



Optimization in Machine Learning: Lecture 8

Duality

by Xiaolin Huang

xiaolinhuang@sjtu.edu.cn

SEIEE 2-429

Institute of Image Processing and Pattern Recognition

<http://www.pami.sjtu.edu.cn/>



上海交通大學

SHANGHAI JIAO TONG UNIVERSITY

1

Langrage Duality

2

Optimality Condition

3

Support Vector Machine



1

Langrage Duality

2

Optimality Condition

3

Support Vector Machine



Lagrangian



- consider the standard form

$$\begin{array}{ll}\min & f_0(x) \\ \text{s. t.} & \underline{f_i(x) \leq 0, i = 1, \dots, m} \\ & h_i(x) = 0, i = 1, \dots, p\end{array}$$



- its Lagrangian is

$$L(x, \lambda, v) = f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p v_i h_i(x)$$

- weighted sum of the objective and the constraint functions
- $L: \mathbf{R}^n \times \mathbf{R}^m \times \mathbf{R}^p \rightarrow \mathbf{R}$
- $\text{dom } L: D \times \mathbf{R}^m \times \mathbf{R}^p$

Lagrange dual function



- Lagrange dual function $g: \mathbf{R}^m \times \mathbf{R}^p \rightarrow \mathbf{R}$

$$\begin{aligned} g(\lambda, v) &= \inf_{x \in D} L(x, \lambda, v) \\ &= \inf_{x \in D} \left(f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p v_i h_i(x) \right) \end{aligned}$$

- lower bound: if $\lambda \geq 0$, then $g(\lambda, v) \leq f^*$
 - proof : suppose \tilde{x} is feasible, i.e., $f_i(\tilde{x}) \leq 0, h_i(\tilde{x}) = 0$

then with $\lambda \geq 0$, we have




$$\begin{aligned} f_0(\tilde{x}) &\geq f_0(\tilde{x}) + \underbrace{\sum_{i=1}^m \lambda_i f_i(\tilde{x}) + \sum_{i=1}^p v_i h_i(\tilde{x})}_{\geq 0} \\ &\geq \inf_{x \in D} \left(f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p v_i h_i(x) \right) = g(\lambda, v) \end{aligned}$$

Lagrange dual and conjugate function



$$\begin{array}{ll} \min & f_0(x) \\ \text{s. t.} & x = 0 \end{array}$$

- conjugate function $f^*(y) = \sup_x (y^\top x - f(x))$
- Lagrangian: $L(x, \lambda, v) = f_0(x) + v^\top x$ 
- Lagrange dual function

$$g(v) = \inf_{x \in D} (f_0(x) + v^\top x) = -f_0^*(-v)$$

Lagrange dual and conjugate function



$$\begin{array}{ll}\min & f_0(x) \\ \text{s. t.} & Ax \leq b \\ & Cx = d\end{array}$$

- conjugate function $f^*(y) = \sup_x (y^\top x - f(x))$
- Lagrangian: $L(x, \lambda, \nu) = f_0(x) + \lambda^\top (Ax - b) + \nu^\top (Cx - d)$
- Lagrange dual function

$$\begin{aligned}g(\lambda, \nu) &= \inf_{x \in D} (f_0(x) + \lambda^\top (Ax - b) + \nu^\top (Cx - d)) \\ &= \inf_{x \in D} (f_0(x) + (\lambda^\top A + \nu^\top C)x - \lambda^\top b - \nu^\top d) \\ &= -f_0^*(-A^\top \lambda - C^\top \nu) - \lambda^\top b - \nu^\top d\end{aligned}$$

Dual problem



- $g(\lambda, v)$ could give a lower bound for the primary problem
- could we get a better, or even tightest, lower bound?

$$\begin{aligned} \max \quad & g(\lambda, v) \\ \text{s. t.} \quad & \lambda \geq 0 \end{aligned}$$

- it is convex, even the primal problem is not

▪ LP $f(x) = f_0^\top x \rightarrow f^*(y) = \sup_{x \in \text{dom } f} (x^\top (y - f_0)) = \begin{cases} 0, & y = f_0 \\ +\infty, & y \neq f_0 \end{cases}$

$\begin{aligned} \min \quad & f_0^\top x \\ \text{s. t.} \quad & Ax \leq b \\ & Cx = d \end{aligned}$		$\begin{aligned} \max \quad & -\lambda^\top b - v^\top d \\ \text{s. t.} \quad & A^\top \lambda + C^\top v = f_0 \\ & \lambda \geq 0 \end{aligned}$
-------------------------------------------------------------------------------------------------------	-------------------------------------------------------------------------------------	-----------------------------------------------------------------------------------------------------------------------------------------------------

$$g(\lambda, v) = -f_0^*(-A^\top \lambda - C^\top v) - \lambda^\top b - v^\top d$$

Dual problem example: SVM



- SVM in primal space

$$\begin{aligned} \min_{x,z} \quad & \frac{1}{2} \|x\|_2^2 + C \sum_i \rho_i \\ \text{s. t.} \quad & b_i(x^\top a_i + z) \geq 1 - \rho_i \\ & \rho_i \geq 0 \end{aligned}$$

- the corresponding Lagrangian

$$\begin{aligned} L(x, z, \rho; \lambda, \nu) = & \frac{1}{2} \|x\|_2^2 + C \sum_i \rho_i \\ & + \sum_i \lambda_i (1 - \rho_i - b_i(x^\top a_i + z)) - \nu^T \rho \end{aligned}$$

Dual problem example: SVM



- to obtain the $g(\lambda, \nu) = \inf_{x, z, \rho} L(x, z, \rho; \lambda, \nu)$

$$L(x, z, \rho; \lambda, \nu) = \frac{1}{2} \|x\|_2^2 + C \sum_i \rho_i + \sum_i \lambda_i (1 - \rho_i - b_i(x^\top a_i + z)) - \nu^\top \rho$$

$$\frac{\partial L}{\partial x} = x - \sum_i b_i \lambda_i a_i = 0$$

$$\frac{\partial L}{\partial z} = \sum_i \lambda_i b_i = 0$$

$$\frac{\partial L}{\partial \rho_i} = C - \lambda_i - \nu_i = 0$$

$$\lambda_i \geq 0, \nu_i \geq 0$$



$$\begin{aligned} \min_{\lambda} \quad & \sum_i \sum_j \lambda_i b_i a_i^\top a_j b_j \lambda_j - \sum_i \lambda_i \\ \text{s. t.} \quad & \sum_i \lambda_i b_i = 0 \\ & 0 \leq \lambda_i \leq C \end{aligned}$$

Dual problem example: LP



- two products I and II with different profit
- three resources A, B, and C with different inventories

	Product I	Product II	inventory
Resource A	0	5	15
Resource B	6	2	24
Resource C	1	1	5
Profit	2	1	

$$\begin{aligned} \max \quad & 2x_1 + x_2 \\ \text{s. t.} \quad & 5x_2 \leq 15 \\ & 6x_1 + 2x_2 \leq 24 \\ & x_1 + x_2 \leq 5 \\ & x_1, x_2 \geq 0 \end{aligned}$$

Weak duality



■ Primal

$$\begin{array}{ll}\max & 2x_1 + x_2 \\ \text{s. t.} & 5x_2 \leq 15 \\ & 6x_1 + 2x_2 \leq 24 \\ & x_1 + x_2 \leq 5 \\ & x_1, x_2 \geq 0\end{array}$$

Dual

$$\begin{array}{ll}\min & 15y_1 + 24y_2 + 5y_3 \\ \text{s. t.} & 6y_2 + y_3 \geq 2 \\ & 5y_1 + 2y_2 + y_3 \geq 1 \\ & y_1, y_2, y_3 \geq 0\end{array}$$

■ Weak duality

$$2x_1 + x_2 \leq 15y_1 + 24y_2 + 5y_3$$

for any feasible solution

Strong duality



■ Primal

$$\begin{array}{ll}\max & 2x_1 + x_2 \\ \text{s. t.} & 5x_2 \leq 15 \\ & 6x_1 + 2x_2 \leq 24 \\ & x_1 + x_2 \leq 5 \\ & x_1, x_2 \geq 0\end{array}$$

Dual

$$\begin{array}{ll}\min & 15y_1 + 24y_2 + 5y_3 \\ \text{s. t.} & 6y_2 + y_3 \geq 2 \\ & 5y_1 + 2y_2 + y_3 \geq 1 \\ & y_1, y_2, y_3 \geq 0\end{array}$$

■ Strong duality

$$2x_1^* + x_2^* = 15y_1^* + 24y_2^* + 5y_3^*$$

- for the optimal solution

Complementary slackness



■ Primal

$$\begin{array}{ll}\max & 2x_1 + x_2 \\ \text{s. t.} & 5x_2 \leq 15 \\ & 6x_1 + 2x_2 \leq 24 \\ & x_1 + x_2 \leq 5 \\ & x_1, x_2 \geq 0\end{array}$$

Dual

$$\begin{array}{ll}\min & 15y_1 + 24y_2 + 5y_3 \\ \text{s. t.} & 6y_2 + y_3 \geq 2 \\ & 5y_1 + 2y_2 + y_3 \geq 1 \\ & y_1, y_2, y_3 \geq 0\end{array}$$

■ Complementary slackness

$$5x_2^* < 15 \rightarrow y_1^* = 0$$

- if there are redundant resource A for the optimal product plan, giving up A will not affect our profile, so the shadow price is zero.
- instrict inequality/inactive constraint correspond to zero dual variable

Primal-dual relationship



■ Primal

$$\begin{aligned} \min \quad & f_0(x) \\ \text{s. t.} \quad & f_i(x) \leq 0, i = 1, \dots, m \\ & h_i(x) = 0, i = 1, \dots, p \end{aligned}$$

Dual

$$\begin{aligned} \max \quad & g(\lambda, v) = \inf_{x \in D} (f_0(x) + \lambda^\top f(x) + v^\top h(x)) \\ \text{s. t.} \quad & \lambda \geq 0 \end{aligned}$$

$$g^* = g(\lambda^*, v^*)$$

$$f^* = f_0(x^*)$$

■ weak duality $f(x) \geq g(\lambda, v), \quad f^* \geq g^*$

■ strong duality $f^* = g^*$

■ slackness condition $\lambda_i^* f_i(x^*) = 0$

$$\lambda_i^* > 0 \rightarrow f_i(x^*) = 0$$

$$f_i(x^*) < 0 \rightarrow \lambda_i^* = 0$$

Duality gap:

$$\min f(x) - g(\lambda, v)$$

Slater's constraint qualification



- for the primal problem

$$\begin{array}{ll}\min & f_0(x) \\ \text{s. t.} & f_i(x) \leq 0, i = 1, \dots, m \\ & h_i(x) = 0, i = 1, \dots, p\end{array}$$

Slater's constraint qualification requires the problem is strictly feasible:

$$\exists x \in \text{int } D, f_i(x) < 0, h_i(x) = 0$$

- if the problem is convex and the Slater's qualification satisfied, then there is

strong duality

- there are many other qualifications

Strong duality proof



- convexity, $\text{int } D \neq \emptyset$, A is full ranked, and the optimal value is f^*

$$\begin{array}{ll}\min & f_0(x) \\ \text{s. t.} & f_i(x) \leq 0, i = 1, \dots, m \\ & Ax - b = 0, i = 1, \dots, p\end{array}$$

- consider a set

$$C = \{(u, v, t) : \exists x \in D, f_i(x) \leq u_i, \forall i, Ax - b = v_i, \forall i, f_0(x) \leq t\}$$

- it is convex and do not have joint point with the following convex set

$$D = \{(0, 0, s) \in \mathbf{R}^m \times \mathbf{R}^p \times \mathbf{R} : s < f^*\}$$

