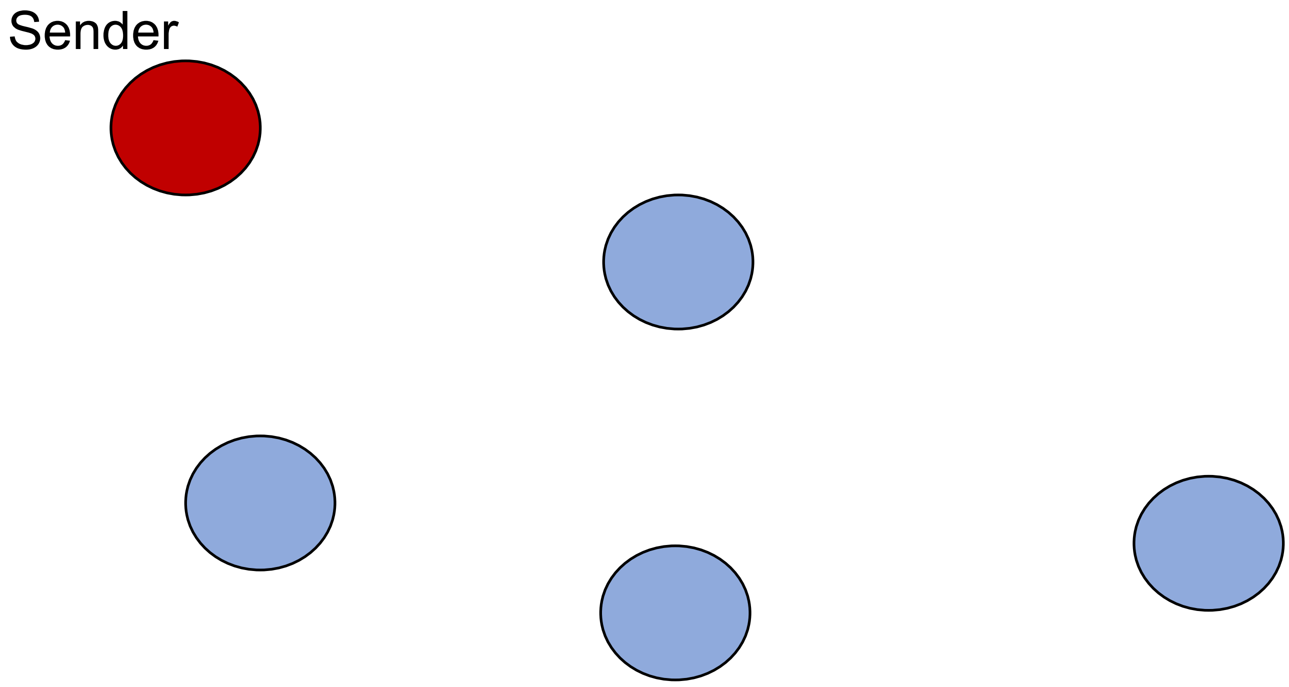
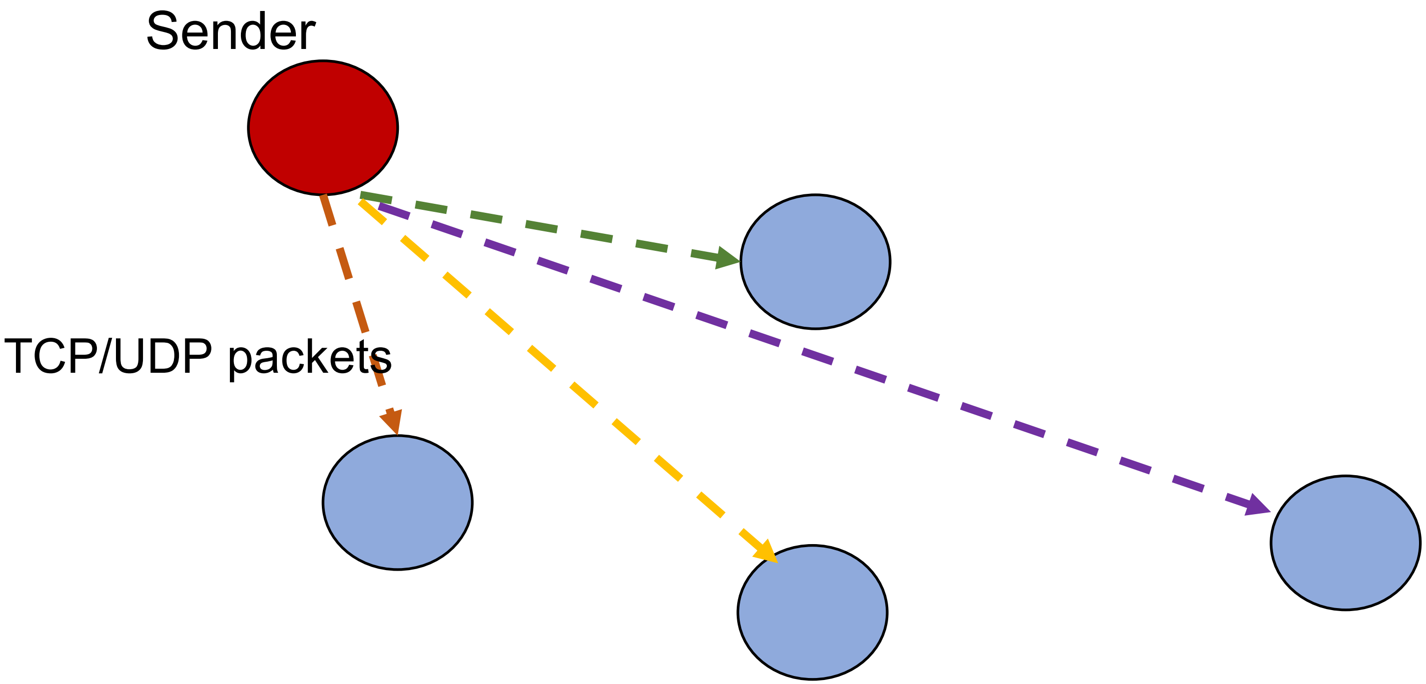
More efficient multicast mechanisms

- Our focus so far has been on the application-level semantics of multicast.
- *What are some of the more efficient underlying mechanisms for a B-multicast?*

