

# Announcements

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- ❖ Mid-term exam
  - ❖ Closed-book, no electronic device
  - ❖ 2 A4 cheatsheets (i.e., 4 pages in total), handwritten (your own writing), stapled, photocopy not allowed
- ❖ Nov. 4 class moved to Nov. 7 10-11:40am (D205)
- ❖ Mid-term course evaluation
- ❖ OH today 1pm-2pm



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# Ve492: Introduction to Artificial Intelligence

## Mid-term Review

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Slides adapted from <http://ai.berkeley.edu>, AIMA, UM, CMU

# Quiz: Search

- ❖ Consider a graph search problem where for every action, the cost is at least  $\epsilon$ , with  $\epsilon > 0$ . Assume the used heuristic is consistent.
  - ❖ Depth-first graph search is guaranteed to return an optimal solution.
  - ❖ Breadth-first graph search is guaranteed to return an optimal solution.
  - ❖ Uniform-cost graph search is guaranteed to return an optimal solution.
  - ❖ Greedy graph search is guaranteed to return an optimal solution.
  - ❖ A\* graph search is guaranteed to return an optimal solution.
  - ❖ A\* graph search is guaranteed to expand no more nodes than depth-first graph search.
  - ❖ A\* graph search is guaranteed to expand no more nodes than uniform-cost graph search.

# Quiz: A\*

- ❖ Assume we are running A\* graph search with a consistent heuristic  $h$ . Assume the optimal cost path to reach a goal has a cost  $c^*$ .
- ❖ All nodes  $n$  reachable from the start state satisfying  $g(n) < c^*$  will be expanded during the search
- ❖ All nodes  $n$  reachable from the start state satisfying  $f(n) = g(n) + h(n) < c^*$  will be expanded during the search
- ❖ All nodes  $n$  reachable from the start state satisfying  $h(n) < c^*$  will be expanded during the search

# Quiz: $A^*$ heuristics

- ❖ Let  $h_1(s)$  be an admissible heuristic. Let  $h_2(s) = 2h_1(s)$ .
  - ❖ The solution found by  $A^*$  tree search with  $h_2$  is guaranteed to be an optimal solution.
  - ❖ The solution found by  $A^*$  tree search with  $h_2$  is guaranteed to have a cost at most twice as much as the optimal path.
  - ❖ The solution found by  $A^*$  graph search with  $h_2$  is guaranteed to be an optimal solution.

# Quiz: A\* Heuristics

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- ❖ Let  $H_1$  and  $H_2$  both be admissible heuristics.
  - ❖  $\max(H_1, H_2)$  is necessarily admissible
  - ❖  $\min(H_1, H_2)$  is necessarily admissible
  - ❖  $(H_1 + H_2)/2$  is necessarily admissible
  - ❖  $\max(H_1, H_2)$  is necessarily consistent

# Quiz: A\* Heuristics

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- ❖ Let  $H_1$  be an admissible heuristic, and let  $H_2$  be an inadmissible heuristic.
  - ❖  $\max(H_1, H_2)$  is necessarily admissible
  - ❖  $\min(H_1, H_2)$  is necessarily admissible
  - ❖  $(H_1 + H_2)/2$  is necessarily admissible
  - ❖  $\max(H_1, H_2)$  is necessarily consistent

# Quiz: Search under Uncertainty

- ❖ You are given a game tree for which you are the maximizer, and in the nodes in which you don't get to make a decision an action is chosen uniformly at random amongst the available options. Your objective is to maximize the probability you win \$10 or more (rather than the usual objective to maximize your expected value).
  - ❖ Running expectimax will result in finding the optimal strategy to maximize the probability of winning \$10 or more.
  - ❖ Running minimax, where chance nodes are considered minimizers, will result in finding the optimal strategy to maximize the probability of winning \$10 or more.
  - ❖ Running expectimax in a modified game tree where every pay-off of \$10 or more is given a value of 1, and every pay-off lower than \$10 is given a value of 0 will result in finding the optimal strategy to maximize the probability of winning \$10 or more.
  - ❖ Running minimax in a modified game tree where every pay-off of \$10 or more is given a value of 1, and every pay-off lower than \$10 is given a value of 0 will result in finding the optimal strategy to maximize the probability of winning \$10 or more.