Strong duality proof



- using separating hyperplane theorem for

$$C = \{(u, v, t): \exists x \in D, f_i(x) \leq u_i, \forall i, Ax - b = v_i, \forall i, f_0(x) \leq t\}$$

$$E = \{(0, 0, t) \in \mathbf{R}^m \times \mathbf{R}^p \times \mathbf{R}: t < f^*\}$$

- there exists nonzero $(\tilde{\lambda}, \tilde{v}, \mu) \in \mathbf{R}^m \times \mathbf{R}^p \times \mathbf{R}$ and α , such that

$$\tilde{\lambda}^\top u + \tilde{v}^\top v + \mu t \geq \alpha, \quad \forall (u, v, t) \in C$$

$$\searrow \quad \tilde{\lambda} \geq 0, \mu \geq 0$$

$$\tilde{\lambda}^\top u + \tilde{v}^\top v + \mu t \leq \alpha, \quad \forall (u, v, t) \in E \quad \left. \vphantom{\tilde{\lambda}^\top u + \tilde{v}^\top v + \mu t \leq \alpha} \right\} \quad \mu f^* \leq \alpha$$

$$\searrow \quad \mu t \leq \alpha, \forall t \leq f^*$$

Strong duality proof



- together with $\tilde{\lambda}^\top u + \tilde{v}^\top v + \mu t \geq \alpha, \forall (u, v, t) \in C$ and $\mu f^* \leq \alpha$, we have

$$C = \{(u, v, t): \exists x \in D, f_i(x) \leq u_i, \forall i, Ax - b = v, \forall i, f_0(x) \leq t\}$$

$$\sum_{i=1}^m \tilde{\lambda}_i f_i(x) + \tilde{v}^\top (Ax - b) + \mu f_0(x) \geq \alpha \geq \mu f^*, \forall x \in D$$

- if $\mu \neq 0$

$$\sum_{i=1}^m \frac{\tilde{\lambda}_i}{\mu} f_i(x) + \frac{\tilde{v}^\top}{\mu} (Ax - b) + f_0(x) \geq f^*, \forall x \in D$$

$$g(\lambda, v) = \inf_{x \in D} (f_0(x) + v^\top f(x) + v^\top h(x)) \geq f^*$$

weak duality: $g(\lambda, v) \leq f^*$

$$\left. \begin{array}{l} g(\lambda, v) = \inf_{x \in D} (f_0(x) + v^\top f(x) + v^\top h(x)) \geq f^* \\ \text{weak duality: } g(\lambda, v) \leq f^* \end{array} \right\} \exists \lambda, v: g(\lambda, v) = f^*$$

Strong duality proof



- together with $\tilde{\lambda}^\top u + \tilde{v}^\top v + \mu t \geq \alpha, \forall (u, v, t) \in C$ and $\mu f^* \leq \alpha$, we have

$$C = \{(u, v, t) : \exists x \in D, f_i(x) \leq u_i, \forall i, Ax - b = v, f_0(x) \leq t\}$$

$$\sum_{i=1}^m \tilde{\lambda}_i f_i(x) + \tilde{v}^\top (Ax - b) + \mu f_0(x) \geq \alpha \geq \mu f^*, \forall x \in D$$

- if $\mu = 0$

$$\sum_{i=1}^m \tilde{\lambda}_i f_i(x) + \tilde{v}^\top (Ax - b) \geq 0, \forall x \in D$$

$$\tilde{v}^\top (Ax - b) \geq 0, \forall x \in D$$

$$\tilde{v}^\top (A\tilde{x} - b) = 0$$

slater's condition: there exist strictly feasible solutions \tilde{x}

$$\sum_{i=1}^m \tilde{\lambda}_i f_i(\tilde{x}) \geq 0 \Rightarrow \tilde{\lambda} = 0$$

$$\mu = 0$$

$(\tilde{\lambda}, \tilde{v}, \mu)$ is nonzero

$$\tilde{v} \neq 0$$

$$\tilde{v}^\top A = 0$$

contradict A is full ranked

Primal-dual relationship



Primal

$$\begin{aligned} \min \quad & f_0(x) \\ \text{s.t.} \quad & f_i(x) \leq 0, i = 1, \dots, m \\ & h_i(x) = 0, i = 1, \dots, p \end{aligned}$$

Dual

$$\begin{aligned} \max \quad & g(\lambda, v) = \inf_{x \in D} (f_0(x) + v^T f(x) + v^T h(x)) \\ \text{s.t.} \quad & \lambda \geq 0 \end{aligned}$$

$$g^* = g(\lambda^*, v^*)$$

$$f^* = f_0(x^*)$$

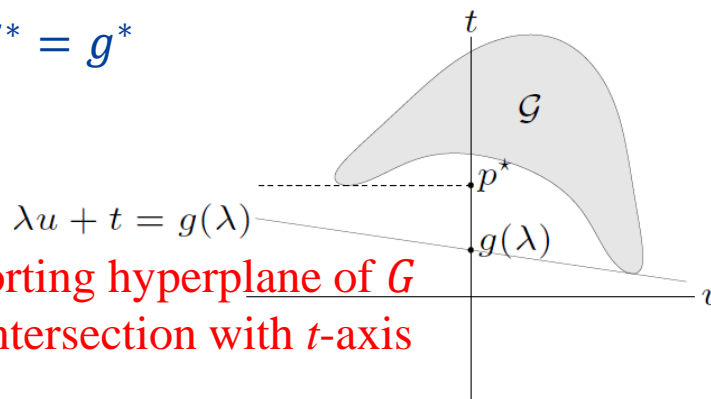
$$\min f_0(x), \text{ s.t. }, f_1(x) \leq 0$$

$$G = \{(f_1(x), f_0(x)), x \in D\}$$

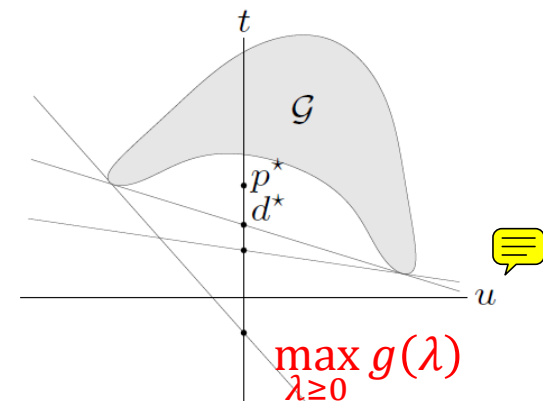
$$g(\lambda) = \inf_{(t,u) \in D} t + \lambda u$$

weak duality $f(x) \geq g(\lambda, v), \quad f^* \geq g^*$

strong duality $f^* = g^*$



$t + \lambda u = g(\lambda)$ is a supporting hyperplane of G
 $g(\lambda)$ is the value of the intersection with t -axis
 (since $\lambda \geq 0$)



Primal-dual relationship



■ Primal

$$\begin{aligned} \min \quad & f_0(x) \\ \text{s. t.} \quad & f_i(x) \leq 0, i = 1, \dots, m \\ & h_i(x) = 0, i = 1, \dots, p \end{aligned}$$

Dual

$$\begin{aligned} \max \quad & g(\lambda, v) = \inf_{x \in D} (f_0(x) + v^\top f(x) + v^\top h(x)) \\ \text{s. t.} \quad & \lambda \geq 0 \end{aligned}$$

$$g^* = g(\lambda^*, v^*)$$

$$f^* = f_0(x^*)$$

■ weak duality $f(x) \geq g(\lambda, v), \quad f^* \geq g^*$

■ strong duality $f^* = g^*$



- convex problems with further condition
- only convexity is not sufficient
- convexity is not necessary: strong duality holds for some non-convex problems

Complementary slackness



- when the strong duality holds

$$\begin{aligned} f_0(x^*) &= g(\lambda^*, v^*) = \inf_x (f_0(x) + \lambda^{*\top} f(x) + v^{*\top} h(x)) \\ &\leq f_0(x^*) + \lambda^{*\top} f(x^*) + v^{*\top} h(x^*) \\ &\leq f_0(x^*) \end{aligned}$$

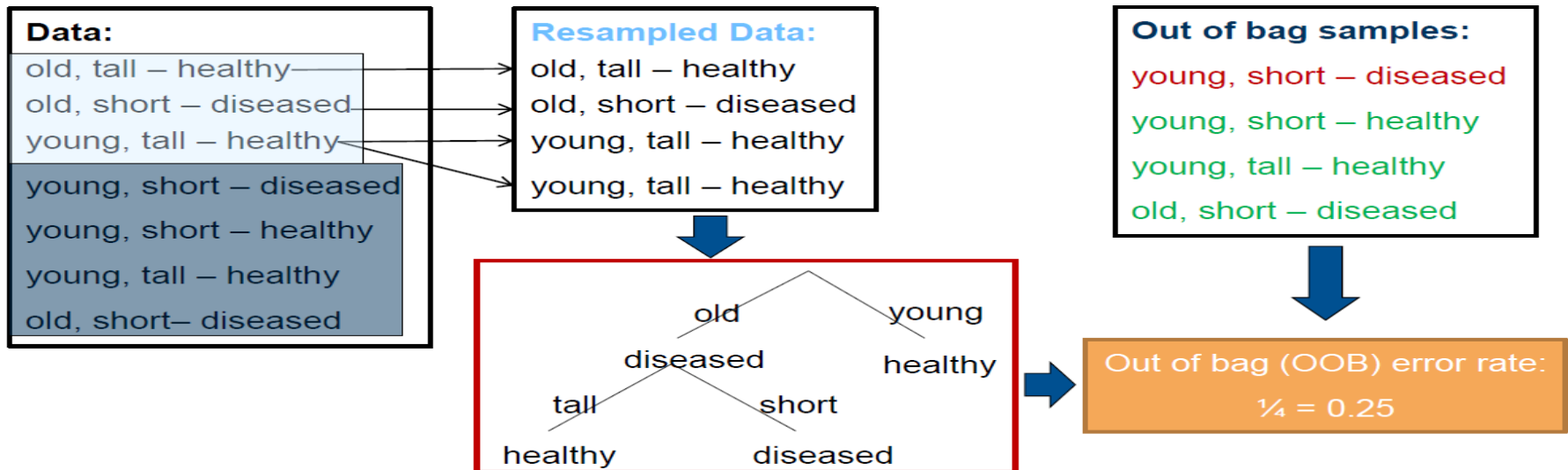
- “=” should be true for the last inequality

$$\begin{aligned} \sum_{i=1}^m \lambda_i^* f_i(x^*) = 0 &\quad \Rightarrow \quad \lambda_i^* f_i(x^*) = 0 \\ &\quad \swarrow \quad \searrow \\ \lambda_i^* > 0 \rightarrow f_i(x^*) = 0 &\quad f_i(x^*) < 0 \rightarrow \lambda_i^* = 0 \end{aligned}$$