B-Multicast

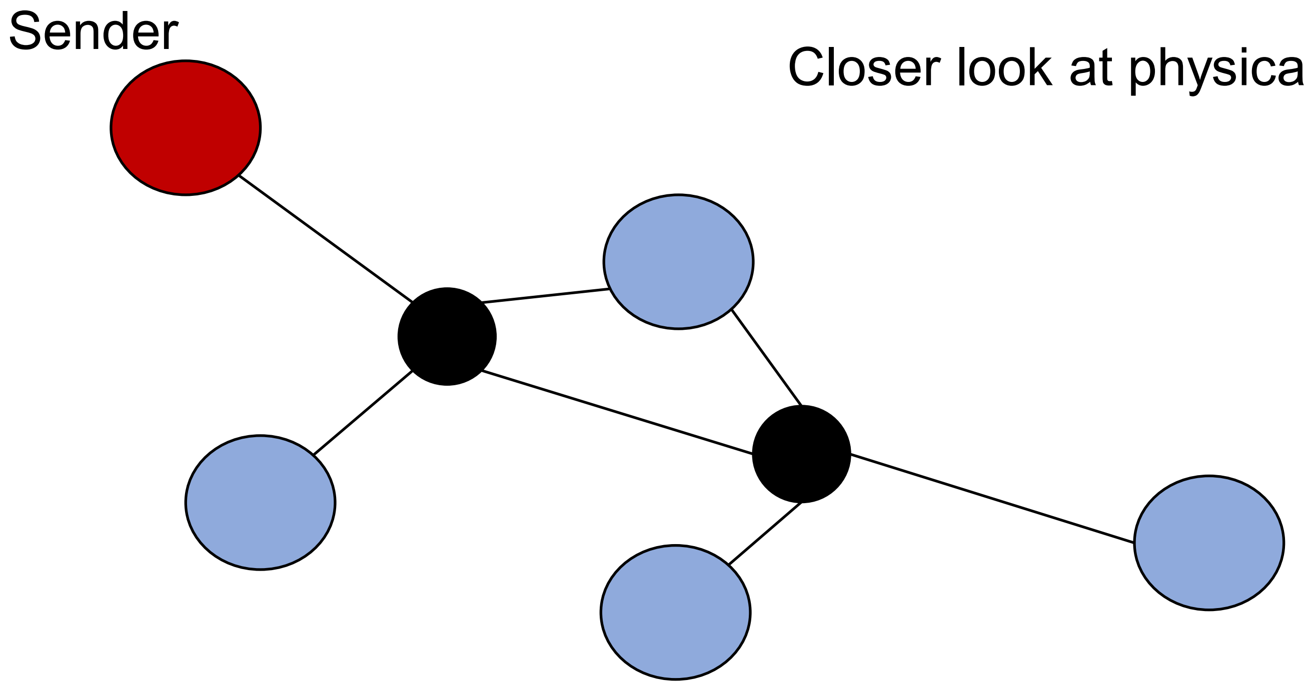


B-Multicast using unicast sends



B-Multicast using unicast sends

Closer look at physical network paths.



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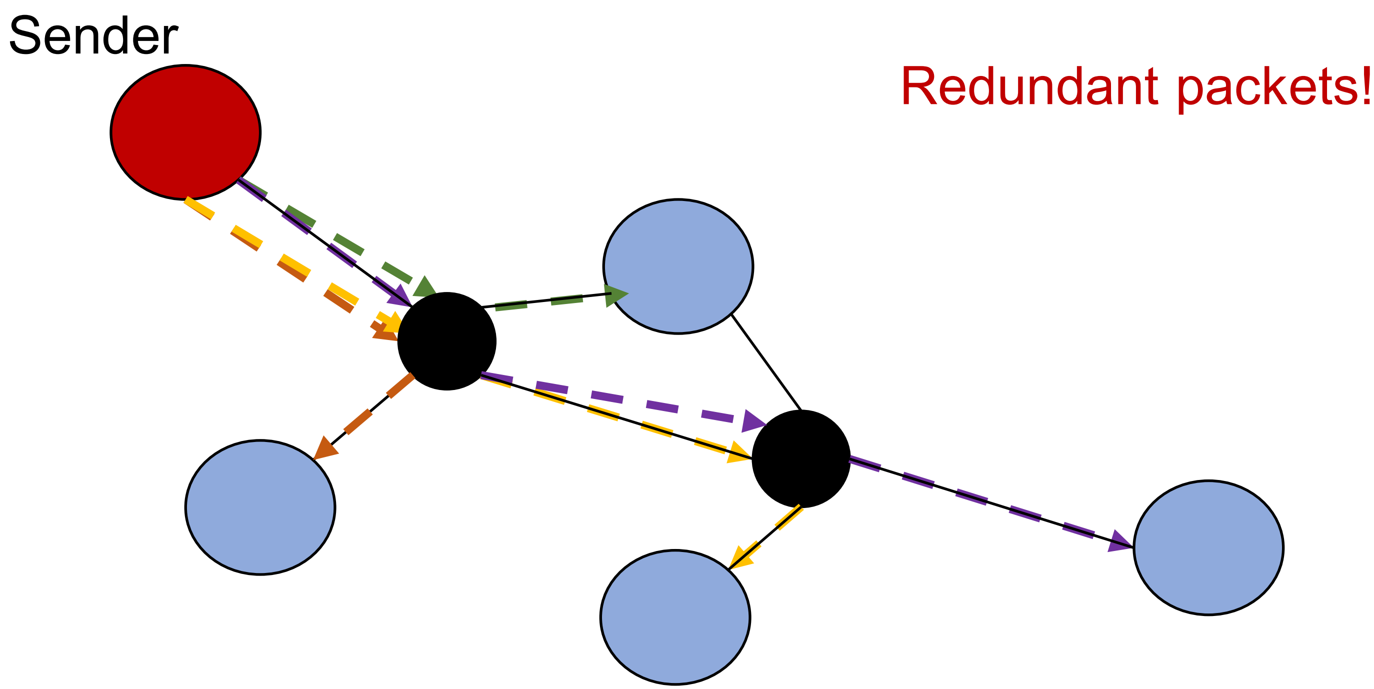