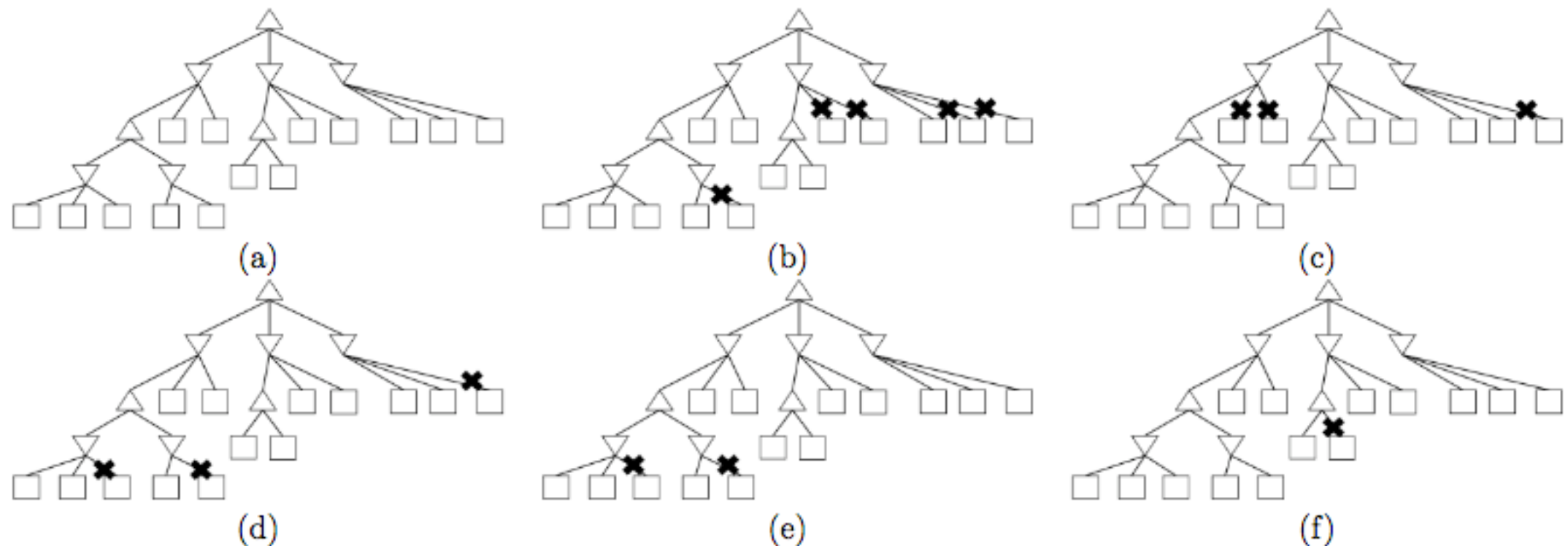


# Quiz: Adversarial Search

- ❖ In the context of adversarial search,  $\alpha$ - $\beta$  pruning
  - ❖ can reduce computation time by pruning portions of the game tree
  - ❖ is generally faster than minimax, but loses the guarantee of optimality
  - ❖ always returns the same value as minimax for the root of the tree
  - ❖ always returns the same value as minimax for all nodes on the leftmost edge of the tree, assuming successor game states are expanded from left to right
  - ❖ always returns the same value as minimax for all nodes of the tree

# Quiz: Adversarial Search

- ❖ Assume we run  $\alpha$ - $\beta$  pruning expanding successors from left to right on a game with tree as shown in Figure (a).



- ❖ For some choice of pay-off values, no pruning will be achieved (shown in Figure (a)).
- ❖ For some choice of pay-off values, the pruning shown in Figure (b) will be achieved.
- ❖ For some choice of pay-off values, the pruning shown in Figure (c) will be achieved.
- ❖ For some choice of pay-off values, the pruning shown in Figure (d) will be achieved.
- ❖ For some choice of pay-off values, the pruning shown in Figure (e) will be achieved.
- ❖ For some choice of pay-off values, the pruning shown in Figure (f) will be achieved.

# Quiz: MDP

- ❖ For Markov Decisions Processes (MDPs), we have that:
  - ❖ A small discount (close to 0) encourages shortsighted, greedy behavior.
  - ❖ A large, negative living reward (  $\ll 0$  ) encourages shortsighted, greedy behavior.
  - ❖ A negative living reward can always be expressed using a discount  $< 1$ .
  - ❖ A discount  $< 1$  can always be expressed as a negative living reward.