Perturbation



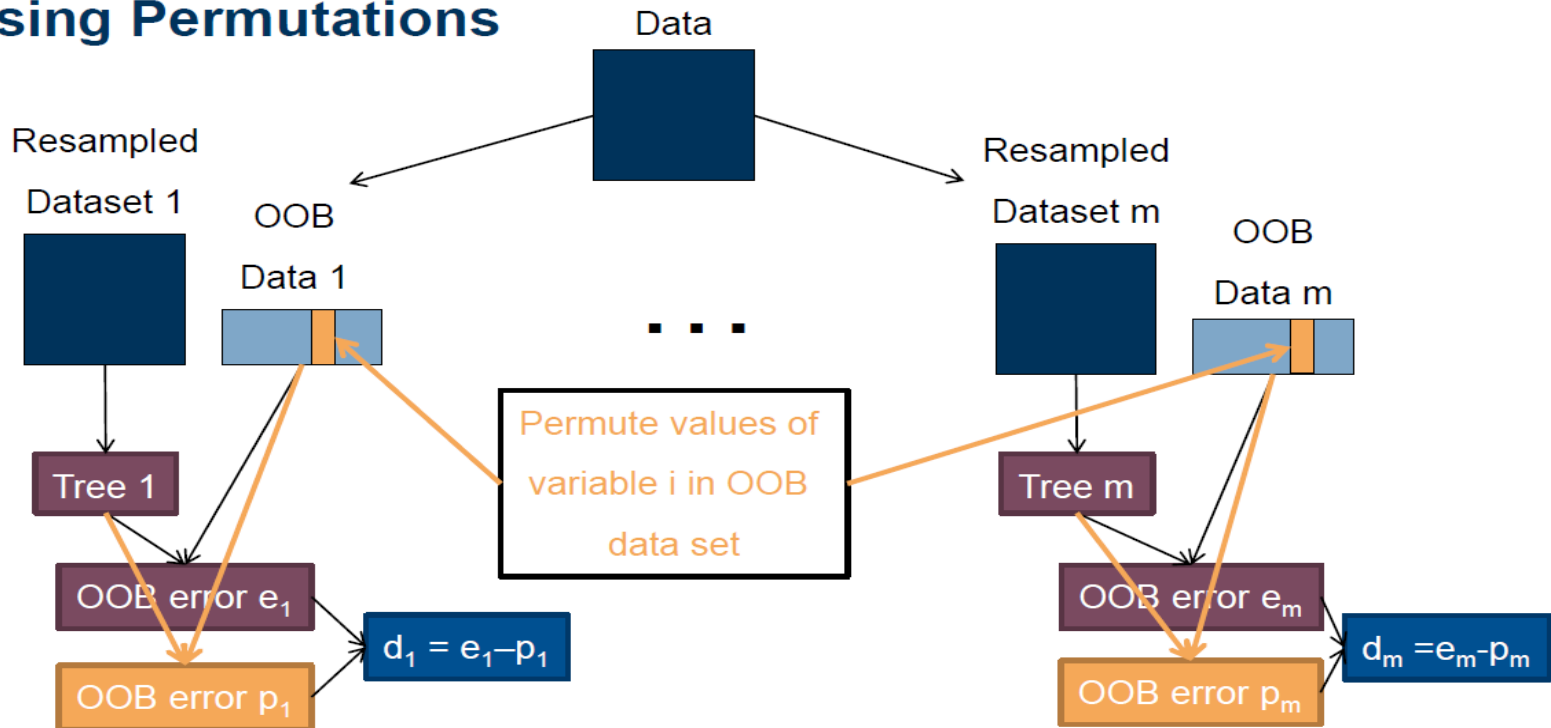
- perturbation: an intuitive way for sensitivity analysis
 - e.g., in random forest, one can permute one contribution and see the difference on the error, e.g., out-of-bag (OOB) error



Perturbation



- perturbation: an intuitive way for sensitivity analysis
using **Permutations**



$$\bar{d} = \frac{1}{m} \sum_{i=1}^m d_i$$

$$s_d^2 = \frac{1}{m-1} \sum_{i=1}^m (d_i - \bar{d})^2$$

$$v_i = \frac{\bar{d}}{s_d}$$

Perturbation



- unperturbed optimization problem and its dual

$$\begin{array}{ll}\min & f_0(x) \\ \text{s. t.} & f_i(x) \leq 0, i = 1, \dots, m \\ & h_i(x) = 0, i = 1, \dots, p\end{array}$$

$$\begin{array}{ll}\max & g(\lambda, v) \\ \text{s. t.} & \lambda \geq 0\end{array}$$

- perturbed optimization problem and its dual

$$\begin{array}{ll}\min & f_0(x) \\ \text{s. t.} & f_i(x) \leq u_i, i = 1, \dots, m \\ & h_i(x) = v_i, i = 1, \dots, p\end{array}$$

$$\begin{array}{ll}\max & g(\lambda, v) - u^\top \lambda - v^\top v \\ \text{s. t.} & \lambda \geq 0\end{array}$$

$f^*(u, v)$ how about the optimal value changes as a function of u and v

Global sensitivity



- apply weak duality to the perturbed problem

$$f^*(u, v) \geq \underbrace{g(\lambda^*, v^*)}_{\text{assume the strong duality holds}} - u^\top \lambda^* - v^\top v^* = f^*(0, 0) - u^\top \lambda^* - v^\top v^*$$

assume the strong duality holds

- perturbed optimization problem and its dual

$$\begin{array}{ll} \min & f_0(x) \\ \text{s. t.} & f_i(x) \leq u_i, i = 1, \dots, m \\ & h_i(x) = v_i, i = 1, \dots, p \end{array}$$

$$\begin{array}{ll} \max & g(\lambda, v) - u^\top \lambda - v^\top v \\ \text{s. t.} & \lambda \geq 0 \end{array}$$

$f^*(u, v)$ how about the optimal value changes as a function of u and v

Global sensitivity



- apply weak duality to the perturbed problem

$$f^*(u, v) \geq \underbrace{g(\lambda^*, v^*)}_{\text{assume the strong duality holds}} - u^\top \lambda^* - v^\top v^* = f^*(0, 0) - u^\top \lambda^* - v^\top v^*$$

assume the strong duality holds

- sensitivity interpretation
 - if λ_i^* is large, f^* increases (becomes worse) greatly when we tighten constraint ($u_i < 0$)
 - if λ_i^* is large and positive,

$$\begin{array}{ll} \min & f_0(x) \\ \text{s. t.} & f_i(x) \leq u_i, i = 1, \dots, m \\ & h_i(x) = v_i, i = 1, \dots, p \end{array}$$

Local sensitivity



- apply weak duality to the perturbed problem

$$f^*(u, v) \geq g(\lambda^*, v^*) - u^\top \lambda^* - v^\top v^* = f^*(0, 0) - u^\top \lambda^* - v^\top v^*$$

- if $f^*(u, v)$ is differentiable at the original, then

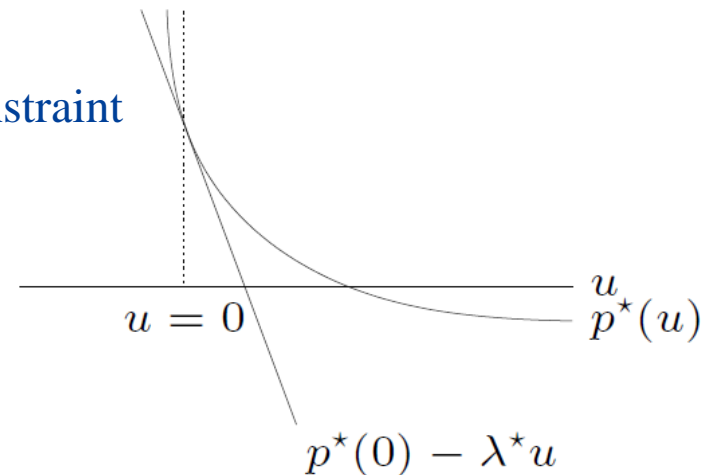
$$\lambda_i^* = - \frac{df^*(u, v)}{du_i} \Big|_{u=0, v=0} \quad v_i^* = - \frac{df^*(u, v)}{dv_i} \Big|_{u=0, v=0}$$

$f^*(u)$ for a problem with one inequality constraint



if the inequality constraint is linear

if the inequality becomes equation



Problem reformulation



- equivalent formulations may have very different dual
- reformulation can be useful when the dual is difficult or uninteresting
- for example

$$\min_x f_0(Ax + b) \quad \longrightarrow \quad g = \inf_x L(x) = \inf_x f_0(Ax + b)$$



$$\begin{aligned} \min_{x,z} f_0(z) &\quad \longrightarrow \quad g(v) = \inf_{x,z} f_0(z) - v^\top (Ax + b - z) \\ \text{s.t. } Ax + b - z &= 0 \end{aligned}$$
$$= \begin{cases} -f_0^*(v) + b^\top v & A^\top v = 0 \\ -\infty & \text{otherwise} \end{cases}$$

1

Langrange Duality

2

Optimality Condition

3

Support Vector Machine



Karush-Kuhn-Tucker conditions



suppose all these functions
are differentiable

$$\begin{array}{ll}\min & f_0(x) \\ \text{s. t.} & f_i(x) \leq 0, i = 1, \dots, m \\ & h_i(x) = 0, i = 1, \dots, p\end{array}$$

- x^* is optimal to the primal problem and u^*, v^* to the dual problem

primal feasible

$$f_i(x^*) \leq 0$$

$$h_i(x^*) = 0$$

dual feasible

$$\lambda_i^* \geq 0$$

complementary slackness

$$\lambda_i^* f_i(x^*) = 0$$



$$\nabla f_0(x^*) + \sum_i \lambda_i^* \nabla f_i(x^*) + \sum_i v_i^* \nabla h_i(x^*) = 0$$

gradient of Lagrangian

Optimality



- for convex problems

$L(x, \lambda, v) = f_0(x) + \lambda^\top f(x) + v^\top h(x)$ is convex

- $L(x, \lambda, v)$ achieves the minimum when x^* satisfying

$$\nabla f_0(x^*) + \sum_i \lambda_i^* \nabla f_i(x^*) + \sum_i v_i^* h(x^*) = 0$$



$$g(\lambda^*, v^*) = f_0(x^*) + \lambda^{*\top} f(x^*) + v^{*\top} h(x^*) = f_0(x^*)$$

duality gap is zero, and they are optimal

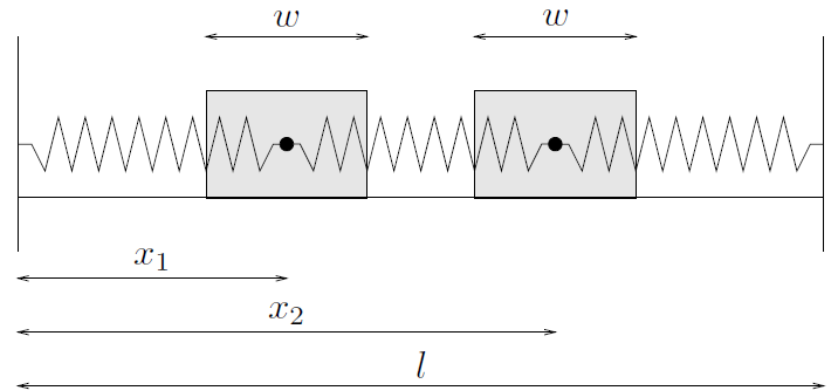
Optimality



- for convex problems
 - if Slater's condition is satisfied, KKT is sufficient and necessary
 - if not, KKT is necessary, but not sufficient
- for non-convex problems
 - KKT is necessary



Interpretation



- the energy

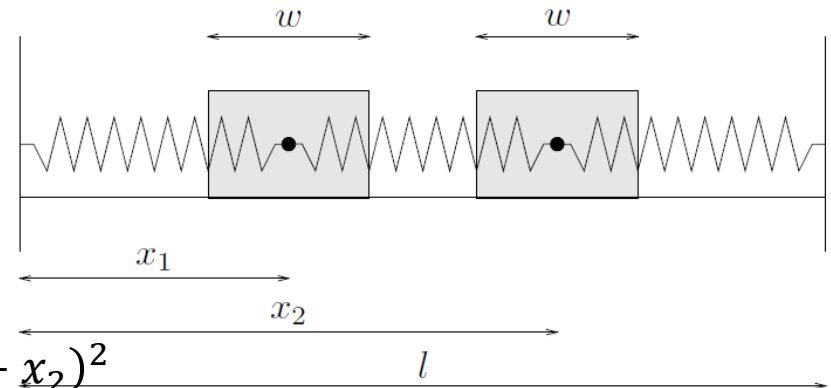
$$f_0(x_1, x_2) = \frac{1}{2}k_1x_1^2 + \frac{1}{2}k_2(x_2 - x_1)^2 + \frac{1}{2}k_3(l - x_2)^2$$