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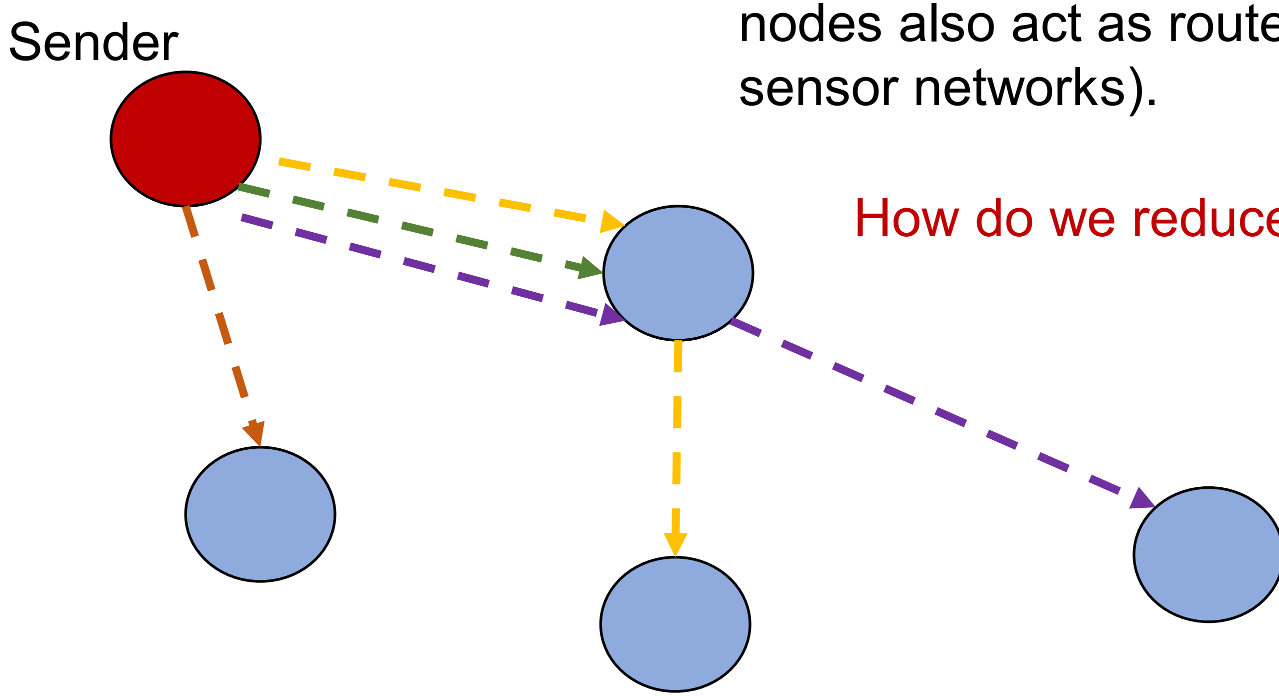
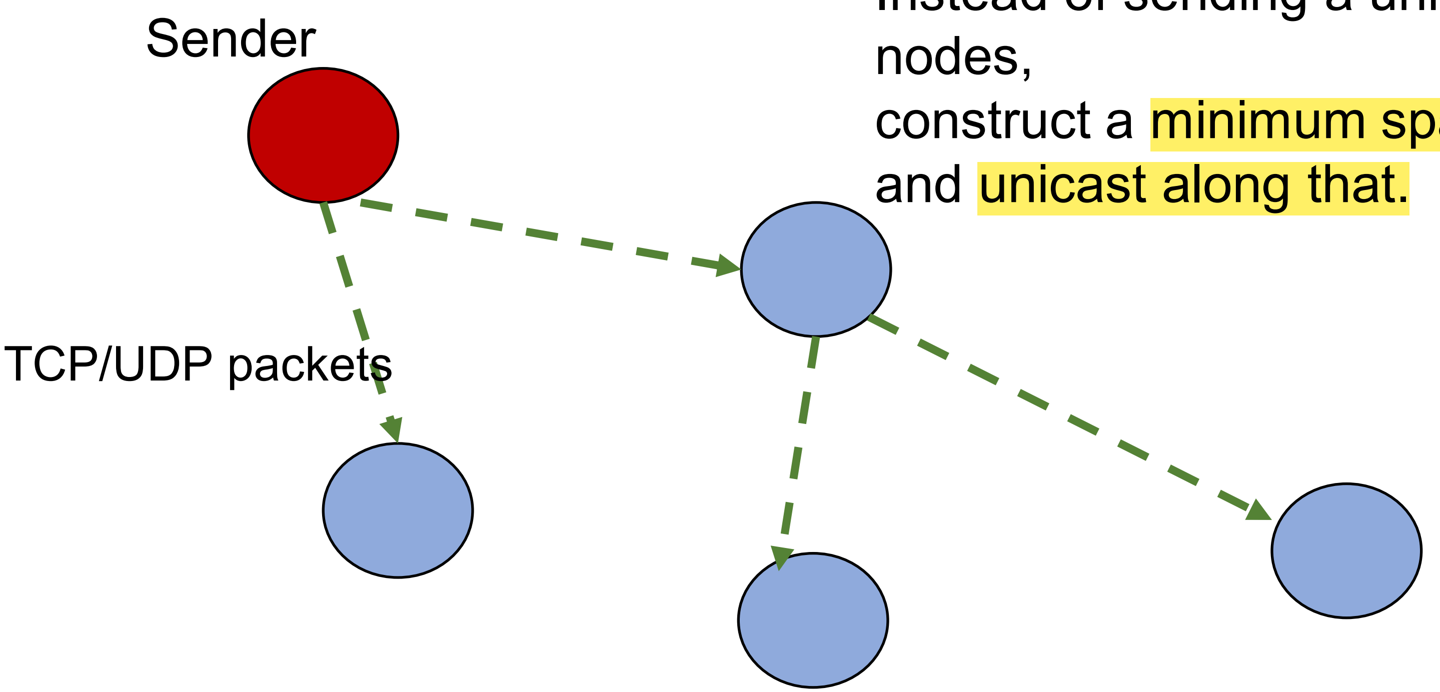


Similar redundancy when individual nodes also act as routers (e.g. wireless sensor networks).

How do we reduce the overhead?