# Quiz: MDP

- ❖ Assume given an MDP  $\mathcal{M} = (S, A, T, R, \gamma)$ . Define a new MDP  $\mathcal{M}' = (S, A, T, \alpha R + \beta, \gamma)$  where  $\alpha > 0$  and  $\beta \in \mathbb{R}$ .
  - ❖ An optimal policy in  $\mathcal{M}$  is optimal in  $\mathcal{M}'$ .
  - ❖ An optimal policy in  $\mathcal{M}'$  is optimal in  $\mathcal{M}$ .
  - ❖ If  $\pi \succeq \pi'$  in  $\mathcal{M}$ , then  $\pi \succeq \pi'$  in  $\mathcal{M}'$
  - ❖ If  $\pi \succeq \pi'$  in  $\mathcal{M}'$ , then  $\pi \succeq \pi'$  in  $\mathcal{M}$

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# Quiz: MDP

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- ❖ Value iteration can converge only if the discount factor ( $\gamma$ ) satisfies  $0 < \gamma < 1$ .
- ❖ Policies found by value iteration may be superior to policies found by policy iteration.
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# Quiz: Reinforcement Learning

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- ❖ Assume that the agent observes the true reward with some Gaussian noise  $\mathcal{N}(0,1)$ , Q-learning would still converge
- ❖ Q-learning can learn the optimal Q-function  $Q^*$  without ever executing the optimal policy.
- ❖ If an MDP has a transition model  $T$  that assigns non-zero probability for all triples  $T(s, a, s')$  then Q-learning will fail.
- ❖ In Q-learning, we decide to explore every  $k$  steps, i.e., if  $t = 0 [k]$  we choose a random action with a uniform distribution, otherwise we choose the greedy action. This version would still converge.

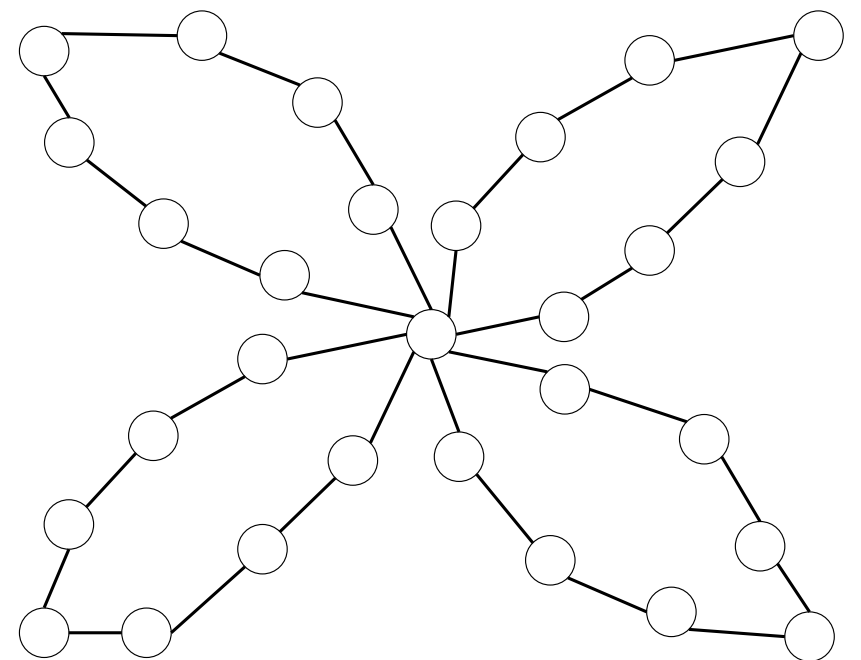
# Quiz: CSP

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- ❖ The most-constrained variable heuristic provides a way to select the next variable to assign in a backtracking search for solving a CSP.
- ❖ By using the most-constrained variable heuristic and the least-constraining value heuristic we can solve every CSP in time linear in the number of variables.
- ❖ CSP problems are always solved faster with arc consistency than with forward checking.
- ❖ When enforcing arc consistency, the values that remain in each domain depend on the order in which the arcs are considered.