- physical constraints: $\frac{w}{2} - x_1 \leq 0, w + x_1 - x_2 \leq 0, \frac{w}{2} - l + x_2 \leq 0$

- the equilibrium could be achieved

$$\begin{aligned}
 \min \quad & \frac{1}{2}k_1x_1^2 + \frac{1}{2}k_2(x_2 - x_1)^2 + \frac{1}{2}k_3(l - x_2)^2 \\
 \text{s. t.} \quad & \frac{w}{2} - x_1 \leq 0 \\
 & w + x_1 - x_2 \leq 0 \\
 & \frac{w}{2} - l + x_2 \leq 0
 \end{aligned}$$

Interpretation



$$\min \quad \frac{1}{2} k_1 x_1^2 + \frac{1}{2} k_2 (x_2 - x_1)^2 + \frac{1}{2} k_3 (l - x_2)^2$$

$$\text{s. t.} \quad \frac{w}{2} - x_1 \leq 0$$

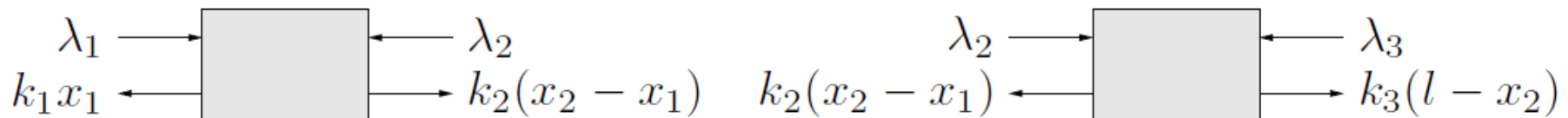
$$w + x_1 - x_2 \leq 0$$

$$\frac{w}{2} - l + x_2 \leq 0$$

$$\lambda_1, \lambda_2, \lambda_3 \geq 0$$

■ KKT $\lambda_1 \left(\frac{w}{2} - x_1 \right) = 0, \lambda_2 (w + x_1 - x_2) = 0, \lambda_3 \left(\frac{w}{2} - l + x_2 \right) = 0$

$$\begin{bmatrix} k_1 x_1 - k_2 (x_2 - x_1) \\ k_2 (x_2 - x_1) - k_3 (l - x_2) \end{bmatrix} + \lambda_1 \begin{bmatrix} -1 \\ 0 \end{bmatrix} + \lambda_2 \begin{bmatrix} 1 \\ -1 \end{bmatrix} + \lambda_3 \begin{bmatrix} 0 \\ 1 \end{bmatrix} = 0$$



Newton's method with equality constraints

$$\min_x f(x) \quad \text{s.t. } Ax = b$$

- if the x^k is feasible, we need to guarantee that $A\Delta x_{\text{nt}} = 0$
- optimality condition

$$Ax^* = b, \nabla f(x^*) + A^\top v^* = 0$$



$$A(x^k + \Delta x_{\text{nt}}) = b, \nabla f(x^k + \Delta x_{\text{nt}}) + A^\top v \approx \nabla f(x^k) + \nabla^2 f(x^k)\Delta x_{\text{nt}} + A^\top v = 0$$



$$Ax^k = b$$

$$A\Delta x_{\text{nt}} = 0, \quad \nabla^2 f(x^k)\Delta x_{\text{nt}} + A^\top v = -\nabla f(x^k)$$

Newton's method with equality constraints

$$\min_x f(x) \quad \text{s. t. } Ax = b$$

- the Newton's direction is obtained by

$$\begin{bmatrix} \nabla^2 f(x^k) & A^\top \\ A & 0 \end{bmatrix} \begin{bmatrix} \Delta x_{\text{nt}} \\ v \end{bmatrix} = \begin{bmatrix} -\nabla f(x^k) \\ 0 \end{bmatrix}$$

given starting point $x \in \text{dom } f$ with $Ax = b$, tolerance $\epsilon > 0$.

repeat

1. Compute the Newton step and decrement $\Delta x_{\text{nt}}, \lambda(x)$.
 2. *Stopping criterion.* **quit** if $\lambda^2/2 \leq \epsilon$.
 3. *Line search.* Choose step size t by backtracking line search.
 4. *Update.* $x := x + t\Delta x_{\text{nt}}$.
-

Newton's method with equality constraints

$$\min_x f(x) \quad \text{s.t.} \quad Ax = b$$

- if the x^k is infeasible, we need to first let the solution goes to feasible set

$$A(x^k + \Delta x_{\text{nt}}) = b, \nabla f(x^k + \Delta x_{\text{nt}}) + A^T v \approx \nabla f(x^k) + \nabla^2 f(x^k) \Delta x_{\text{nt}} + A^T v = 0$$

$$\Downarrow \quad Ax^k \neq b$$

$$A\Delta x_{\text{nt}} = -(Ax^k - b), \quad \nabla^2 f(x^k) \Delta x_{\text{nt}} + A^T v = -\nabla f(x^k)$$

$$\Downarrow$$

$$\begin{bmatrix} \nabla^2 f(x^k) & A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} \Delta x_{\text{nt}} \\ v \end{bmatrix} = - \begin{bmatrix} \nabla f(x^k) \\ Ax^k - b \end{bmatrix}$$

primal-dual interpretation

Inequality constrained problem



$$\begin{array}{ll}\min & f_0(x) \\ \text{s. t.} & f_i(x) \leq 0, i = 1, \dots, m \\ & Ax = b\end{array}$$

suppose all these functions are
twice continuously differentiable

- convex problem, A is full-ranked, f^* is finite and attained
- we assume the problem is strictly feasible
(and so strong duality holds and dual optimum is attained)
- using indicator function to reformulate

$$\begin{array}{ll}\min & f_0(x) + \sum_{i=1}^m I_-(f_i(x)) \\ \text{s. t.} & Ax = b\end{array}$$

$$I_-(u) = \begin{cases} 0, & u \leq 0 \\ \infty, & u > 0 \end{cases}$$

Logarithmic barrier



$$\begin{aligned}
 \min \quad & f_0(x) + \sum_{i=1}^m I_{-}(f_i(x)) \\
 \text{s. t.} \quad & Ax = b
 \end{aligned}$$

- the indicator function is not continuous
- use logarithmic function to approach the indicator

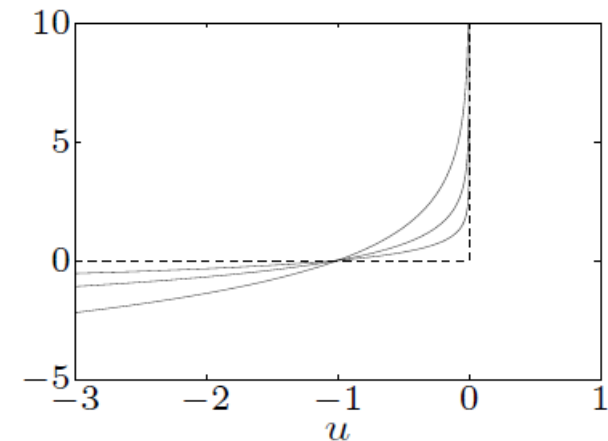
$$I_{-}(u) \approx -\frac{1}{t} \log(-u), t \rightarrow \infty$$

- logarithmic barrier function

$$\phi(x) = -\sum_{i=1}^m \log(-f_i(x))$$

$$\nabla \phi(x) = \sum_{i=1}^m \frac{1}{-f_i(x)} \nabla f_i(x)$$

$$\nabla^2 \phi(x) = \sum_{i=1}^m \frac{1}{f_i(x)^2} \nabla f_i(x) \nabla f_i(x)^{\top} + \sum_{i=1}^m \frac{1}{-f_i(x)} \nabla^2 f_i(x)$$



Logarithmic barrier



$$\begin{array}{ll} \min & f_0(x) + \sum_{i=1}^m I_-(f_i(x)) \\ \text{s. t.} & Ax = b \end{array}$$

- the indicator function is not continuous
- use logarithmic function to approach the indicator

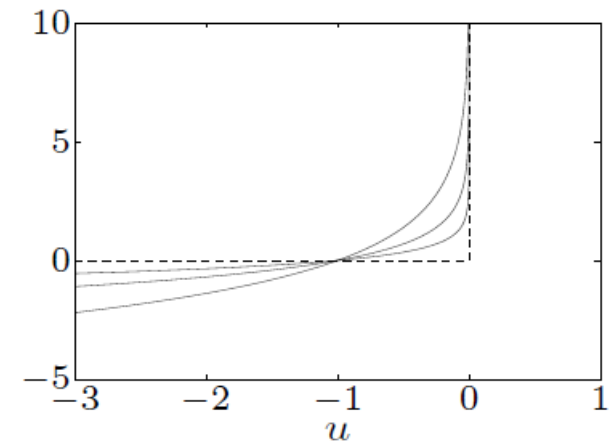
$$I_-(u) \approx -\frac{1}{t} \log(-u), t \rightarrow \infty$$

- logarithmic barrier function

$$\phi(x) = -\sum_{i=1}^m \log(-f_i(x))$$

- with increasing t , we can approach the original problem

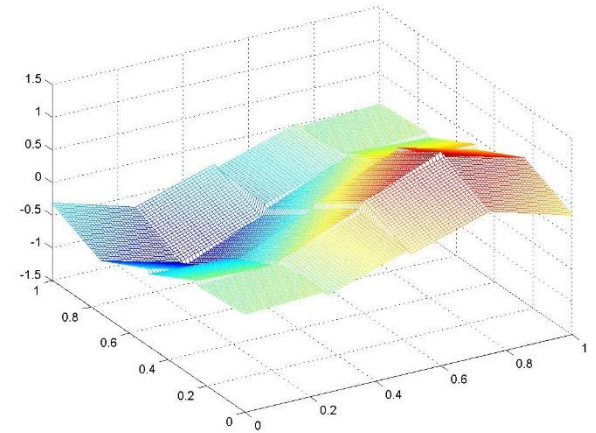
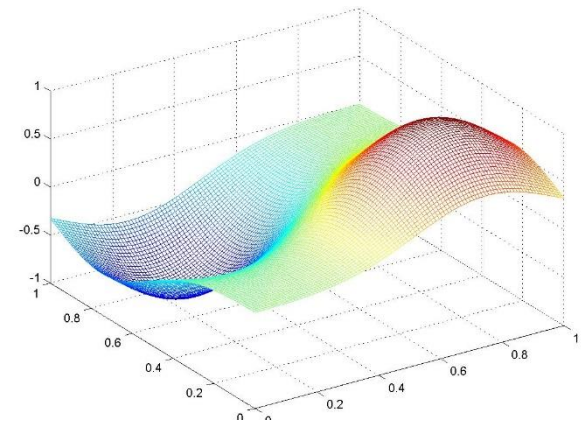
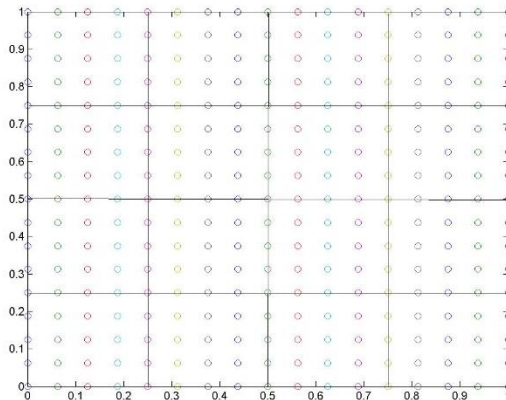
$$\min_x t f_0(x) + \phi(x), \text{ s. t. } Ax = b$$



Nonlinearity: subregion



- Takagi-Sugeno Model:
 - divide the domain into subregions
 - locally train a linear model