Tree-based multicast

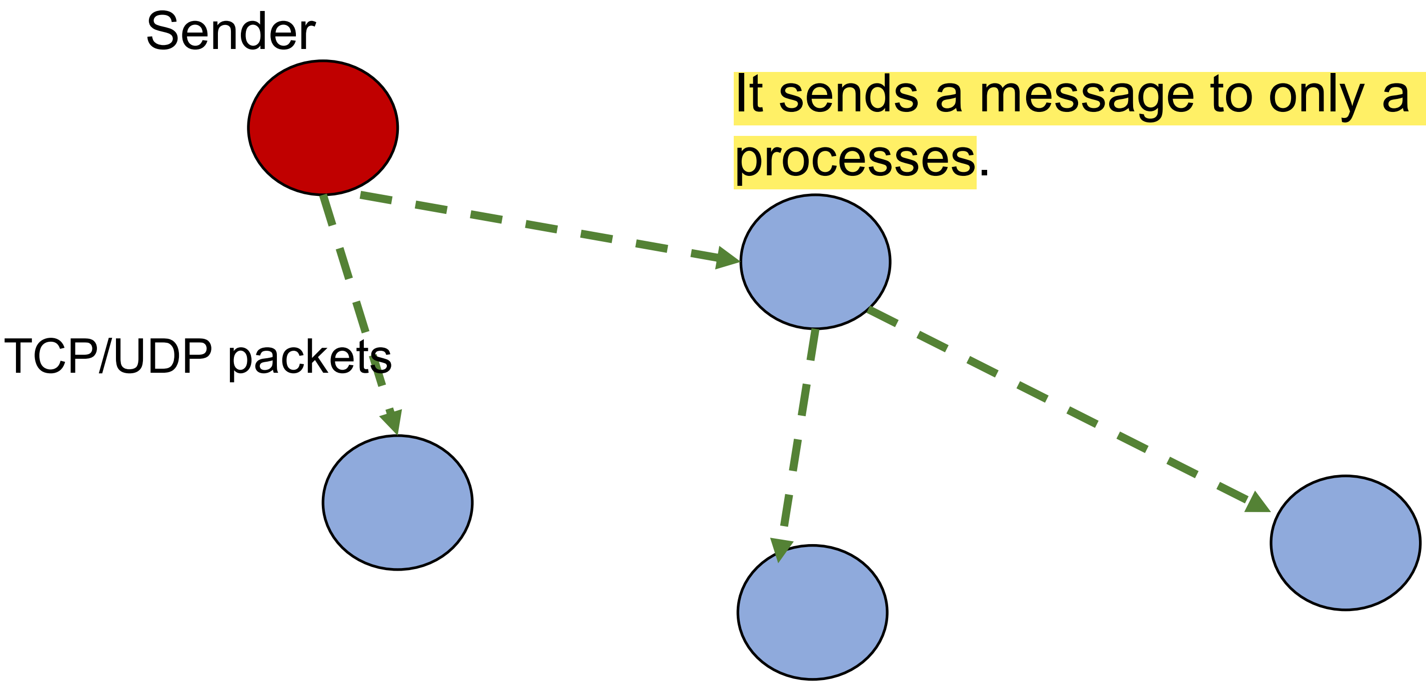


Instead of sending a unicast to all nodes,
construct a minimum spanning tree and unicast along that.

Tree-based multicast

A process does not directly send messages to *all* other processes in the group.

It sends a message to only a subset of processes.



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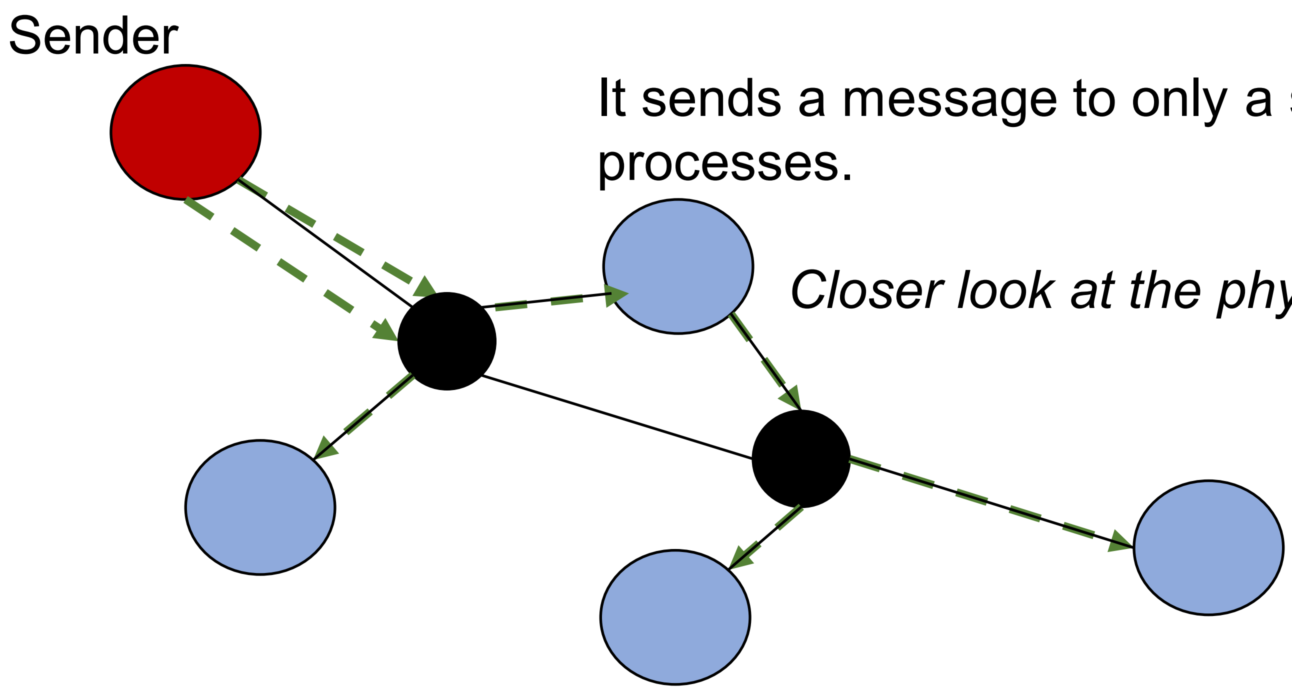
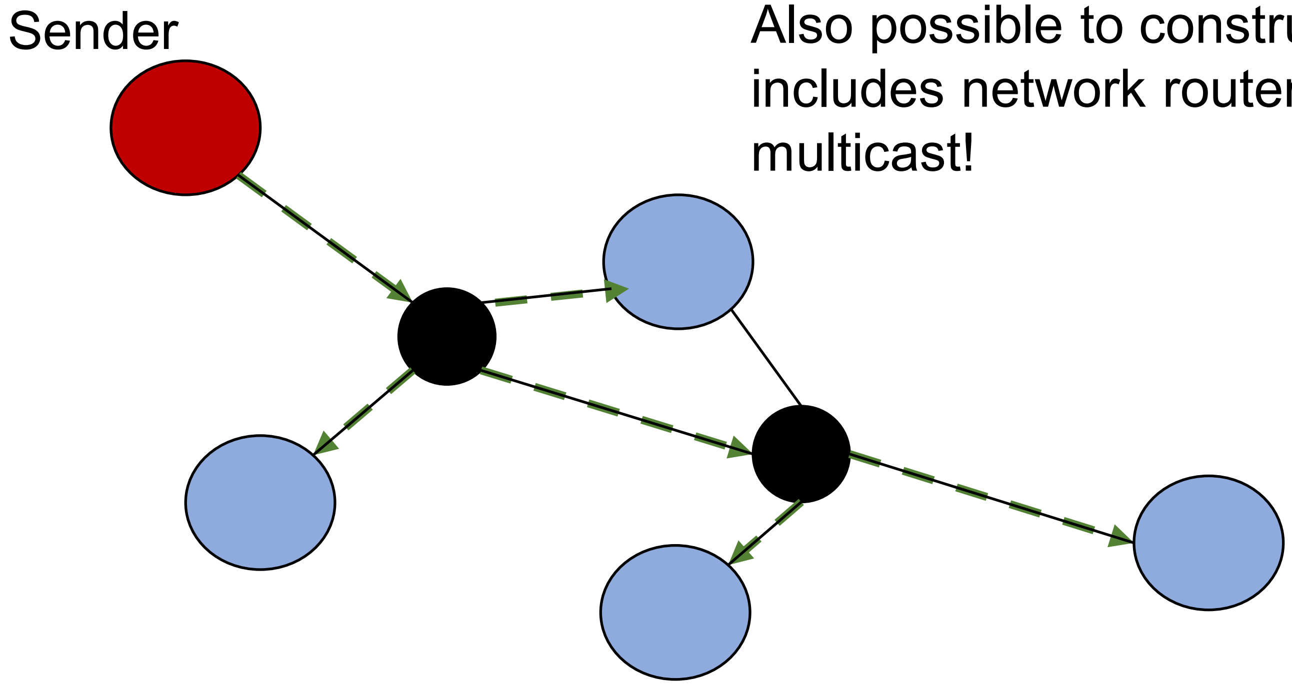
Closer look at the physical network.