# Quiz: CSP

- ❖ Assume given a CSP whose constraint graph is given below and that all the variables have the same domain.
- ❖ What is the complexity of solving it with a direct application of backtracking search?
- ❖ Which efficient strategy could you apply to solve it? What would be the complexity?





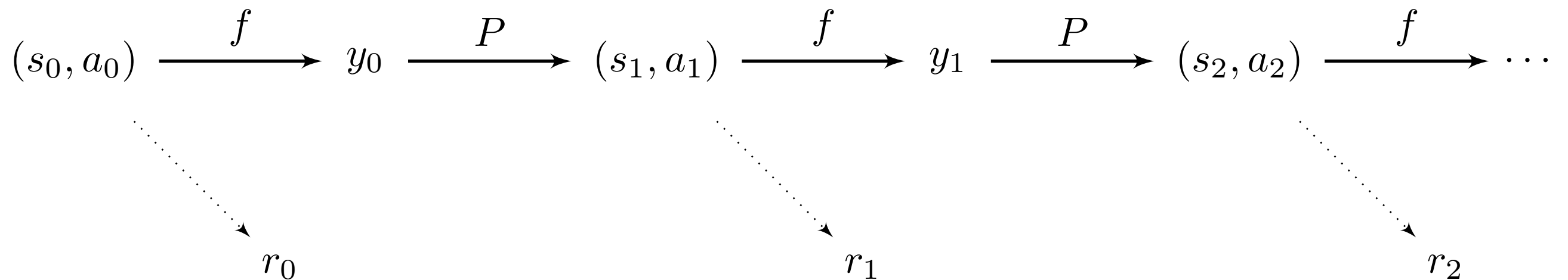
# CSP Problem: Job Scheduling

## ❖ When can I move in?

Task	Description	Duration	Predecessor
a	Erecting walls	7	none
b	Carpentry for roof	3	a
c	Roof	1	b
d	Installations	8	a
e	Facade painting	2	c & d
f	Windows	1	c & d
g	Garden	1	c & d
h	Ceilings	3	a
i	Painting	2	f & h
j	Moving in	1	i

# MDP Problem: Post-Decision State

- ❖ Consider an infinite-horizon, discounted MDP  $(S, A, T, R, \gamma)$  where  $T(s,a,s') = P(s' | f(s,a))$ ,  $R(s,a,s') = R(s,a)$  and  $f$  is some deterministic function mapping  $S \times A \rightarrow Y$ , where  $Y$  is a set of states called post-decision states.



- ❖  $V^\pi(s_0) = \mathbb{E}[R(s_0, \pi(s_0)) + \gamma R(s_1, \pi(s_1)) + \gamma^2 R(s_2, \pi(s_2)) + \dots]$
- ❖  $W^\pi(y_0) = \mathbb{E}[R(s_1, \pi(s_1)) + \gamma R(s_2, \pi(s_2)) + \gamma^2 R(s_3, \pi(s_3)) + \dots]$

# MDP Problem: Post-Decision State

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- ❖ Write  $W^*$  in terms of  $V^*$
- ❖ Write  $V^*$  in terms of  $W^*$
- ❖ Provide the equivalent of the Bellman equation for  $W^*$
- ❖ Fill in the blanks
  - ❖ Initialize policy  $\pi_1$  arbitrarily
  - ❖ For all  $i = 1, 2, \dots$ 
    - ❖ Compute  $W^{\pi_i}(y) \quad \forall y$
    - ❖ Compute new policy  $\pi_{i+1}$  from  $\pi_{i+1}(s) = \arg \max \_\_\_\_\_\_ \quad \forall s$
    - ❖ If  $\_\_\_\_\_\_ \quad \forall s$ , return  $\pi_i$

# Game Theory Problem

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- ❖ Two players choose simultaneously a coin of 10 cents, 50 cents or 1 dollar, which they show to each other.
- ❖ If they chose the same coin, player I wins. Otherwise, player II wins.
- ❖ Write this game in normal form. Is there any pure NE?
- ❖ Express a system of inequalities to find a mixed NE.

# Search: The wolf, goat, cabbage problem

- ❖ You are on the bank of a river with a boat, a cabbage, a goat, and a wolf. Your task is to get everything to the other side. Rules:
  - ❖ Only you can handle the boat
  - ❖ When you're in the boat, there is only space for one more item
  - ❖ You can't leave the goat alone with the wolf, nor with the cabbage (or something will be eaten)