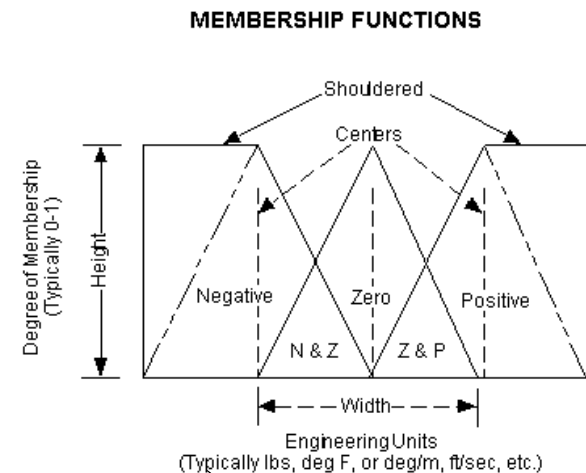
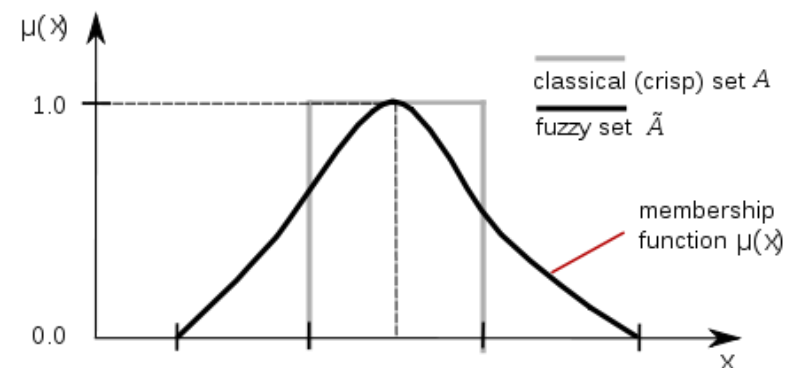
Fuzzy systems



- how to deal with the discontinuity
 - local linear functions $f_i(x), \forall x \in S_i$
 - simply sum them with an indicator function $I(A) = 1$ iff A is true

$$F(x) = \sum_{i=1} I(x \in S_i) f_i(x)$$

- we could replace the indicator function by a *membership* functions
- ANFIS (Adaptive neuro fuzzy inference system)



Central path



- central path

$$x^*(t) = \operatorname{argmin}_x t f_0(x) + \phi(x), \text{ s.t. } Ax = b$$

- for a given t , there exists a $w(t)^*$ such that

$$t \nabla f_0(x^*(t)) + \sum_{i=1}^m \frac{1}{-f_i(x^*(t))} \nabla f_i(x^*(t)) + A^\top w(t)^* = 0$$

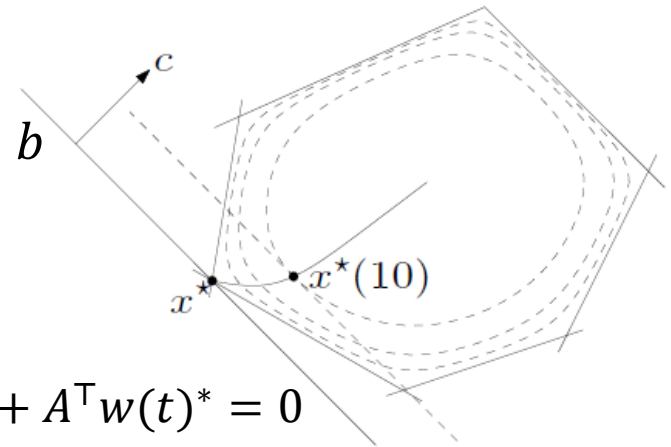
$$\lambda^*(t) = 1/(-t f_i(x^*(t)))$$

$$v^*(t) = w(t)^*/t$$

- the following Lagrangian is minimized at $x^*(t)$

$$L(x, \lambda_i^*(t), v^*(t)) = f_0(x) + \sum_{i=1}^m \lambda_i^*(t) f_i(x) + v^*(t)^\top (Ax - b)$$

$$f^* \geq g(\lambda_i^*(t), v^*(t)) = L(x^*(t), \lambda^*(t), v^*(t)) = f_0(x^*(t)) - m/t$$



KKT conditions



- consider $x^*(t), \lambda^*(t), v^*(t)$

$$x^*(t) = \underset{x}{\operatorname{argmin}} \quad t f_0(x) + \phi(x), \text{ s. t. } Ax = b$$

$$\lambda^*(t) = 1/(-t f_i(x^*(t)))$$

$$v^*(t) = w(t)^*/t$$

$$\begin{aligned} \min \quad & f_0(x) \\ \text{s. t.} \quad & f_i(x) \leq 0, i = 1, \dots, m \\ & Ax = b \end{aligned}$$

primal feasible

$$f_i(x^*(t)) \leq 0$$

$$h_i(x^*(t)) = 0$$

dual feasible

$$\lambda_i^*(t) \geq 0$$

complementary slackness

$$\lambda_i^*(t) f_i(x^*(t)) = -1/t$$

$$\nabla f_0(x^*(t)) + \sum_i \lambda_i^* \nabla f_i(x^*(t)) + \sum_i v_i^* \nabla h_i(x^*(t)) = 0 \quad \text{gradient of Lagrangian}$$

Barrier method



given strictly feasible x , $t := t^{(0)} > 0$, $\mu > 1$, tolerance $\epsilon > 0$.

repeat

1. *Centering step.* Compute $x^*(t)$ by minimizing $tf_0 + \phi$, subject to $Ax = b$.
 2. *Update.* $x := x^*(t)$.
 3. *Stopping criterion.* **quit** if $m/t < \epsilon$.
 4. *Increase t .* $t := \mu t$.
-

- two loops: outer *iteration*, and *centering*
- outer iteration
 - a larger μ leads to faster convergence
- centering
 - standard analysis for unconstrained problems
 - a larger μ leads to slower convergence

Feasibility and phase I



- interior-point method needs a strictly feasible solution
- if not, solve a feasibility problem

$$\begin{array}{ll}\text{find} & x \\ \text{s. t.} & f_i(x) \leq 0, i = 1, \dots, m \\ & Ax = b\end{array}$$



$$\begin{array}{ll}\min & s \\ \text{s. t.} & f_i(x) \leq s, i = 1, \dots, m \\ & Ax = b, s \geq 0\end{array}$$

$$\begin{array}{ll}\min & \sum s_i \\ \text{s. t.} & f_i(x) \leq s_i, i = 1, \dots, m \\ & Ax = b, s_i \geq 0\end{array}$$

1

Langrange Duality

2

Optimality Condition

3

Support Vector Machine



Dual problem example of SVM



- SVM in primal space

$$\begin{aligned} \min_{x,z} \quad & \frac{1}{2} \|x\|_2^2 + C \sum_i \rho_i \\ \text{s. t.} \quad & b_i(x^\top a_i + z) \geq 1 - \rho_i \\ & \rho_i \geq 0 \end{aligned}$$

- the corresponding Lagrangian

$$\begin{aligned} L(x, z, \rho; \lambda, \nu) = & \frac{1}{2} \|x\|_2^2 + C \sum_i \rho_i \\ & + \sum_i \lambda_i (1 - \rho_i - b_i(x^\top a_i + z)) - \nu^\top \rho \end{aligned}$$

Dual problem example of SVM



- to obtain the $g(\lambda, \nu) = \inf_{x, z, \rho} L(x, z, \rho; \lambda, \nu)$

$$L(x, z, \rho; \lambda, \nu) = \frac{1}{2} \|x\|_2^2 + C \sum_i \rho_i + \sum_i \lambda_i (1 - \rho_i - b_i(x^\top a_i + z)) - \nu^\top \rho$$

$$\frac{\partial L}{\partial x} = x - \sum_i b_i \lambda_i a_i = 0$$

$$\frac{\partial L}{\partial z} = \sum_i \lambda_i b_i = 0$$

$$\frac{\partial L}{\partial \rho_i} = C - \lambda_i - \nu_i = 0$$

$$\lambda_i \geq 0, \nu_i \geq 0$$



$$\begin{aligned} \min_{\lambda} \quad & \sum_i \sum_j \lambda_i b_i a_i^\top a_j b_j \lambda_j - \sum_i \lambda_i \\ \text{s. t.} \quad & \sum_i \lambda_i b_i = 0 \\ & 0 \leq \lambda_i \leq C \end{aligned}$$

Primal-dual relation



- SVM in primal space

$$\begin{aligned} \min_{x,z} \quad & \frac{1}{2} \|x\|_2^2 + C \sum_i \rho_i \\ \text{s. t.} \quad & b_i(x^\top a_i + z) \geq 1 - \rho_i \\ & \rho_i \geq 0 \end{aligned}$$

- SVM in dual space

$$\begin{aligned} \min_{\lambda} \quad & \sum_i \sum_j \lambda_i b_i a_i^\top a_j b_j \lambda_j - \sum_i \lambda_i \\ \text{s. t.} \quad & \sum_i \lambda_i b_i = 0 \\ & 0 \leq \lambda_i \leq C \end{aligned}$$

- strong duality and

$$f(x) = x^\top a + z = \sum_i \lambda_i b_i a_i^\top a + z$$

Support vector



- SVM in dual space $f(x) = \sum_i \lambda_i b_i a_i^T a + z$

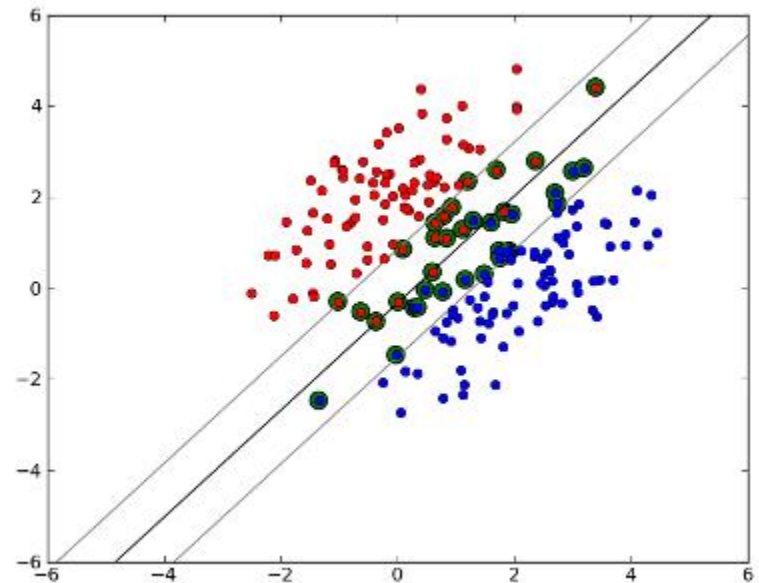
$$\begin{aligned} \min_{\lambda} \quad & \sum_i \sum_j \lambda_i b_i a_i^T a_j b_j \lambda_j - \sum_i \lambda_i \\ \text{s. t.} \quad & \sum_i \lambda_i b_i = 0 \\ & 0 \leq \lambda_i \leq C \end{aligned}$$

- there are only a part of samples

$$\lambda_i \neq 0$$

“Support Vector”

which one is support vector?