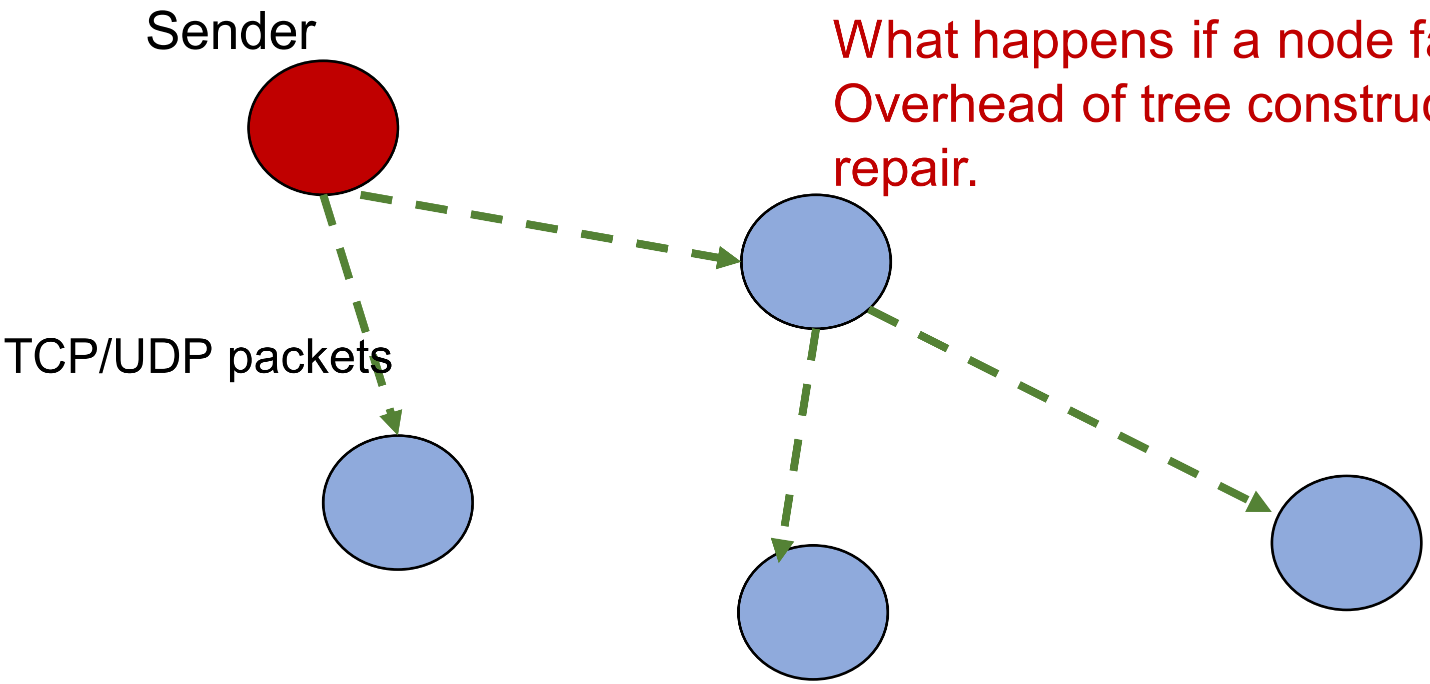


Tree-based multicast



Also possible to construct a tree that includes network routers. IP multicast!

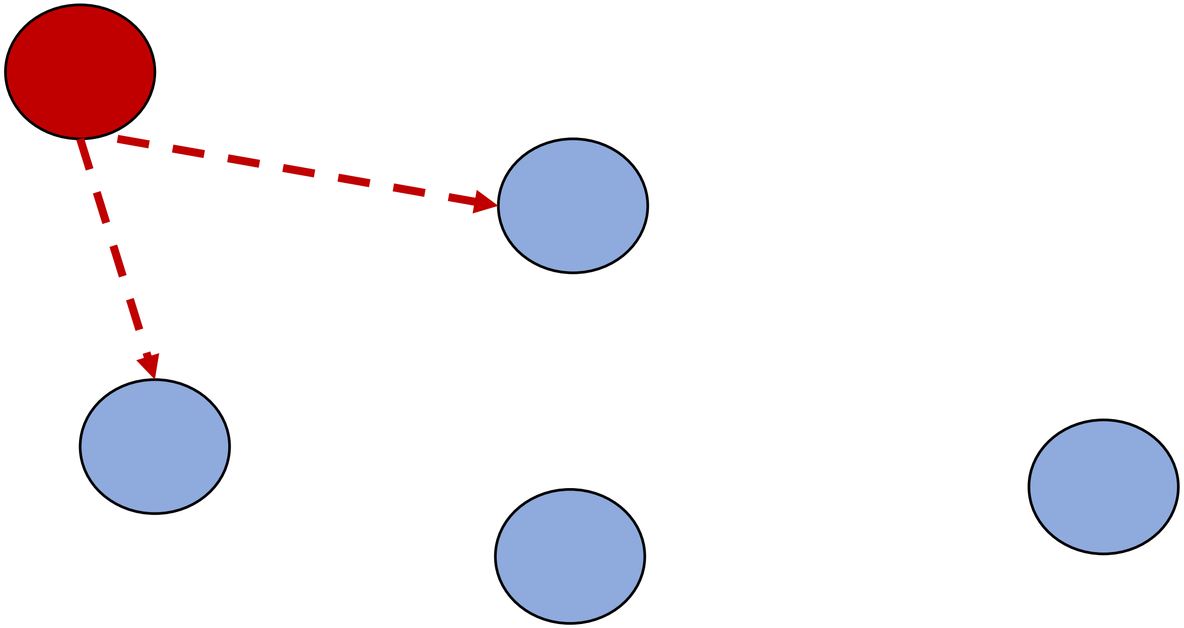
Tree-based multicast



What happens if a node fails?
Overhead of tree construction and repair.

Third approach: Gossip

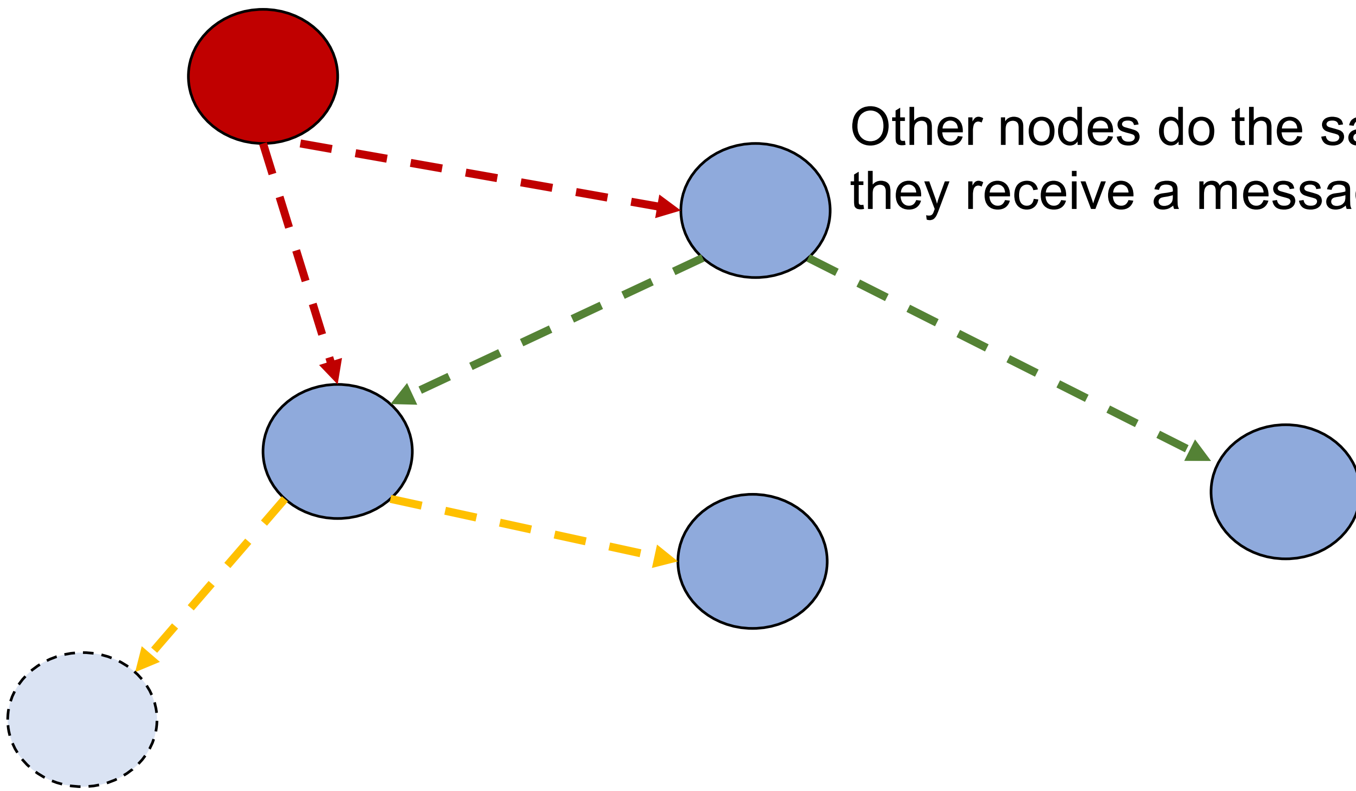
Transmit to b random targets.



Third approach: Gossip

Transmit to b random targets.

Other nodes do the same when they receive a message.



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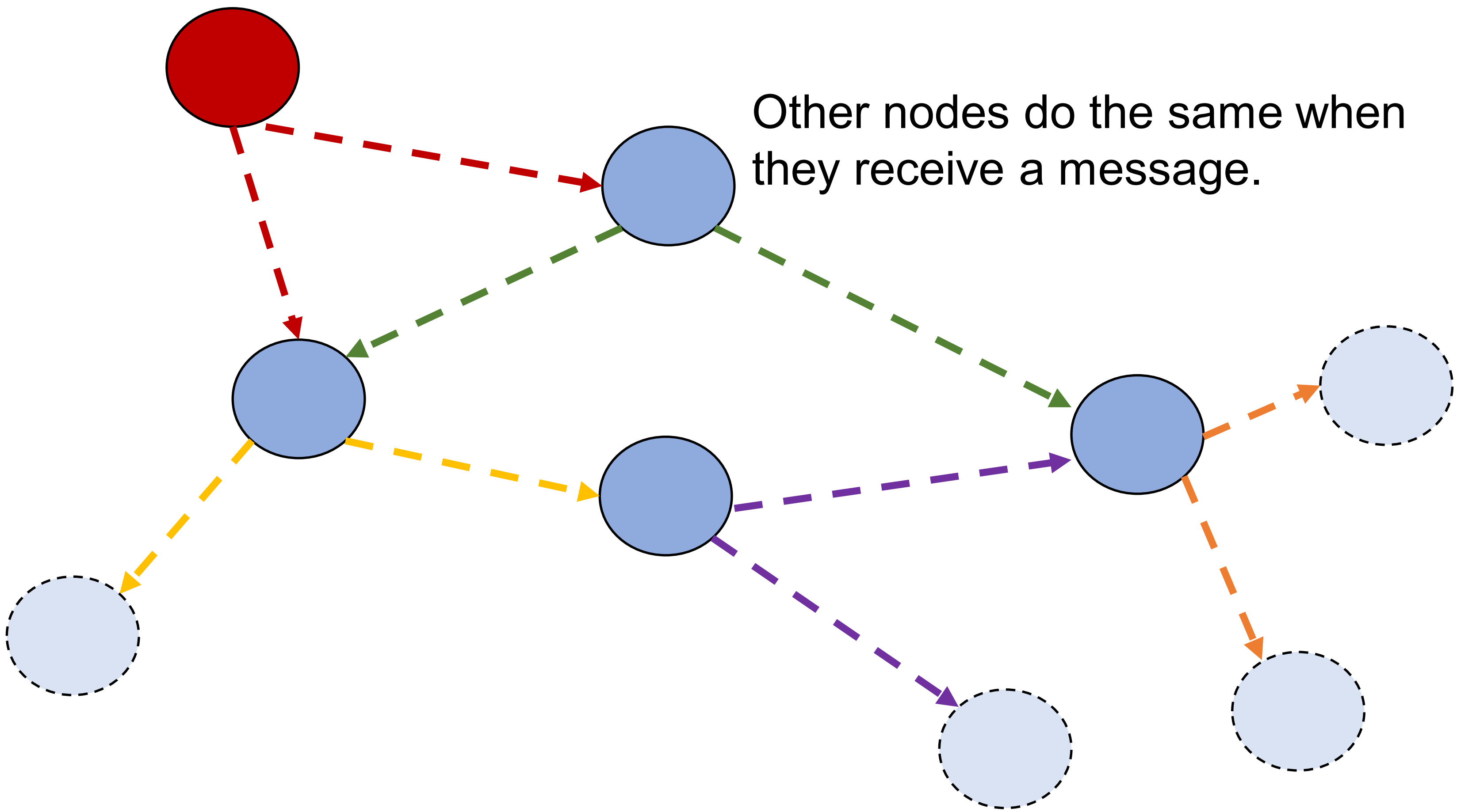
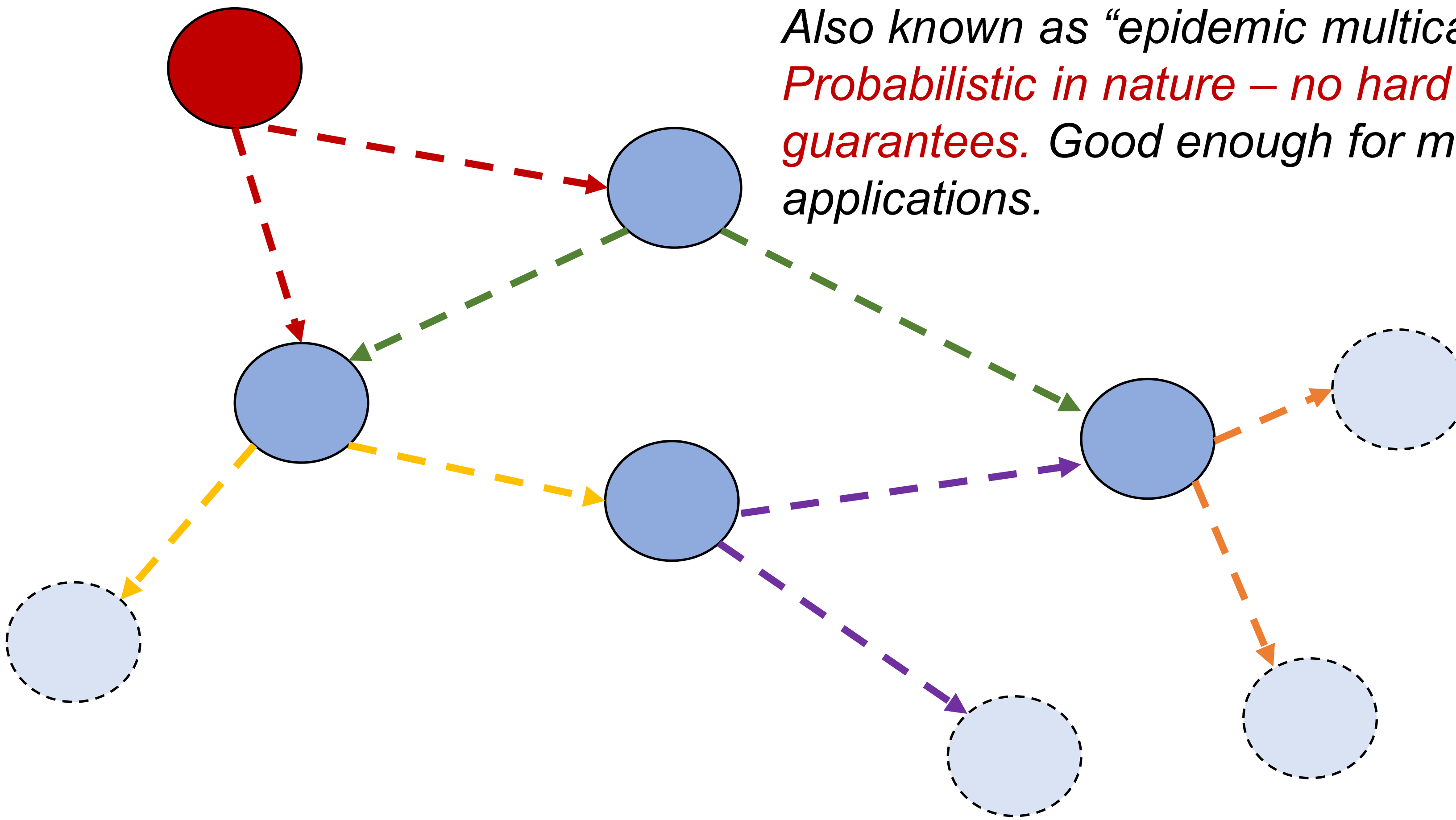