Support vector

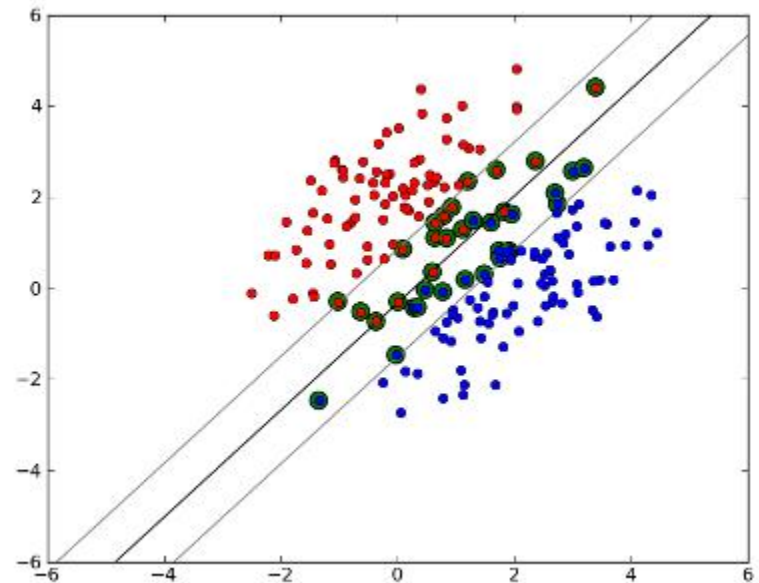


- complementary slackness

$$\lambda_i = C \begin{cases} \lambda_i > 0 \longrightarrow 1 - \rho_i - b_i(x^\top a_i + z) = 0 \\ \nu_i = 0 \longrightarrow \rho_i \geq 0 \end{cases}$$



$$b_i(x^\top a_i + z) \leq 1$$



Support vector



- complementary slackness

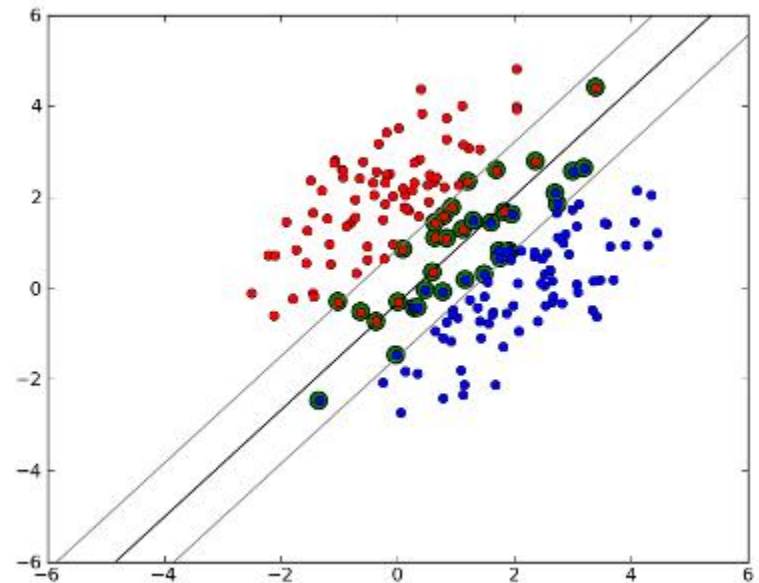
$$0 < \lambda_i < C \begin{cases} \lambda_i > 0 \longrightarrow 1 - \rho_i - b_i(x^\top a_i + z) = 0 \\ \nu_i > 0 \longrightarrow \rho_i = 0 \end{cases}$$



$$b_i(x^\top a_i + z) = 1$$



determine z



Kernel trick



- to introduce non-linearity, usually a non-linear mapping is needed:

$$\phi(a): \mathbf{R}^n \rightarrow \mathbf{R}^d$$

- SVM in primal space

$$\begin{aligned} \min_{x,z} \quad & \frac{1}{2} \|x\|_2^2 + C \sum_i \rho_i \\ \text{s. t.} \quad & b_i(x^\top \phi(a_i) + z) \geq 1 - \rho_i, \\ & \rho_i \geq 0 \end{aligned}$$

- SVM in dual space

$$\begin{aligned} \min_{\lambda} \quad & \sum_i \sum_j \lambda_i b_i \phi(a_i)^\top \phi(a_j) b_j \lambda - \sum_i \lambda_i \\ \text{s. t.} \quad & \sum_i \lambda_i b_i = 0 \\ & 0 \leq \lambda_i \leq C \end{aligned}$$

Kernel trick



- the discriminant function is :

$$f(x) = \sum_i \lambda_i b \boxed{\phi(a_i)^\top \phi(a)} - z$$

- kernel trick: we do not need to know the formulation of $\phi(x)$, instead, we only need to know the inner product

- kernel functions: $K(u, v): \mathbf{R}^n \times \mathbf{R}^n \rightarrow R$

describe the relationship of the two samples

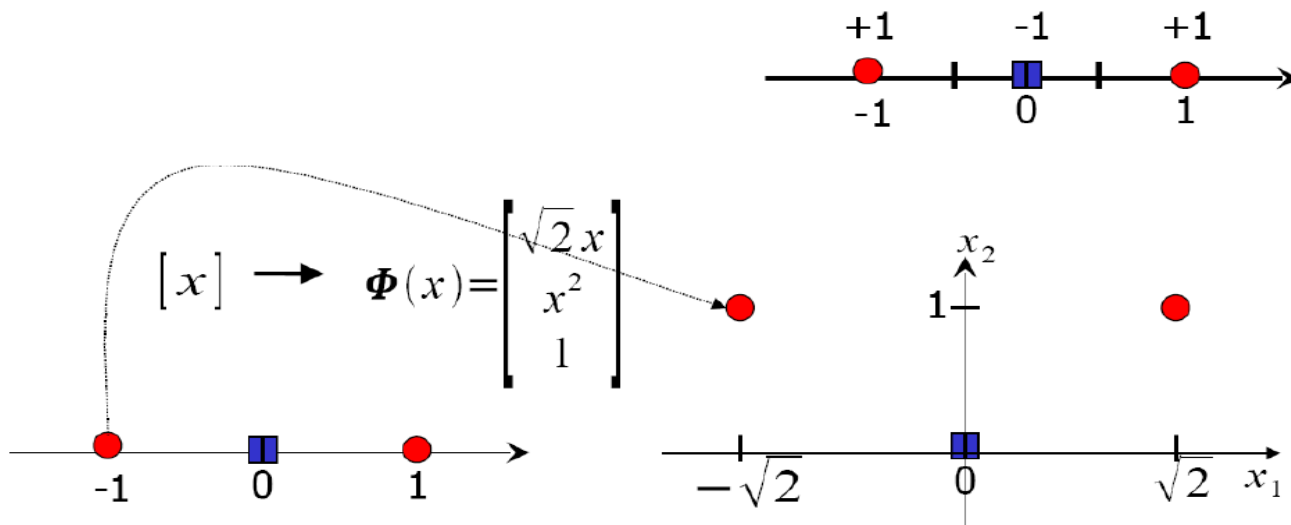
Polynomial kernel



- polynomial kernel:

$$K(u, v) = (u^T v + c)^d$$

- when $c = 0, d = 1$, it reduces to *linear kernel*
- for a one-dimensional, two-order polynomial kernel



Polynomial kernel

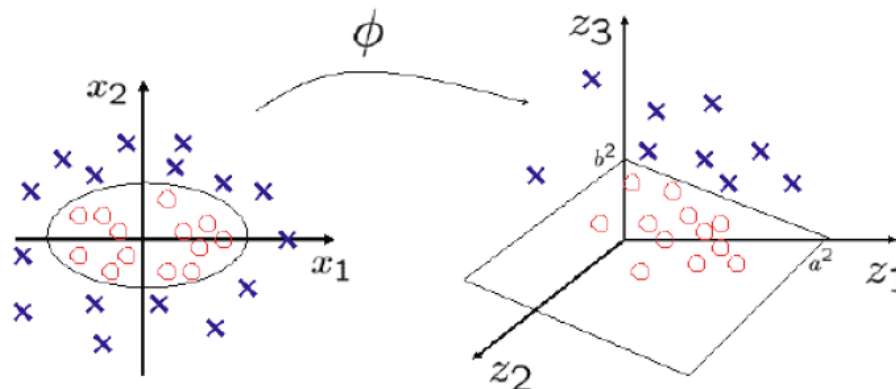


- polynomial kernel:

$$K(u, v) = (u^T v + c)^d$$

- when $c = 0, d = 1$, it reduces to *linear kernel*
- for a two-dimensional, two-order polynomial kernel

$$\phi(\mathbf{u}) = [u_1^2, \sqrt{2}u_1u_2, u_2^2];$$




RBF kernel



- Radial basis function (Gaussian) kernel

$$\mathcal{K}(\mathbf{u}, \mathbf{v}) = \exp \left(-\frac{\|\mathbf{u} - \mathbf{v}\|_2^2}{2\sigma^2} \right)$$

- even for one-dimensional case

$$\phi(u) = \exp \left(-\frac{u^2}{2} \right) \left[1, \sqrt{2}u, \sqrt{\frac{1}{2!}}u^2, \sqrt{\frac{1}{3!}}u^3, \dots \right]^T$$


an indefinite dimensional mapping

Mercer kernel



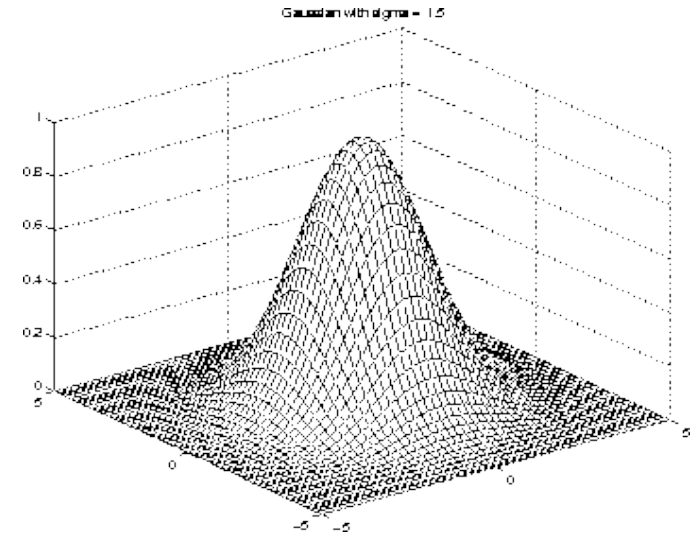
- radial basis function (Gaussian) kernel

$$\mathcal{K}(\mathbf{u}, \mathbf{v}) = \exp \left(-\frac{\|\mathbf{u} - \mathbf{v}\|_2^2}{2\sigma^2} \right)$$

- it is a similarity/dissimilarity measure
- many similarity matrix can be used if:
- Mercer's Theorem:

the matrix introduced by K is positive-semidefinite

$$K_{ij} = K(a_i, a_j) = \phi(a_i)^T \phi(a_j)$$



Solving algorithms



- solve SVM from primal or dual?

$$\begin{aligned} \min_{x,z} \quad & \frac{1}{2} \|x\|_2^2 + C \sum_i \rho_i \\ \text{s.t.} \quad & b_i(x^\top a_i + z) \geq 1 - \rho_i \\ & \rho_i \geq 0 \end{aligned}$$

$$\begin{aligned} \min_{\lambda} \quad & \sum_i \sum_j \lambda_i b_i a_i^\top a_j b_j \lambda - \sum_i \lambda_i \\ \text{s.t.} \quad & \sum_i \lambda_i b_i = 0 \\ & 0 \leq \lambda_i \leq C \end{aligned}$$

independent of m

- primal space: $n + 1$ unknown variables ($2m + 1$ constraints)
- dual space: m unknown variables and $m + 1$ constraints
- big data problem usually solved from primal

independent of n

SMO



- consider the dual problem

$$\begin{aligned} \min_{\lambda} \quad & \sum_i \sum_j \lambda_i b_i K_{ij} b_j \lambda_j - \sum_i \lambda_i \\ \text{s. t.} \quad & \sum_i \lambda_i b_i = 0 \\ & 0 \leq \lambda_i \leq C \end{aligned}$$

- choose only a small number of variables
- the smallest number is 2
- Sequential Minimization Optimization (SMO)
 - this idea is not restricted to SVM