Third approach: Gossip

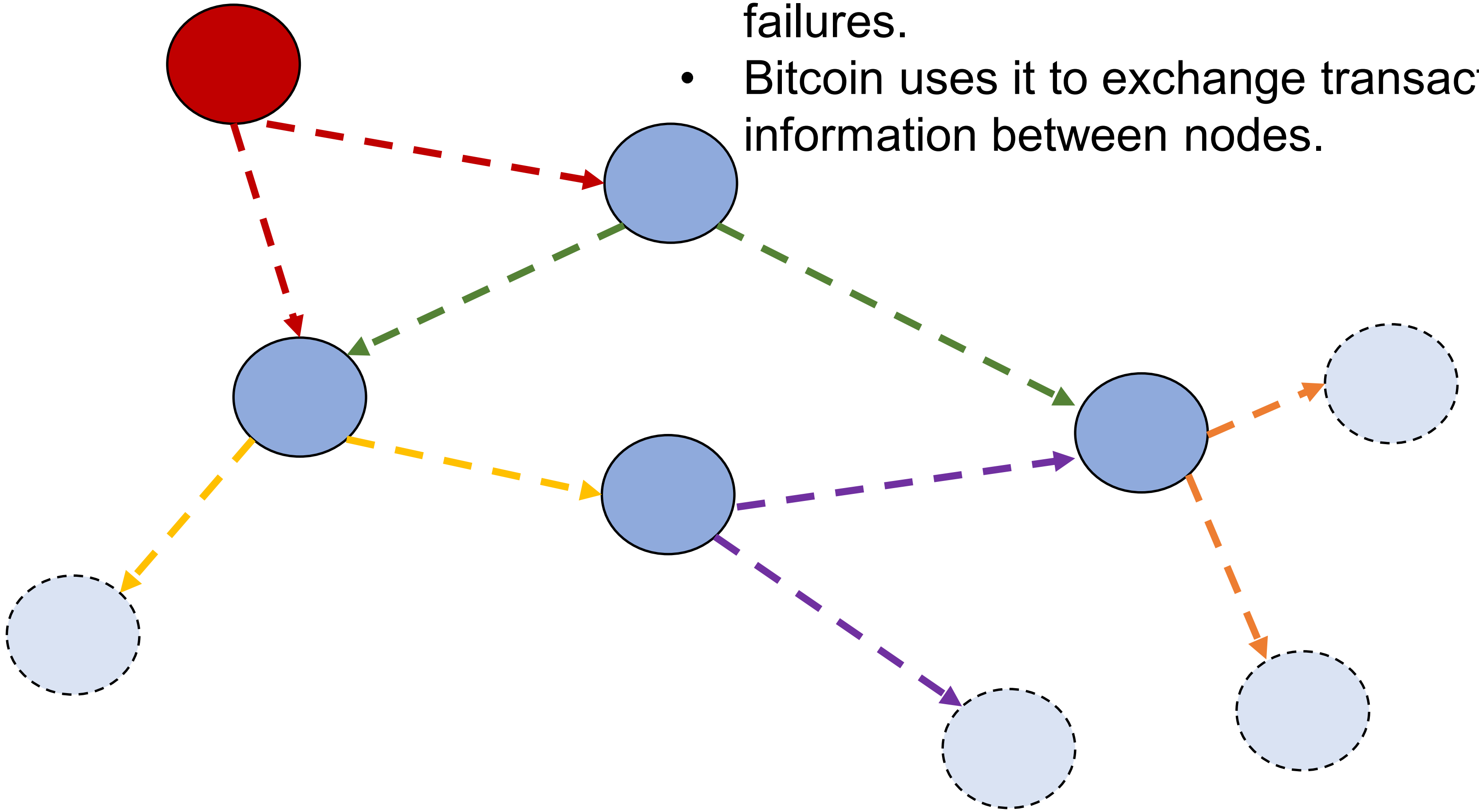
No "tree-construction" overhead.
More efficient than unicasting to all receivers.
Also known as "epidemic multicast".
Probabilistic in nature – no hard guarantees. Good enough for many applications.