SMO

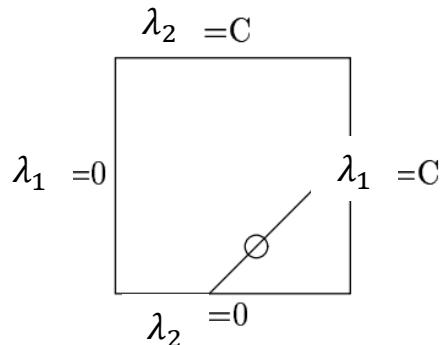


- consider the dual problem

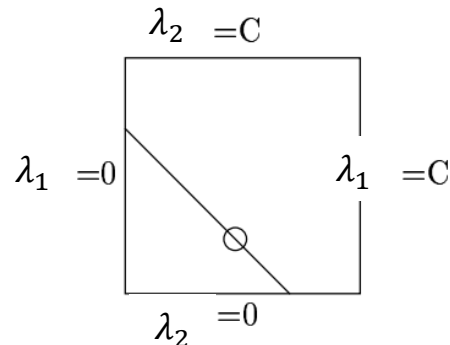
$$\begin{aligned} \min_{\lambda} \quad & \sum_i \sum_j \lambda_i b_i K_{ij} b_j \lambda_j - \sum_i \lambda_i \\ \text{s. t.} \quad & \sum_i \lambda_i b_i = 0 \\ & 0 \leq \lambda_i \leq C \end{aligned} \quad \rightarrow \quad \begin{aligned} \min_{\lambda_1, \lambda_2} \quad & K_{11}\lambda_1^2 + 2b_1b_2K_{12}\lambda_1\lambda_2 + K_{22}\lambda_2^2 - \lambda_1 - \lambda_2 \\ \text{s. t.} \quad & b_1\lambda_1 + b_2\lambda_2 = 0 \\ & 0 \leq \lambda_1, \lambda_2 \leq C \end{aligned}$$



- according to different b_1, b_2



$$a_1 \neq a_2 \rightarrow \lambda_1 - \lambda_2 = 0$$



$$a_1 = a_2 \rightarrow \lambda_1 + \lambda_2 = 0$$

SMO



- λ_1, λ_2 can be optimally updated, actually with analytic expressions

$$\min_{\lambda_1, \lambda_2} \begin{aligned} & K_{11}\lambda_1^2 + 2b_1b_2K_{12}\lambda_1\lambda_2 + K_{22}\lambda_2^2 \\ & -\lambda_1 - \lambda_2 \end{aligned}$$

$$\begin{aligned} \text{s. t. } & b_1\lambda_1 + b_2\lambda_2 = 0 \\ & 0 \leq \lambda_1, \lambda_2 \leq C \end{aligned}$$

- the remaining question is how to select the variables to be update
 - purely random: the improvement may be small, but no time required
 - the largest improvement pair: a 2-D loop
 - the largest two variables violating optimality condition: 1-D loop

Another example: one-bit CS



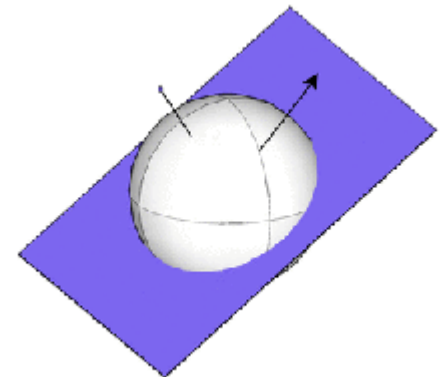
- compressive sensing

$$\min_{x,z} \mu \|x\|_1 + \frac{1}{m} \sum (b_i - a_i^\top x)^2$$

- in real situation, the observations (actually all variables) are quantized
- the extreme case, we only have one-bit information

$$b_i = \text{sign}(a_i^\top x + \varepsilon)$$

- is that possible to also recover the signal?
- but norm information is needed



Another example: one-bit CS



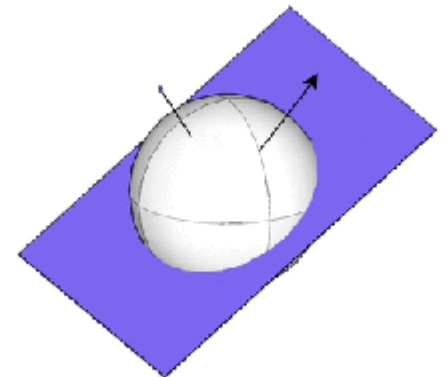
- one-bit compressive sensing

$$\begin{aligned} \min_{x,z} \quad & \mu \|x\|_1 + \frac{1}{m} \sum \max\{0, 1 - b_i(a_i^\top x)\} \\ \text{s. t.} \quad & \|x\|_2 = 1 \end{aligned}$$

- in real situation, the observations (actually all variables) are quantized
- the extreme case, we only have one-bit information

$$b_i = \text{sign}(a_i^\top x + \varepsilon)$$

- is that possible to also recover the signal?
- but norm information is needed



Another example: one-bit CS



- one-bit compressive sensing

$$\begin{array}{ll} \min_{x,z} & \mu \|x\|_1 + \frac{1}{m} \sum \max\{0, 1 - b_i(a_i^\top x)\} \\ \text{s. t.} & \|x\|_2 = 1 \end{array}$$

- relaxation

$$\begin{array}{ll} \min_{x,z} & \mu \|x\|_1 + \frac{1}{m} \sum \max\{0, 1 - b_i(a_i^\top x)\} \\ \text{s. t.} & \|x\|_2 \leq 1 \end{array}$$

- reformulation (could have different dual)

$$\begin{array}{ll} \min_{x,y,z} & \mu \|y\|_1 + \frac{1}{m} \sum \max\{0, 1 + z_i\} + I_{\|x\|_2 \leq 1}(x) \\ \text{s. t.} & x = y, z_i = -b_i(a_i^\top x) \end{array}$$

Another example: one-bit CS



$$\begin{aligned} \min_{x,y,z} \quad & \mu \|y\|_1 + \frac{1}{m} \sum \max\{0, 1 + z_i\} + I_{\|x\|_2 \leq 1}(x) \\ \text{s. t.} \quad & x = y, z_i = -b_i(a_i^\top x) \end{aligned}$$

- Lagrangian

$$L(x, y, z; \lambda, v) = \mu \|y\|_1 + \frac{1}{m} \sum \max\{0, 1 + z_i\} + I_{\|x\|_2 \leq 1}(x) + \lambda^\top (x - y) + v^\top (-b \cdot Ax - z)$$

- minimization over primal variables

$$\min_x L(x, y, z; \lambda, v) = \min_x I_{\|x\|_2 \leq 1}(x) + \lambda^\top x - v^\top (b \cdot Ax) = -\|\sum v_i b_i a_i - \lambda\|_2$$

$$\min_y L(x, y, z; \lambda, v) = \min_y \mu \|y\|_1 - \lambda^\top y = \begin{cases} 0, & \|\lambda\|_\infty \leq \mu \\ -\infty, & \text{otherwise} \end{cases}$$

$$\min_{z_i} L(x, y, z; \lambda, v) = \min_{z_i} \frac{1}{m} \max\{0, 1 + z_i\} - v_i z_i = \begin{cases} v_i, & |v_i| \leq 1/m \\ -\infty, & \text{otherwise} \end{cases}$$

Another example: one-bit CS



$$\begin{aligned} \max_{\lambda, v} \quad & \sum v_i - \|\sum v_i b_i a_i - \lambda\|_2 \\ \text{s. t.} \quad & \|\lambda\|_\infty \leq \mu, \quad \|v\|_\infty \leq 1/m \end{aligned}$$

separable
dual coordinate ascent

- Langrangian

$$L(x, y, z; \lambda, v) = \mu \|y\|_1 + \frac{1}{m} \sum \max\{0, 1 + z_i\} + I_{\|x\|_2 \leq 1}(x) + \lambda^\top (x - y) + v^\top (-b \cdot Ax - z)$$

- minimization over primal variables

$$\min_x L(x, y, z; \lambda, v) = \min_x I_{\|x\|_2 \leq 1}(x) + \lambda^\top x - v^\top (b \cdot Ax) = -\|\sum v_i b_i a_i - \lambda\|_2$$

$$\min_y L(x, y, z; \lambda, v) = \min_y \mu \|y\|_1 - \lambda^\top y = \begin{cases} 0, & \|\lambda\|_\infty \leq \mu \\ -\infty, & \text{otherwise} \end{cases}$$

$$\min_{z_i} L(x, y, z; \lambda, v) = \min_{z_i} \frac{1}{m} \max\{0, 1 + z_i\} - v_i z_i = \begin{cases} v_i, & |v_i| \leq 1/m \\ -\infty, & \text{otherwise} \end{cases}$$

Principle component analysis



- for finding the principle axes of a dataset

$$\max_{x^T x = 1} x^T C x \quad C = \frac{1}{n-1} A^T A$$

it should be zero-mean,
otherwise, the covariance
matrix will be ?

- this optimization problem can be solved as the following:
 1. compute the mean
 2. compute the covariance
 3. find the principle axes
 4. project data onto the eigenvectors



Principle component analysis



- for finding the principle axes of a dataset

$$\max_{x^T x = 1} x^T C x \quad C = \frac{1}{n-1} A^T A$$

- the Lagrangian is

$$L(x, \lambda) = x^T C x - \lambda(x^T x - 1)$$

- from the KKT condition

$$\frac{\partial L(x, \lambda)}{\partial x} = Cx - \lambda x = 0$$



$$Cx = \lambda x$$

Other principle components



- for the second principle axis:
 - maximize the variance
 - uncorrelated (orthogonal) with $x_1^\top A$

$$\text{cov}(x_1^\top A, x_2^\top A) = x_1^\top A x_2 = x_2^\top A x_1 = \lambda x_1^\top x_2 = 0$$

- then we are going to solve

$$\max_{x^\top x=1, x_1^\top x_2=0} x^\top A x$$

- similarly, consider its Langrangian

$$L(x, \lambda) = x^\top A x - \lambda(x^\top x - 1) - \beta x_1^\top w$$

Other principle components



- consider the derivatives of $L(x, \lambda) = x^T A x - \lambda(x^T x - 1) - \beta x_1^T x$

$$\frac{\partial L(x, \lambda)}{\partial x} = Cx - \lambda x - \beta x_1 = 0$$

- multiply by x_1^T on the left, we have

$$\boxed{x_1^T Cx - \lambda x_1^T x} - \beta \boxed{x_1^T x_1} = 0 \rightarrow \beta = 0$$

both equal zero,
because of the
un-correlation

nonzero

$$Cx = \lambda x$$

Nonlinear PCA

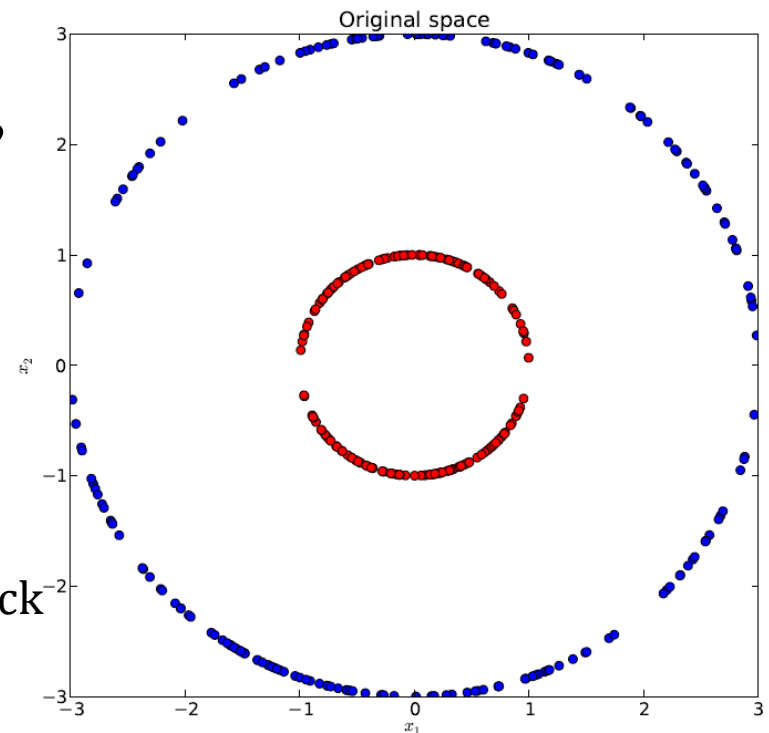


- PCA is to find linear subspace

- can we extend PCA to nonlinear problems?
 - the previous PCA is in primal
 - can we go to dual space and use *kernel trick*?

$$C = \frac{1}{n-1} \boxed{A^T A}$$

it is possible to use kernel trick



PCA in terms of dot products



- the eigenvectors lie in the span of a_1, a_2, \dots, a_m


- Proof.**

$$Cx = \frac{1}{m} \sum_{j=1}^m a_j a_j^T x = \lambda x$$

Therefore,

$$\begin{aligned} x &= \frac{1}{\lambda x} \sum_{j=1}^m a_j a_j^T x \\ &= \frac{1}{\lambda x} \sum_{j=1}^m (a_j \cdot x) a_j \end{aligned}$$

$(aa^T)x = (a \cdot x)a$





Show that $(xx^T)v = (x \cdot v)x$

$$\begin{aligned}(xx^T)v &= \begin{pmatrix} x_1x_1 & x_1x_2 & \dots & x_1x_M \\ x_2x_1 & x_2x_2 & \dots & x_2x_M \\ \vdots & \vdots & \ddots & \vdots \\ x_Mx_1 & x_Mx_2 & \dots & x_Mx_M \end{pmatrix} \begin{pmatrix} v_1 \\ v_2 \\ \vdots \\ v_M \end{pmatrix} \\ &= \begin{pmatrix} x_1x_1v_1 + x_1x_2v_2 + \dots + x_1x_Mv_M \\ x_2x_1v_1 + x_2x_2v_2 + \dots + x_2x_Mv_M \\ \vdots \\ x_Mx_1v_1 + x_Mx_2v_2 + \dots + x_Mx_Mv_M \end{pmatrix}\end{aligned}$$



$$\begin{aligned} &= \begin{pmatrix} (x_1 v_1 + x_2 v_2 + \dots + x_M v_M) x_1 \\ (x_1 v_1 + x_2 v_2 + \dots + x_M v_M) x_2 \\ \vdots \\ (x_1 v_1 + x_2 v_2 + \dots + x_M v_M) x_M \end{pmatrix} \\ &= \begin{pmatrix} x_1 v_1 + x_2 v_2 + \dots + x_M v_M \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_M \end{pmatrix} \\ &= (x \cdot v) x \end{aligned}$$

□

Nonlinear feature mapping



- we now can apply a nonlinear feature mapping $\phi(a)$
- then matrix $\Phi = \phi(A)$ (and assume it is centered)
- its principle component can be calculated as linear PCA:

$$Cx = \frac{1}{m} \sum_{j=1}^m \phi(a_j) \phi(a_j)^\top x = \lambda x$$

$$x = \sum_{i=1}^m \alpha_i \phi(a_i)$$

as showed previously, the solutions lie in **the span of $\phi(a_i)$** :

$$\frac{1}{m} \sum_{j=1}^m \phi(a_j) \phi(a_j)^\top \sum_{i=1}^m \alpha_i \phi(a_i) = \lambda \sum_{i=1}^m \alpha_i \phi(a_i)$$

Nonlinear feature mapping



- Kernel trick:

$$\sum_{j=1}^m \phi(a_j) \phi(a_j)^T \sum_{i=1}^m \alpha_i \phi(a_i) = m\lambda \sum_{i=1}^m \alpha_i \phi(a_i)$$

kernel trick

- Again, we do not need to know the feature mapping:

eigenvector

$$\phi(a)^T x = \phi(a)^T \sum_{i=1}^m \alpha_i \phi(a_i) = \sum_{i=1}^m \alpha_i K(a, a_i)$$

Kernel PCA



- calculate the kernel matrix K
- centralize the kernel matrix

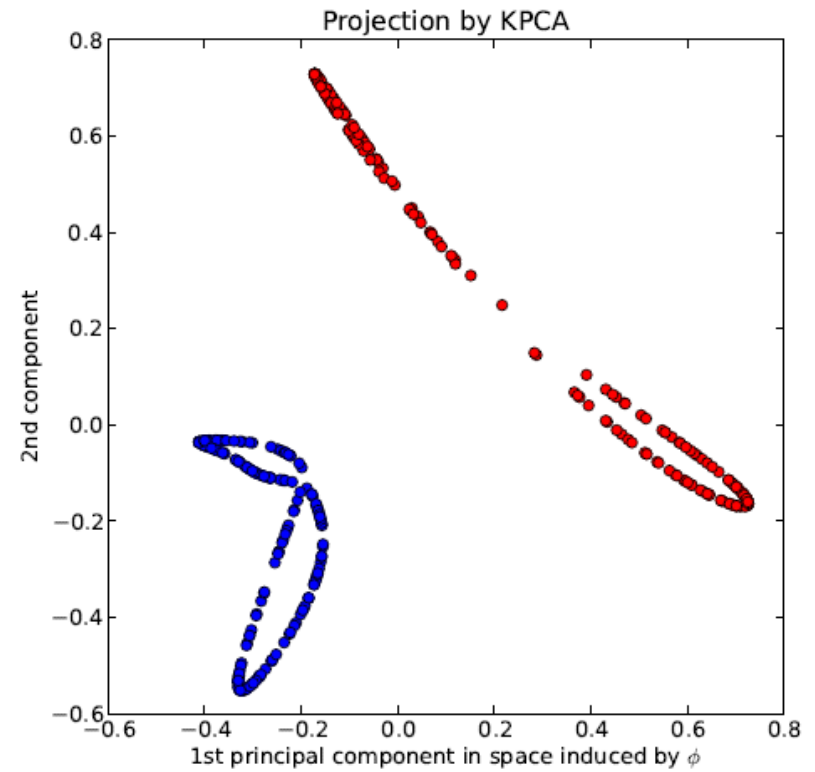
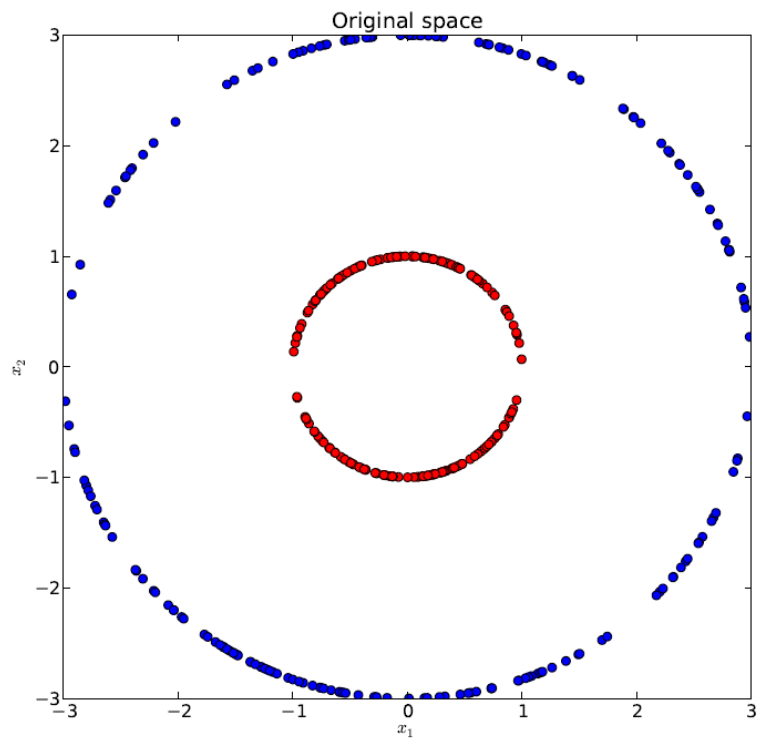
$$\hat{K} = K - \frac{1}{n} \mathbf{1} \mathbf{1}^T K - \frac{1}{n} K \mathbf{1} \mathbf{1}^T + \frac{\mathbf{1}^T K \mathbf{1}}{n^2} \mathbf{1} \mathbf{1}^T$$

- eigen-value decomposition: $[U, V] = \text{eig}(\hat{K})$
- find dual variables: $\alpha_i = \lambda_j^{-\frac{1}{2}} v_j$
- projection onto subspace:

$$\sum \alpha_{ji} K(a_i, a)$$



Kernel PCA



Kernel PCA and PCA



- If we choose linear kernel: $C = A^T A$, $A = AA^T$

- PCA:

$$\boxed{C}w = \lambda w$$

$n \times n$

- KPCA:

$$\boxed{K}\beta = \mu\beta$$

$m \times m$

- the projected data

$$Ax = \sum \alpha_{ji} K(a_i, a)$$

Recall the PSD condition



- before we always requires the kernel matrix is PSD
- are the previous algorithm applicable?
 - SMO for SVM

$$\begin{aligned} \min_{\alpha} \quad & \sum_i \sum_j \alpha_i b_i K_{ij} b_j \alpha_j - \sum_i \alpha_i \\ \text{s. t.} \quad & \sum_i \alpha_i b_i = 0 \\ & 0 \leq \alpha_i \leq C \end{aligned}$$

- inverse problem for LS-SVM
- eigenvalue for kPCA

$$K\alpha = \mu\alpha$$

Recall the PSD condition



- before we always requires the kernel matrix is PSD
- are the previous algorithm applicable?
 - SMO for SVM

$$\begin{aligned} \min_{\alpha} \quad & \sum_i \sum_j \alpha_i b_i K_{ij} b_j \alpha_j - \sum_i \alpha_i \\ \text{s. t.} \quad & \sum_i \alpha_i b_i = 0 \\ & 0 \leq \alpha_i \leq C \end{aligned}$$

yes, but local optimality

- inverse problem for LS-SVM
- eigenvalue for kPCA

yes, solvable

yes, solvable

$$K\alpha = \mu\alpha$$

Recall the PSD condition



- before we always requires the kernel matrix is PSD
- does the primal-dual relationship exist?
 - SVM

$$\begin{aligned} \min_{x,z} \quad & \frac{1}{2} \|x\|_2^2 + C \sum_i \rho_i \\ \text{s. t.} \quad & b_i (x^\top \phi(a_i) + z) \geq 1 - \rho_i, \\ & \rho_i \geq 0 \end{aligned}$$

$$\begin{aligned} \min_{\alpha} \quad & \sum_i \sum_j \alpha_i b_i K_{ij} b_j \alpha_j - \sum_i \alpha_i \\ \text{s. t.} \quad & \sum_i \alpha_i b_i = 0 \\ & 0 \leq \alpha_i \leq C \end{aligned}$$

- if not, what is the relationship between them?

1

Langrage Duality

2

Optimality Condition

3

Support Vector Machine



Homework



- 5.29
- Consider the following problem

$$\begin{aligned} \min_{x,z} \quad & \frac{1}{2} \|x\|_2^2 - \nu \xi + C \sum_i \rho_i \\ \text{s. t.} \quad & b_i(x^\top \phi(a_i) + z) \geq \xi - \rho_i, \forall i, \\ & \xi \geq 0, \rho_i \geq 0, \forall i \end{aligned}$$

you are asked to prove that ν is an upper bound on the fraction of margin errors, i.e., the number of samples falling in the margin is less than νm , i.e.,

$$\#\{i: y_i f(x_i) < \rho\} \leq \nu m$$

THANKS