Third approach: Gossip

Used in many real-world systems:

- Facebook's distributed datastore uses it to determine group membership and failures.
- Bitcoin uses it to exchange transaction information between nodes.



Multicast Summary

- Multicast is important for applications in distributed systems.
- Applications may have different requirements:
  - Basic
  - Reliable
  - Ordering: FIFO, Causal, Total
  - Combinations of the above.
- Underlying mechanisms to spread the information:
  - Unicast to all receivers.
  - Tree-based multicast, and gossip: sender unicasts messages to only a subset of other processes, and they spread the message further.
  - Gossip is more scalable and more robust to failures.

Today's agenda

- Wrap up Multicast
  - Chapter 15.4
  - Tree-based multicast and Gossip

- Mutual Exclusion
  - Chapter 15.2

- Goal: reason about ways in which different processes in a distributed system can safely manipulate shared resources.

Why Mutual Exclusion?

- Bank's Servers in the Cloud: Two of your customers make simultaneous deposits of $1,000 into your bank account, each from a separate ATM.
  - Both ATMs read initial amount of $10,000 in your account concurrently from the bank's cloud server
  - Both ATMs add $1,000 to this amount (locally at ATM)
  - Both write the final amount to the server
  - What's wrong?

Why Mutual Exclusion?

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  - Both ATMs read initial amount of $10,000 in your account concurrently from the bank's cloud server
  - Both ATMs add $1,000 to this amount (locally at ATM)
  - Both write the final amount to the server
  - You lost $1,000!
- The ATMs need *mutually exclusive* access to your account entry at the server
  - or, mutually exclusive access to executing the code that modifies the account entry.

More uses of mutual exclusion

- Distributed file systems
  - Locking of files and directories
- Accessing objects in a safe and consistent way
  - Ensure at most one server has access to object at any point of time
- In industry
  - Chubby is Google’s locking service

Problem statement for mutual excls.

- *Critical Section* Problem:
  - Piece of code (at all processes) for which we need to ensure there is at most one process executing it at any point of time.

- Each process calls three functions
  - enter() in order to enter the critical section (CS)
  - AccessResource() to run the critical section code
  - exit() to exit the critical section

Our bank example

ATM1:

```
enter();
  // AccessResource()
obtain bank amount;
add in deposit;
update bank amount;
  // AccessResource() end
exit();
```

ATM2:

```
enter();
  // AccessResource()
obtain bank amount;
add in deposit;
update bank amount;
  // AccessResource() end
exit();
```

Mutual exclusion for a single OS

- If all processes are running in one OS on the same machine (or VM):
  - Semaphores
  - Mutexes
  - Condition variables
  - Monitors
  - …

Processes sharing an OS: Semaphores

- Semaphore == an integer that can only be accessed via two special functions
- Semaphore S=1; // Max number of allowed accessors.

wait(S) (or P(S) or down(S)):

```
while(1) { // each execution of the while loop is atomic
  if (S > 0) {
        S--;
        break;
    }
}
```

enter()

signal(S) (or V(S) or up(s)):

```
S++; // atomic
```

exit()

Atomic operations are supported via hardware instructions such as compare-and-swap, test-and-set, etc.

Our bank example

ATM1:

```
enter();
  // AccessResource()
obtain bank amount;
add in deposit;
update bank amount;
  // AccessResource() end
exit();
```

ATM2:

```
enter();
  // AccessResource()
obtain bank amount;
add in deposit;
update bank amount;
  // AccessResource() end
exit();
```

Our bank example

Semaphore S=1; // shared

ATM1:

```
wait(S);   //enter
  // AccessResource()
obtain bank amount;
add in deposit;
update bank amount;
  // AccessResource() end
signal(S); // exit
```

ATM2:

```
wait(S);  //enter
  // AccessResource()
obtain bank amount;
add in deposit;
update bank amount;
  // AccessResource() end
signal(S); // exit
```

Mutual exclusion in distributed systems

- Processes communicating by passing messages.
- Cannot share variables like semaphores!
- *How do we support mutual exclusion in a distributed system?*

Mutual exclusion in distributed systems

- Our focus today: Classical algorithms for mutual exclusion in distributed systems.
  - Central server algorithm
  - Ring-based algorithm
  - Ricart-Agrawala Algorithm
  - Maekawa Algorithm

Mutual Exclusion Requirements

- Need to guarantee 3 properties:
  - Safety (essential):
    - At most one process executes in CS (Critical Section) at any time.
  - Liveness (essential):
    - Every request for a CS is granted eventually.
  - Ordering (desirable):
    - Requests are granted in the order they were made.

System Model

- Each pair of processes is connected by reliable channels (such as TCP).
- Messages sent on a channel are eventually delivered to recipient, and in FIFO order.
- Processes do not fail.
  - Fault-tolerant variants exist in literature.

Mutual exclusion in distributed systems

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Central Server Algorithm

- Elect a central server (or leader)
- Leader keeps
  - A queue of waiting requests from processes who wish to access the CS
  - A special token which allows its holder to access CS
- Actions of any process in group:
  - enter()
    - Send a request to leader
    - Wait for token from leader
  - exit()
    - Send back token to leader

Central Server Algorithm

- Leader Actions:
  - On receiving a request from process P*i*

    if (leader has token)

    Send token to P*i*

    else

    Add P*i* to queue

  - On receiving a token from process P*i*

    if (queue is not empty)

    Dequeue head of queue (say P*j*), send that process the token

    else

    Retain token

Analysis of Central Algorithm

- Safety – at most one process in CS
  - Exactly one token
- Liveness – every request for CS granted eventually
  - With *N* processes in system, queue has at most *N* processes
  - If each process exits CS eventually and no failures, liveness guaranteed
- Ordering:
  - FIFO ordering guaranteed in order of requests received at leader
  - Not in the order in which requests were sent or the order in which processes enter CS!

Analysis of Central Algorithm

- Safety – at most one process in CS
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To be continued in next class

- Metrics for analyzing performance of mutual exclusion algorithms.

- Other algorithms for mutual exclusion in distributed systems.
  - Central server algorithm
  - Ring-based algorithm
  - Ricart-Agrawala Algorithm
  - Maekawa Algorithm